

CS 5100 01/23

Last time

heuristic function

$h(n)$ : estimated cost from node  $n$  to a goal node

Previously in UCS

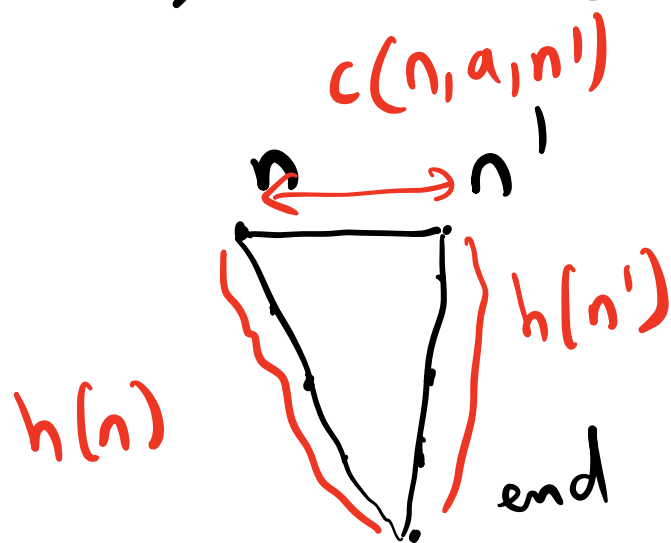
$g(n)$ : total path cost from the start node to node  $n$

$$f(n) = g(n) + h(n)$$

total estimated cost from start node to a goal node while going through node  $n$

Conditions on a heuristic function to get optimality

- 1) Admissibility: never overestimate the cost to reach a goal node
- 2) Consistency (triangle inequality)



$$\boxed{h(n)} \leq \boxed{h(n') + c(n, a, n')} \\ \forall n, n'$$

Consistency  $\Rightarrow$  admissibility  
(by induction)

\* The values of  $f(n)$  are non-decreasing along any path.

Take a node  $n'$  which is a successor of node  $n$ , we need to have

$$\underline{f(n')} \geq \underline{f(n)} = \underline{g(n)} + \underline{h(n)}$$

(Show this as exercise)

