

## DS 筆記

### 1-1. 遞迴的基本概念

① 把問題縮的越來越小

② 問題變小還是同一個問題，所以可以用同一個 function

⇒ 精簡化程式，讓別人 好理解，debug

### 1-2. 反向印出字串

\* `string.substr(index, length)`

Key: print = 問題簡化

Ex: 字串的字數漸漸減一 ⇒ 問題簡化

最後有空字串 ⇒ 終止條件 (base case)

\* 遞迴出來的結果相同，但程式者需思考何者更適合使用者

### 1-5 GCD

```
gcd1(int x, int y) {
```

```
    if (y == 0)
```

```
        return x;
```

```
    else if (y > x)
```

```
        return gcd1(x, y % x);
```

```
    else
```

```
        return gcd1(y, x % y);
```

```
}
```

```
gcd2(int x, int y) {
```

```
    if (!(x % y)) return y;
```

```
    else return gcd2(y, x % y);
```

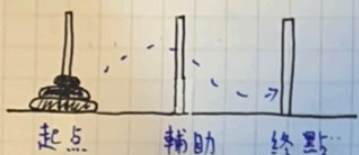
```
}
```

### 1-7 找第 K 大 (小)

① 用類 binary search → 只處理其中一邊 *more efficient*

② Linear Recursion = 線性遞迴 (只擇一遞迴)

### 1-8 河内塔介绍



Key point:  $f(n) \rightarrow f(n-1)$

```
solveTowers (count, source, destination, spare)
```

```
if (count is 1)
```

Move a disk directly from source to destination  $\Rightarrow$  把最后一层解决

```
else {
```

```
solveTowers (count - 1, source, spare, destination);  $\Rightarrow$  把终点当辅助
```

```
solveTowers (1, source, destination, spare);
```

```
solveTowers (count - 1, spare, destination, source);  $\Rightarrow$  把起点当辅助
```

```
}
```

### 1-10 Binary Recursion

```
drawTicks (length)
```

```
if (length > 0)
```

```
drawTics (length - 1)
```

draw tick of the length

```
drawTics (length - 1)
```

```
int sum (int a, int b)
```

```
{ ... a + sum (a + 1, b);
```

```
} // 8 call
```

vs

```
int sum (int a, int n)
```

```
{ ... sum (a, n/2) +
```

```
sum (a +  $\frac{n}{2}$ , n -  $\frac{n}{2}$ );
```

```
} // 15 call
```

### 1-12 Multiplying Rabbit (Fibonacci)

可用第 0 项

Better:

```
if k=1 return (k, 0)
```

```
else
```

```
(i, j) = linearFibonacci (k-1)
```

```
return (i+j, i)
```



1-14 指數

迴圈

```
double power1 (double x, int n)
```

```
double result = 1;
```

```
while (n > 0)
```

```
result = x * x;
```

```
n--;
```

```
return result;
```

```
double power2 (double x, int n)
```

```
if (n == 1)
```

```
return x;
```

```
else
```

```
return x * power2(x, n-1);
```

1-17 n 選 k 個組合數

Base cases:  $c(k, k) = 1$  or  $c(n, 0)$

If  $k > n \Rightarrow c(n, k) = 0$

有幾種選法 = 有多少 Base cases

For example:  $c(4, 2)$

Base Cases =  $c(2, 0)$ ,  $c(1, 0)$ ,  $c(1, 1)$ ,  $c(2, 2)$ ,  $c(2, 2)$ ,  $c(2, 2)$

Non-Base Cases =  $c(n, k) - 1$

Tail Recursion = 可以轉為迴圈的 Recursion

$\Rightarrow$  有時用 loop 相較於 Recursion 更加有效率

## 2-1 物件導向概念

Encapsulation (封裝) :

hides inner details

有一群物件(相同、相似的)

用 Encapsulation → 更有效率, 更不受外界影響

Inheritance (繼承) :

reused

把之前寫過得東西(程式)

能用的直接 copy 過來

Polymorphism (多型) :

在執行時, 根據 Input 的不同

做出符合資料類別的結果

## 2-2 物件導向程式設計介紹

\* Operation contracts :

Exceptions :

① Assume they never happen

② Ignore invalid situations

③ Return a value that signals a problem

④ Throw an exception

① Purpose    ② Assumptions    ③ Input    ④ Output



### 2-3 資料抽象化原理

\* ADT = Abstract Data Type

Modularity (模組化):

① 將程式切成一小塊一小塊

② Isolates errors

③ Eliminates redundancies (冗長)

→ Easier write, read, modify (修改)

內聚 = 一個 function 別做太多事 (越少越好)

耦合 = function 間傳遞的參數越少越好

Data abstraction:

不需要去管如何達成目的

只需要把要做什麼講清楚就好

補充:

```
reverse (in aList = List, out source = boolean)
```

```
for (i = 1 to aList.getLength() - 1) {  
    aList.retrieve (1, dataItem, success); → 檢索  
    aList.remove (1, success);  
    aList.insert (aList.getLength() - i + 2, dataItem, success);  
} // for
```

### 2-8 C++ 實作

① Data members should be private.

