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Dhirubhai Ambani Institute of Information & Communication Technology End Semester Examination, Winter Semester 2021-2022

Course TitleIT215 System SoftwareMax Marks74Date6-May-2022Time120 mins

Instructions:

- All questions are compulsory.
- The answer to each question must be written in the space provided after the question in the question paper only (No need for Answer Booklet).
- Write your answers in brief and with clarity. Writing long answers does not fetch a higher score.

Q1 Concurrent Programming:

[16 Marks]

Consider the code1, code2, code3 and code4 below. For each of these codes, what are all possible counter variables values that can be printed as the first value on stdout? Briefly explain your answer.

```
// code1
                                           // code2
int counter = 2;
                                           int counter = 2:
void foo() {
                                           void foo() {
 counter++;
                                            counter++;
 printf("%d", counter);
                                            printf("%d", counter);
int main() {
                                           int main() {
 pthread_t tid[2];
                                            pthread t tid[2];
 int i;
                                            int i;
 for (i = 0; i < 2; i++) {
                                            for (i = 0; i < 2; i++) {
  pthread_create(&tid[i], 0, foo, 0);
                                             pthread_create(&tid[i], 0, foo, 0);
                                             pthread_join(tid[i], 0);
 counter++;
 printf("%d", counter);
                                            counter++;
                                            printf("%d", counter);
// code3
                                           // code4
int counter = 2;
                                           int counter = 2;
void foo() {
                                           void foo() {
 pthread_mutex_lock(&m);
                                            counter++;
                                            printf("%d", counter);
 counter++;
 printf("%d", counter);
 pthread_mutex_unlock(&m);
                                           int main() {
```

```
pthread_t tid[2];
int main() {
                                           int i;
 pthread_t tid[2];
                                           for (i = 0; i < 2; i++) {
 int i;
                                            if (fork() == 0) {
 pthread_mutex_t m;
                                             foo();
 ret=pthread_mutex_init(&m, NULL);
                                            }
 for (i = 0; i < 2; i++) {
  pthread_create(&tid[i], 0, foo, 0);
                                           counter++;
                                           printf("%d", counter);
 counter++;
 printf("%d", counter);
 ret=pthread_mutex_destroy(&m);
```

Answer For Code1:

Answer For Code2:

Answer For Code3:

Answer For Code4:

Q2 Synchronization with Concurrent Programming

a. Consider the thread_fn function as shown below. Assume that two threads are allowed to call the function, which is controlling the execution of the critical section part of the code using a shared variable. Do you see any problem with this code? Briefly explain the reason. If you see a problem, rewrite the code to fix it. [2 Marks] int shared=0;

```
void *thread_fn(void *arg) {
  if (shared == 0)
    shared = 1;
  else return;
  // critical section code goes here
  shared = 0;
}
```

Answer:

b. For each situation, state the one primitive that, when used correctly around the relevant critical section, prevents race conditions and results in the most concurrency. When more than one primitive will work with equal concurrency, give the primitive that is simplest, as defined below. If your answer is a semaphore, you must specify its initial value.

Your response to each question will be exactly one of the following primitives, listed here in order from simplest to most complex: none needed, mutex, semaphore(n), rwlock.

You may assume that all relevant information has been given to you. For example, if it is not explicitly stated that a thread writes to a variable, then there is no possible race condition involving writes. No additional logic or variables may be added to the programs; you are only wrapping critical sections with concurrency primitives.

Consider the following situations:

[5 Marks]

	onsider the following situations.	[5 Warks]
Sit	tuation:	Answer:
Α.	Two threads read from a global	
	variable.	
B.	Two threads increment a global	
	counter.	
C.	Two hundred threads search through a	
	regular linked list of integers.	
D.	Two hundred threads search through a	
	regular linked list of integers; one	
	thread occasionally removes and frees	
	nodes from the list.	
E.	At most seven threads may be within	
	the critical section simultaneously.	
F.	One thread waits, blocked, for events	
	that may occur at any time and are	
	inserted into a queue when they do	
	occur; it is unacceptable for any event	
	to be missed.	
G.	The operating system maintains the	
	process table which can be read by	
	several threads but only the main	
	kernel thread is allowed to create a	
	new process or remove the finished	
.	process in the process table.	
н.	Your system has four USB ports that	
	need to be shared by a maximum of	
	four threads.	
I.	In the producer-consumer problem, a	
	producer generated an item and places	
	it in the circular queue if the queue is not full, otherwise, it waits for the	
	queue to have at least one free space.	
	Assume that queue can have a	
	maximum of 100 items.	
J.	In the producer-consumer problem, a	
0.	consumer consumes an item from a	
	·	
	•	
	•	
	consumer consumes an item from a circular queue if the queue has at least one item, otherwise, it waits for the queue to have at least one item. Assume that queue can have a maximum of 100 items.	

