#### Computer Organization and Architecture

## 1 Introduction

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# **Computer is Everywhere**













**Change or not?** 



# **Review: What is Computer?**



## **Review: What is Computer?**

- 计算机是指 "通用电子数字计算机 (general-purpose electronic digital computer)"
  - 通用: 不是一种专用设备
    - 所有计算机在给予足够时间和容量存储器的条件下,都可以完成同样的计算
    - 当希望完成新的计算时,不需要对计算机重新设计
  - 电子(非机械): 采用电子元器件
  - 数字(非模拟):信息采用数字化的形式表示
- 计算机系统
  - 硬件: 处理器、存储器、外部设备
  - 软件:程序,文档



## **Architecture and Organization**

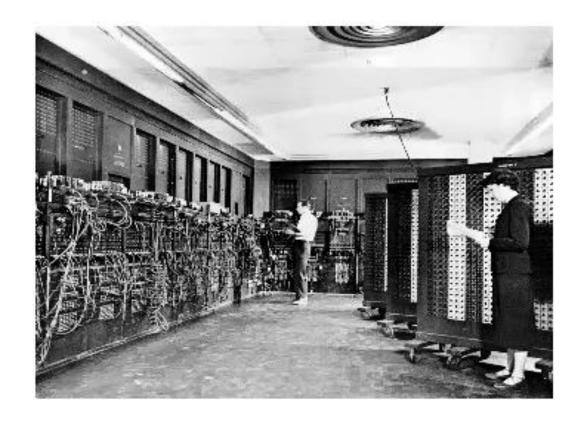
- Architecture (visible to programmer)
  - Its attributes have a direct impact on the logical execution of a program
  - Instruction set, the number of bits to represent data types, ...
    - E.g.: Is there a multiply instruction?
- Organization (opaque/encapsulated to programmer)
  - The operational units and their interconnections
  - Control signals, memory technology, ...
    - E.g.: Implement multiply by a hardware unit or repeated addition?



## **A Brief History of Computers**

- First generation: Vacuum tubes (1946-1957)
  - ENIAC (1946-1955): decimal, manually programming

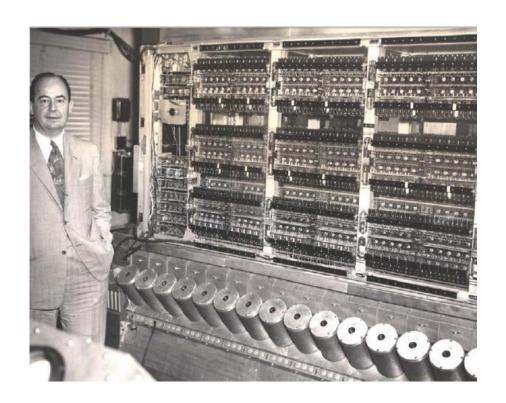






## **A Brief History of Computers**

- First generation: Vacuum tubes (1946-1957)
  - IAS (1946-1952\*): binary, stored program





#### **Review: von Neumann Machine**



#### **Review: von Neumann Machine**

#### • 组成部分

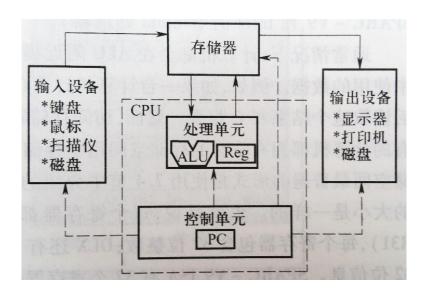
• 存储器: 地址和存储的内容

• 处理单元: 执行信息的实际处理

• 控制单元: 指挥信息的处理

• 输入设备:将信息送入计算机中

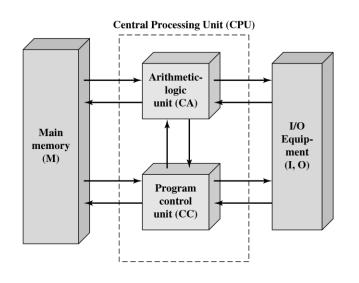
• 输出设备: 将处理结果以某种形式显示在计算机外





## **A Brief History of Computers**

- First generation: Vacuum tubes (1946-1957)
  - IAS (1946-1952\*): the von Neumann machine
    - Idea: main memory stores programs and data
    - Prototype of all subsequent computers
      - Central Arithmetical (CA)
      - Central Control (CC)
      - Memory (M)
      - Input (I) / Output (O)



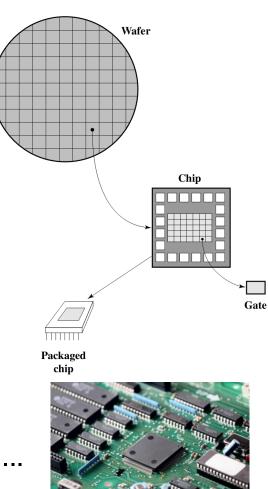


- Second generation: Transistors (1958-1964)
  - NCR and RCA, IBM 7000: transistors are smaller, cheaper, and dissipate less heat; they also can be used as same as vacuum tubes
  - Introduction of more complex arithmetic and logic units and control units, the use of high-level programming languages and the provision of system software.





