



Extension of the CKKS Cryptosystem: Part 2

Introduction to Homomorphic Cryptosystems - Lecture 7

THE BINARY STEP FUNCTION



Definition

The binary step function (also called Heaviside step function) is zero for negative and one for positive arguments:

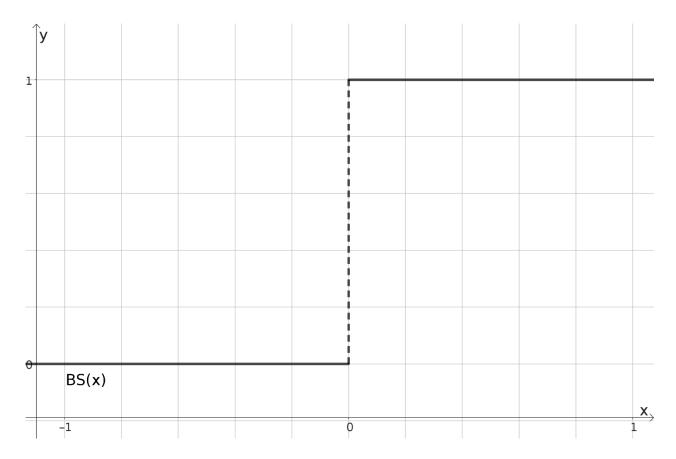
$$BS(x) = \begin{cases} 0 & \text{if } x < 0 \\ 1 & \text{if } x > 0 \end{cases}$$

What about the zero argument?

Depends on the use case. For our use case we define $BS(0) = \frac{1}{2}$.

Use Cases

- Activation function for neural networks
- > Filter



Mathematical definition

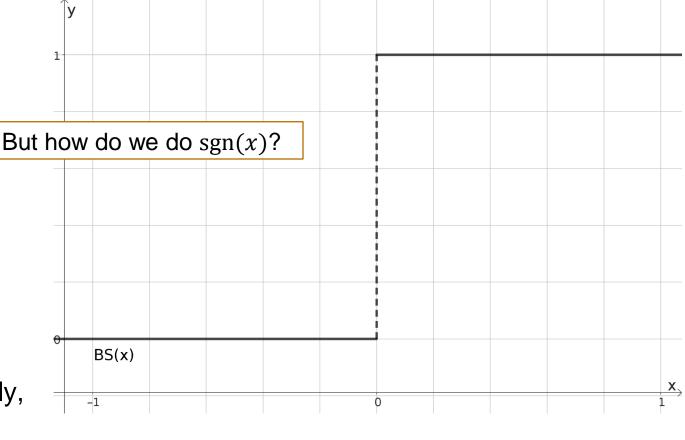
We want to calculate the BS function homomorphically, so we need a definition based on our limited functions:

$$BS(x) = \frac{\operatorname{sgn}(x) + 1}{2}$$
$$\operatorname{sgn}(x) = \frac{x}{\sqrt{x^2}}$$

This leaves us with this definition:

$$BS(x) = \frac{1}{2} * \left(x * \frac{1}{\sqrt{x^2}} + 1 \right)$$

Which we can implement homomorphically, as we have already extended CKKS by division and square root.

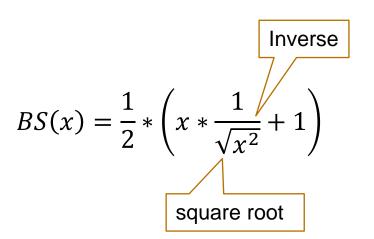


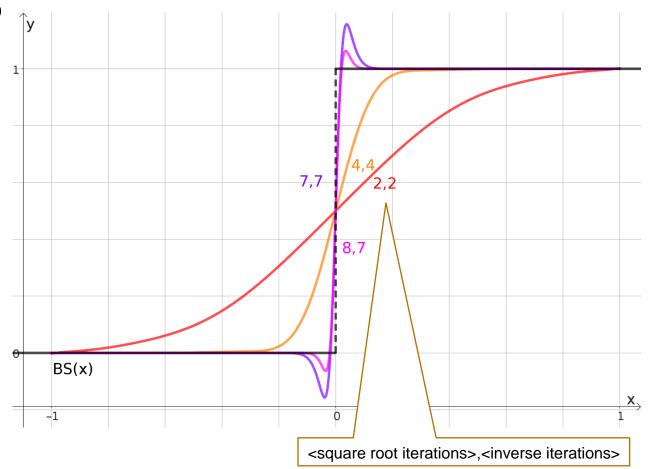


Homomorphic Implementation

Depending on the number of iterations we do for the division and square root, we achieve varying levels of precision when approximating the BS function homomorphically.

The function contains two complex operations (which we already implemented):







How can we use the Binary Step Function?

Example Task

From a given set of tuples:

$$\{(\lambda_1, L_1), (\lambda_2, L_2), \dots (\lambda_n, L_n)\}, \lambda_i \in \mathbb{R}, L_i \in \mathbb{R}$$

Find the λ_x who's corresponding L_x is the minimum of all L_i , $0 < i \le n$.

This is obviously easy on unencrypted numbers:

```
def find_smallest_L(lambdas,L):
    min = -1
    min_L = MAX

for i in range(0, len(L))
    if L[i] < min_L
    min_L = L[i]
    min = i

But we cannot do this on encrypted numbers!</pre>
def find_smallest_L(lambdas,L):
    min_L = min(L)

for i in range(0, len(L))
    if L[i] == min_L
    eturn lambdas[i]

But we cannot do this on encrypted numbers!
```



How can we use the Binary Step Function?