$n \quad n'$

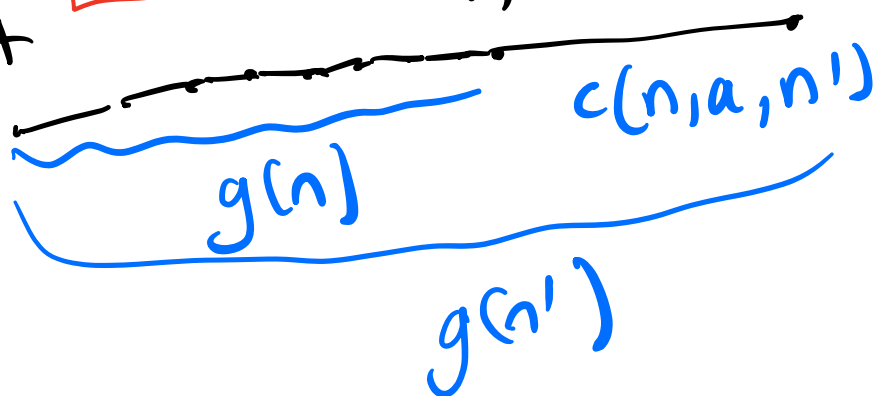
$$f(n') = \underline{g(n')} + h(n')$$

$$\rightarrow = \underline{g(n)} + \boxed{c(n, a, n')} + h(n')$$

by consistency of  $h$

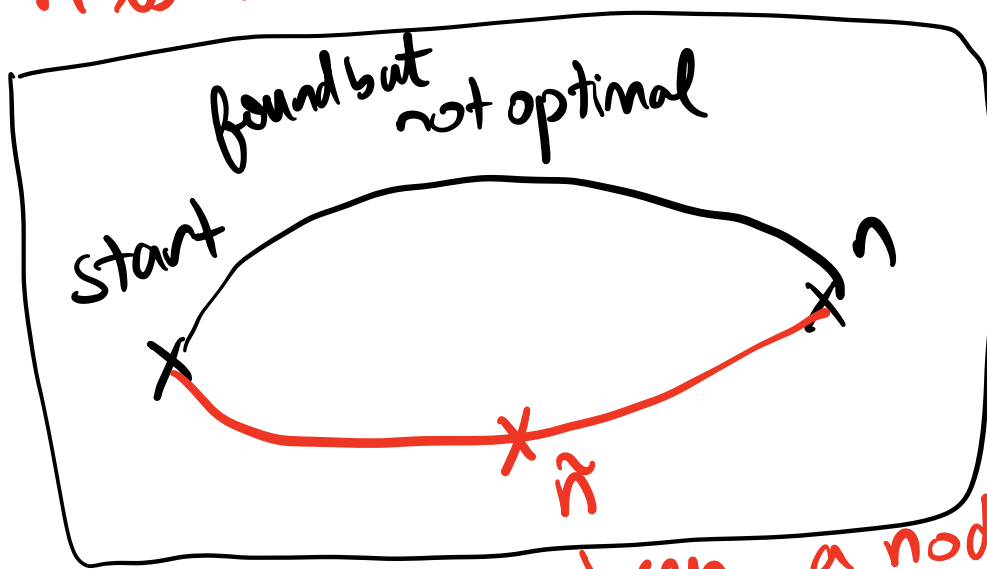
$$\rightarrow \geq \underline{g(n)} + \boxed{h(n)} = f(n)$$

start



\* Show that whenever a node is selected for expansion (using priority queue with  $f(n) = g(n) + h(n)$ ), the optimal path to node  $n$  has been found

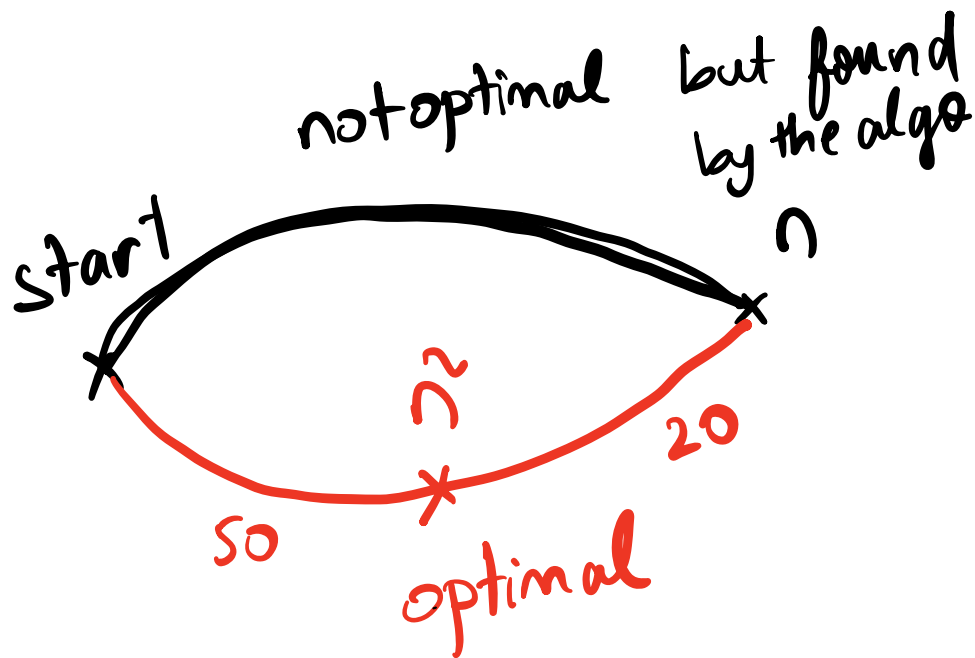
Assume that the path found to node  $n$  is not the optimal path



