Big Data Systems

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Lecture 2 – A Primer on DBMS Storage and Indexing

Outline

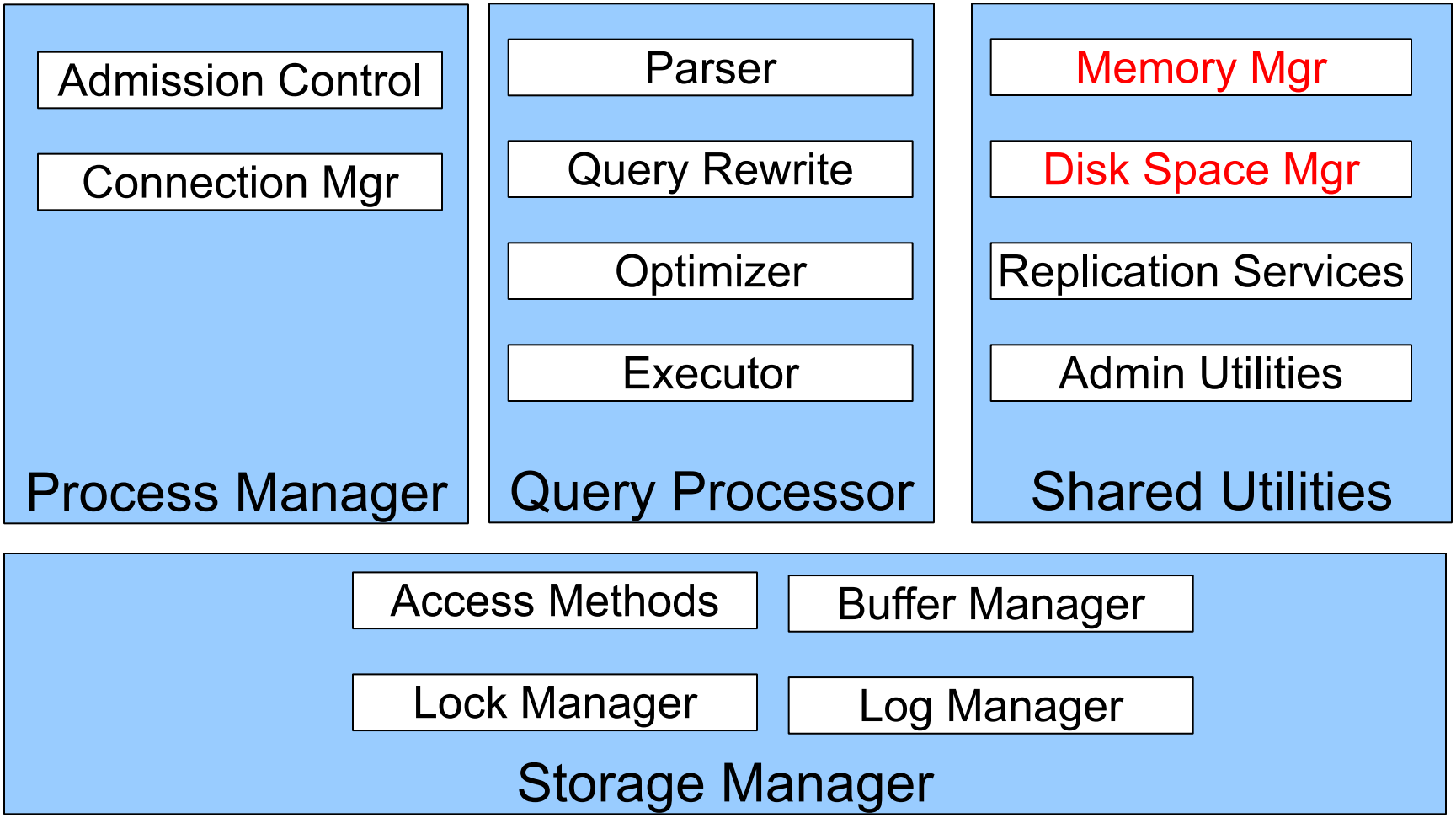
- **Data storage**
  - Disk and files
  - Operations on files
- **Indices**
  - Index structures
  - Hash-based indices
  - B+ trees

Why is this important?

- **Data storage**

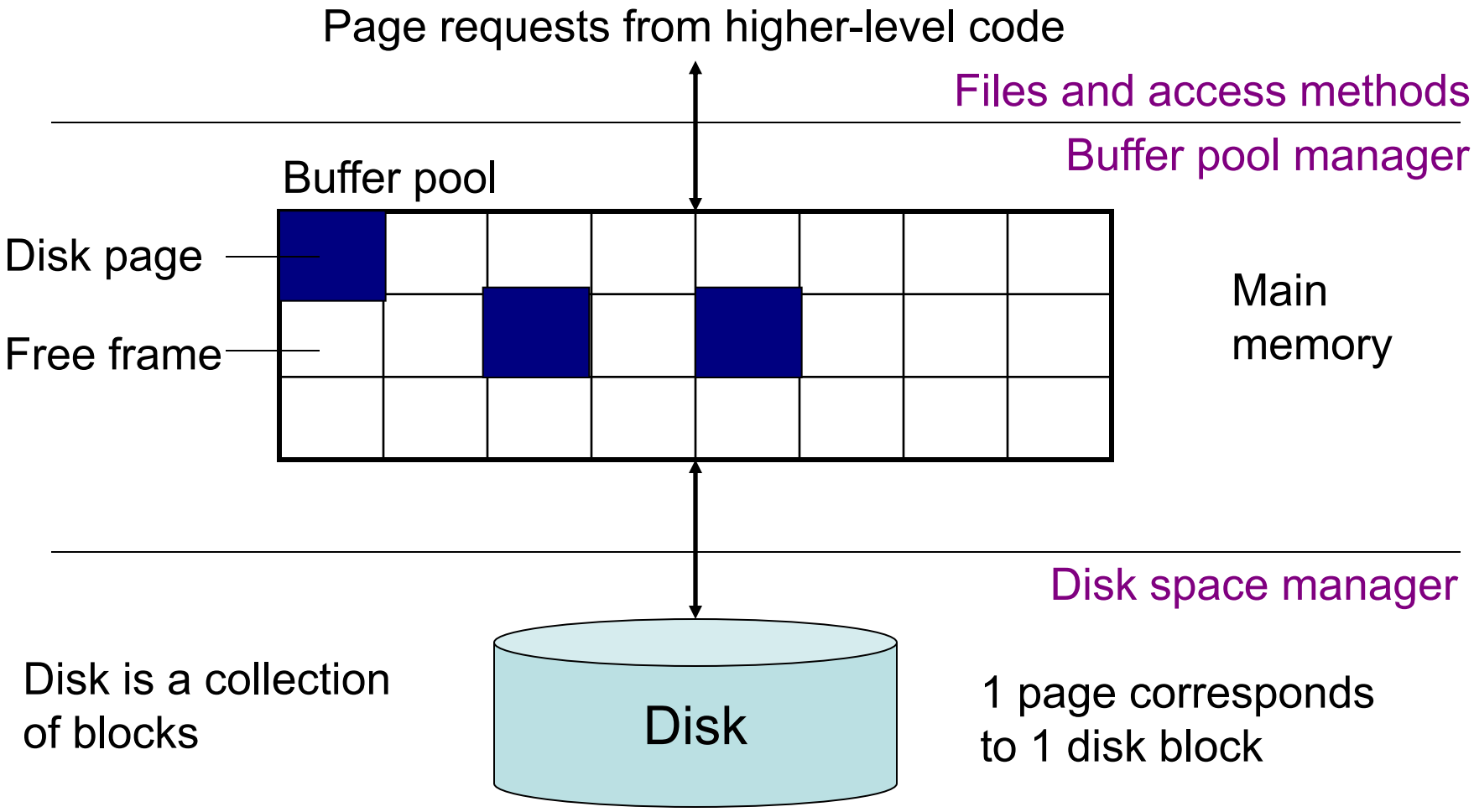
- **Indices**

DBMS Architecture



[Anatomy of a Db System. J. Hellerstein & M. Stonebraker. Red Book. 4ed.]

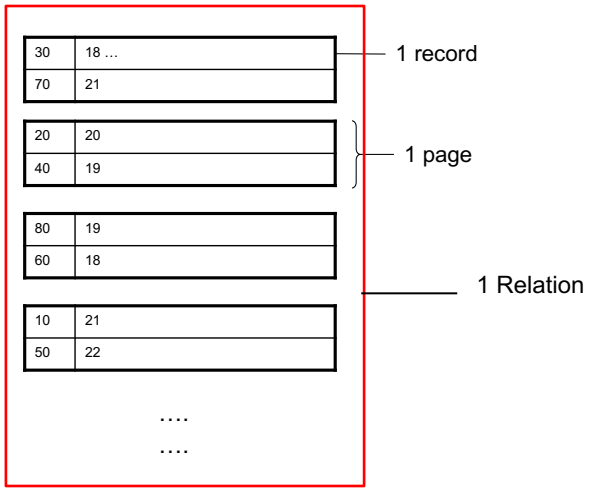
Buffer Manager



Page Formats

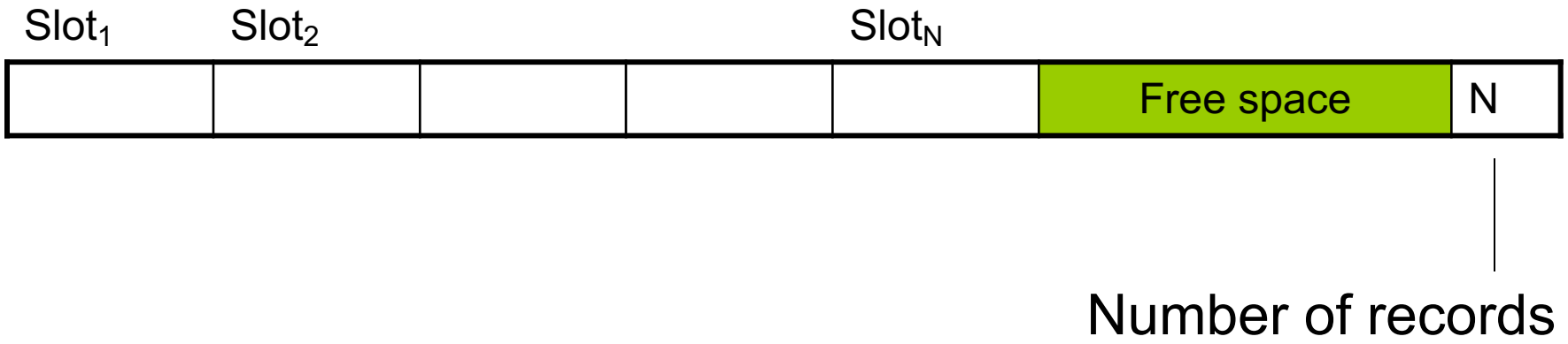
Basic abstraction

- Typically, 1 relation = 1 file
- A file consists of one or more pages
- 1 page = 1 disk block = fixed size (e.g. 4KB)
  - Disk block? The smallest unit of data that can be read from or written to a disk drive
- Records:
  - Fixed length
  - Variable length
- Record id = RID
  - Typically RID = (PageID, SlotNumber)