c. Let us consider a 1GB of computer memory is divided into 1024 blocks, each with 1MB size. You are writing an operating system task that will track the occupied and free blocks. The two functions occupy_block is called when a process requests to occupy the available (free) memory block to the kernel and free_up_block is called when a process wants to free up the already occupied block. Both occupy_block and free_up_block must run atomically to ensure no two process requests are handled at the same time. Assume that initially, all the blocks are free. Partial code is provided; please fill in the required code in the space provided. [10 Marks]

```
#include <pthread.h>
#include <semaphore.h>
int blocks[1024] = \{0\};
int read_index=0, write_index=0;
sem_t occupied, free;
pthread mutex t mutex;
int main() {
 pthread_t tid;
 pthread_create(&tid, NULL, free_up_block, NULL);
 pthread create(&tid, NULL, occupy block, NULL);
return;
/* occupy_block thread */
void* occupy block(void* start block addr) {
 while(1){
```

```
block[write_index] = (int *)start_block_addr;
  write_index = (write_index + 1) % 1024;
 sleep(rand()%5); /* wait for up to 5 sec */
 return NULL;
/* free_up_block thread */
void * free_up_block(void* vargp) {
 while(1){
 int *start_block_addr = block[read_index];
  read_index = (read_index + 1) % 1024;
 sleep(rand()%5); /* wait for up to 5 sec */
  }
      return (void *)start_block_addr;
}
```

Q3 GCC Compilation and Makefiles

a. We have the following content in the makefile where myapp.c uses functions defined in abc.c and def.c from the dynamic library. You need to find all the errors in the makefile. Write the new makefile with all corrections.
 [5 Marks]

CC=gcc %.o: %.c \$(CC) -o \$@ \$^ libmylib.so: abc.o def.o ar rs \$@ \$< myapp.out: myapp.o \$(CC) -o \$@ \$^

Answer:

b. Let us say you are building libraries for an application used in a time-sensitive (real-time) environment. Which type of libraries, static or dynamic, will you provide? Why?
 [2 Marks]

Answer:

c. Match tool (1,2,3,4) on the left side to its functionality (A,B,C,D) on the right from the table below. **[2 Marks]**

1. ldd	A. Links multiple object and library files to generate the final executable
2. ld	B. Converts source code to assembly code
3. cc	C. Converts assembly code to object code
4. as	D. Used to check unreachable libraries

Answer:

Q4 Socket Programming

a. Let's assume that we have a server computer with 2 ethernet connections on 2 different network interface cards, therefore each connection having different IPv4 addresses (191.168.1.1, 191.168.1.2). We want to have 4 instances of server code running that will use two IP addresses each with 15000, 15001 port numbers. Each of the 4 server instances should support a maximum of 500 client connections. In other words, server 1 will run on 191.16.1.1:15000, server 2 on 191.16.1.1:15001, server 3 on 191.16.1.2:15000 and server 4 on 191.16.1.2:15001 each accepting maximum of 500 clients' connection. Write the server code using the template code provided below for this implementation. [5 Marks]

main()	
{ getaddrinfo(
&hints, &listp);	,
	= socket(listp->ai_family, listp->ai_socktype, listp-
>ai_protocol);	
bind(, listp->ai_addr, listp->ai_addrlen);
listen();
}	

Answer: (Additional space is available on the next page)

b. Consider the server-side code of network communication using socket between server and client. Assume that there are multiple clients requesting for connection at a time. Assuming all system calls succeed and therefore the error handling code is never executed.
 [5 Marks]

```
#define BUFF SIZE 512
#define SERVER_PORT 15213
char buffer[BUFF_SIZE];
int main(){
  int server_sock, recvSize;
  struct sockaddr in serverAddr, clientAddr;
  /*ignore the SIGPIPE signal*/
  signal(SIGPIPE,SIG_IGN);
  /*open server socket */
  if((server_sock = socket(AF_INET,SOCK_STREAM,IPPROTO_TCP))<0){
         exit(-1);
  serverAddr.sin_addr.s_addr = htonl(INADDR_ANY);
  serverAddr.sin_port = htons(SERVER_PORT);
  serverAddr.sin_family = AF_INET;
  if(bind(server_sock,(struct sockaddr *)&serverAddr,sizeof(struct sockaddr)<0)){
    /*handle bind failing*/
             exit(-1);
  if(listen(server_sock,15)<0){
    /*handle listen failing*/
             exit(-1);
  }
  while(1){
    int client socket:
    size t clientLen = sizeof(struct sockaddr);
    if((client_socket = accept(server_sock,(struct sockaddr
*)&clientAddr,&clientLen))<0){
       /*handle failing of accept*/
```

```
continue;
}
do{
    if((recvSize = recv(socket,buffer,BUFF_SIZE,0))<0){
        break;
        }
    if(send(socket,buffer,recvSize,0)<0){
        break;
        }
}while(recvSize >0);
/*once the code reaches this point, we have received 0 bytes from the recv* call*/close(socket);
```

There are two bugs in this code. Please locate the 2 logic bugs in this code and describe them. A logic bug is one where the programmer misunderstood how their program will execute and will produce unwanted behavior under certain input conditions. Show the code changes required to fix these bugs.