there must have been a node  $\hat{n}$  on the optimal path from start to  $n$

$$f(\hat{n}) \leq f(n) \quad (\text{From the previous exercise})$$

$\hat{n}$  must be popped first from the priority queue



$$f(\hat{n}) \leq f(n)$$

priority queue

100

~~50~~

$\{ \underline{n}, \hat{n} \}$

\* pop  $\hat{n}$  before  $n$

\* explore successors of  $\hat{n}$  → find  $n$

70

70

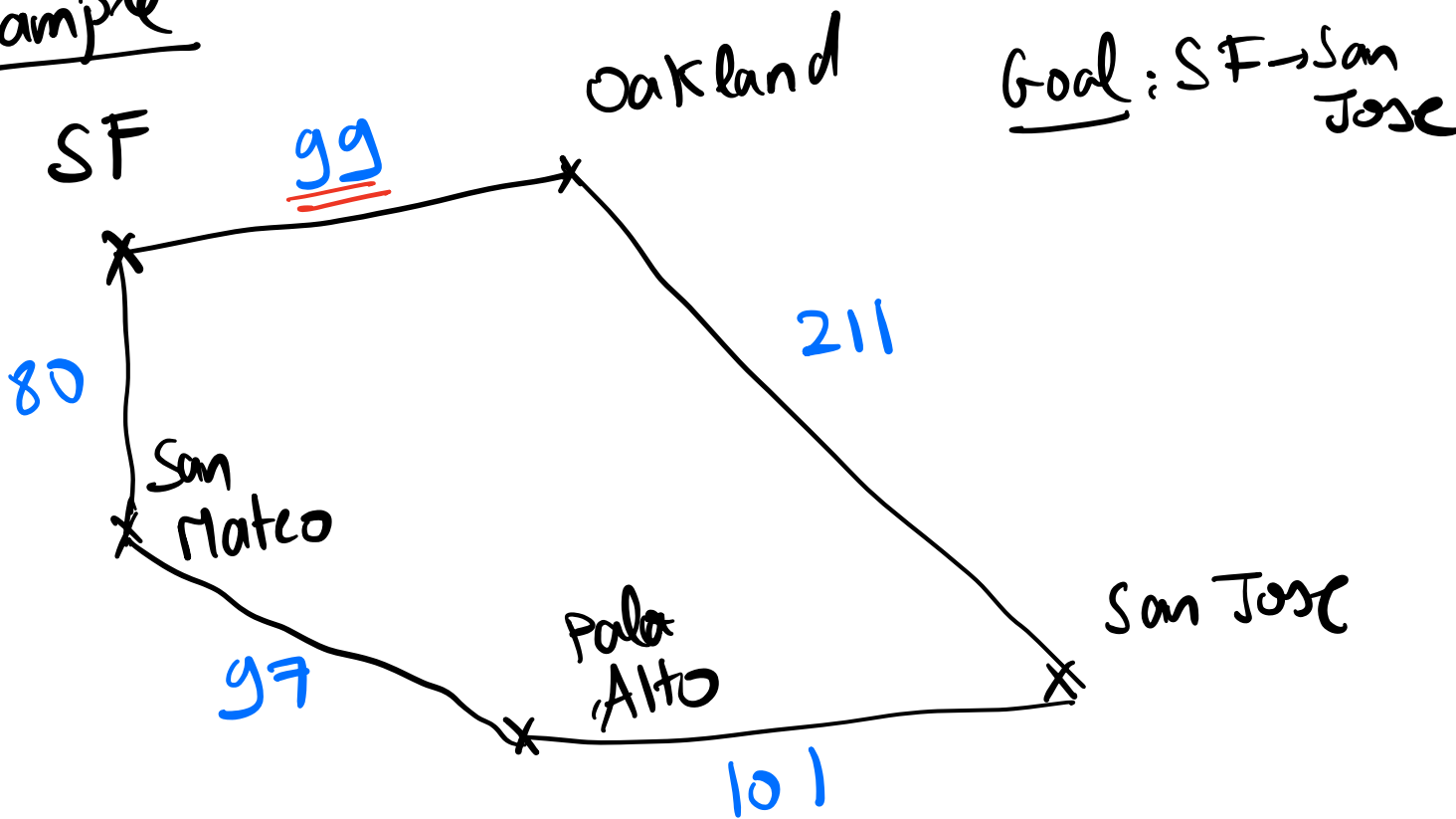
$\{ n \}$

Algorithm : A\* algorithm

(identical to UCS

but we're using the  
Function  $f(n) = g(n) + h(n)$   
when considering the priority  
queue instead of only  
using  $g(n)$

Example



Use straight line distance heuristic

$$h_{SF \rightarrow SJ} = 200$$

$$\rightarrow h_{SM \rightarrow SJ} = 120$$

$$\underline{h_{OA \rightarrow SJ}} = 150$$

$$h_{PA \rightarrow SJ} = 70$$

node = SF

frontier = { SF }

explored = {}

loop

node popped = SF (not goal)

explored = { SF }

loop actions

child = San Mateo

$80 + 120 = 200$

frontier = { San Mateo }

child = Oakland

$99 + 150 = 249$

frontier = { San Mateo, Oakland }

node popped = San Mateo (not goal)

explored = { SF, San Mateo }

loop actions

child = Palo Alto

$247 = 80 + 97 + 70$

frontier = { Oakland, Palo Alto }

node popped = Palo Alto (not goal)

explored = { SF, San Mateo, Palo Alto }

loop actions

child = San Jose (not goal)  
249 278+0

frontier = { Oakland, San Jose }

node popped = Oakland

explored = { SF, San Mateo, PA, Oakland }

loop actions

child = San Jose

don't do anything

Node popped = San Jose (goal node)