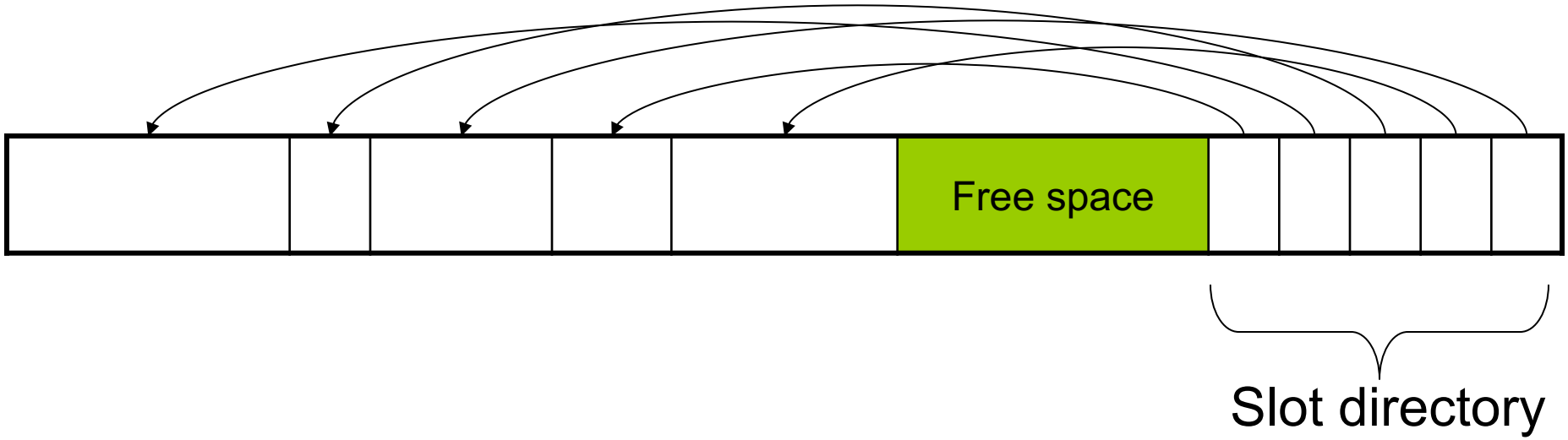


Page Format Approach 1

Fixed-length records: packed representation



Page Format Approach 2



Each slot contains
<record offset, record length>

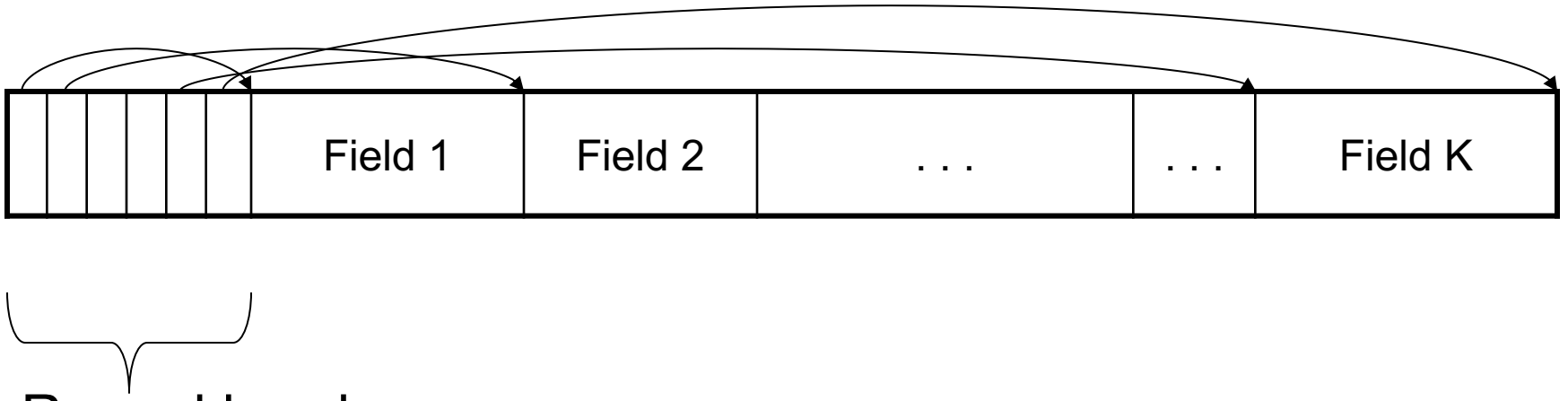
Can handle variable-length records

Record Formats

Fixed-length records $\rightarrow$ Each field has a fixed length
(i.e., it has the same length in all the records)

Record Formats

Variable length records



Remark: NULLS require no space at all (why ?)

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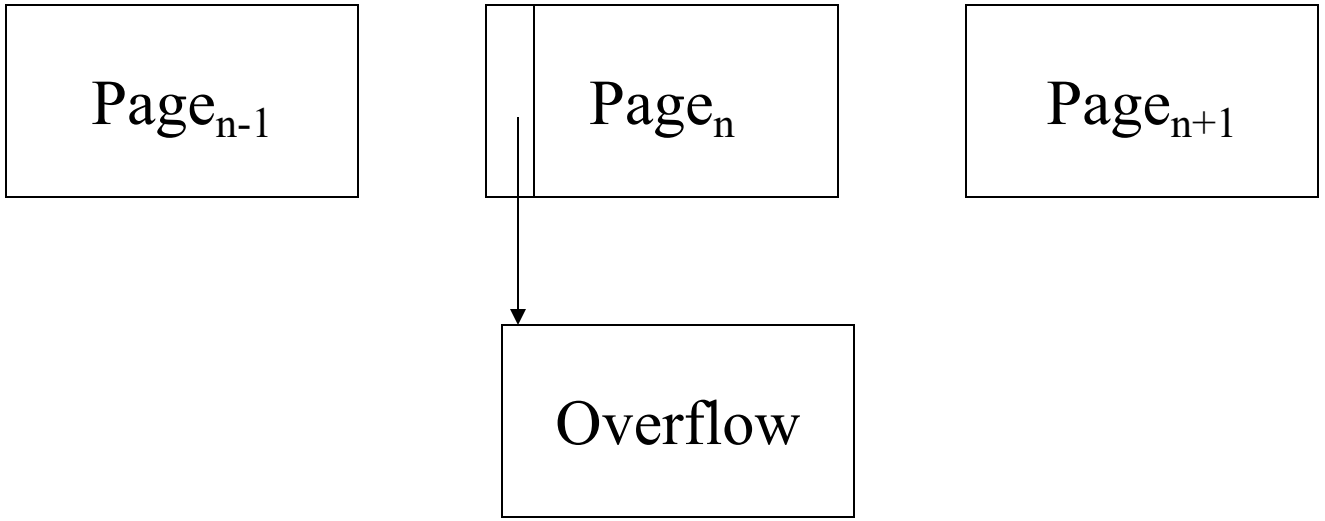
  - B+ trees

Modifications: Insertion

- File is unsorted (= ***heap file***)
  - add it wherever there is space (easy ☺)
- File is sorted (clustered index)
  - Is there space on the right page ?
    - Yes: we are lucky, store it there
  - Is there space in a neighboring page ?
    - Look 1-2 pages to the left/right, shift records
  - If anything else fails, create ***overflow page***

Overflow Pages



- After a while the file starts being dominated by overflow pages: time to reorganize

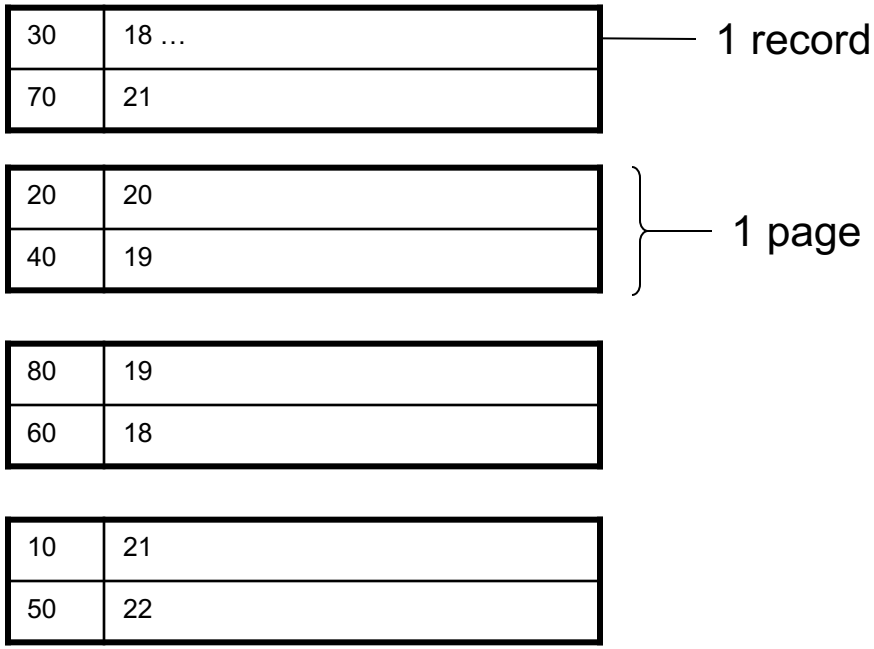
Modifications: Updates

- If new record is shorter than previous, easy ☺
- If it is longer, need to shift records
  - May have to create overflow pages

Searching in a Heap File

File is not sorted on any attribute

`Student(sid: int, age: int, …)`



Heap File Search Example

- 10,000 students
- 10 student records per page
- Total number of pages: 1,000 pages
- Find student whose sid is 80
  - Must read on average 500 pages
- Find all students older than 20
  - Must read all 1,000 pages
- Can we do better?

Sequential File

File sorted on an attribute, usually on primary key

`Student(sid: int, age: int, …)`

10	21 …
20	20

30	18
40	19

50	22
60	18

70	21
80	19

Sequential File Example

- Total number of pages: 1,000 pages
- Find student whose sid = 80
  - Could do binary search, read $\log_2(1{,}000) \approx 10$ pages
- Find all students older than 20
  - Must still read all 1,000 pages
- Can we do better?

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Indexes

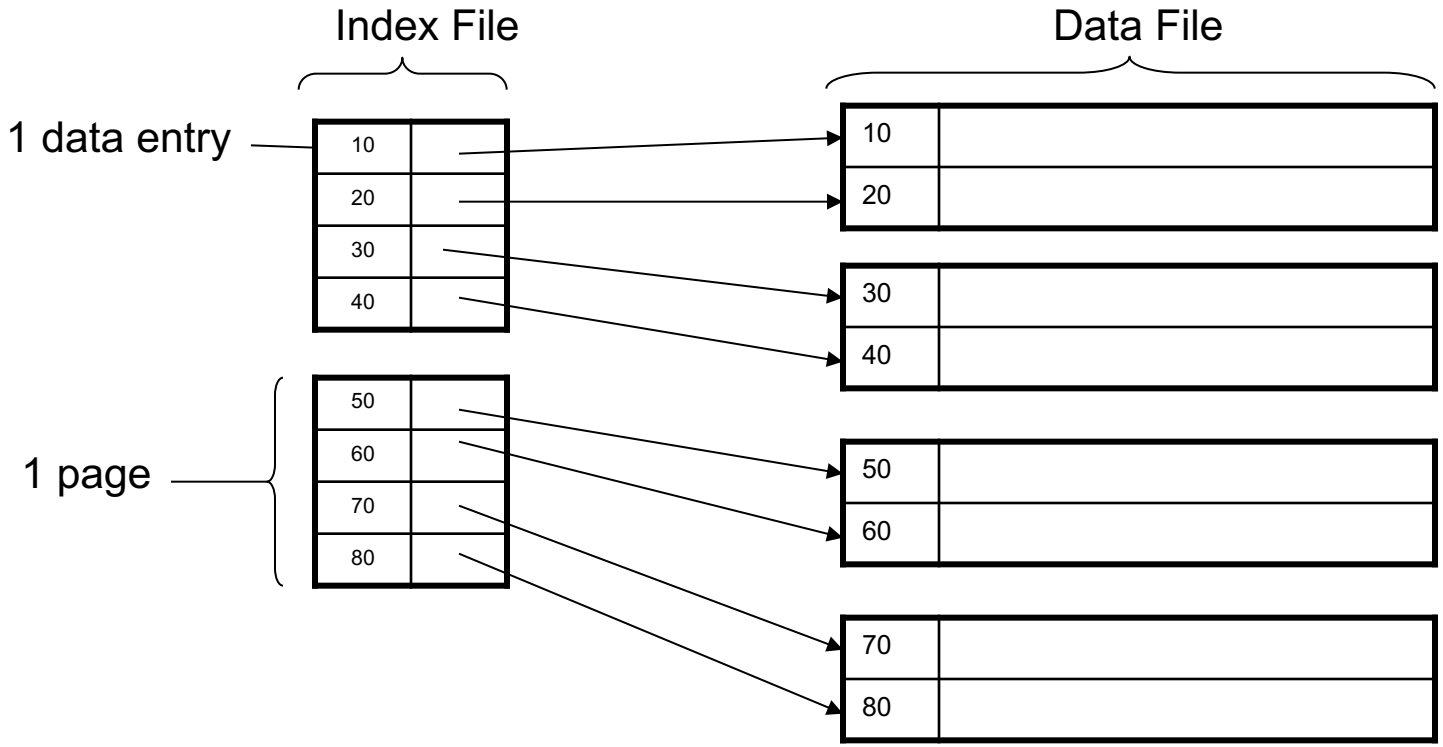
- **Index:** data structure that organizes data records on disk to optimize selections on the ***search key fields*** for the index
- An index contains a collection of ***data entries***, and supports efficient retrieval of all data entries with a given search key value **k**
- **Search key** = can be any set of fields
  - not the same as the primary key, nor a key
- **Data entry** for key k can be:
  - The actual record with key k
    - In this case, **the index is also a special file organization**
    - This type of index is also called the **primary index** of a file
  - (k, RID: Record ID)
  - (k, list-of-RIDs)

Index Classification Overview

- Primary/secondary
  - Primary = determines the location of indexed records
  - Secondary = does not directly determine data location
- Dense/sparse
  - Dense = every key in the data appears in the index (Unordered)
  - Sparse = the index contains only some keys (Ordered)
- Clustered/unclustered
  - Clustered = records close in index are close in data. The data is stored on disk in the same order as the index
  - Unclustered = records close in index may be far in data
- B+ tree / Hash table / …

