

CS/ECE 3810: Computer Organization

Lecture 5: Arithmetic for Computers

Anton Burtsev
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Arithmetic for Computers

- Operations on integers
 - Addition and subtraction
 - Multiplication and division
 - Dealing with overflow
- Floating-point real numbers
 - Representation and operations

Integer addition

- Example 7 + 6

$$0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0111_{\text{two}} = 7_{\text{ten}}$$

[illegible]

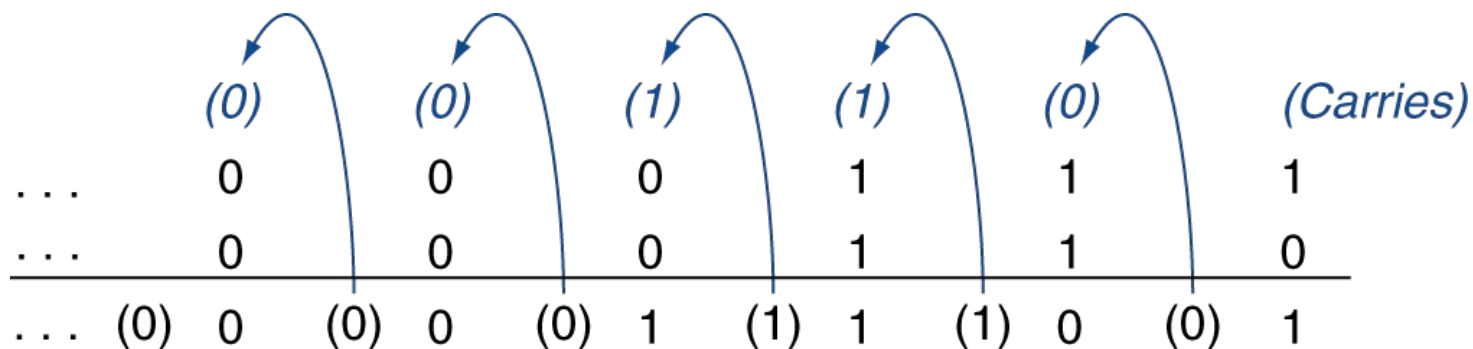
Integer addition

- Example 7 + 6

$$\begin{array}{r}
 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0111_{\text{two}} = 7_{\text{ten}} \\
 + \quad 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0110_{\text{two}} = 6_{\text{ten}} \\
 \hline
 = \quad 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 \ 1101_{\text{two}} = 13_{\text{ten}}
 \end{array}$$

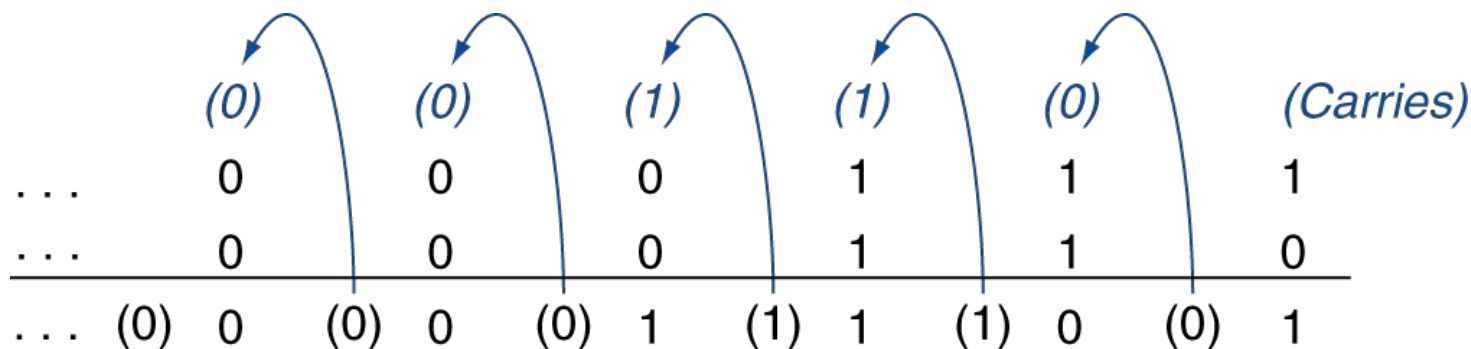
Integer Addition

- Example: $7 + 6$



Integer Addition

- Example: $7 + 6$



- Overflow if result out of range
 - Adding + (positive) and – (negative) operands, no overflow
 - Adding two + (positive) operands
 - Overflow if result sign is 1
 - Adding two –(negative) operands
 - Overflow if result sign is 0