## Exercise (code on canvas)

start at value 1

actions  either add 1  
multiply by 2

Goal: value  $n$

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Games like chess, whenever I am at some state, my action will really depend on what my opponent will also be doing. (the action that I choose will give me an advantage, some positive utility)

\* Initial state

\* Player(s)

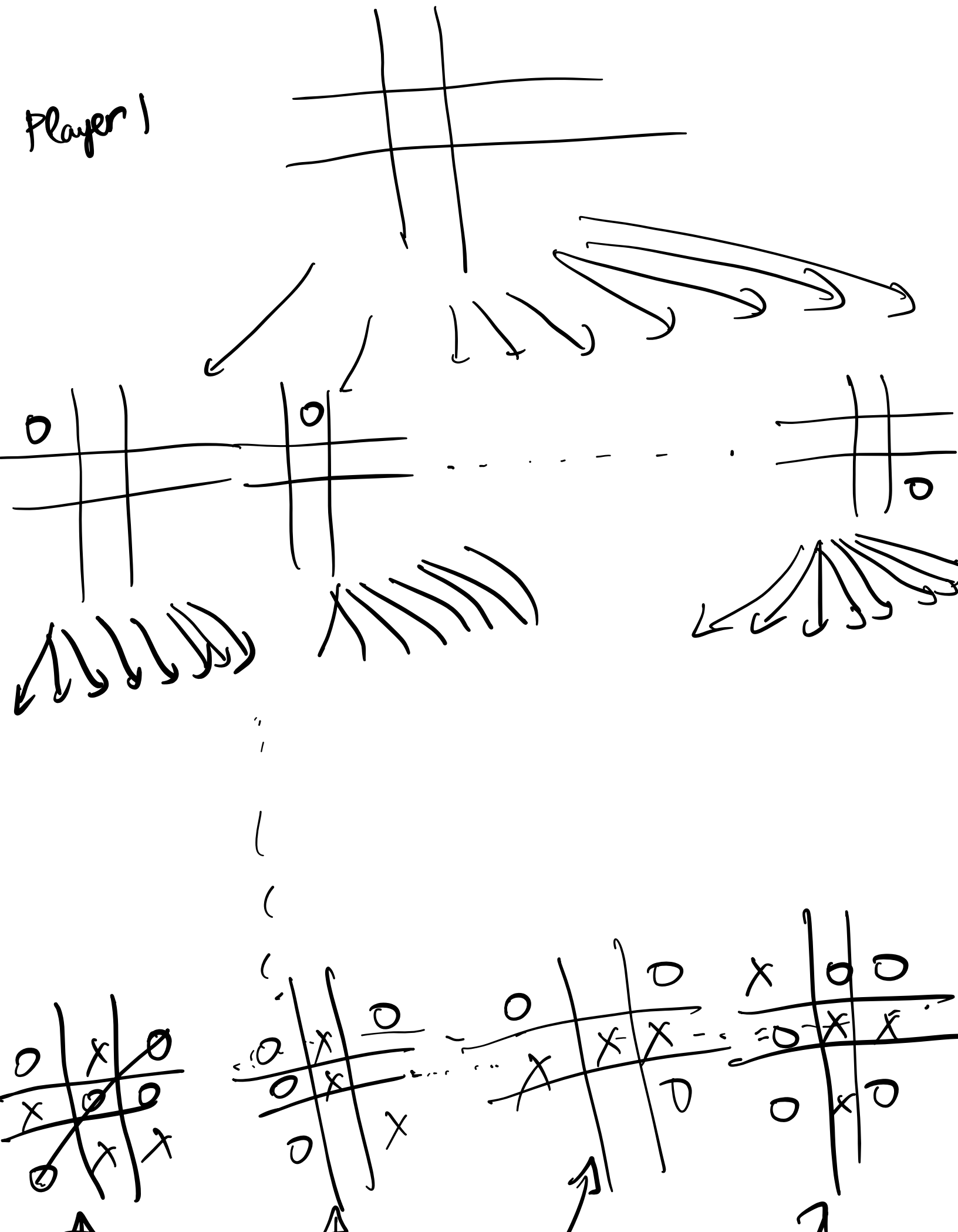
\* Actions(s)

\* Result(s,a) (transition model)  
    → state

\* Utility(s,p) (positive if great for me  
and negative if bad for me)  
    (good for my opponent)