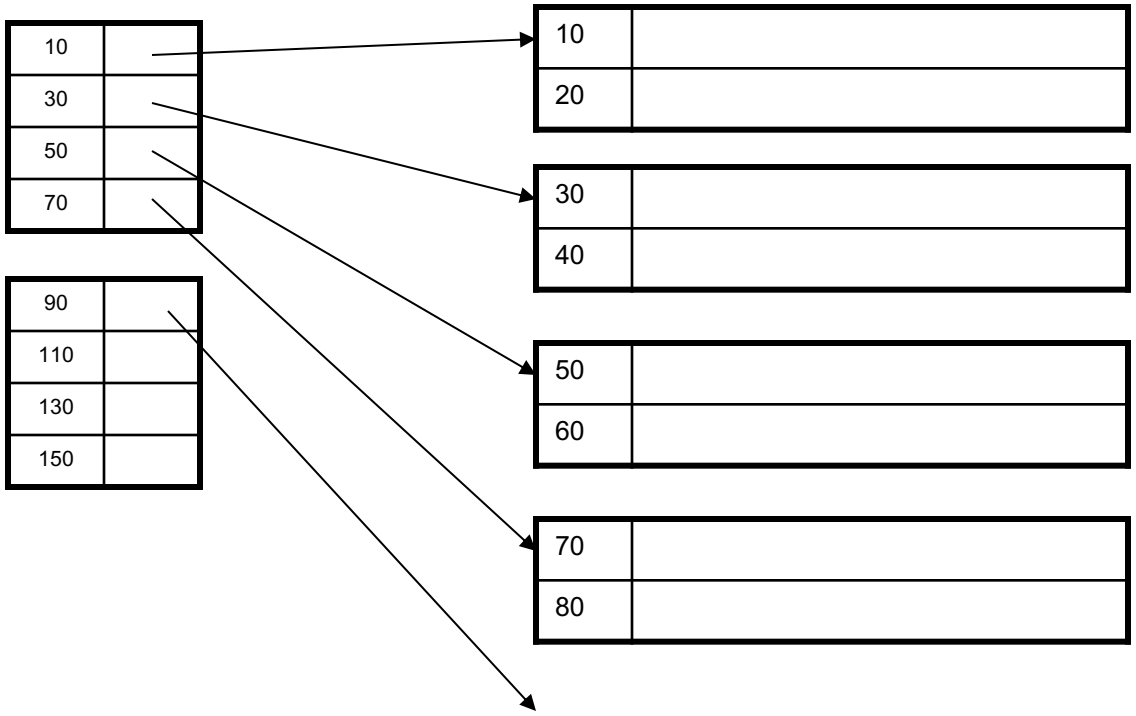
Primary Index

- Index determines the location of indexed records
- *Dense* index: sequence of (key,pointer) pairs



Primary Index

- *Sparse* index



Primary Index Example

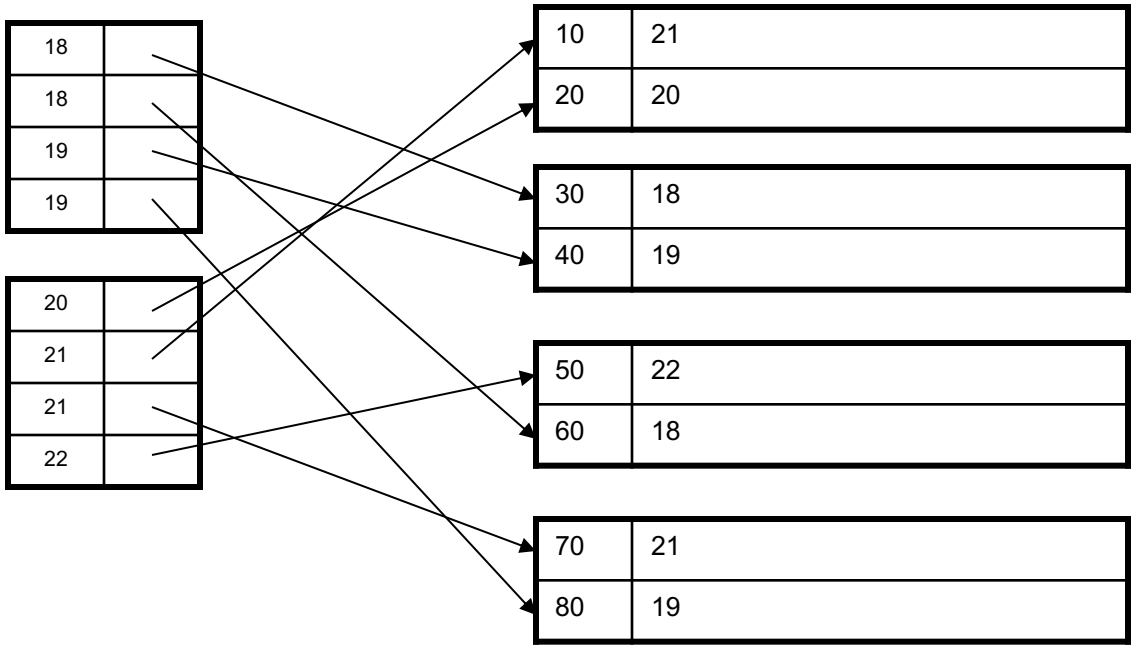
- Let's assume all pages of index fit in memory

- Find students whose sid is 80
  - Index (dense or sparse) points directly to the page
  - Only need to read 1 page from disk.
- Find all students older than 20
  - Must still read all 1,000 pages.

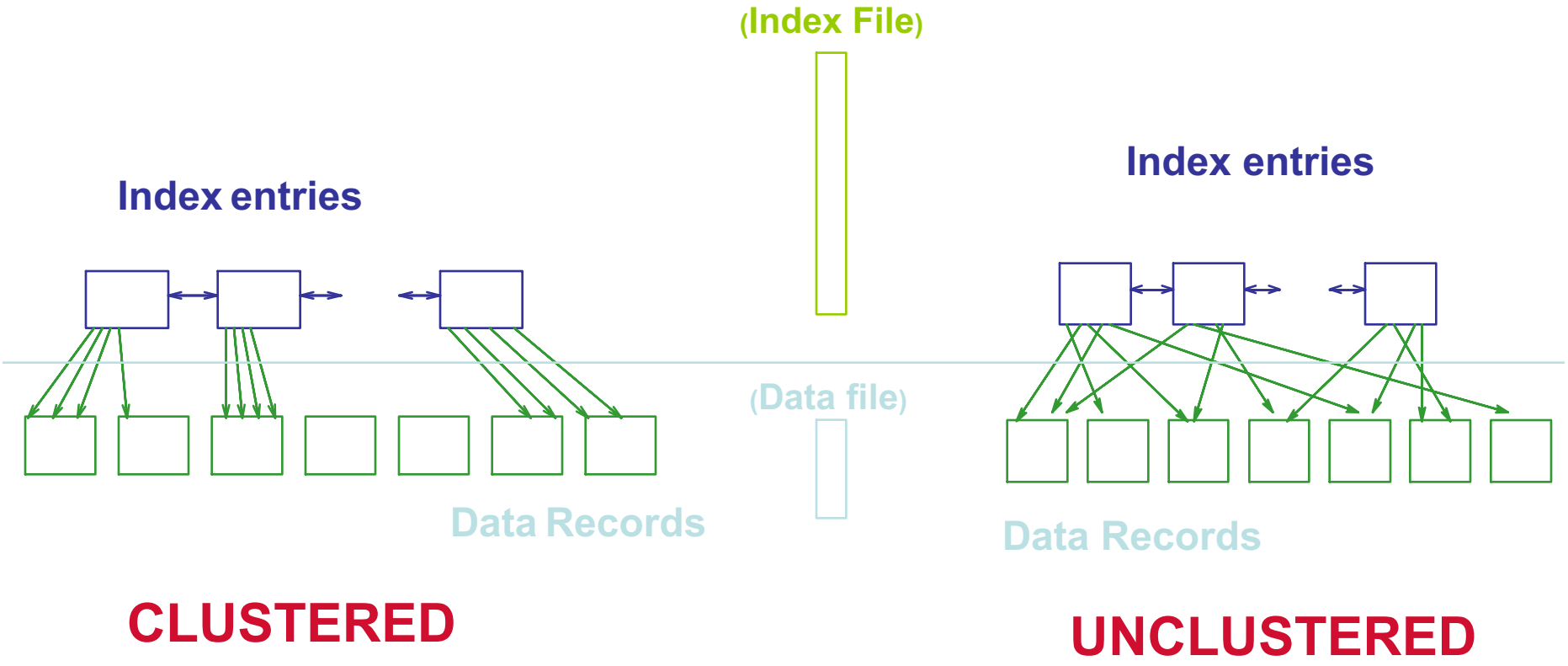
- … we handled point queries but not range.
- How can we make *both* queries fast?

Secondary Indexes

- To index other attributes than primary key
- Always dense (why not sparse?)



Clustered vs. Unclustered Index



Clustered = records close in index are close in data

Clustered/Unclustered

- Primary index = clustered by definition
- Secondary indexes = usually unclustered

Secondary Indexes

- Applications
  - Index other attributes than primary key
  - Index unsorted files (heap files)

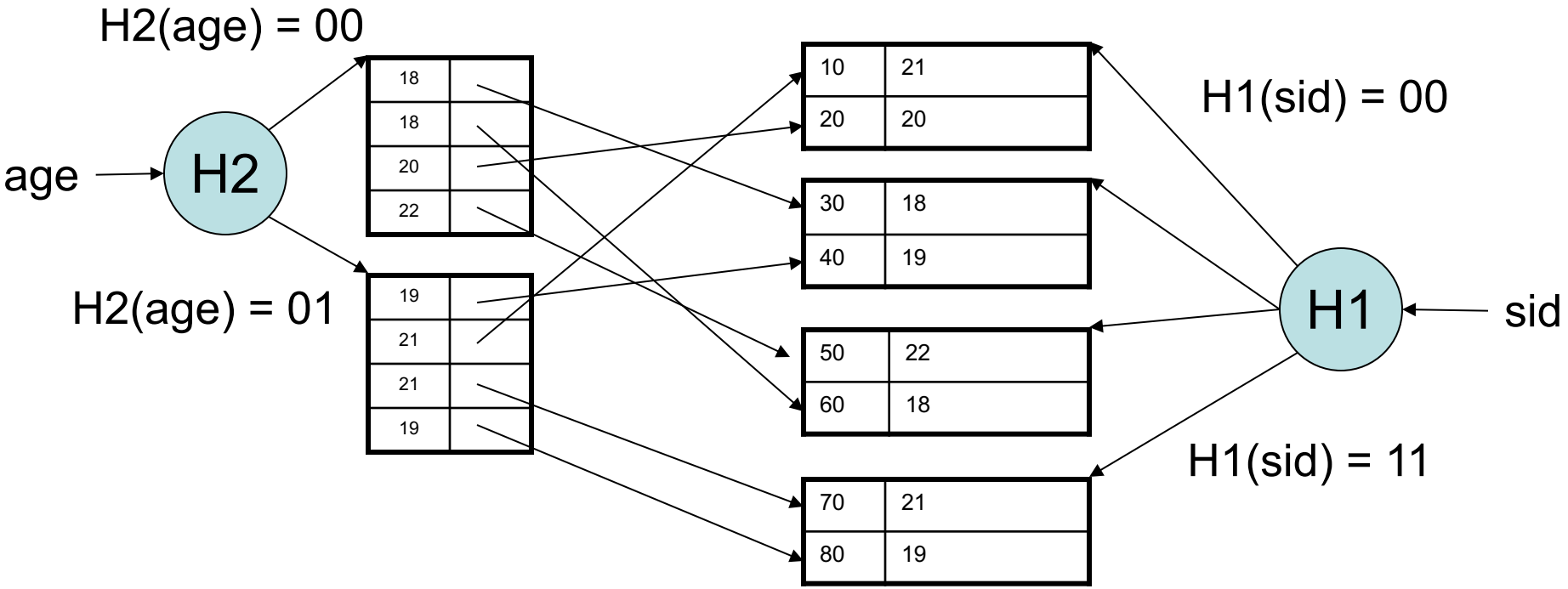
Large Indexes

- What if index does not fit in memory?
- Would like to index the index itself
  - Hash-based index
  - Tree-based index

Hash-Based Index

Another example of secondary index

Another example of primary index



Good for point queries but not range queries

Tree-Based Index

- How many index levels do we need?
- Can we create them automatically? Yes!
- Can do something even more powerful!

