

# Chapter 8

## The Disjoint Set ADT

# Equivalence Class

## 1) Definition of Equivalence Class:

Suppose we have a set  $U=\{1,2,\dots,n\}$  of  $n$  elements and a set  $R=\{(i_1,j_1), (i_2,j_2),\dots, (i_r,j_r)\}$  of  $r$  relations. The relation  $R$  is an equivalence relation iff the following conditions are true (symbol ' $\equiv$ ' represent the equivalence relation on sets,  $x,y,z$  are elements in set ) :

- Reflexive  $x \equiv x$ .
- Symmetric  $x \equiv y, y \equiv x$
- Transitive  $x \equiv y$  and  $y \equiv z$ , then  $x \equiv z$

# Equivalence Class

例如:

判别3个数a, b, c能否构成三角形的三条边?

能构成三角形的等价类:

$\{ (3,4,5), (4,5,6), (5,6,7), \dots \}$

不能构成三角形的等价类:

$\{ (1,2,3), (2,3,5), \dots \}$

# Equivalence Class

## 2) Example:

set  $s = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11\}$

pairs of equivalence:

$(0\ 4), (3\ 1), (6\ 10), (8\ 9), (7\ 4), (6\ 8), (3\ 5), (2\ 11),$   
 $(11\ 0)$

Initial:  $\{0\}, \{1\}, \{2\}, \{3\}, \{4\}, \{5\}, \{6\}, \{7\}, \{8\},$   
 $\{9\}, \{10\}, \{11\}$

$0 \equiv 4$   $\{0, 4\}, \{1\}, \{2\}, \{3\}, \{5\}, \{6\}, \{7\}, \{8\}, \{9\},$   
 $\{10\}, \{11\}$

$3 \equiv 1$   $\{0, 4\}, \{1, 3\}, \{2\}, \{5\}, \{6\}, \{7\}, \{8\}, \{9\},$   
 $\{10\}, \{11\}$

# Equivalence Class

$6 \equiv 10$  {0, 4}, {1, 3}, {2}, {5}, {6, 10}, {7}, {8},  
{9}, {11}

$8 \equiv 9$  {0, 4}, {1, 3}, {2}, {5}, {6, 10}, {7},  
{8, 9}, {11}

$7 \equiv 4$  {0, 4, 7}, {1, 3}, {2}, {5}, {6, 10}, {8, 9}, {11}

$6 \equiv 8$  {0, 4, 7}, {1, 3}, {2}, {5}, {6, 8, 9, 10}, {11}

$3 \equiv 5$  {0, 4, 7}, {1, 3, 5}, {2}, {6, 8, 9, 10}, {11}

$2 \equiv 11$  {0, 4, 7}, {1, 3, 5}, {2, 11}, {6, 8, 9, 10}

$11 \equiv 0$  {0, 4, 7, 2, 11}, {1, 3, 5}, {6, 8, 9, 10}

# Equivalence Class

## 3) Online equivalence class operation

- **Combine(a,b) : combine the equivalence classes that contains elements a and b into a single class**
- **Find(e) : determine the class that currently contains element e.**