\* terminal\_test(s)  
    ↓  
    allowing me to check  
    whether the game is done

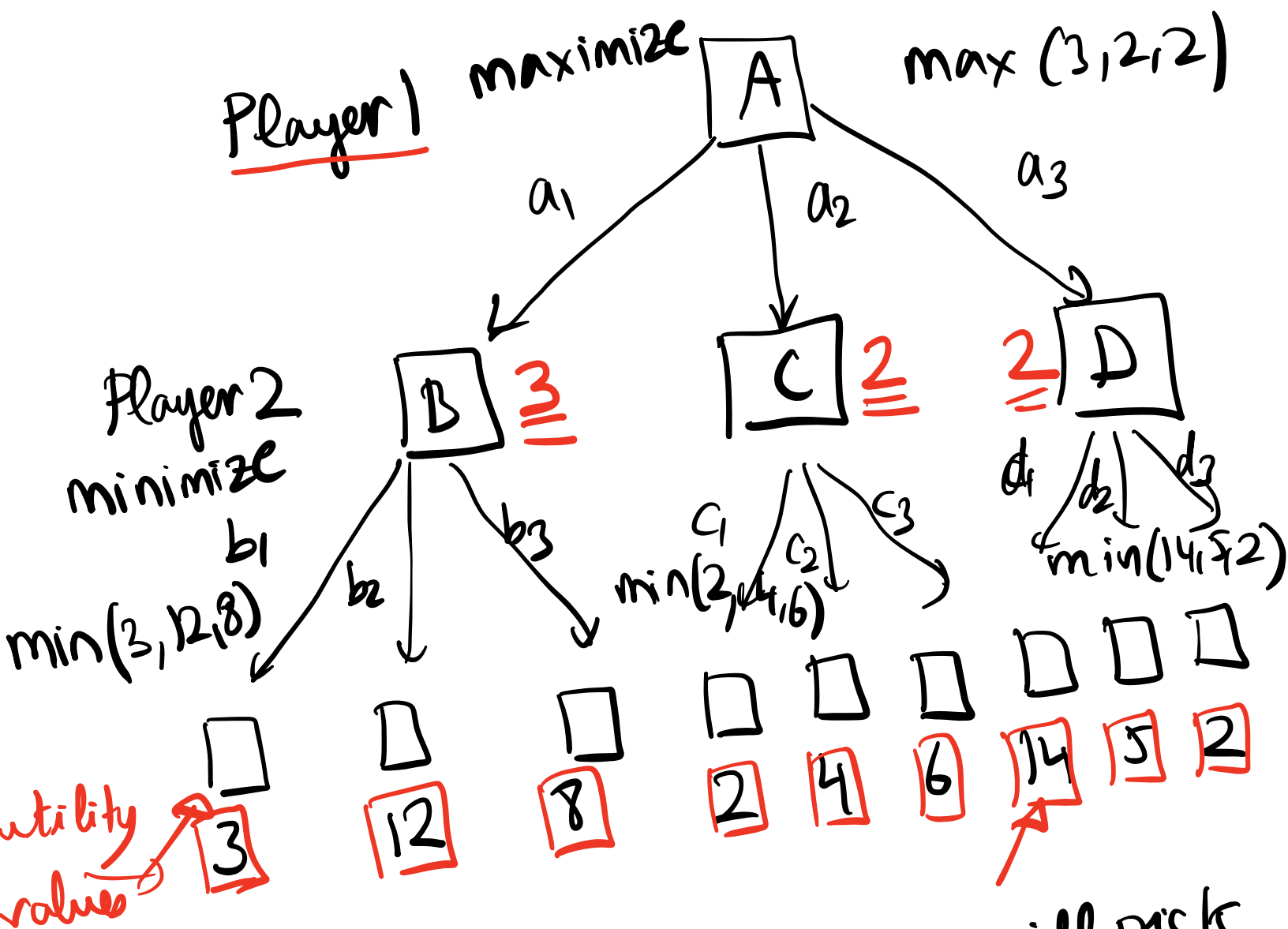
Player 1



↑  
1

↑  
-1

↑  
tie  
0



If I pick  $a_1 \rightarrow$  Xiang ying will pick  $b_1$  to minimize utility

If I pick  $a_2 \rightarrow$  Xiangying will pick  $c_1$

If I pick  $a_3 \rightarrow$  Xiangying pick  $d_3$

		0
	X	

$a_1$   
 $b_1$

$\rightarrow 3$