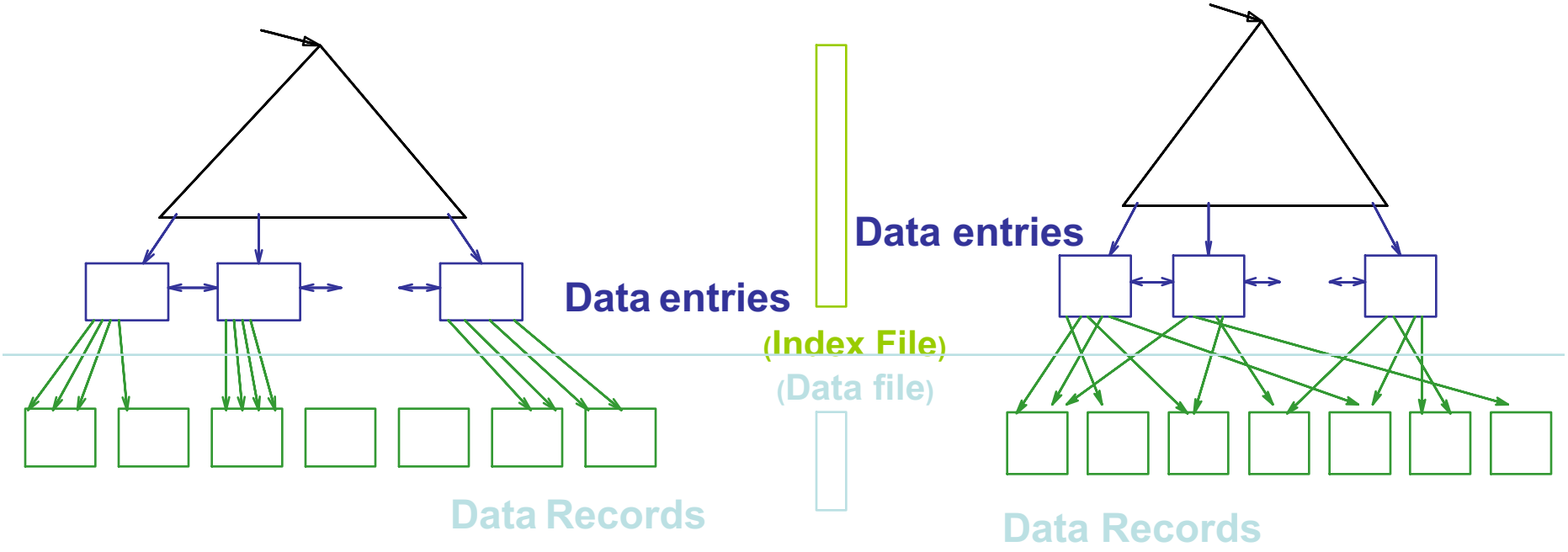
B+ Trees

- Search trees

- Idea in B Trees (Search Tree)
  - Make 1 node = 1 page (= 1 block)
  - Keep tree balanced in height

- Idea in B+ Trees
  - Make leaves into a linked list : facilitates range queries

B+ Trees



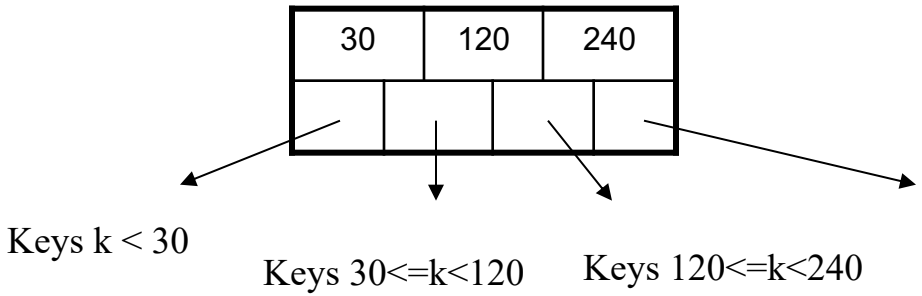
CLUSTERED

UNCLUSTERED

Note: can also store data records directly as data entries (primary index)

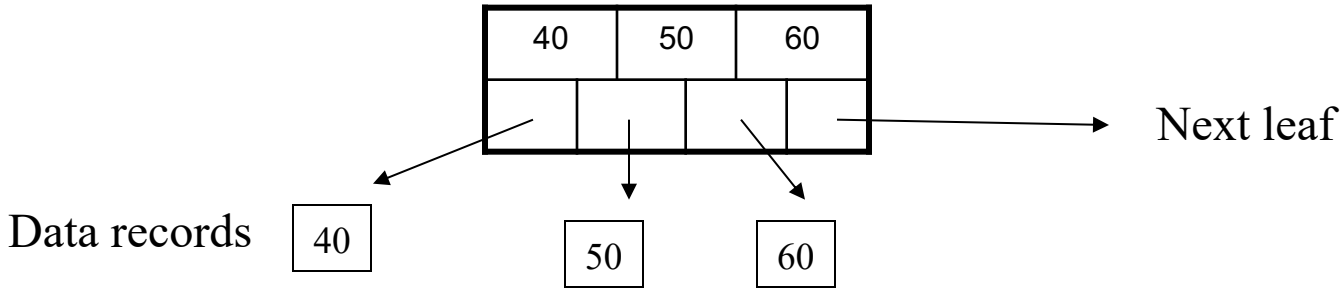
B+ Trees Basics

- Parameter d = the *degree*
- Each node has **d <= m <= 2d keys** (except root)



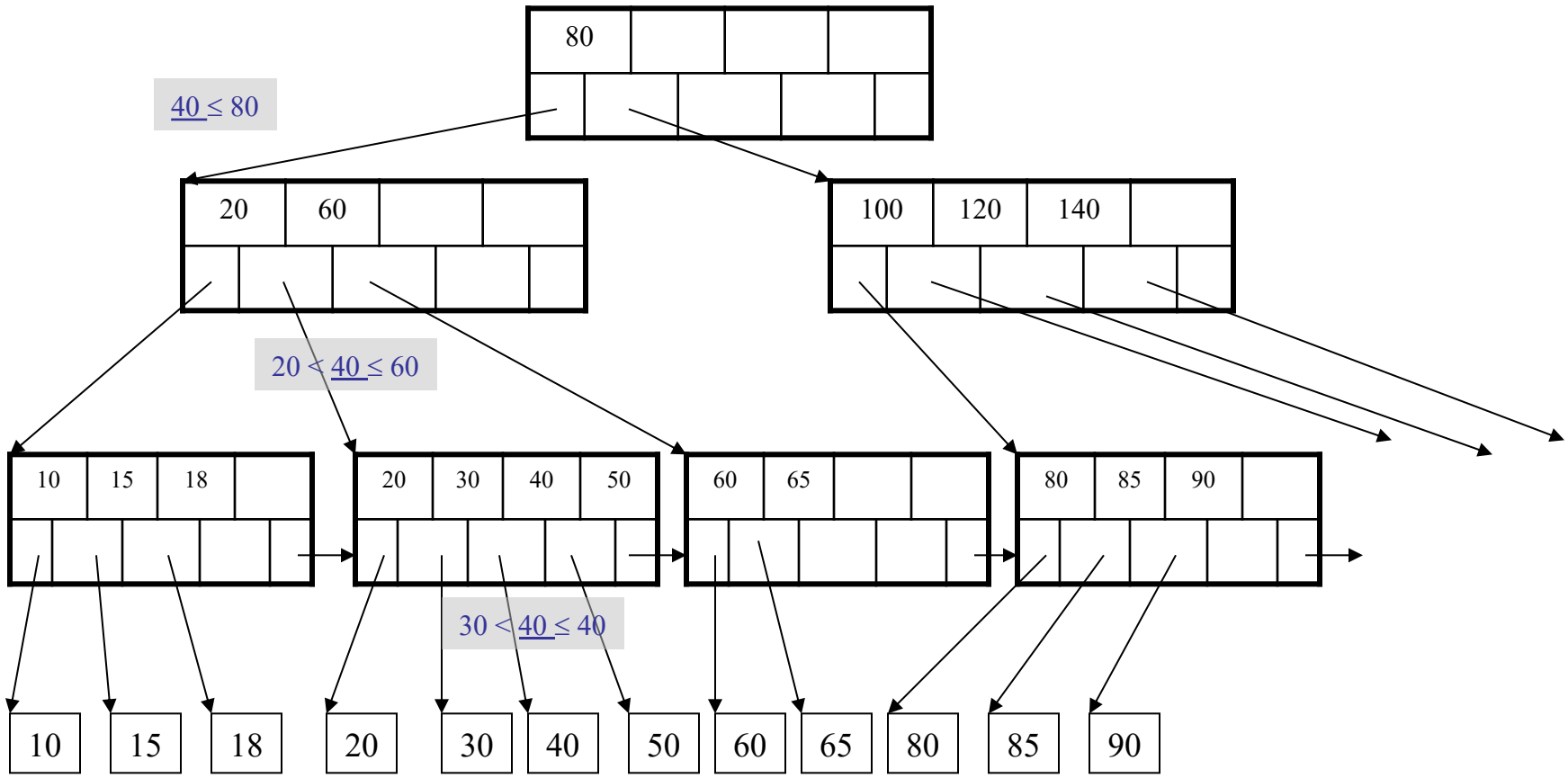
Each node also has **m+1 pointers**

- Each leaf has **d <= m <= 2d keys**:



B+ Tree Example

d = 2



Searching a B+ Tree

- Exact key values:
  - Start at the root
  - Proceed down, to the leaf
- Range queries:
  - Find lowest bound as above
  - Then sequential traversal

```
Select name
From Student
Where age = 25
```

```
Select name
From Student
Where 20 <= age
  and  age <= 30
```

B+ Tree Design

- How to choose the degree d ?
- Example:
  - Key size = 4 bytes
  - Pointer size = 8 bytes
  - Block size = 4096 bytes
- 2d x 4 + (2d+1) x 8 <= 4096
- d = 170

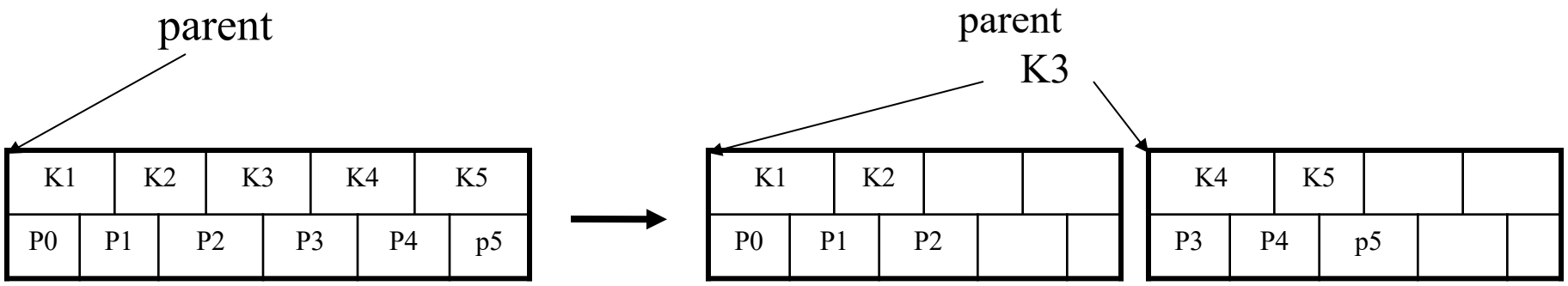
B+ Trees in Practice

- Typical order: 100. Typical fill-factor: 67%.
  - average fanout = 133
- Typical capacities
  - Height 4: 133^4 = 312,900,700 records
  - Height 3: 133^3 = 2,352,637 records
- Can often hold top levels in buffer pool
  - Level 1 = 1 page = 8 Kbytes
  - Level 2 = 133 pages = 1 Mbyte
  - Level 3 = 17,689 pages = 133 Mbytes

Insertion in a B+ Tree

Insert (K, P)

