

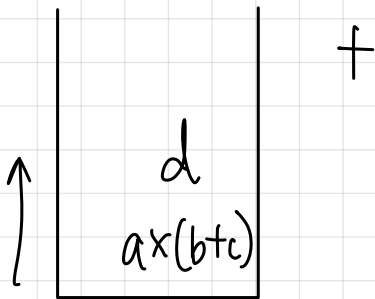
# stack

ADT stack:

{  
  isEmpty()  
  push()  
  pop()    *pop 不传参数*  
  getTop()  
  pop()  
}

postfix Evaluator ✓

abc + x d +



step in infix expression:

1. Append an operand to the end of an initially empty string postfixExpr

2. push ( onto a stack

3. push an operator onto stack if stack is empty; otherwise pop operators and append them to postfixExpr as long as they have a precedence higher than or equal to  $\geq$  that of the operator in the infix expression.

4. pop operators from stack and 運算

心得: 在用 stack 做 prefix, infix, postfix 時理解了很久, 一直看不太懂何時要 pop 何時要 push, 最後多看了幾次影片後就有比懂, 其中在表達式中有括號處理起來最容易。

## Queue

佇列 = 排隊

### ADT Queue

is Empty

enqueue 新增

dequeue 移除

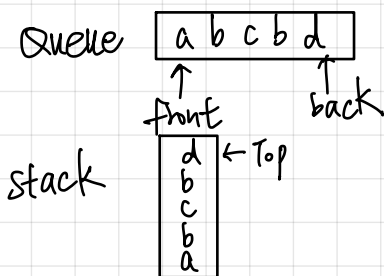
getFront 擷取

dequeue(out queueFront: QueueItem) 擷取後移除

### Recognize Palindromes 迴文判斷

Using stack & queue

queue 的前端 vs. stack 的 Top



```
void Queue::enqueue(const QueueItemType & newItem) {
```

```
    QueueNode* newPtr = new QueueNode;
```

```
    newPtr->item = newItem;
```

```
    if (isEmpty()) newPtr->next = newPtr;
```

```
    else {
```

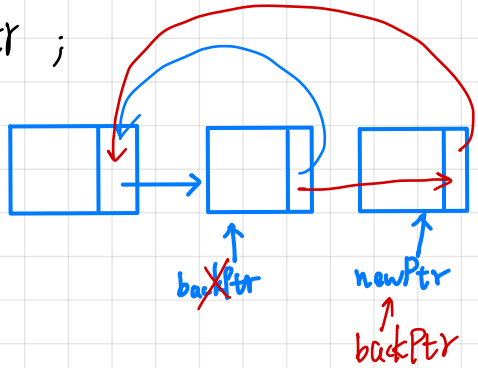
$\text{newPtr} \rightarrow \text{next} = \text{backPtr} \rightarrow \text{next};$

$\text{backPtr} \rightarrow \text{next} = \text{newPtr};$

} // else

$\text{backPtr} = \text{newPtr};$

} // end enqueue



`void Queue :: dequeue() throw (QueueException) {`

`if (isEmpty()) throw...;`

`else {`

`QueueNode * tempPtr = backPtr->next;`

`if (backPtr->next == backPtr)`

`backPtr = NULL;`

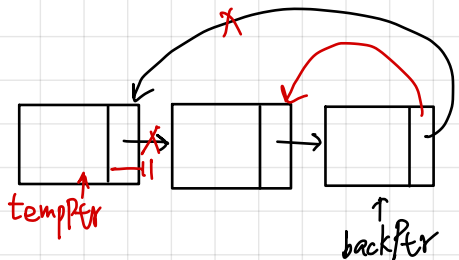
`else backPtr->next = tempPtr->next;`

`tempPtr->next = NULL;`

`delete tempPtr;`

`{ // else`

`} // end dequeue.`



心得: Queue 跟前面學過的 stack 滿像的, 只是存資料跟取資料的順序有點不同, 實作起來也不會很難.

排序:

stable sort vs. unstable sort

bubble

selection

insertion

quick

merge

heap

radix

2 5<sub>a</sub> 5<sub>b</sub> 6 8<sub>b</sub> 8<sub>a</sub> 14 26 28 29

順序可能被對調

selection sort

單筆資料太大 → 搬動費時 才會選 selection sort

Mergesort 排序快, 缺點需要額外的陣列存放排好的結果

quick sort 不是 stable

```
void radixsort(int A[], int first, int last) {  
    for (int base = 1; (maxData/base) > 0; base *= 10) {
```

```
        ...  
        for (i = first; i <= last; i++)  
            bucket[(A[i]/base)%10]++;  
        // 算各餘數有幾個
```

```
        for (i = 1; i < 10; i++)  
            bucket[i] += bucket[i-1];
```

```
        for (i = first; i <= last; i++)  
            temp[bucket[(A[i]/base)%10]++] = A[i];
```

```
        for (i = first; i <= last; i++)  
            A[i] = temp[i];
```

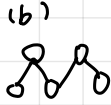
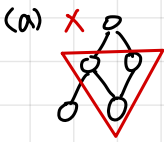
心得:

這單元教了很多排序法, 也在作業中確實看到每個不同排序法效率上的差異, 每個排序法都要想一下才可以理解. 也在實作中看到 stable & unstable 分別是哪些

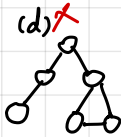
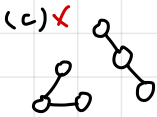
# Tree

有大小階級關係可用 Tree.

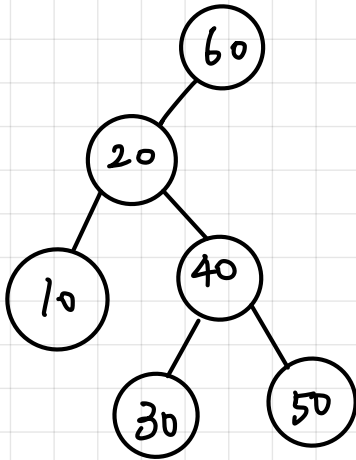
不是 binary tree:



旋轉看是否是 binary tree



order:



preorder: 60, 20, 10, 40, 30, 50, 70

inorder: 10, 20, 30, 40, 50, 60, 70

postorder: 10, 30, 50, 40, 20, 70, 60

要保持原樣  $\Rightarrow$  用 preorder 存檔案.

preorder  
+  
postorder

無法還原一顆唯一的 binary search tree

最高比較次數  $\Rightarrow$  樹高.

心得:

大一的計概有寫過 preorder, inorder, postorder 的 code

那時候理解了很久才寫出來, 第二次遇到

binary tree 我覺得理解起來有輕鬆一點