

Application of Erasure Codes

Extra Lecture Notes

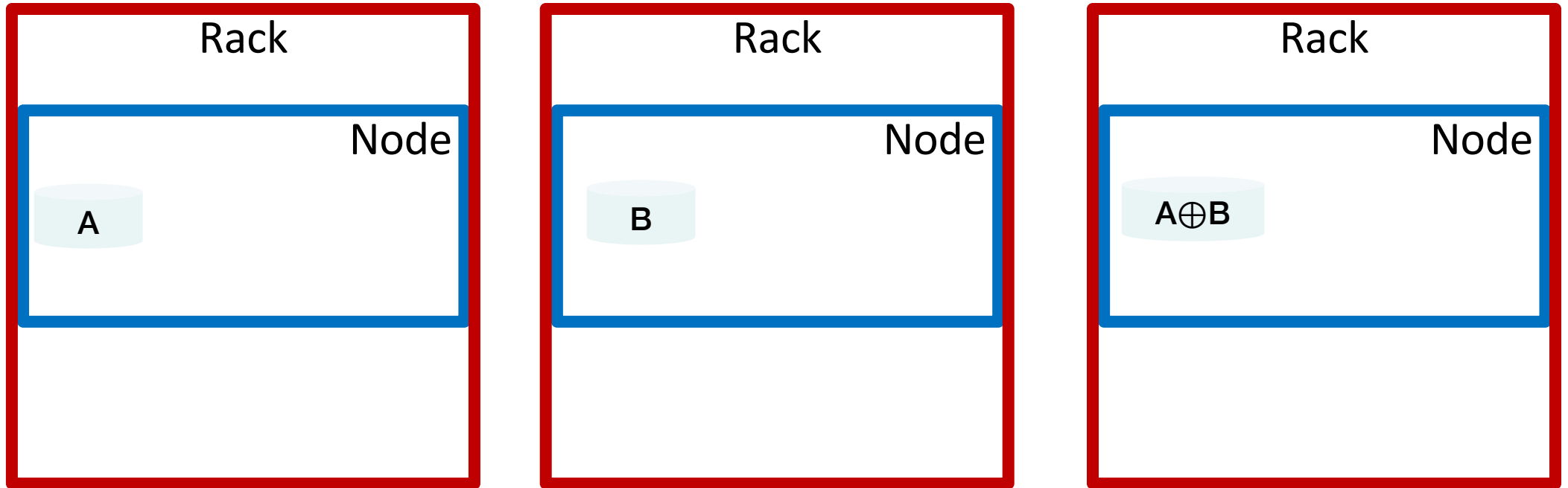
Mitigate Hardware Failure

- Replication
 - Store multiple copies across multiple machines and data centers
 - Simple to implement (unless consistency is an issue)
 - Significant large storage overhead (2 times and more the size of the entire data)
- Erasure coding
 - Store "codes data" to recover lost blocks in case of a failure.
 - For, example parity codes (🤔 we heard this before)
 - Less overhead
 - Combined with more reliable hardware, it can be a prime choice to balance cost, fault tolerance, and availability

/foo.csv

Block A
Block B

Erasure Codes using Bitwise XOR



Erasure Codes using Bitwise XOR

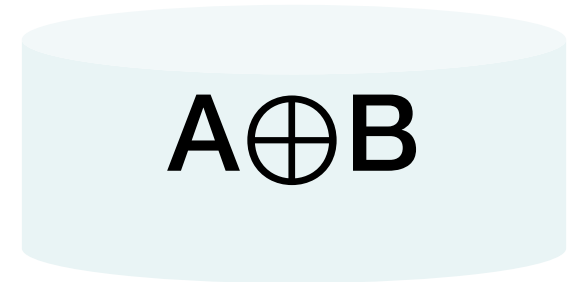
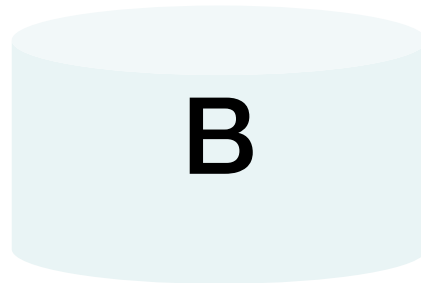
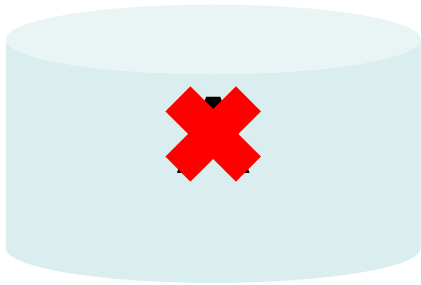


A

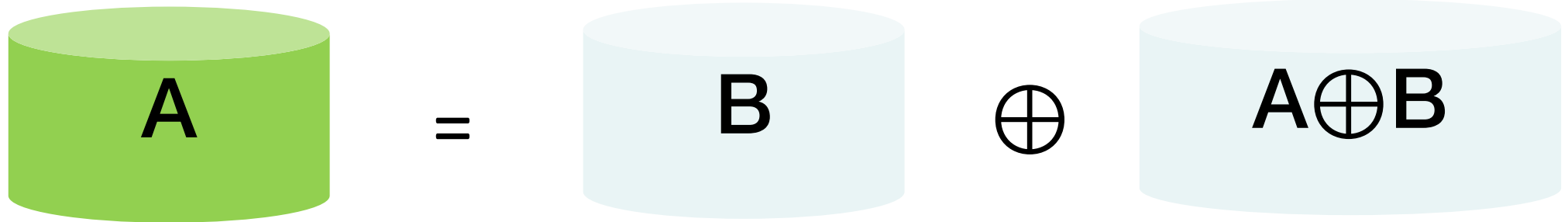
B

$A \oplus B$

Erasure Codes using Bitwise XOR



Erasure Codes using Bitwise XOR



Reed-Solomon

- Given a specific setup:
 - N : the number of data blocks for a given file
 - K : the number coding blocks
 - M = N+K total blocks
- Characteristics:
 - Recover **any 1** block from **any N** other blocks
 - Can withstand **up to K** concurrent failures
 - Applicable for any values of N and K
- Notation: RS(N,K)
 - (6,3): 6 data blocks, 3 coding blocks
- Technique:
 - Galois Field arithmetic. $GF(2^w)$

Reed-Solomon (4,2)

A

B

C

D

1

2

Reed-Solomon (4,2)



B

C

D

1

2

Reed-Solomon (4,2)



$$A = B + C + D + 1$$

Reed-Solomon (4,2)



$$\boxed{A} = \boxed{B} + \boxed{C} + \boxed{D} + \boxed{1}$$

⋮

Any N Blocks

$$\boxed{A} = \boxed{C} + \boxed{D} + \boxed{1} + \boxed{2}$$