Example Task

From a given set of tuples:

$$\{(\lambda_1, L_1), (\lambda_2, L_2), \dots (\lambda_n, L_n)\}, \lambda_i \in \mathbb{R}, L_i \in \mathbb{R}$$

Find the λ_{min} who's corresponding L_{min} is the minimum of all L_i , $0 < i \le n$.

Homomorphic Solution

With the help of the BS function, we can solve the task even on encrypted numbers. We need two steps:

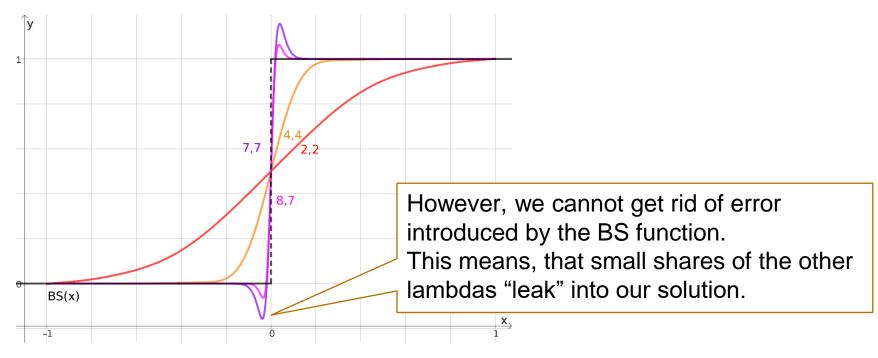
- 1. Find $L_{min} = \min_{i} L_{i} = \min(\min(\min(L_{1}, L_{2}), L_{3}), ...)$ As we already implemented the min function, this step is no problem.
- 2. Calculate λ_{min} :

$$\lambda_{min} = 2 * \sum_{i=1}^{n} BS(L_{min} - L_i) * \lambda_i$$



Scaling Factor

As we saw already, the BS function becomes less accurate when approaching zero:



This is why we normally include a scaling factor to all calculations involving the BS function:

$$\lambda_{min} = 2 * \sum_{i=1}^{n} BS(s(L_{min} - L_i)) * \lambda_i$$



How can we use the Binary Step Function?

Example Task

From a given set of tuples:

$$\{(\lambda_1, L_1), (\lambda_2, L_2), \dots (\lambda_n, L_n)\}, \lambda_i \in \mathbb{R}, L_i \in \mathbb{R}$$

Find the λ_{min} who's corresponding L_{min} is the minimum of all L_i , $0 < i \le n$.

Homomorphic Solution

With the help of the BS function, we can solve the task even on encrypted numbers. We need two steps:

- 1. Find $L_{min} = \min_{i} L_{i} = \min(\min(\min(L_{1}, L_{2}), L_{3}), ...)$ As we already implemented the min function, this step is no problem.
- 2. Calculate λ_{min} :

$$\lambda_{min} = 2 * \sum_{i=1}^{n} BS(s(L_{min} - L_i)) * \lambda_i$$

This is a very interesting solution!
Can we generalise this to get a
function that allows us to do
arbitrary case differentiations?



EQUALS AND SMALLER THAN FUNCTION



Equals Function

Goal

We want to use the BS function to create a conditional assignment. In pseudo code it would look like this:

Solution

$$EQ(a, b, c, d) = c * BS(a - b) * BS(b - a) * 4 + d * (BS(a - b) - BS(b - a))^{2}$$

Properties

$$EQ(a,b,c,d) = \begin{cases} c & \text{if } a = b \\ d & \text{if } a < b \\ d & \text{if } a > b \end{cases}$$



Smaller Than Function

Goal

We want to use the BS function to create a conditional assignment. In pseudo code it would

look like this:

Solution

$$HELP(a, b, c) = \frac{1}{2} \left(c - c \left(BS(a - b) - BS(b - a) \right) \right) - \frac{1}{2} c * BS(a - b)BS(b - a) * 4$$

$$ST(a, b, c, d) = EQ(HELP(a, b, c), 0, d, c)$$

Properties

$$ST(a, b, c, d) = \begin{cases} d & \text{if } a = b \\ c & \text{if } a < b \\ d & \text{if } a > b \end{cases}$$



But: It's more a theory

EQ and ST are demonstrations on how we can work with the limited operations in a FHE cryptosystems.

Their real use depend on different factors:

- Are there different options? (Newton-Raphson, etc..)
- What is the order of magnitude of the numbers?
- What is the distance of the compared numbers?

A big distance can increase the error significantly, as both numbers are part of the equation and the *BS* implementation is not perfect.



Conditional Branching

EQ and ST (and thus, abusing the BS function) enable more complex operations. In some way they resemble conditional branching, as they also allow to change the control flow based on conditions:

```
def example(bool, input):
    if(bool == 0):
        input = <operation1>
        input = <operation2>
        [...]
        return input
    else:
        input = <operation3>
        input = <operation4>
        [...]
        return input

        r = c_1 * <operation1> + c_2 * <operation3>
        r = c_1 * <operation2> + c_2 * <operation4>
        [...]
        return input
```



Summary – What did we learn today?

Binary Step Function

The BS function is zero for negative values and one for positive values. It can be implemented homomorphically, but the error is significant.

Conditional Branching

The BS function can be used to implement conditional branching.

This allows the execution of operations based on a condition.