**Combine(a,b) is equivalent to**

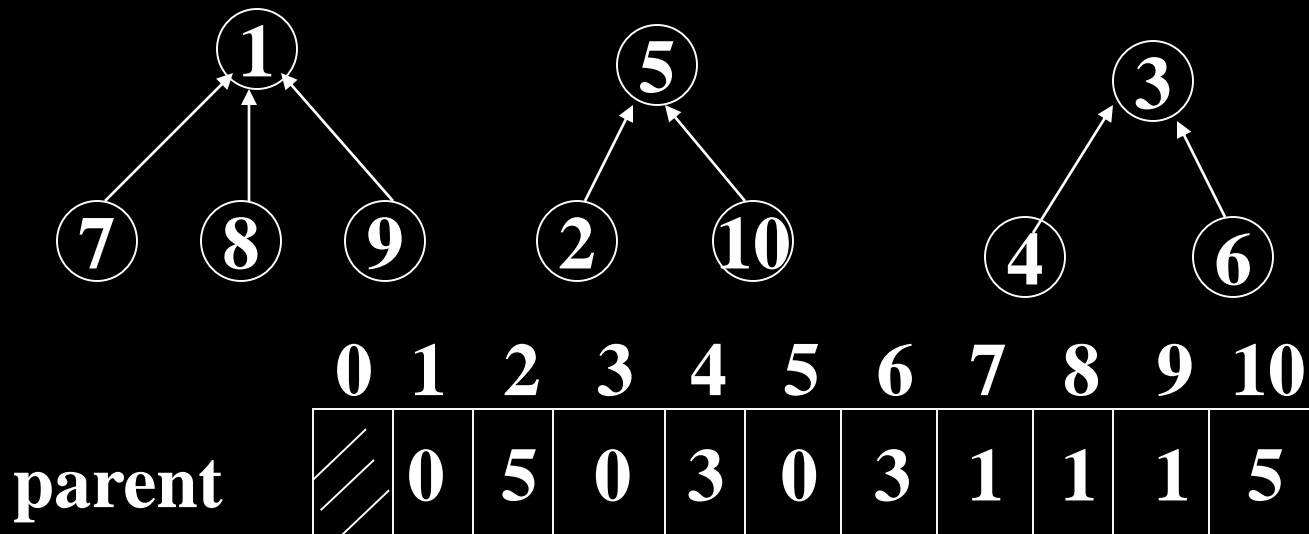
**i=Find(a); j=Find(b); if(i!=j) Union(i,j);**

# Equivalence Class

## 4) Tree Representation(Union-Find sets)

**Example:**

$S_1=\{1,7,8,9\}$ ,  $S_2=\{5,2,10\}$ ,  $S_3=\{3,4,6\}$ , they all belong to  $S=\{1,2,3,\dots,10\}$



# Equivalence Class

simple tree solution to union-find problem

```
void Initialize(int n)
```

```
{  parent=new int[n+1];
```

```
   for(int e=1;e<=n;e++)
```

```
       parent[e]=0;
```

```
}
```

```
int Find(int e)
```

```
{  while(parent[e])
```

```
       e=parent[e];
```

```
   return e;
```

```
}
```

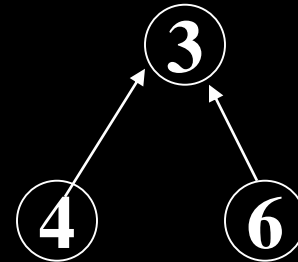
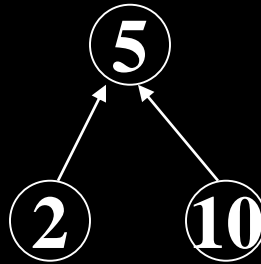
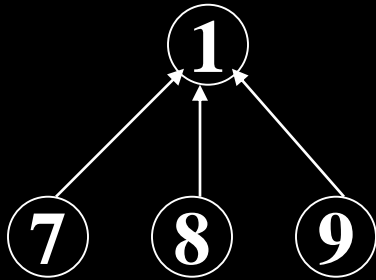
```
void Union(int i, int j)
```

```
{  parent[j]=i;
```

```
}
```



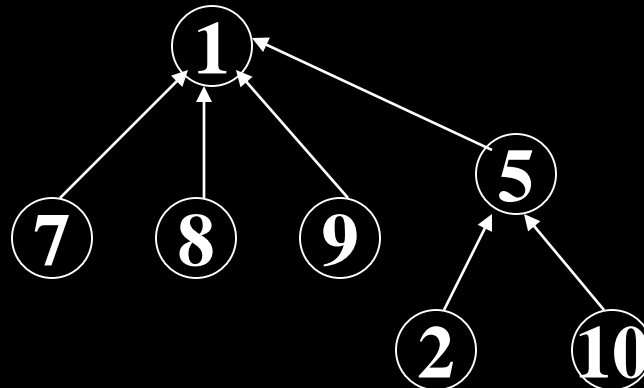
## Equivalence Class



parent

	0	1	2	3	4	5	6	7	8	9	10
parent		0	5	0	3	<del>10</del>	3	1	1	1	5

Union ( 1,5 )



