

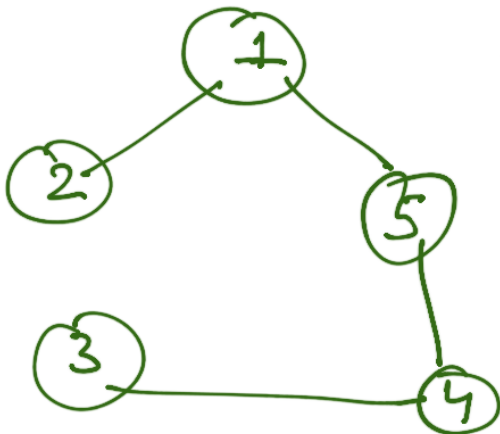
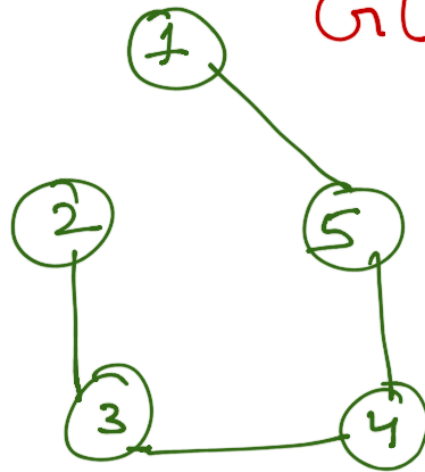
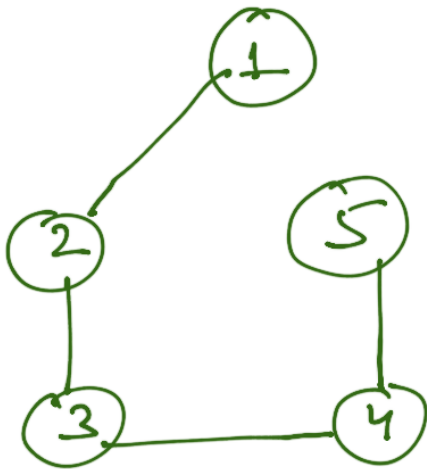
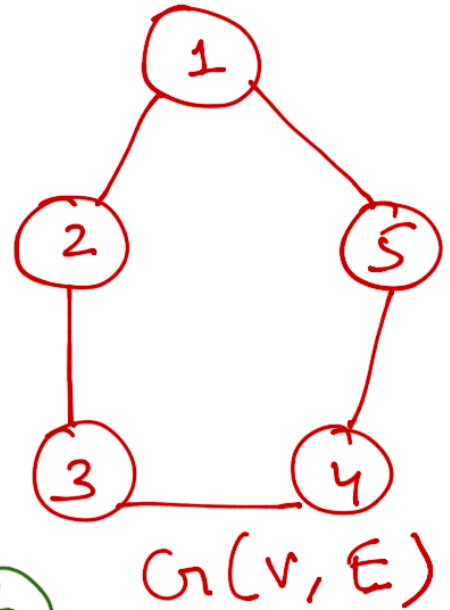
# MST

Let MST be  $G'(V', E')$

$$V' = V$$

$$E' \subset E$$

$$E' = |V| - 1$$



→ If costs is given to these edges, MST will be the one with least cost

& it should be connected No cycle in MST

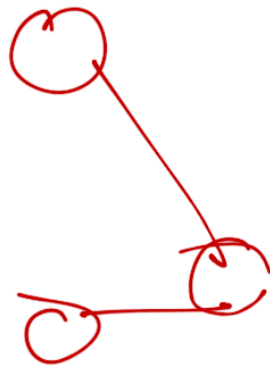
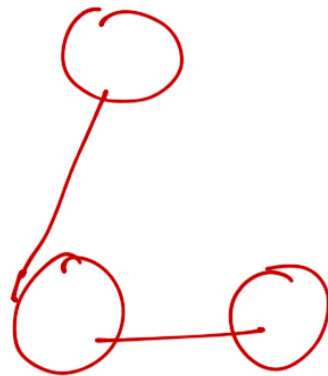
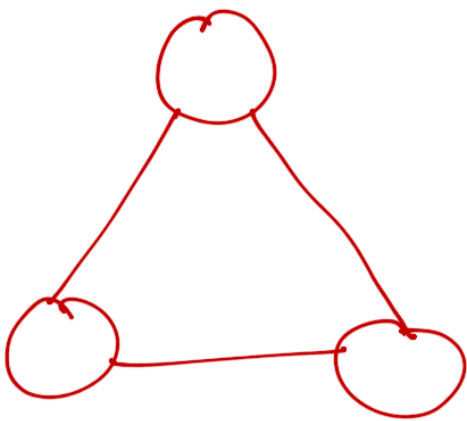
## Properties