- Third to *N* generation: Integrated circuits (1965-now)
  - Idea:
    - fabricate an entire circuit in a piece of silicon rather than assemble discrete components made from separate pieces of silicon
    - These transistors can be connected with a process of metallization to for circuits
  - Scale:
    - small → large → very large → ultra large ...

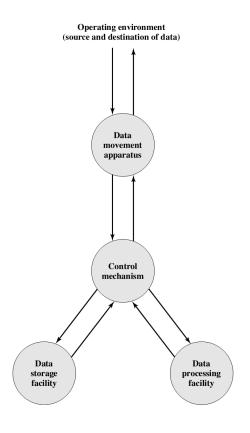


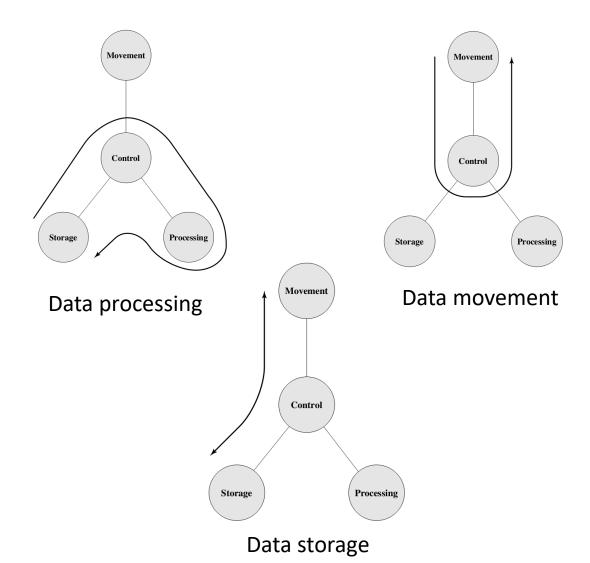


- Moore's law (Gordon Moore, 1965)
  - The number of transistors that could be put on a single chip is doubling every year (1965-1969) / 18 months (1970-now)
  - Consequence
    - The cost of computer logic and memory circuitry has dramatically fallen for the cost of a single chip is unchanged
    - Smaller size leads to flexibility and probability
    - Increased operating speed for shortened electrical path length
    - Reduction in dependency on power and cooling
    - The interconnections on the integrated circuit are more reliable than solder connections and few interchip connections



Constant function







Performance improvements

Generation	Approximate Dates	Technology	Typical Speed (operations per second)
1	1946–1957	Vacuum tube	40,000
2	1958–1964	Transistor	200,000
3	1965–1971	Small and medium scale integration	1,000,000
4	1972–1977	Large scale integration	10,000,000
5	1978–1991	Very large scale integration	100,000,000
6	1991–	Ultra large scale integration	1,000,000,000



## **Computer Performance**

- One of key parameters of computer
  - Performance, cost, size, security, reliability, power consumption, ...
- Sample performance evaluation criteria
  - CPU: speed
  - Memory: capacity, speed
  - I/O: speed, capacity

The main goal / driver is the increase of CPU speed



#### **CPU Performance**

- System clock
  - Clock rate / clock speed (HZ): fundamental rate in cycles per second at which a computer performs its most basic operations
  - Clock cycle / clock tick: a single electronic pulse of a CPU
  - Cycle time (s): the time between pulses
    - Sometimes "clock cycle"



## **CPU Performance (cont.)**

- Instruction execution
  - A processor is driven by a clock with a constant frequency f or, equivalently, a constant cycle time t.
  - Let  $CPI_i$  be the number of cycles required for instruction type i, and  $I_i$  be the number of executed instructions of type i for a given program.
  - The overall can be calculated as follows:

$$CPI = \frac{\sum_{i=1}^{n} (CPI_i \times I_i)}{I_c}, \ I_c = \sum_{i=1}^{n} I_i$$

Process time to execute a given program:

$$T = I_c \times CPI \times t$$
$$T = I_c \times [p + (m \times k)] \times t$$

Transfer data between processor and memory



## **CPU Performance (cont.)**

Million Instructions Per Second (MIPS)

$$MIPS = \frac{I_c}{T \times 10^6} = \frac{f}{CPI \times 10^6}$$

Million Floating Point Operations Per Second (MFLOPS)

$$MFLOPS = \frac{N_{floating-point\ op}}{T \times 10^6}$$



## **CPU Performance (cont.)**

- Benchmarks
  - Measure the performance of systems using a set of benchmark programs
  - Averaging results:
    - Arithmetic mean:  $R_A = \frac{1}{m} \sum_{i=1}^{m} R_i$
    - Harmonic mean:  $R_H = \frac{m}{\sum_{i=1}^{m} \frac{1}{R_i}}$



### Summary

- Concepts
  - organization, architecture
- Computer history
  - The von Neumann machine, Moore's law
- Possible computer operations
- Computer performance
  - CPU performance evaluation



# **Thank You**

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