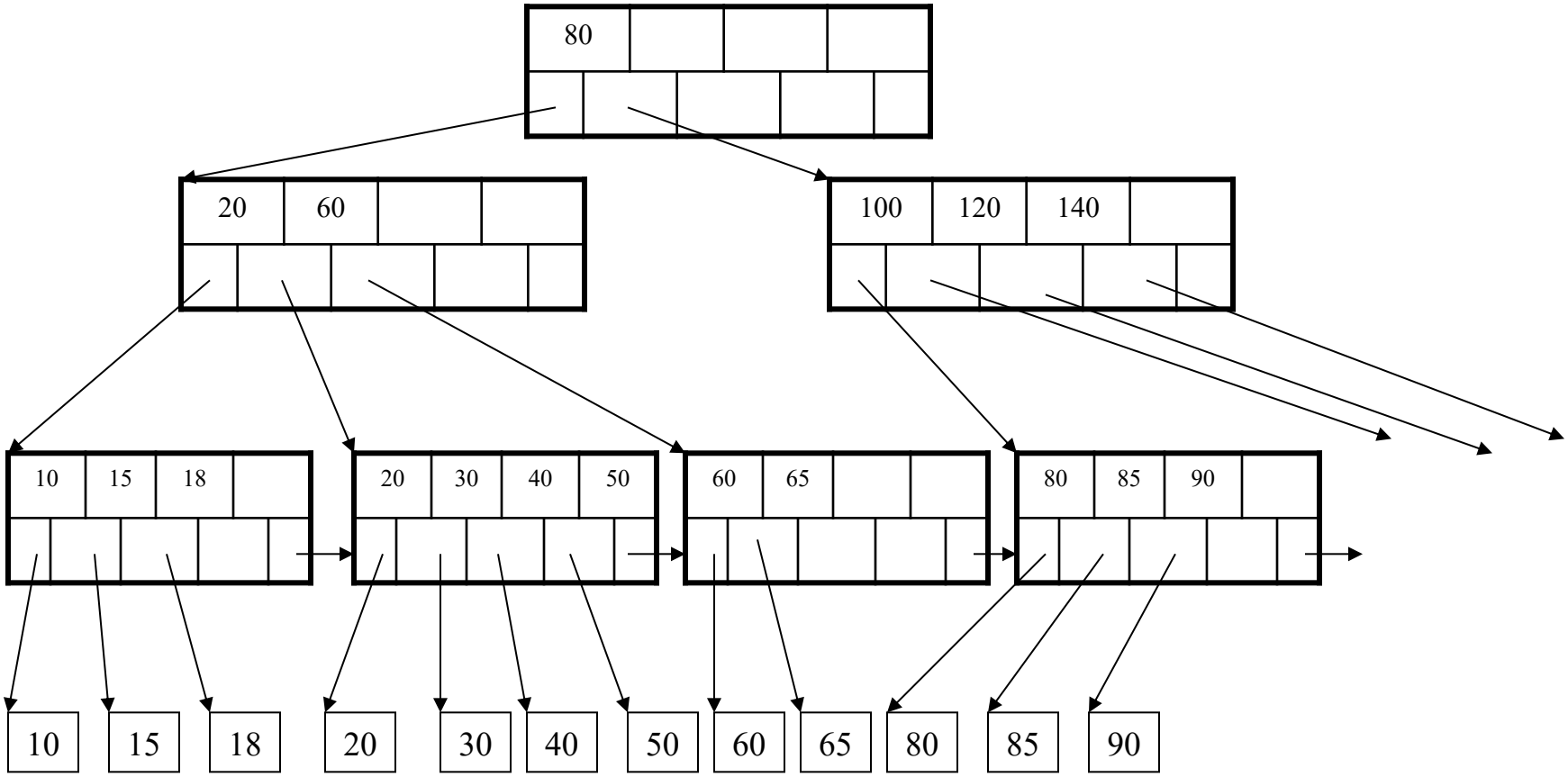
- Find leaf where K belongs, insert
- If no overflow (2d keys or less), halt
- If overflow (2d+1 keys), split node, insert in parent:



- If leaf, also keep K3 in right node
- When root splits, new root has 1 key only