Integer subtraction

- Example 7 – 6
 - Directly:

$$\begin{array}{r} 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0111_{\text{two}} = 7_{\text{ten}} \\ - \quad 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0110_{\text{two}} = 6_{\text{ten}} \\ \hline = \quad 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0001_{\text{two}} = 1_{\text{ten}} \end{array}$$

- Or as adding (-6)

$$\begin{array}{r} 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0111_{\text{two}} = 7_{\text{ten}} \\ + \quad 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1010_{\text{two}} = -6_{\text{ten}} \\ \hline = \quad 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0001_{\text{two}} = 1_{\text{ten}} \end{array}$$

Integer Subtraction

- Add negation of second operand

- Example: $7 - 6 = 7 + (-6)$

+7: 0000 0000 ... 0000 0111

-6: 1111 1111 ... 1111 1010

+1: 0000 0000 ... 0000 0001

- Overflow if result out of range
 - Subtracting two +(positive) or two -(negative) operands, no overflow
 - Subtracting +(positive) from -(negative) operand
 - Overflow if result sign is 0
 - Subtracting -(negative) from +(positive) operand
 - Overflow if result sign is 1

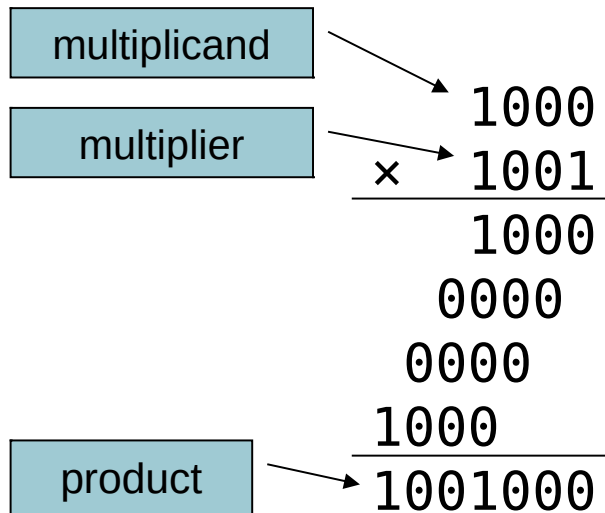
Dealing with Overflow

- Some languages (e.g., C) ignore overflow
 - Use MIPS `addu`, `addui`, `subu` instructions
- Other languages (e.g., Ada, Fortran) require raising an exception
 - Use MIPS `add`, `addi`, `sub` instructions
 - On overflow, invoke exception handler
 - Save PC in exception program counter (EPC) register
 - Jump to predefined handler address
 - `mfc0` (move from coprocessor reg) instruction can retrieve EPC value, to return after corrective action

Multiplication

Multiplication

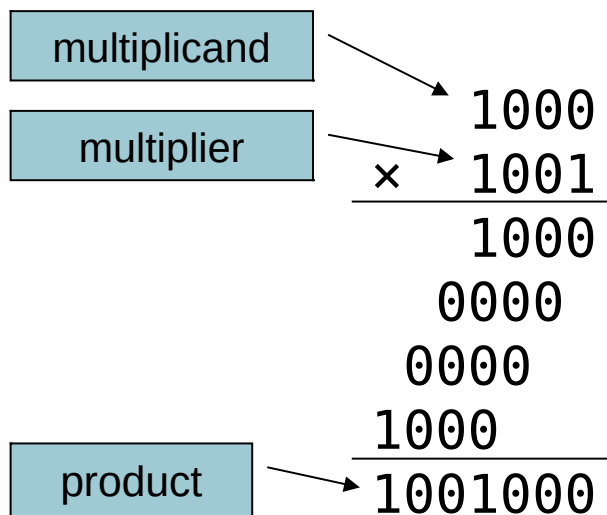
- Start with long-multiplication approach



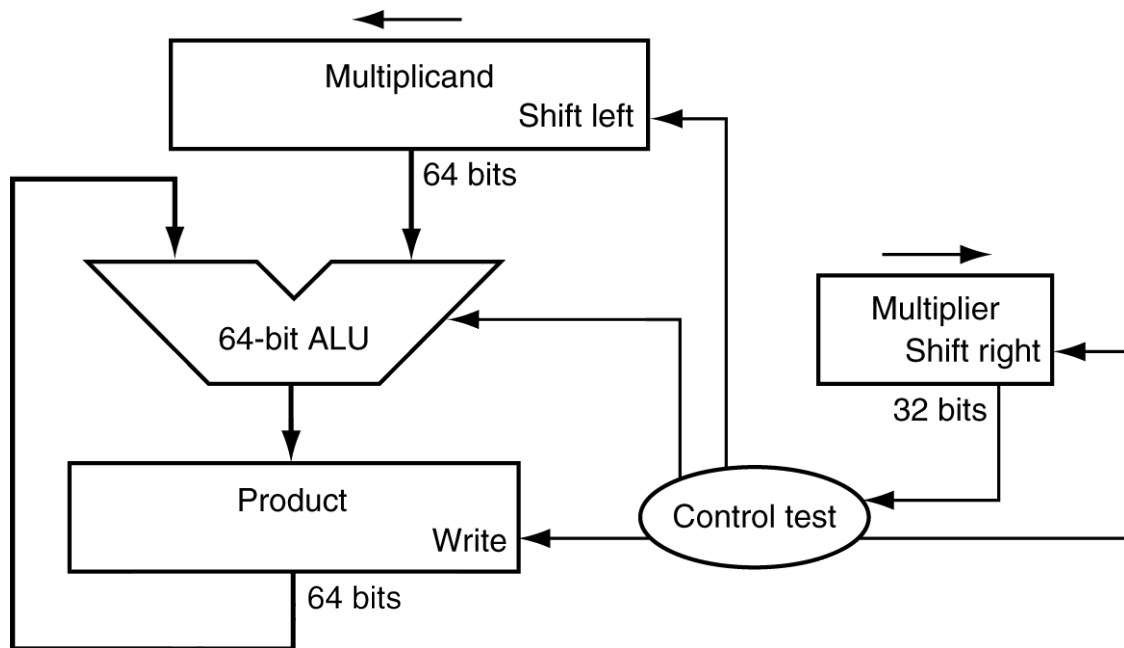
Length of product is
the sum of operand
lengths

Multiplication

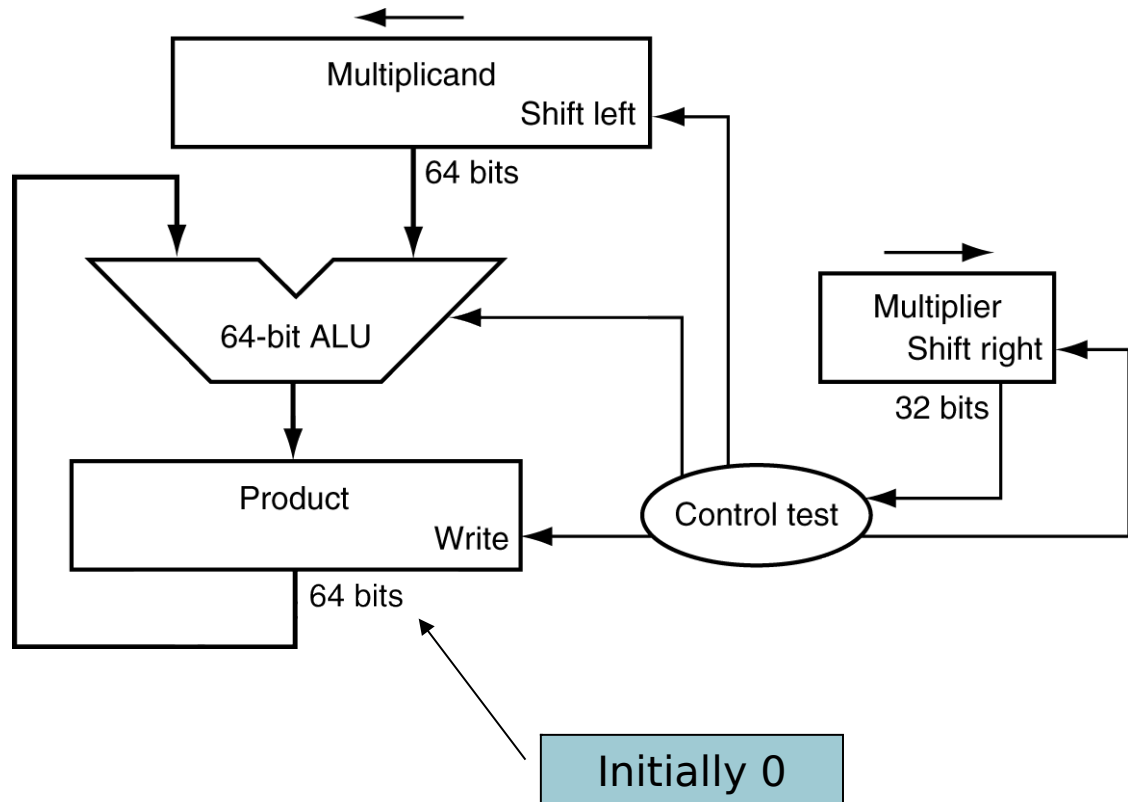
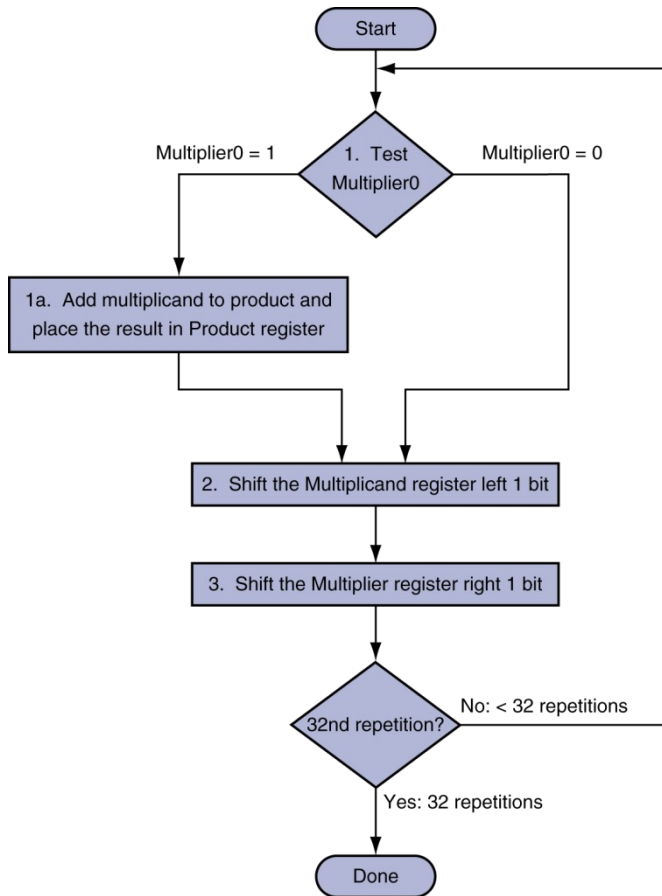
- Start with long-multiplication approach



Length of product is the sum of operand lengths

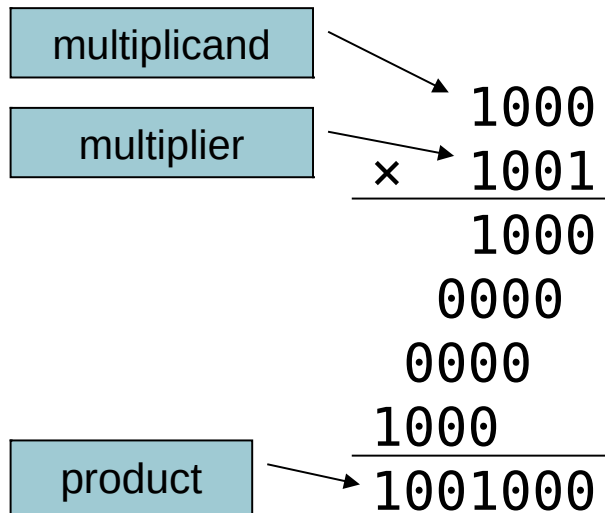


Multiplication Hardware



Can we do it faster?