1. Remove one edge from MST, it will become disconnected
2. Adding one edge will create a loop
3. If each edge has a distinct wt, then only one unique MST
4. A complete (each vertex conn. to every other vertex) undirected graph, can have  $n^{n-2}$  no. of MST's
5. Every connected undirected graph has at least one MST.

→ Disconnected graph has no MST

→ for a complete graph by  
removing  $\max(e - n + 1)$  edges, we  
can construct a MST

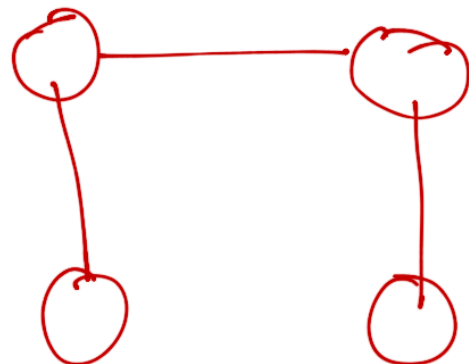
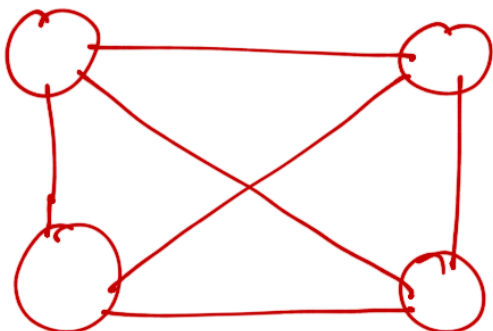
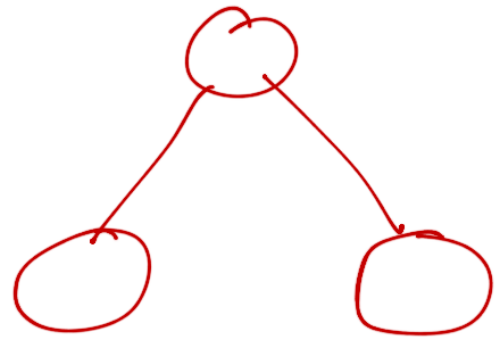
Eg:



$$\max(e - n + 1)$$

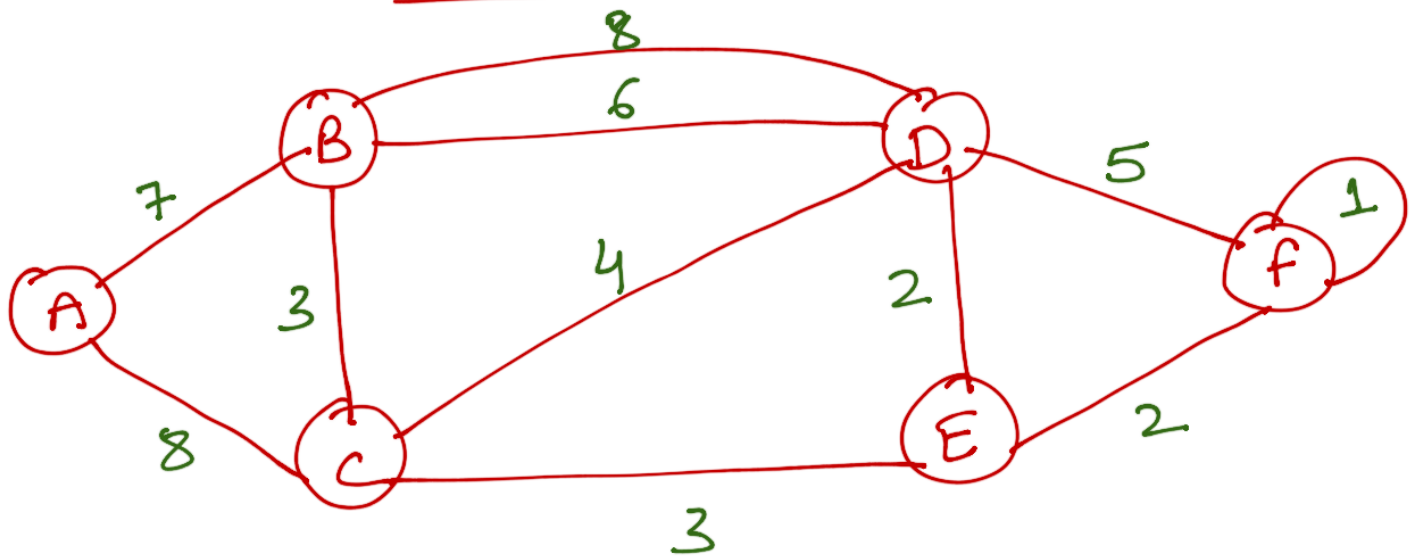
$$(\cancel{3} - \cancel{3} + 1)$$

$$= 1$$

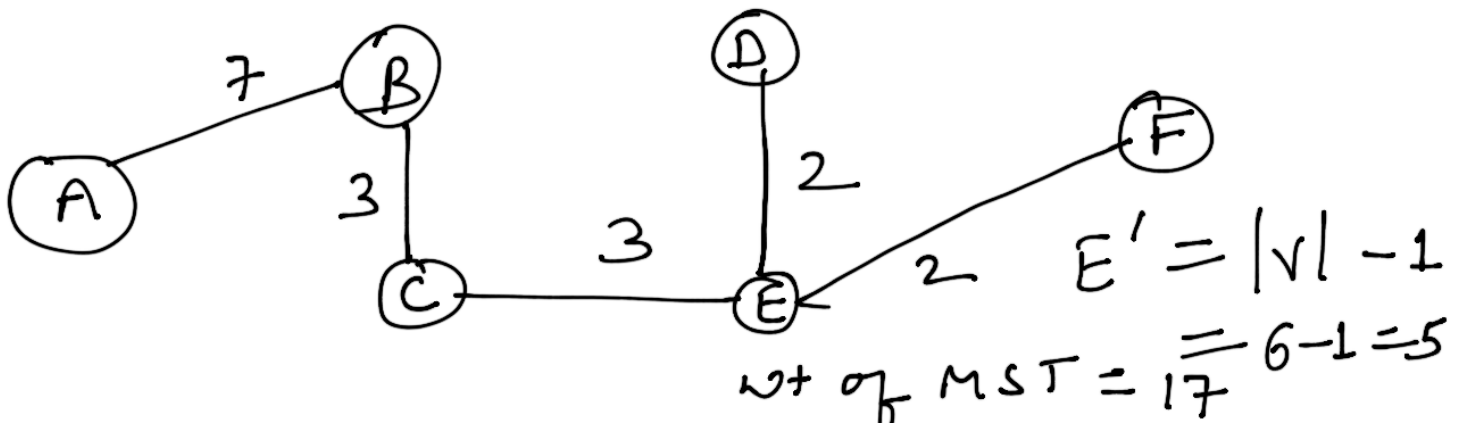


$$\max(6 - 4 + 1) = 3 \text{ edges can be removed}$$

# Prim's Algorithm

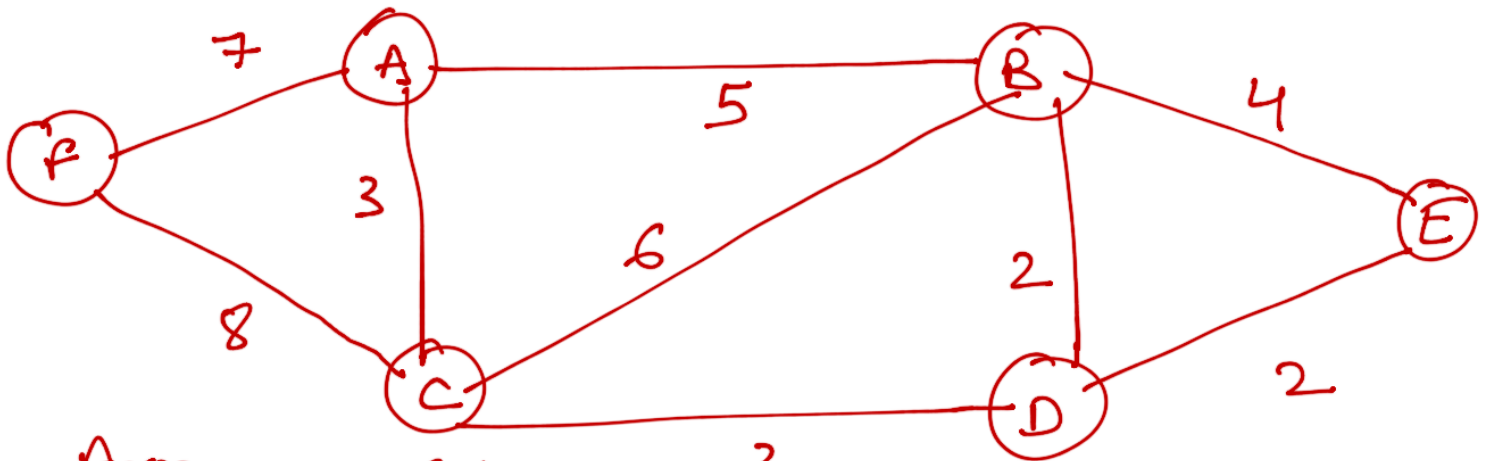
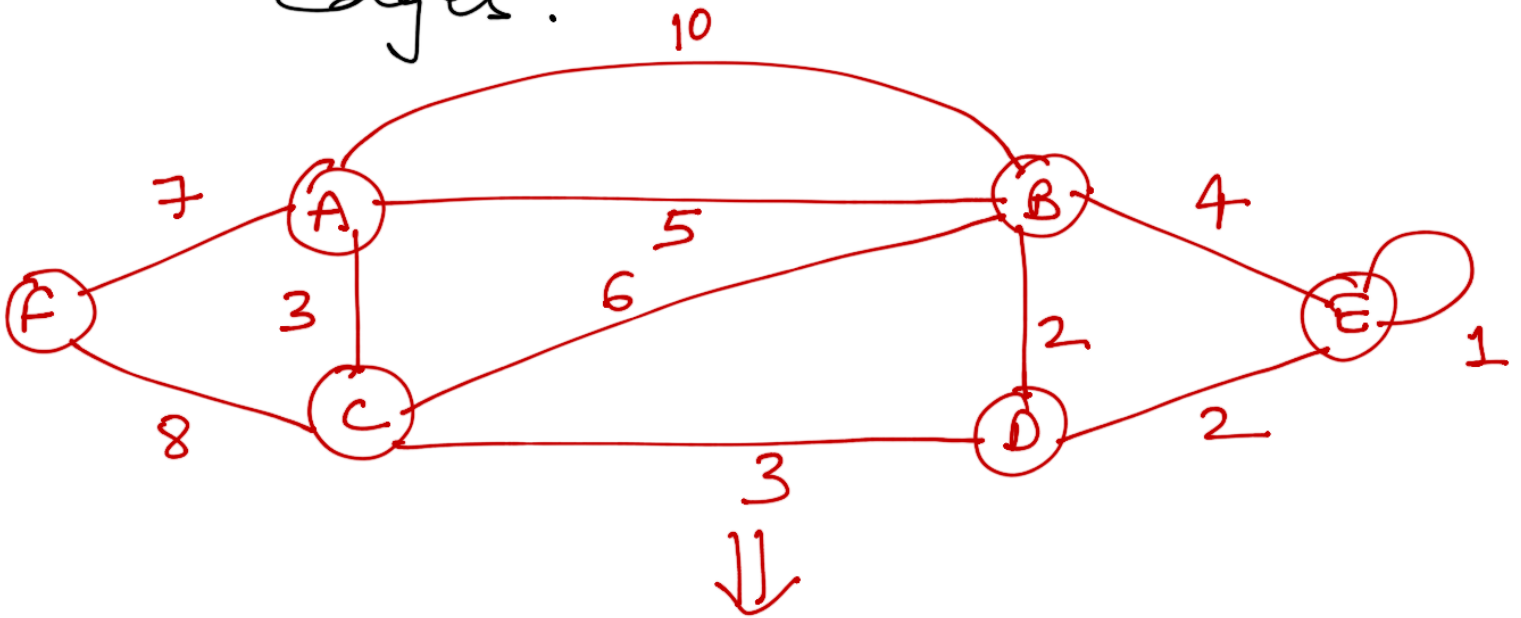


1. Remove all loops & parallel edges (based on cost, remove)
2. Choose any node as starting node,  
check incident edges, &  
choose one with minimum cost.



# Kruskal's Algorithm for MST

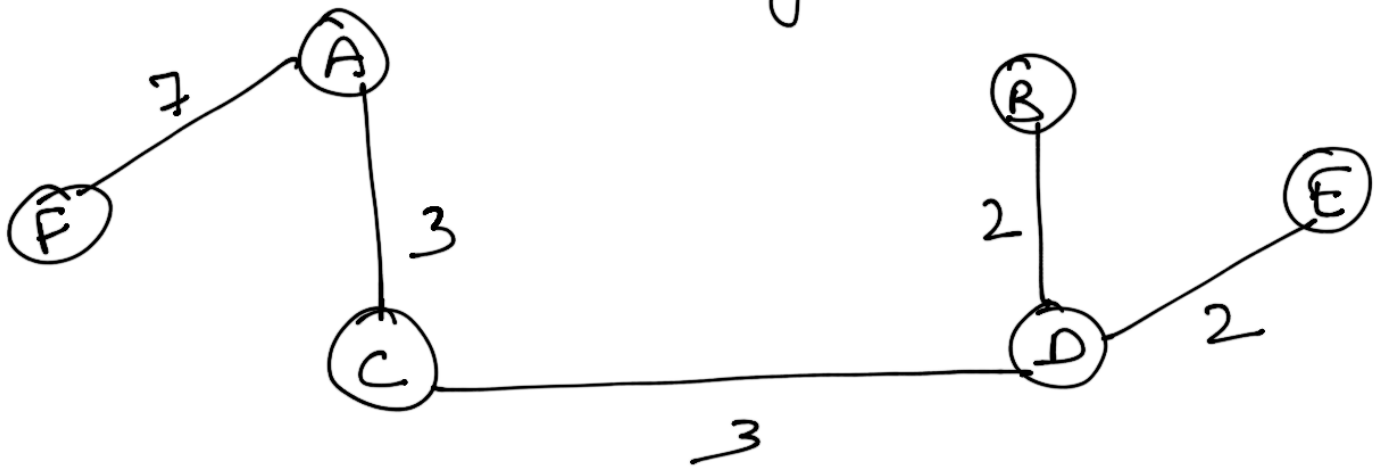
1. Remove all loops & parallel edges.



2. Arrange Edges in order of wts

BD, DE, AC, CD, BE, AB, BC, AF, FC

3. Taking minimum edges, but no cycle should be allowed



$$BD = 2$$

$$DE = 2$$

$$AC = 3$$

$$CD = 3$$

$$\cancel{X} BE = 4$$

$$\cancel{X} AB = 5$$

$$\cancel{X} BC = 6$$

$$AF = 7$$

$$\cancel{X} FC = 8$$

$$\begin{aligned} E' &= |V| - 1 \\ &= 6 - 1 \\ &= 5 \end{aligned}$$