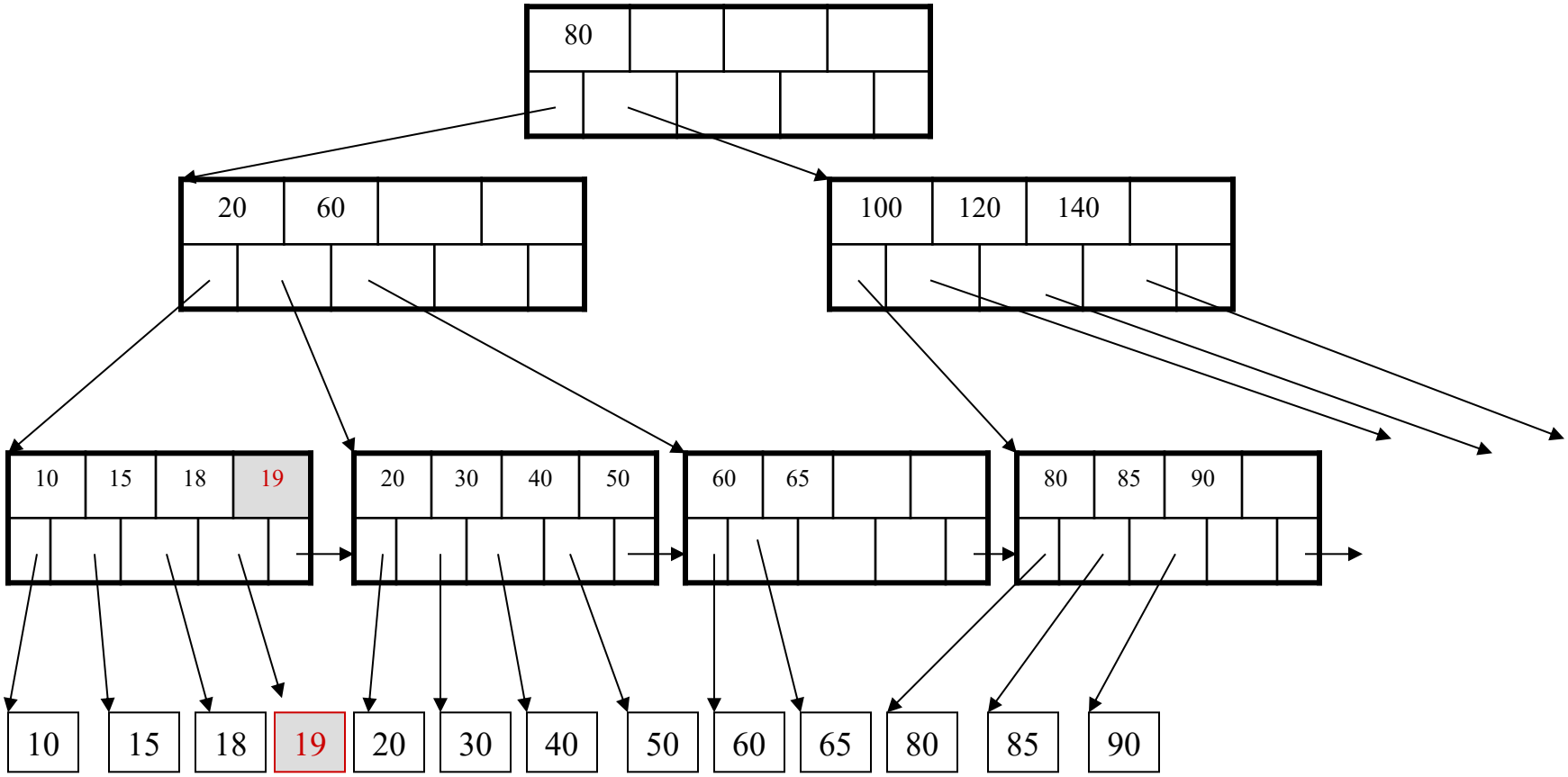
Insertion in a B+ Tree

Insert K=19



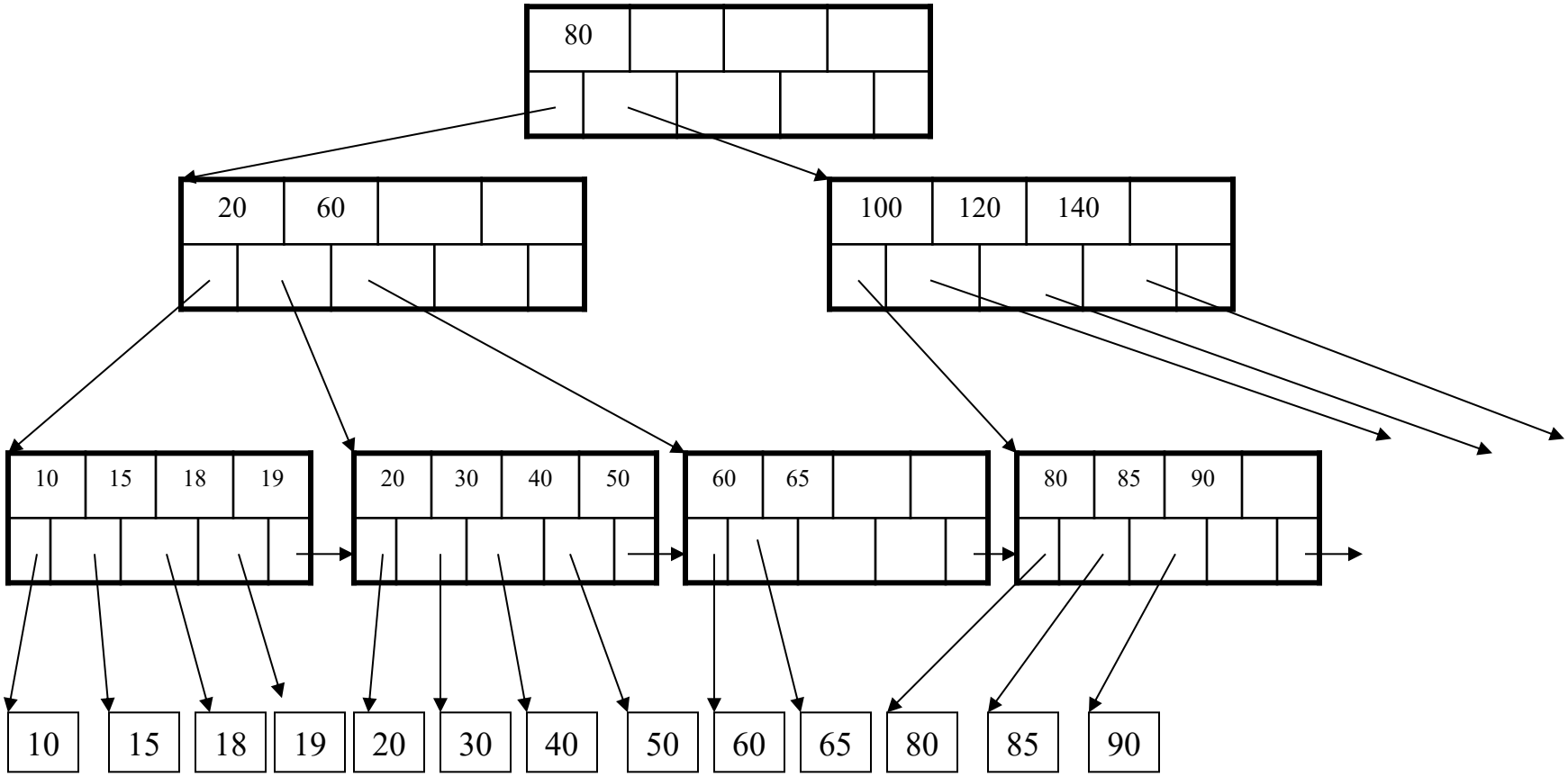
Insertion in a B+ Tree

After insertion



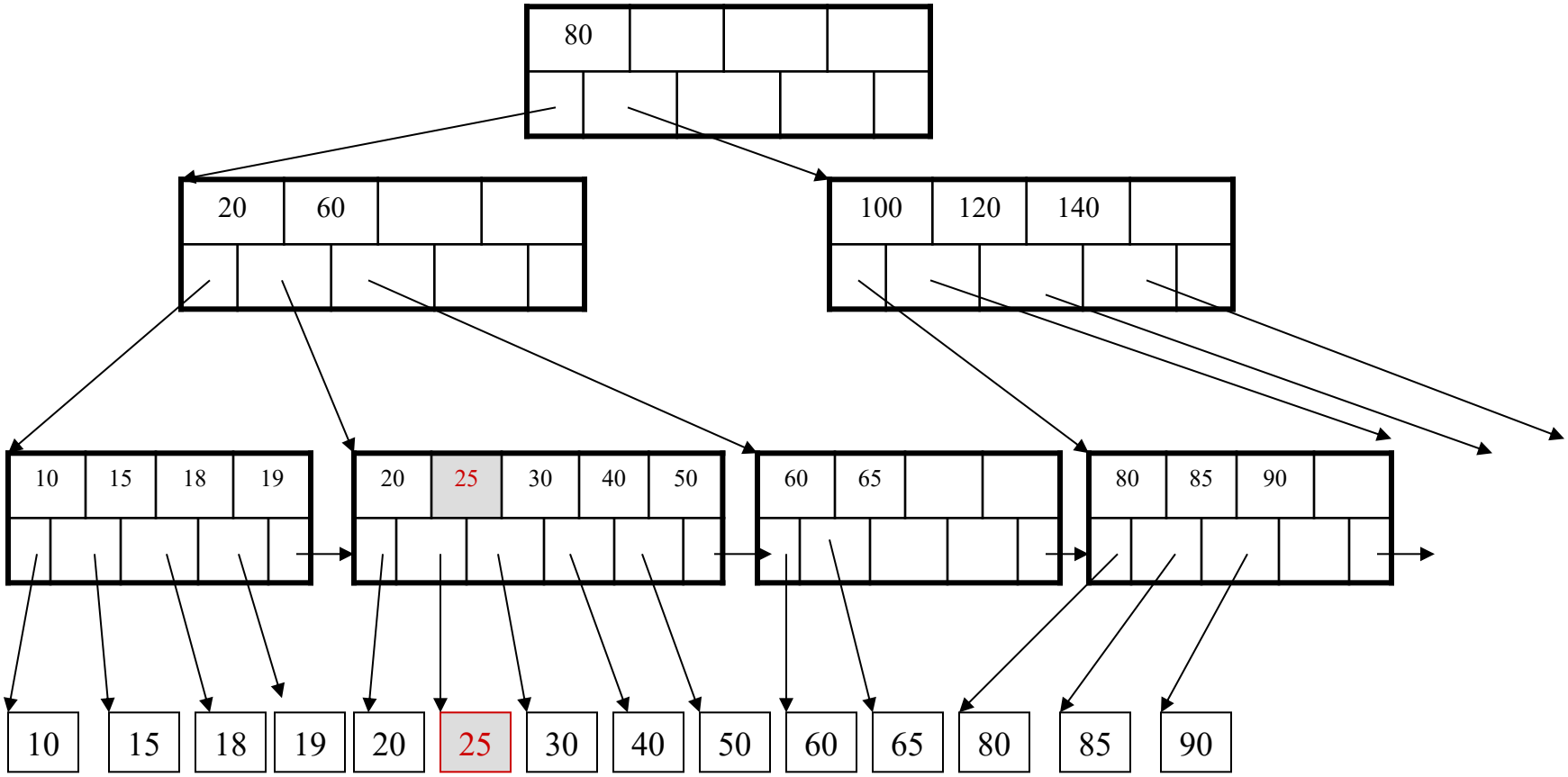
Insertion in a B+ Tree

Now insert 25



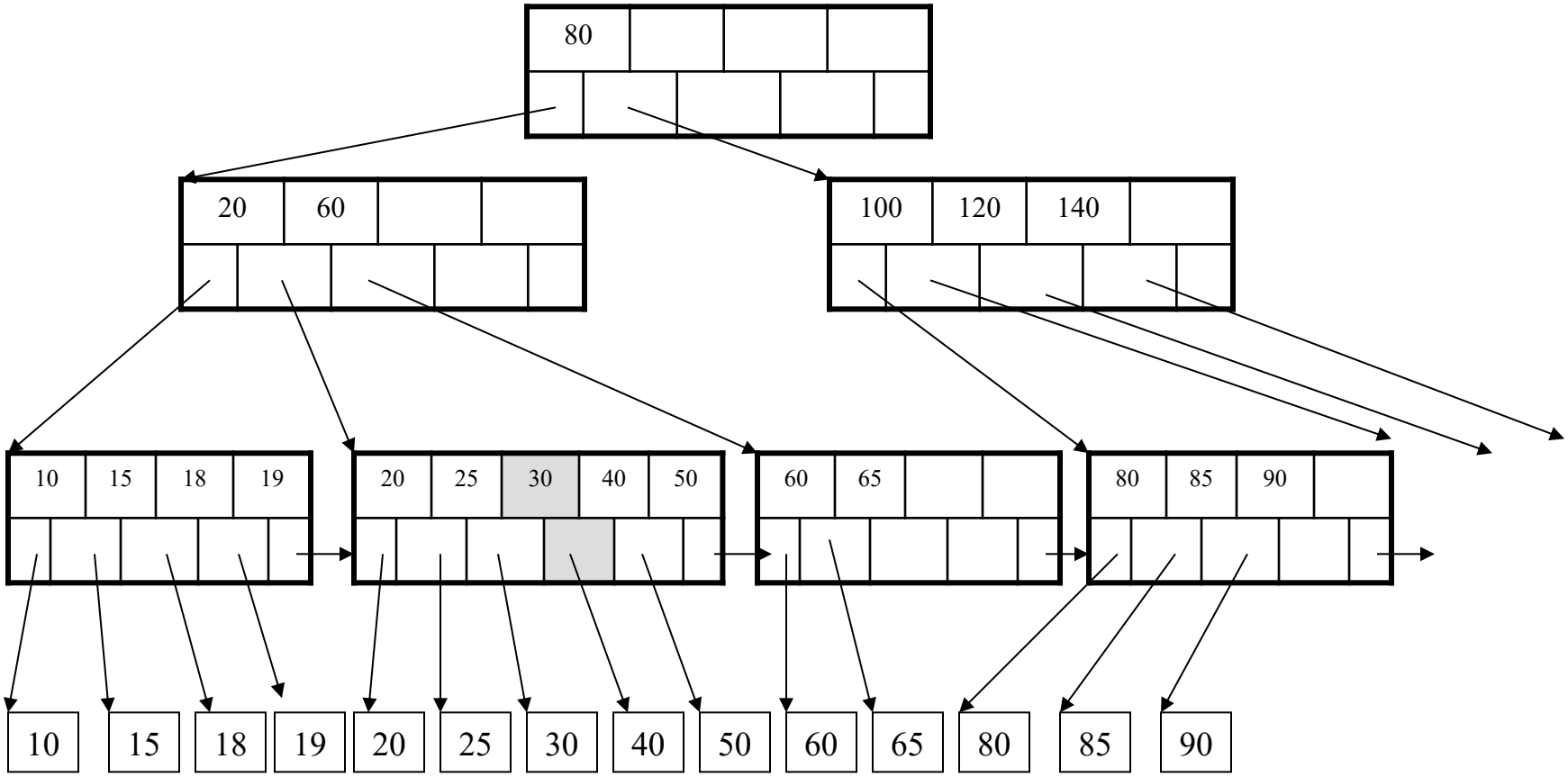
Insertion in a B+ Tree

After insertion



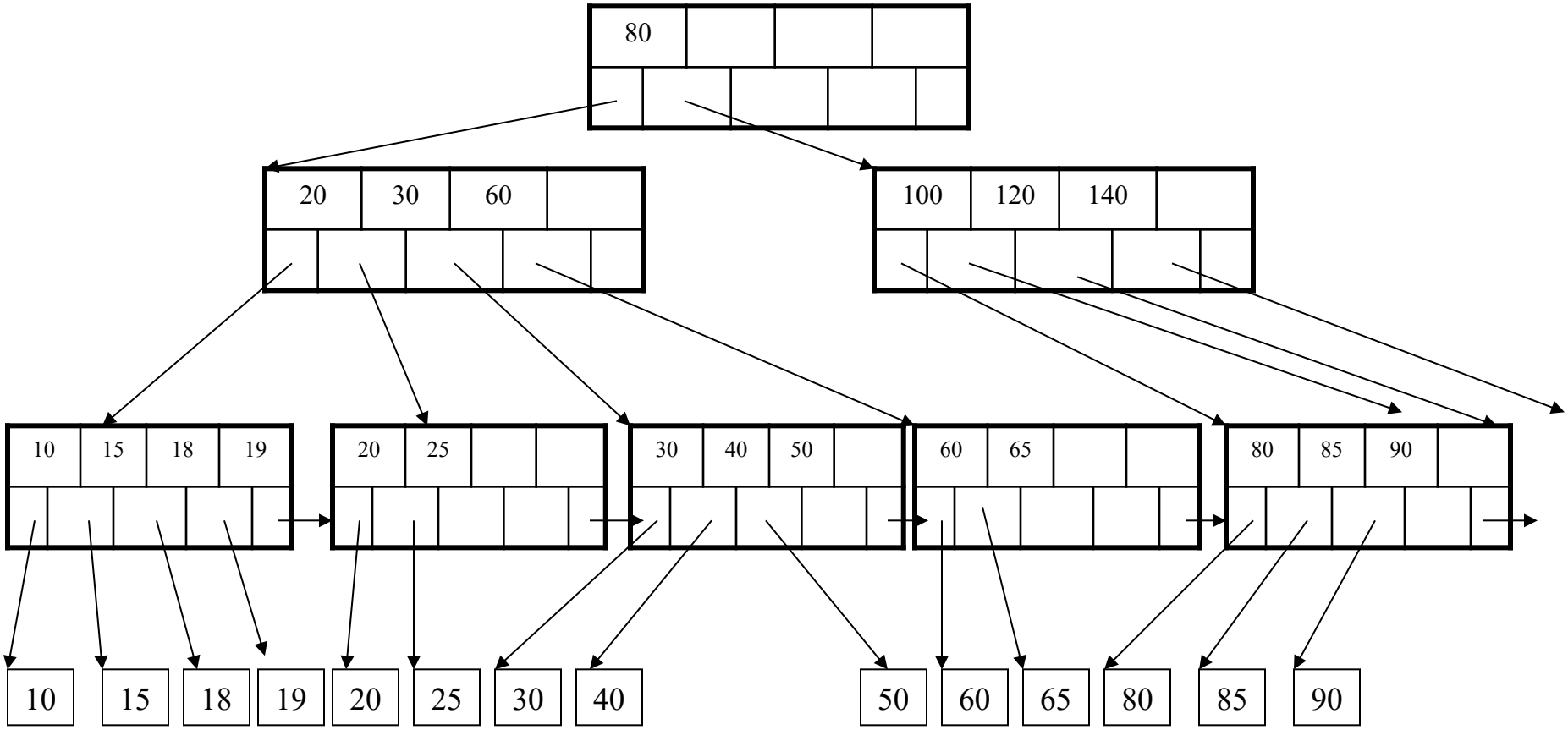
Insertion in a B+ Tree

But now have to split !