# Equivalence Class

Java

```
public class DisjSets
{ public DisjSets( int numElements )
    public void union( int root1, int root2 )
    public int find( int x )
    private int [ ] s;
}
```

```
public DisjSets( int numElements )
{ s = new int [ numElements ];
    for( int i = 0; i < s.length; i++ )
        s[ i ] = -1; //一个根结点
}
```

# Equivalence Class

```
public void union( int root1, int root2 )
```

```
{  s[ root2 ] = root1;
```

```
}
```

```
public int find( int x )
```

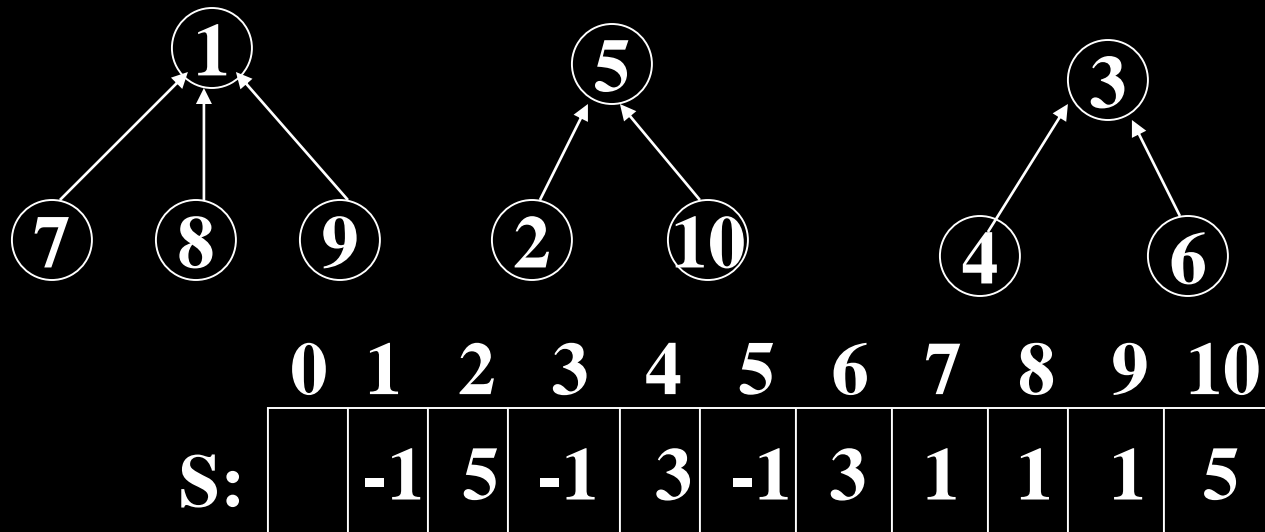
```
{  if( s[x] < 0 )
```

```
    return x;
```

```
    else
```

```
    return find( s[ x ] );
```

```
}
```



# Equivalence Class

## 5) Performance Evaluation

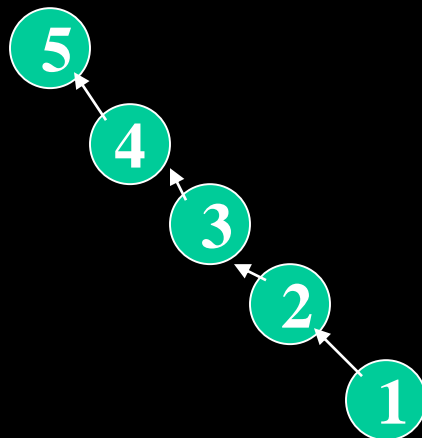
Time complexity: Find--  $O(h)$ ,

Union--  $\theta(1)$

Assume that  $u$  times unions and  $f$  times finds are to be performed,  $f > u$ ,

in the worst case a tree with  $m$  elements can have a height of  $m$ :

Union(2,1), Union(3,2), Union(4,3), Union(5,4)...



# Equivalence Class

## 6) Performance Enhancement

- ★improve Union in order to decrease the time each find take, so that the height of tree will not increase linearly .
- ★.Improvement of *Find* –path compression

# Equivalence Class

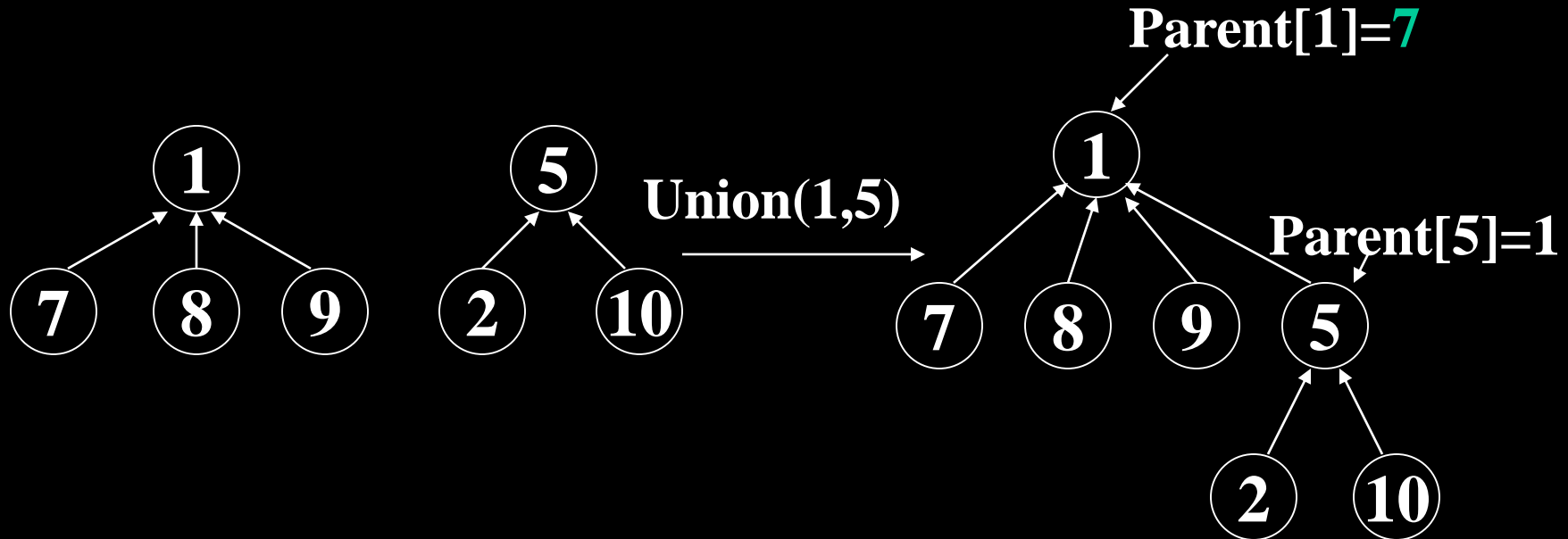
## ★improve Union

two rules:

- **Weight rule:** if the number of nodes in tree  $i$  is less than the number in tree  $j$ , then make  $j$  the parent of  $i$ ; otherwise, make  $i$  the parent of  $j$ .
- **Height rule:** if the height of tree  $i$  is less than that of tree  $j$ , then make  $j$  the parent of  $i$ ; otherwise, make  $i$  the parent of  $j$ .

# Equivalence Class

Let's discuss the weight rule(C++) :



Besides the *parent* field, each node has a boolean field *root*. The *root* field is true iff the node is presently a root node. The *parent* field of each root node is used to keep a count of the total number of nodes in the tree.

两个数组： Parent, root 黑板上具体讲。

# Equivalence Class

Union with the weight rule

```
void Initialize(int n)
{ root=new bool[n+1];
  parent=new int[n+1];
  for(int e=1;e<=n;e++)
  { parent[e]=1;
    root[e]=true;
  }
}

int Find(int e)
{ while(!root[e])
  e=parent[e];
  return e;
}
```



# Equivalence Class

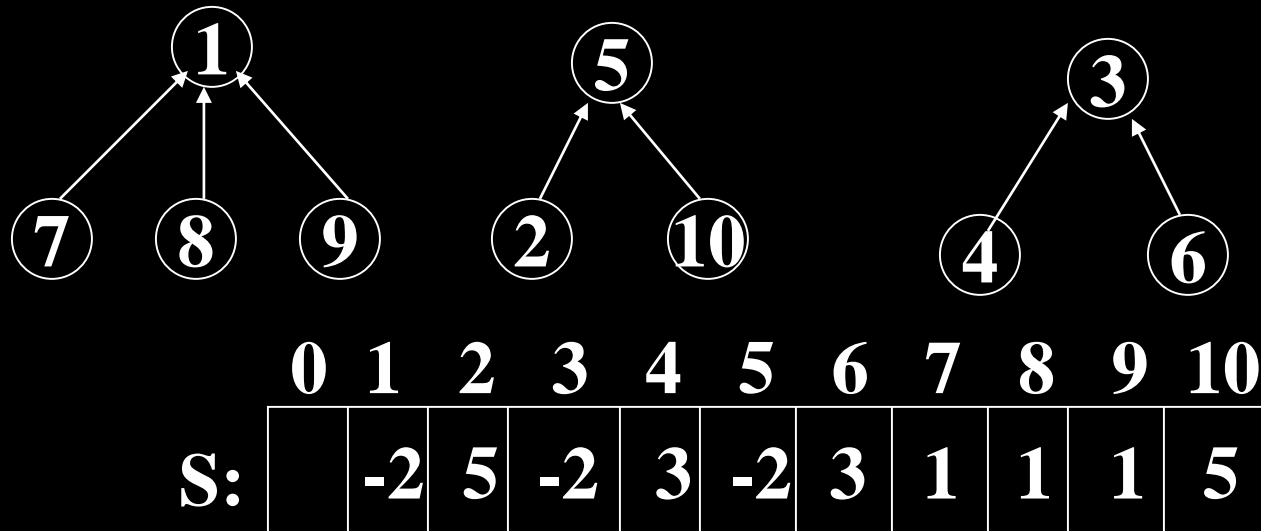
```
void Union(int i, int j)
{ if(parent[i]<parent[j]) // i becomes subtree of j
  { parent[j]=parent[i]+parent[j];
    root[i]=false;
    parent[i]=j;
  }
  else { parent[i]=parent[i]+parent[j];
        root[j]=false;
        parent[j]=i;
      }
}
```

# Equivalence Class

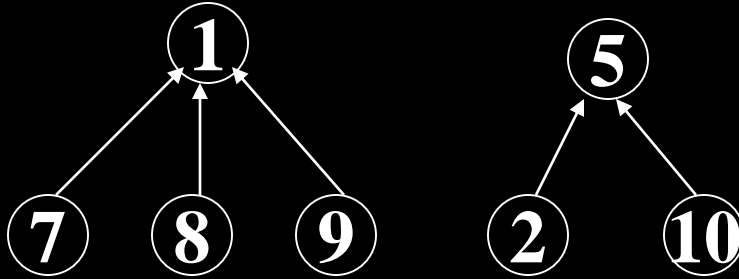
## Java(高度规则)

用一个数组来实现，根结点中放负数，而且是代表高度。

```
public void union( int root1, int root2 )  
{   if( s[ root2 ] < s[ root1 ] )  
      s[ root1 ] = root2;  
    else {   if( s[ root1 ] == s[ root2 ] )  
              s[ root1 ]--;  
      s[ roo2 ] = root1;} }
```



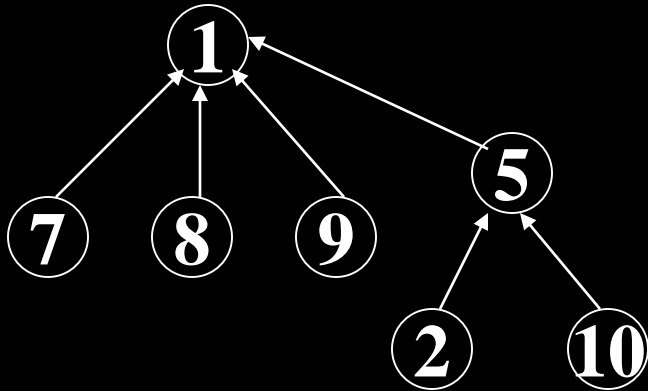
# Equivalence Class



**union ( 1, 5 )**

**s[root1]--**

**s[root2] = root1**

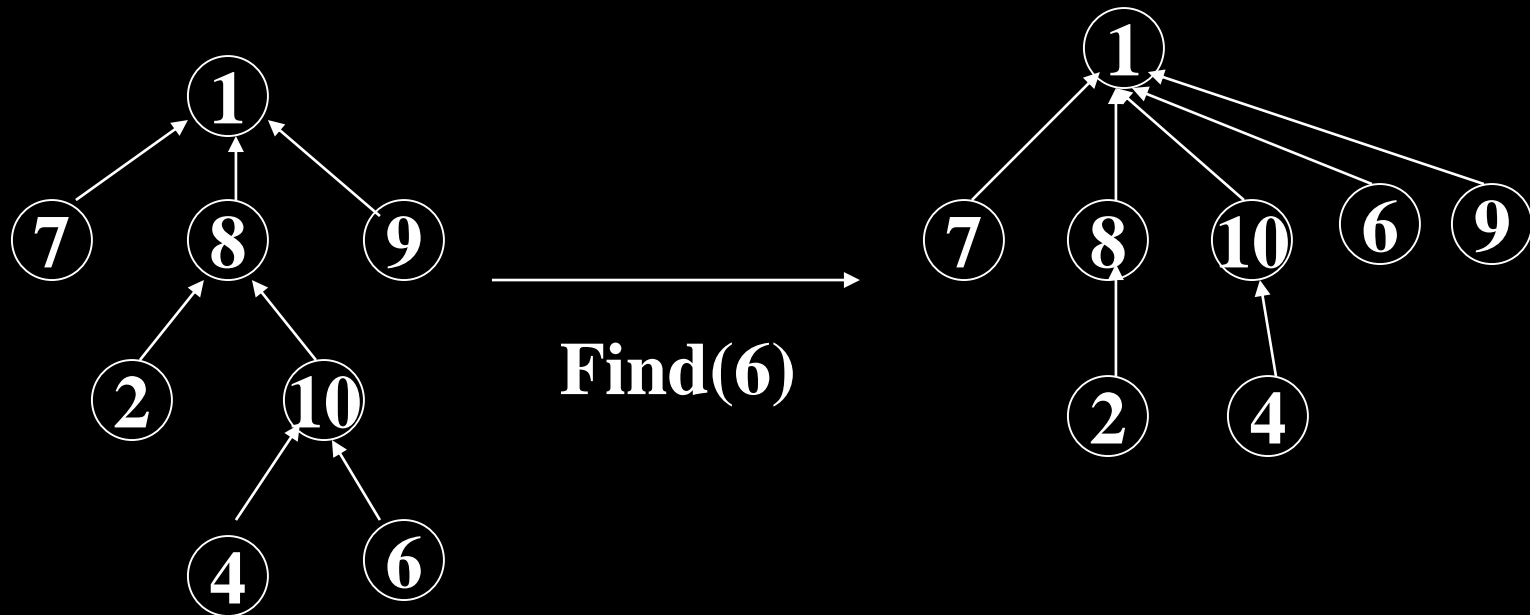


# Equivalence Class

## ★.Improvement of *Find* –path compression

When processing a equivalence pair, we need to operate *Find* twice, *WeightUnion* once.

Example of improvement:



# Equivalence Class

```
int Find( int e) { /* C++ */  
{ int j=e;  
  while(!root[j]) j=parent[j];  
  int f=e;  
  while(f!=j)  
    { int pf=parent[f];  parent[f]=j; f=pf; }  
}
```

# Equivalence Class

## Java

```
public int find( int x )  
{   if( s[ x ] < 0 )  
        return x;  
    else  
        return s[ x ] = find( s[ x ] );  
}
```