→ Constructors can be public

② Public → 提供給外界使用, private → 外界看不到、用不到

③ 子類別可以繼承父類別的所有 methods (private 除外)

④ 外部看到的 position 要減一, 才是 implement 時的 index → 這樣設計才對  
→ 對外部來說, pos 不會為 0

## 2-16 C++ 命名空間, 例外處理

① 可以把一堆 class 放在同一個 namespace 並使用 → 省時省力

② 容易出錯的地方:

try: 去試看看 ⇒ throw ...

catch: 把錯誤訊息保留下來

```
try { ... throw(type); ...
```

```
    catch (type1) {
```

```
        statement(s);
```

```
        ...
```

```
    } // catch(1)
```

```
    catch (type 2) {
```

```
        statement(s);
```

```
        ...
```

```
    } // catch(2)
```

```
    ...
```

```
    ...
```

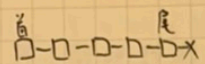
```
    } // try()
```

⇒ 跟 switch 用法很像



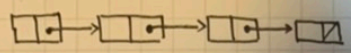
# Chapter 3

## • Link list :



Array  $\leftrightarrow$  Link list

## Preliminaries



## # 指標 pointer

=> 不用移動資料 (array 不行)

## • Pointer :

int \*p declaration of integer pointer  
(initially **undefined**, NOT NULL!)

p = &x (x 的門牌)

p = new int (申請一棟房子)

delete p (歸還房子)

p = NULL (門牌清空)

## • Pointer - Base implementation (實用)

- Public
  - isEmpty
  - getLength
  - insert
  - remove
  - retrieve
- Private
  - find
- Local variables
  - cur
  - pre

## • 動態陣列

int arraysize = 50;

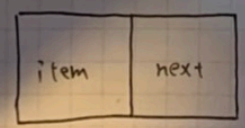
double \*Anarray = new double[arraysize];

delete [] Anarray (歸還)

• A destructor is require for dynamically allocated memory

```

List::~~List()
{
  while (!isEmpty)
    remove();
}
  
```



```

struct Node {
  int item;
  Node *next;
}
  
```

## • Array v.s Pointer

- Size: Linked list grow and shrink as necessary.
- Storage: Array requires less memory.
- retrieval: Array  $\rightarrow$  常數時間 (快), Linked List  $\rightarrow$  線性時間 (慢)
- insert / delete: if 資料多, Array 較慢

## • Restoring Linked list by using File:

- ofstream outFile

Ex: outFile << cur->item << endl;  $\rightarrow$  寫檔

`outFile.close();`

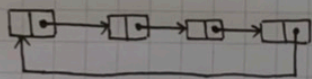
- ifstream inFile

Ex: inFile >> next item;  $\rightarrow$  讀檔

## • Processing Linked list Recursively $\rightarrow$ 利用 head

### • Variation (變化)

- circular linked list



- Dummy (Fake) head Node

$\Rightarrow$  Eliminate the special node!  
insert / delete

- Double Linked list (双向指標)

```
{  
    (初始化 = 指向自己)  
    Node * prev;  
    Node * next;  
}
```

◦ 處理指標順序不宜更改

## • Application (maintaining Inventory)

- Local DVD store
- operation on the inventory



## Chapter 4

### • 定义语言: Defining Language

$x | y$  (x or y)

$xy$  or  $x \cdot y$  (紧接着)

$\langle \text{number} \rangle = \langle \text{digit} \rangle \langle \text{number} \rangle | \langle \text{digit} \rangle$

$\langle \text{digit} \rangle = 0 | 1 | 2 | 3 | 4 | \dots | 9$

$\langle \text{addition} \rangle = \langle \text{digit} \rangle + \langle \text{addition} \rangle | \langle \text{digit} \rangle$

$\geq 2$  位数值以上  $\langle \text{addition} \rangle = \langle \text{number} \rangle + \langle \text{addition} \rangle | \langle \text{number} \rangle$

### • The Basic of grammar $\rightarrow$ recognition algorithm (判别演算法)

is Id # 递归

终止条件: 单一字元 / 递归呼叫: 扣除最后一字元

is Id suffix

### • Palindromes 回文

ex: Anna, 38+83=121

Base

$\langle \text{pal} \rangle = \text{empty string} \langle \text{ch} \rangle | a \langle \text{pal} \rangle a | b \langle \text{pal} \rangle b | \dots | z \langle \text{pal} \rangle z$

$\langle \text{ch} \rangle = a | b | c | \dots | z$

### • $A^n B^n$

$\langle \text{Legal-word} \rangle = \text{empty string} | A \langle \text{Legal-word} \rangle B$

### • Practice

$\langle S \rangle = \langle L \rangle | \langle D \rangle \langle S \rangle \langle S \rangle$

$\langle L \rangle = \langle A \rangle \langle B \rangle$

$\langle D \rangle = 1 | 2$

⑥ 1. write all 3个字元字串? 1AA, 1BB, 1BA, 2AA, 2BB, 2BA  
2. 包含多于3个字元字串? 12AAB (无数组)

• 描述至少一个字母所组成的字串, 第一个字母大写, others 小写

## • Algebraic Expression (代數)

◦ Infix 中序運算式  $ex: A + B$

◦ Prefix 前序 "  $ex: + \cdot AB$

◦ Postfix 後序 "  $ex: AB +$

## ◦ Advantage of prefix, postfix

- No precedence rules 優先權
- No association 結合律
- No parentheses 括弧

• 中序轉前序  $(a+b) \times c$

$\Rightarrow * + abc$

## • Prefix

### ◦ Grammar

$\langle \text{prefix} \rangle = \langle \text{identifier} \rangle \mid \langle \text{operator} \rangle \langle \text{prefix} \rangle \langle \text{prefix} \rangle$

### ◦ recursive recognition

- Base Case: One lowercase is a prefix exp
- Recursive:  $\langle \text{operator} \rangle \langle \text{prefix} \rangle \langle \text{prefix} \rangle$

\* 重要特性 = 一個前序式後面再接上非空字串不一定是前序式

## • Back tracking (回溯)

### ◦ 八皇后問題



- 64 格

- 8 前進方向

→ 避開死棋的位置

### • Recursive

- Base Case: 1 欄填滿

- Recursive Step: 填其他欄位