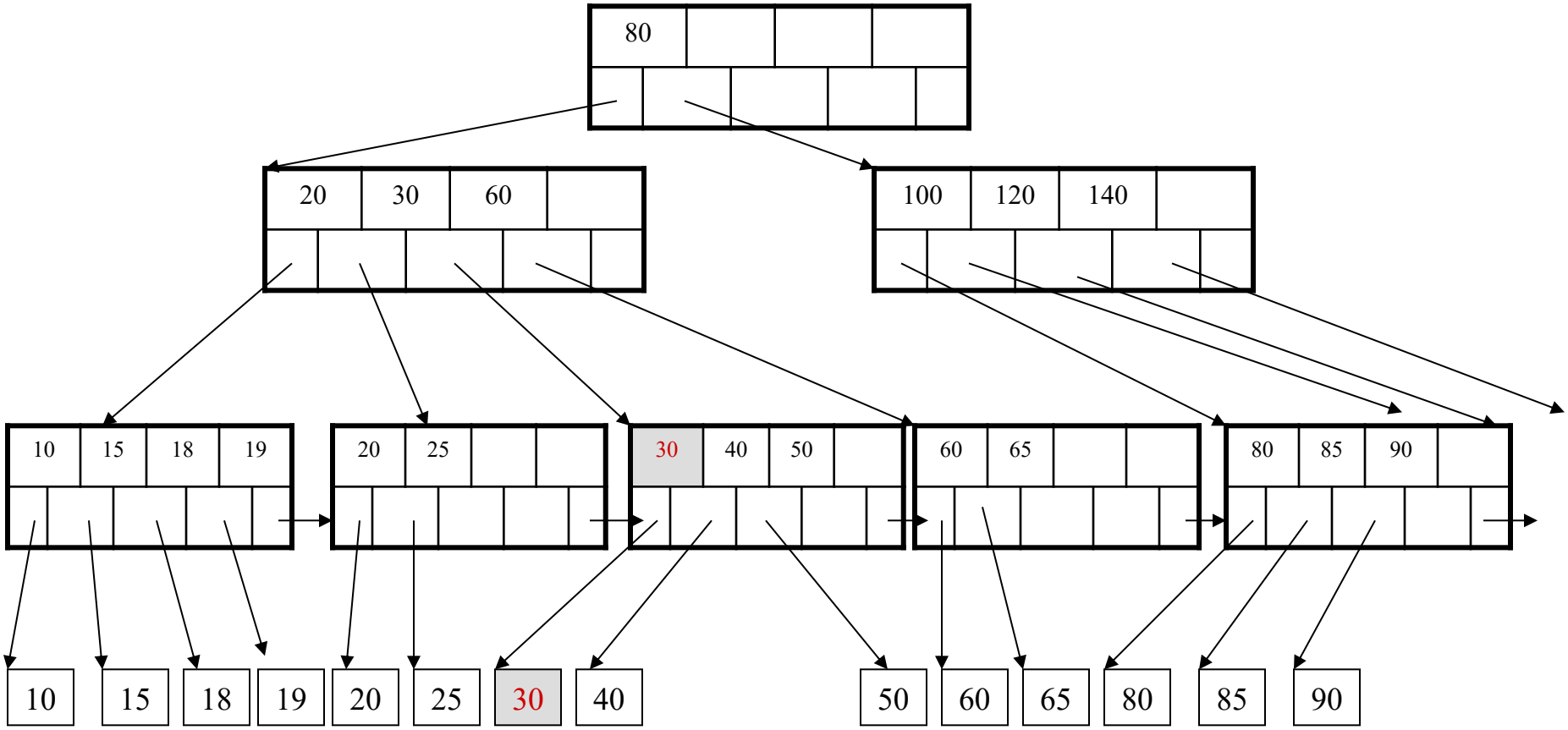
Insertion in a B+ Tree

After the split



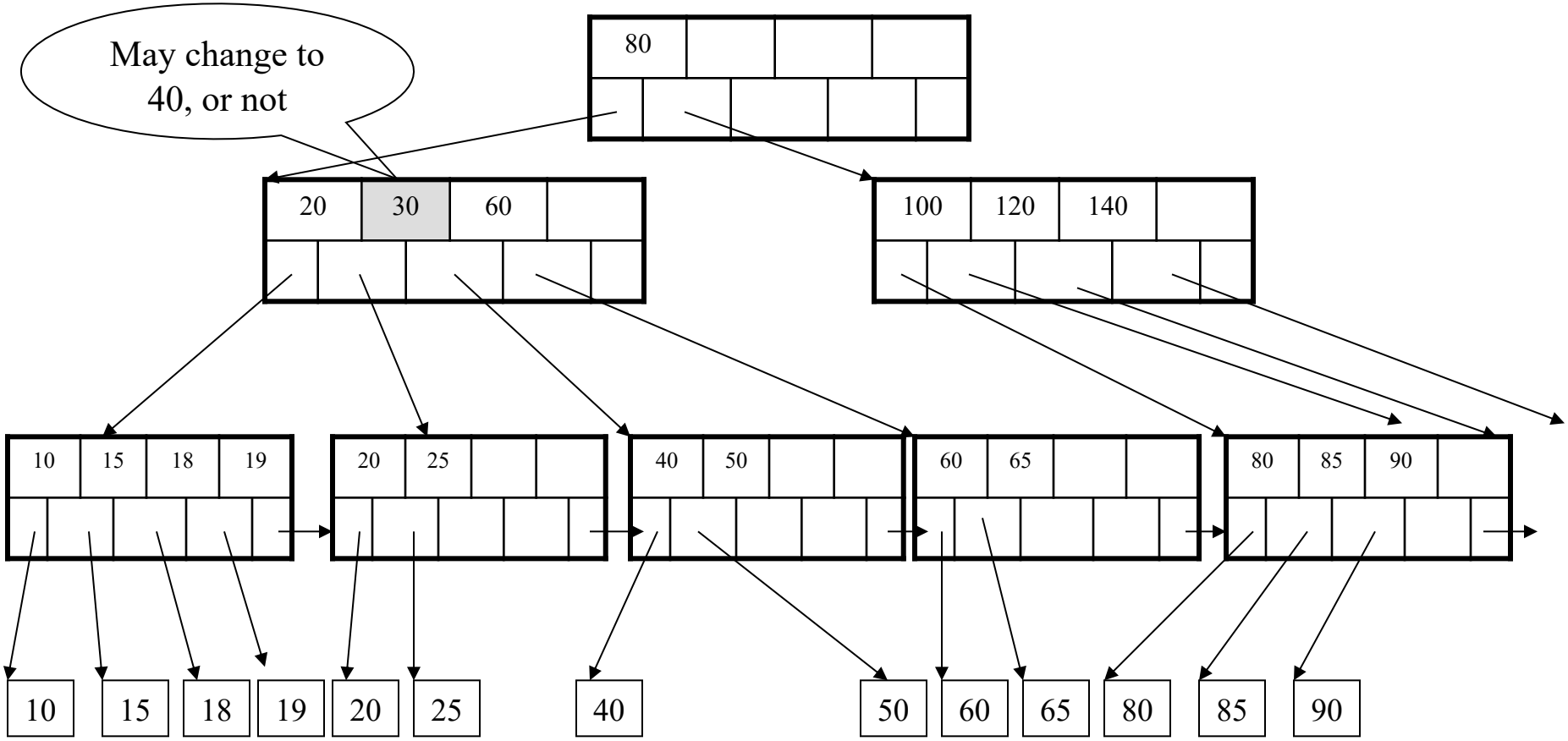
Deletion from a B+ Tree

Delete 30



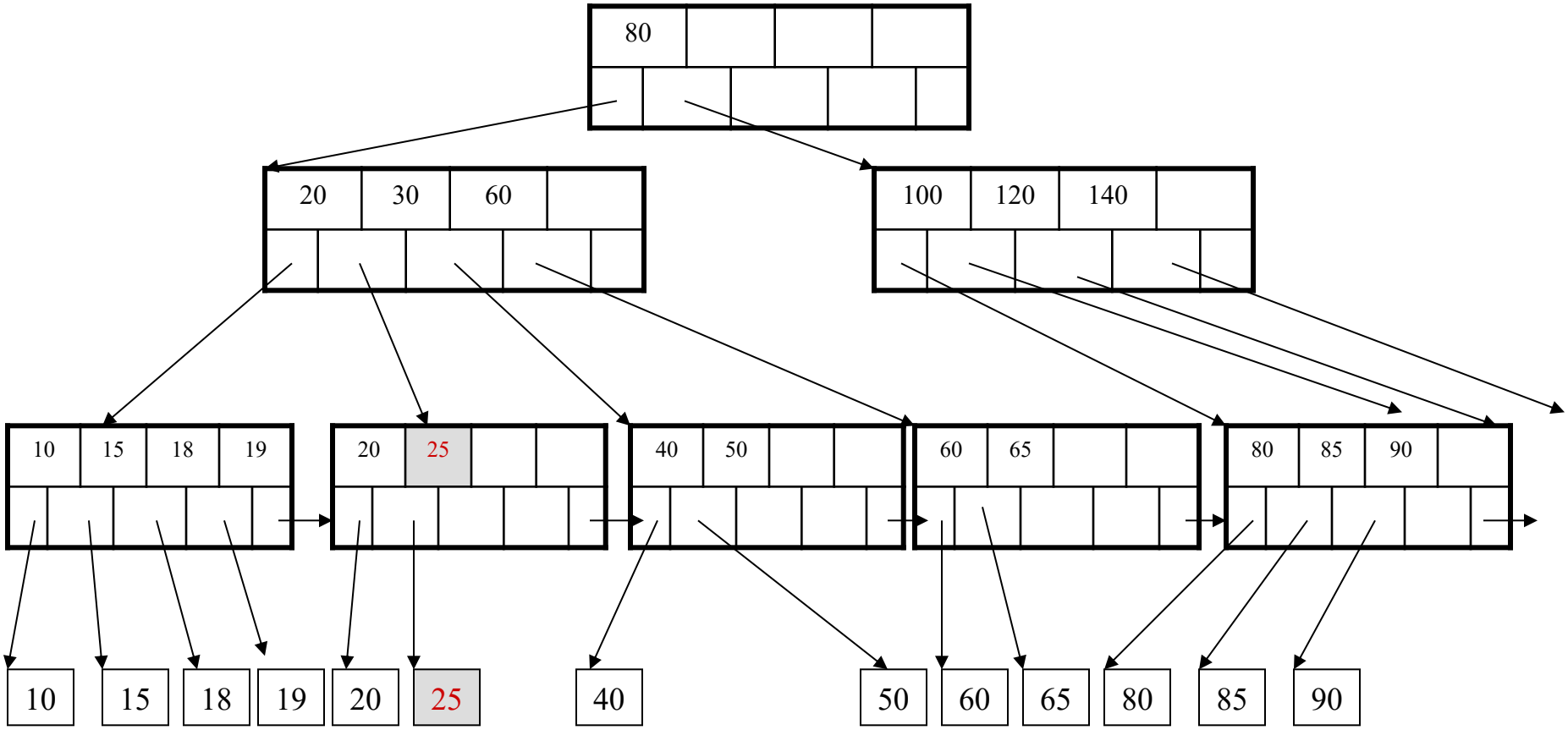
Deletion from a B+ Tree

After deleting 30



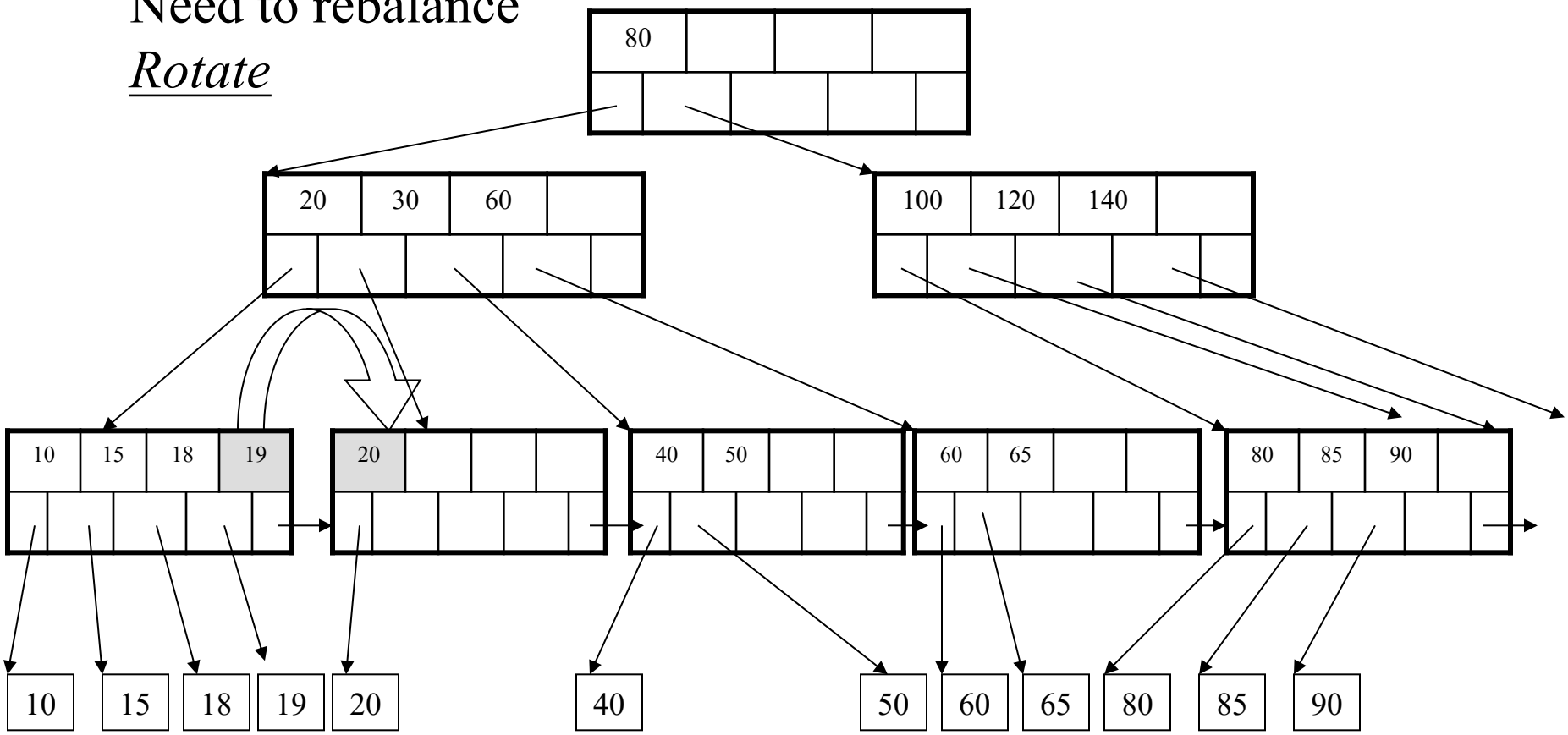
Deletion from a B+ Tree

Now delete 25



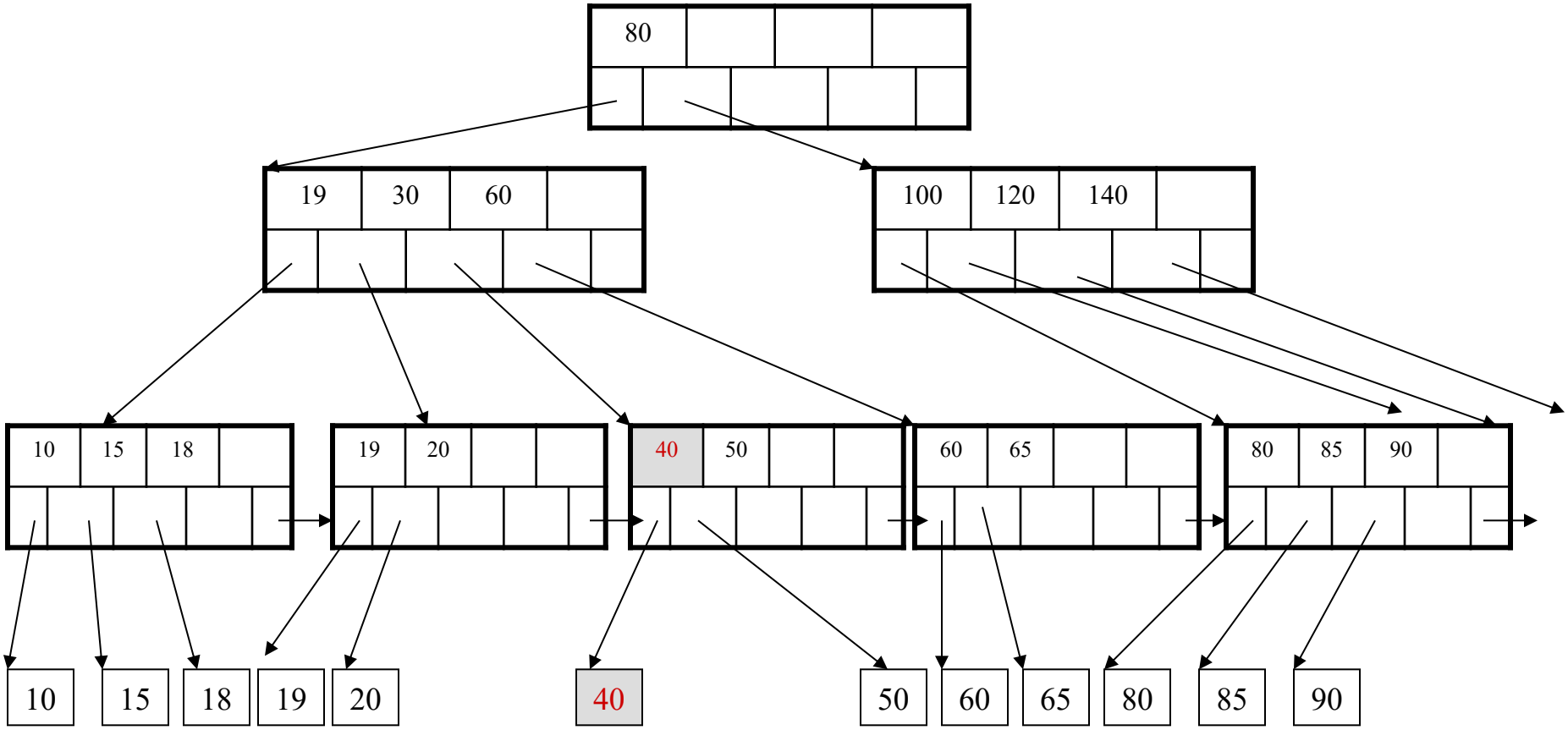
Deletion from a B+ Tree

After deleting 25
Need to rebalance
Rotate



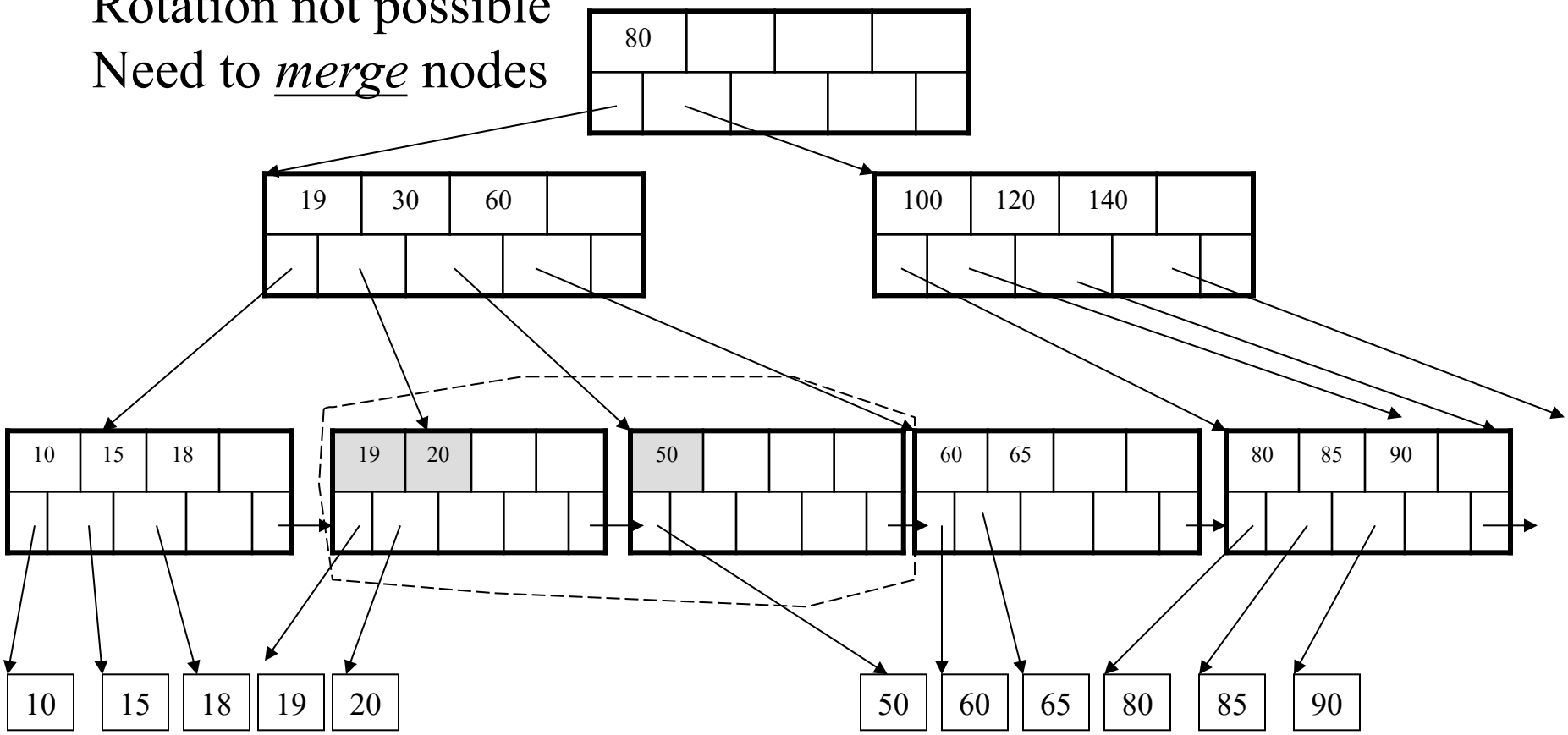
Deletion from a B+ Tree

Now delete 40



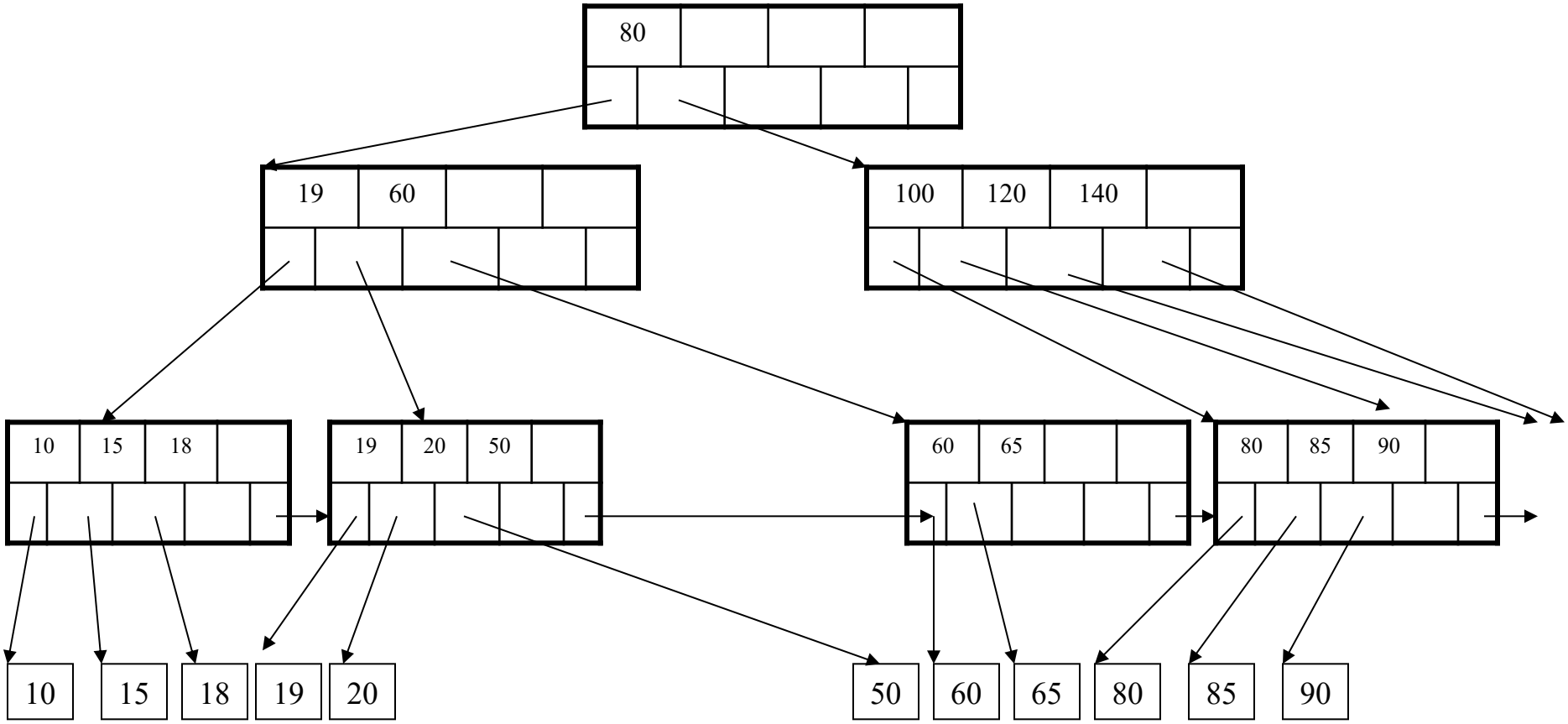
Deletion from a B+ Tree

After deleting 40
Rotation not possible