Answer:

Q5 File IO

Assume that the code below is executed as a program with both files having some data. For each row of the code section, provide details of changes in Process Open File Tables and System-Wide Open File Table. Changes to both tables are shown as an example when line 3 is executed to open file1.txt and return fd1. Please make sure to copy the required information from the previous row to the next row before you update. The last row of tables must have complete information. [10 Marks]

Line	Code	Process Open File Tables	System-Wide Open File Table	
1	char	Index SysFD Ptr	SysFD Offset RefCnt inode Ptr	
	wrt_msg[10]="abcdefghij"	1 fd1=10	10 0 1 1000	
2	char read_msg[10];	2		
3	int fd1 = open("file1.txt",	3		
	"O_RDWR");	4		
4	int fd2 = open("file2.txt",			
	"O_RDWR");			
5	write(fd1, wrt_msg,	Index SysFD Ptr	SysFD Offset RefCnt inode Ptr	
6	strlen(wrt_msg));	1		
7	lseek(fd1, -2, SEEK_END);	2		
8	read(fd2, wrt_msg, 2);	3		
9	write(fd2, read_msg, 3);	4		
	(, = 3, ,,			
10	if $(fork() > 0)$	Parent Process	SysFD Offset RefCnt inode Ptr	
11	{ `	Index SysFD Ptr		
12	int $fd3 = dup(fd1)$;	1		
13	read(fd3, read_msg, 5);	2		
14	int fd4 = dup(fd2);	3		
15	}	4		
16	else	Child Process	SysFD Offset RefCnt inode Ptr	
17	{	Index SysFD Ptr	Oysi D Onset Reform model to	
18	int fd4=open("file1.txt",	1		
10	"O RDWR");	2		
19	read(fd4, read_msg, 5);	3		
20	}	4		
20	•			
	Show Content of Final	Child Process	SysFD Offset RefCnt inode Ptr	
	Tables	Index SysFD Ptr		
		1		
		2		
		3		
		4		
		Parent Process		
		Index SysFD Ptr		
		1		
		2		
		3		
		4		

Q6 Process Management, Signals

Consider the following two different C code. Assume all functions return without error, no signals are sent from other processes, and printf is atomic. [5 Marks]

```
// Code1:
int main() {
    int pid = fork();
    if(pid > 0) {
        kill(pid, SIGKILL);
        // Code2:
    int a = 1;
    void handler(int sig) {
        a = 0;
        }
```

```
printf("a");
                                                void emptyhandler(int sig) {
 } else {
  /* getppid() returns the pid of the parent
                                                int main() {
process */
                                                 signal(SIGINT, handler);
  kill(getppid(), SIGKILL);
                                                 signal(SIGCONT, emptyhandler);
  printf("b");
                                                 int pid = fork();
                                                 if(pid == 0) {
}
                                                   while(a == 1)
                                                    pause();
                                                   printf("a");
                                                 } else {
                                                   kill(pid, SIGCONT);
                                                   printf("b");
                                                   kill(pid, SIGINT);
                                                   printf("c");
```

For each code snippet in the table below write a Y next to an outcome if it could occur, otherwise write N.

Answer:

,		
Code1 Outcome	Possible (Y/N)?	
Nothing is printed.		
"a" is printed.		
"b" is printed.		
"ab" is printed.		
"ba" is printed.		
A process does not terminate.		

Code2 Outcome	Possible (Y/N)?	
Nothing is printed.		
"ba" is printed.		
"abc" is printed.		
"bac" is printed.		
"bca" is printed.		
A process does not terminate.		

Q6 Device Driver

a. What is the purpose of using module_init() and module_exit() system calls? Write the shell command to execute module_init(myinit) but not module_exit(myexit). Please note that myinit() and module_init(myinit) are defined in mymodule.c driver code.
 [2 Mark]

Answer: (Additional space on next page is available)

b. What is the significance of file_operations structure when used as member in cdev structure as shown below? [2 Mark]
struct cdev {
 struct kobject kobj;
 struct module *owner;
 const struct file_operations *ops;
 struct list_head list;
 dev_t dev;
 unsigned int count;
};
Answer:

c. What are the two ways we can identify whether a device is a character device or a block device? [2 Mark]

Answer: