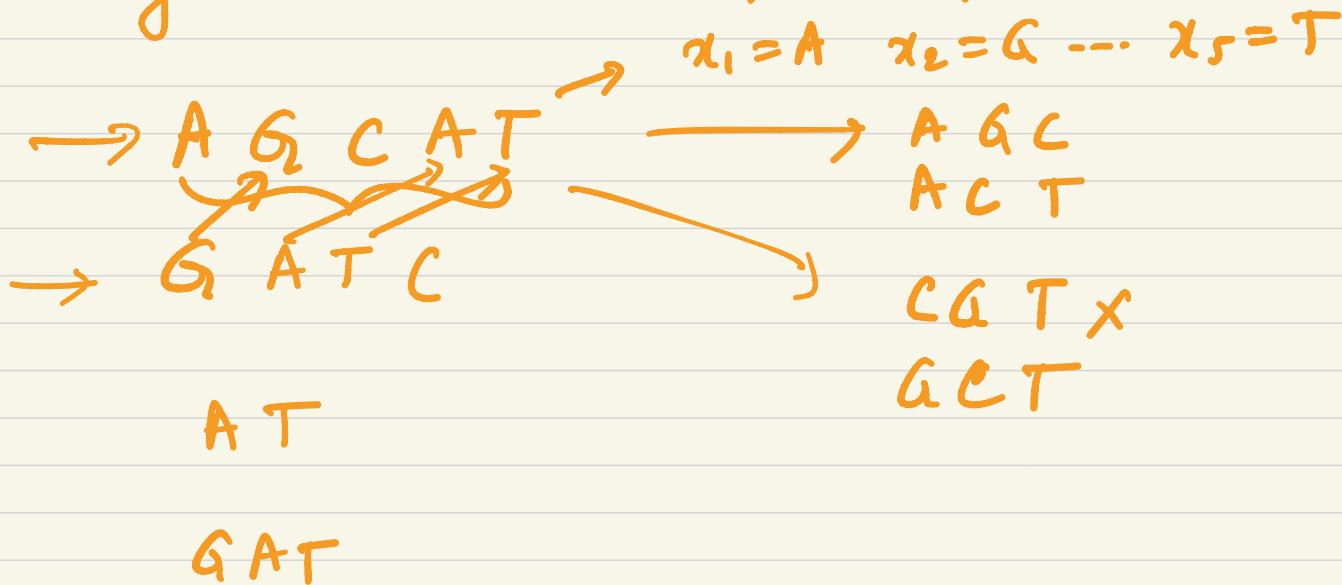


## Longest Common Subsequence Problem.



## Notation :

A sequence  $X = \langle x_1, \dots, x_k \rangle$

$X_i := \langle x_1, \dots, x_i \rangle$  prefix.

A subsequence of  $X$  is a sequence  
of the form  $\langle x_{i_1}, x_{i_2}, \dots, x_{i_r} \rangle$

$$1 \leq i_1 < i_2 < \dots < i_r \leq k$$

Def<sup>n:</sup>

Let  $X$  and  $Y$  be two sequences. We say  
the  $Z$  is a common subsequence of  $X \times Y$

if  $Z$  is subseq of  $X$   
 $\forall Z \text{ `` `` `` } Y$

Defn: We say that  $Z$  is a longest common  
subsequence of  $X$  &  $Y$  if for all common subsequence  
 $W$ ,  $|W| \leq |Z|$

$X = \langle A, T, G, C, A \rangle$  5

$A, T, G, C$  4

$Y = \langle T, G, C, A, T, G, C \rangle$  7

$T, G, C, A$

Thm:- Let  $X = \langle x_1, \dots, x_m \rangle$  and  $Y = \langle y_1, \dots, y_n \rangle$  be two sequences. Let  $Z = \langle z_1, \dots, z_k \rangle$  be an LCS of  $X$  and  $Y$ .

- ① Let  $x_m = y_n$ , then  $z_k = x_m = y_n$  and  $Z_{k-1} = \langle z_1, \dots, z_{k-1} \rangle$  is an LCS of  $X_{m-1} \times Y_{n-1}$ .
- ② If  $x_m \neq y_m$  and  $z_k \neq x_m$  then  $Z$  is an LCS of  $X_{m-1} \times Y$ .
- ③ If  $x_m \neq y_m$  and  $z_k \neq y_n$  then  $Z$  is an LCS of  $X \times Y_{n-1}$ .

Proof:- ① If  $z_k \neq x_m$  then append  $x_m (= y_n)$  at the end of  $Z$  to get a longer common ss. Contra!

Note  $Z_{k-1}$  is a lcs of  $x_{m-1} \& y_{n-1}$

Suppose  $Z_{k-1}$  is not an LCS of " "

Let  $W$  be an LCS of  $x_{m-1} \& y_{n-1}$

$$|Z_{k-1}| < |W|$$

$$\begin{array}{c} Z \\ = \\ \underbrace{z_{k-1} + \langle x_m \rangle}_{\text{LCS}} \quad \underbrace{W + \langle z_m \rangle}_{\text{LCS}} \end{array}$$

$W + \langle x_m \rangle$  will be a  
common ss of  $x \& y$  of

length longer than the  
length of  $Z$

contra!

(2)  $x_m \neq y_n \wedge z_k \neq x_m$ . then

$Z$  is  $\text{css}$  of  $x_{m-1} \wedge Y$

Suppose  $Z$  is not a LCS of  $x_{m-1} \wedge Y$

then HW: contra!

def  $c[i, j]$  = length of an LCS of  
 $X_i \times Y_j$

$$c[i, j] = \begin{cases} 0 & i=0 \text{ or } j=0 \\ c[i-1, j-1] + 1 & i, j > 0, x_m = y_n \\ \max\{c[i, j-1], c[i-1, j]\} & i, j > 0, x_m \neq y_n \end{cases}$$

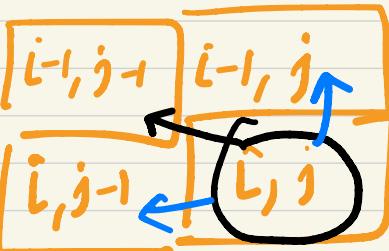
$\Sigma$   
 $\downarrow$   
 string of length 0.

$$Z_{k+1} = \text{LCS}(X_{i-1}, Y_{j-1})$$

$\downarrow$

$$c[i-1, j-1]$$

$i$	$j$	$x$	$\varepsilon$	A	T	G	C	A
0	0	$x_0$	$\varepsilon_0$	0	0	0	0	0
1	1	$x_1$	$\varepsilon_1$	0	0	1	1	1
2	2	$x_2$	$\varepsilon_2$	0	0	1	2	2
3	3	$x_3$	$\varepsilon_3$	0	0	1	2	3
4	4	$x_4$	$\varepsilon_4$	0	1	2	3	4
		T	0					
		G	0					
		C	0					



■  $x_i = y_j$

■  $x_i \neq y_j$

HW:- Pseudocode [CLRS]

## Optimal Binary Search Tree

Keys 1, 2, 3

How many BST's are there with keys 1, 2, 3 ?