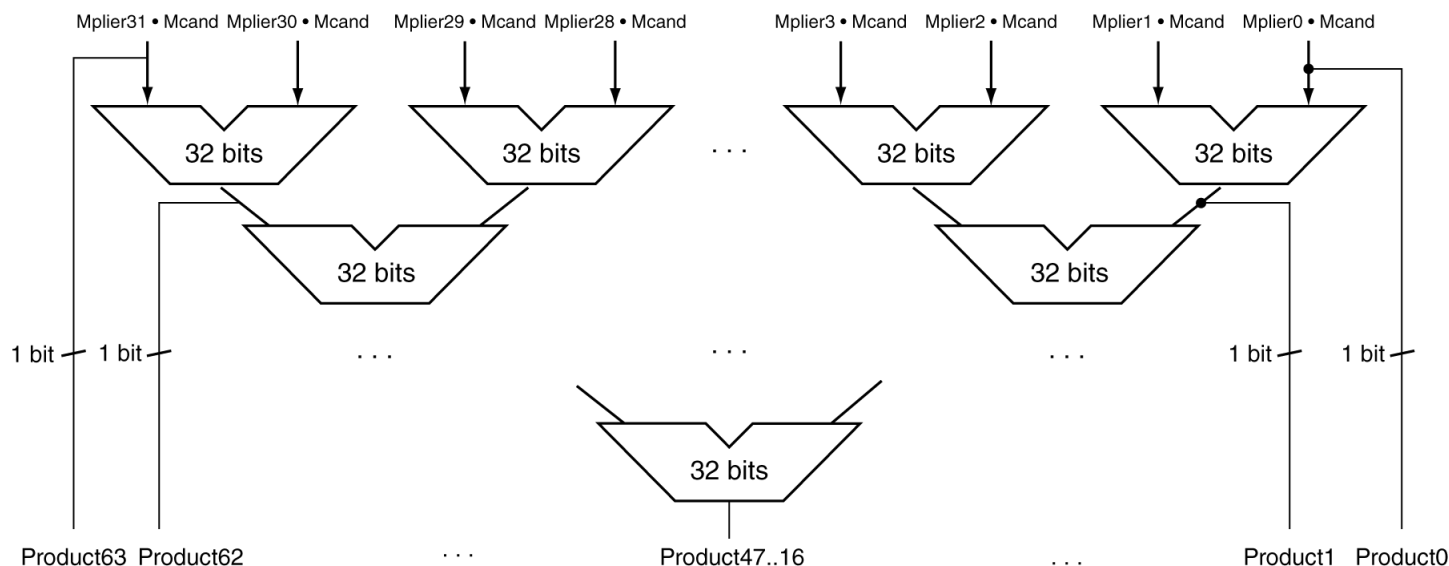
- Start with long-multiplication approach



Length of product is
the sum of operand
lengths

Faster Multiplier

- Uses multiple adders
 - Cost/performance tradeoff



- Can be pipelined
 - Several multiplication performed in parallel

MIPS Multiplication

- Two 32-bit registers for product
 - HI: most-significant 32 bits
 - LO: least-significant 32-bits
- Instructions
 - `mult rs, rt` / `multu rs, rt`
 - 64-bit product in HI/LO
 - `mfhi rd` / `mflo rd`
 - Move from HI/LO to rd
 - Can test HI value to see if product overflows 32 bits
 - `mul rd, rs, rt`
 - Least-significant 32 bits of product → rd

Division

Division

		$\overline{1001}_{\text{ten}}$	Quotient
Divisor	1000_{ten}	1001010_{ten}	Dividend
		$\underline{-1000}$	
		10	
		101	
		1010	
		$\underline{-1000}$	
		10_{ten}	Remainder

At every step,

- shift divisor right and compare it with current dividend
- if divisor is larger, shift 0 as the next bit of the quotient
- if divisor is smaller, subtract to get new dividend and shift 1 as the next bit of the quotient