Need to *merge* nodes



Deletion from a B+ Tree

Final tree



Summary on B+ Trees

- Popular index structure on most DBMSs

- Very effective at answering 'point' queries:
  productName = 'gizmo'

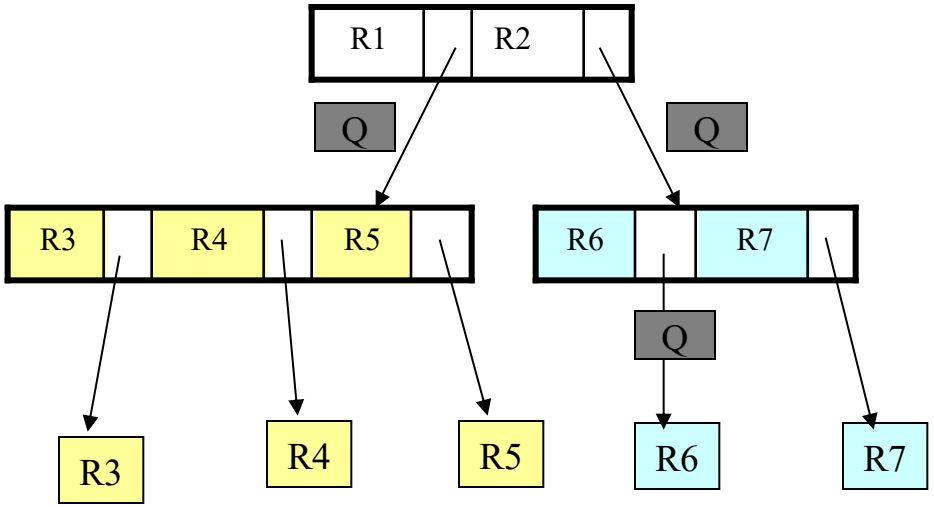
- Effective for range queries:
  50 < price AND price < 100

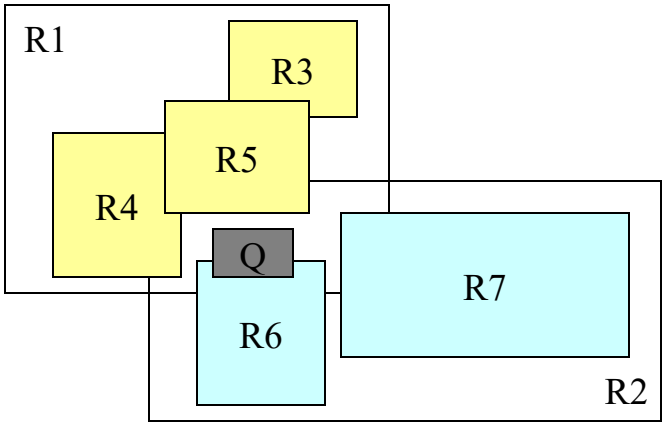
- Less effective for multirange:
  50 < price < 100 AND 2 < quant < 20

R-Tree: a multidimensional B-Tree

Designed for spatial data

Search key values are bounding boxes





For insertion: at each level, choose child whose bounding box needs least enlargement (in terms of area)