Division

	$\begin{array}{r} 1001_{\text{ten}} \\ \hline 1000_{\text{ten}} \end{array}$		1001010_{ten}	Quotient																	
Divisor				Dividend																	
<table border="0" style="width: 100%;"> <tr> <td style="width: 25%;">0001001010</td> <td style="width: 25%;">0001001010</td> <td style="width: 25%;">0000001010</td> <td style="width: 25%;">0000001010</td> </tr> <tr> <td>100000000000</td> <td>0001000000</td> <td>0000100000</td> <td>0000001000</td> </tr> <tr> <td style="text-align: right;">✂</td> <td style="text-align: right;">✂</td> <td style="text-align: right;">✂</td> <td></td> </tr> <tr> <td>Quo: 0</td> <td>000001</td> <td>0000010</td> <td>000001001</td> </tr> </table>						0001001010	0001001010	0000001010	0000001010	100000000000	0001000000	0000100000	0000001000	✂	✂	✂		Quo: 0	000001	0000010	000001001
0001001010	0001001010	0000001010	0000001010																		
100000000000	0001000000	0000100000	0000001000																		
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At every step,

- shift divisor right and compare it with current dividend
- if divisor is larger, shift 0 as the next bit of the quotient
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Divide Example

- Divide 7_{ten} (0000 0111_{two}) by 2_{ten} (0010_{two})

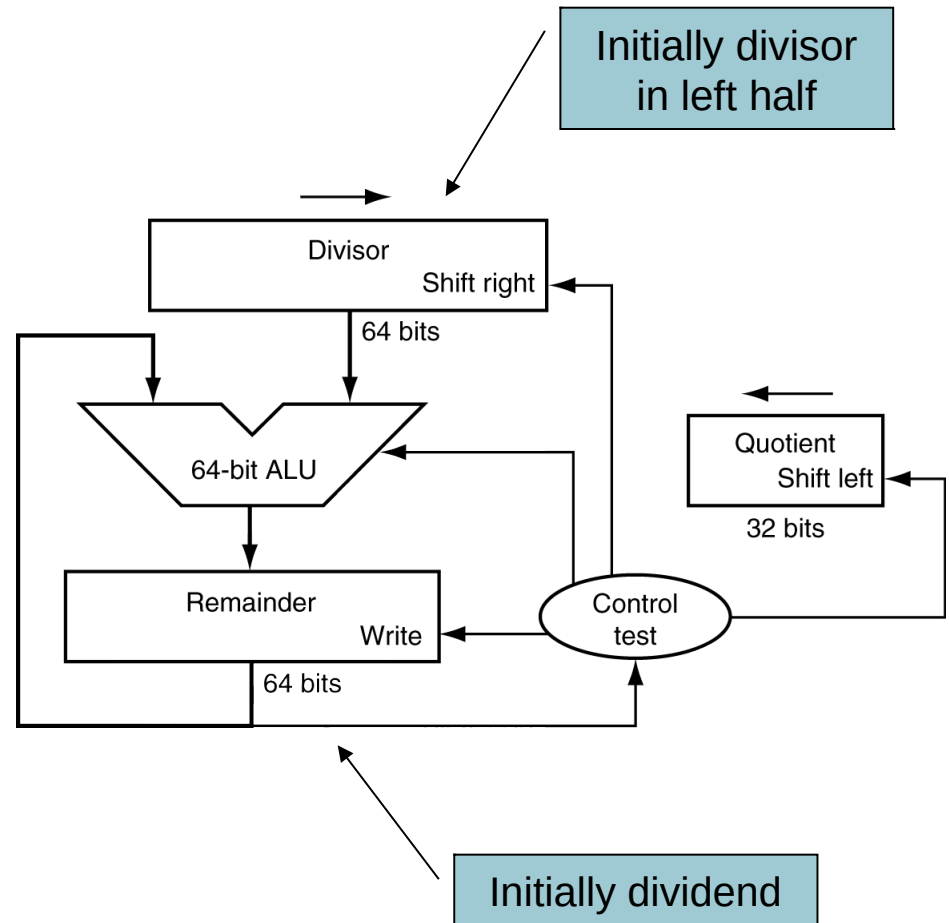
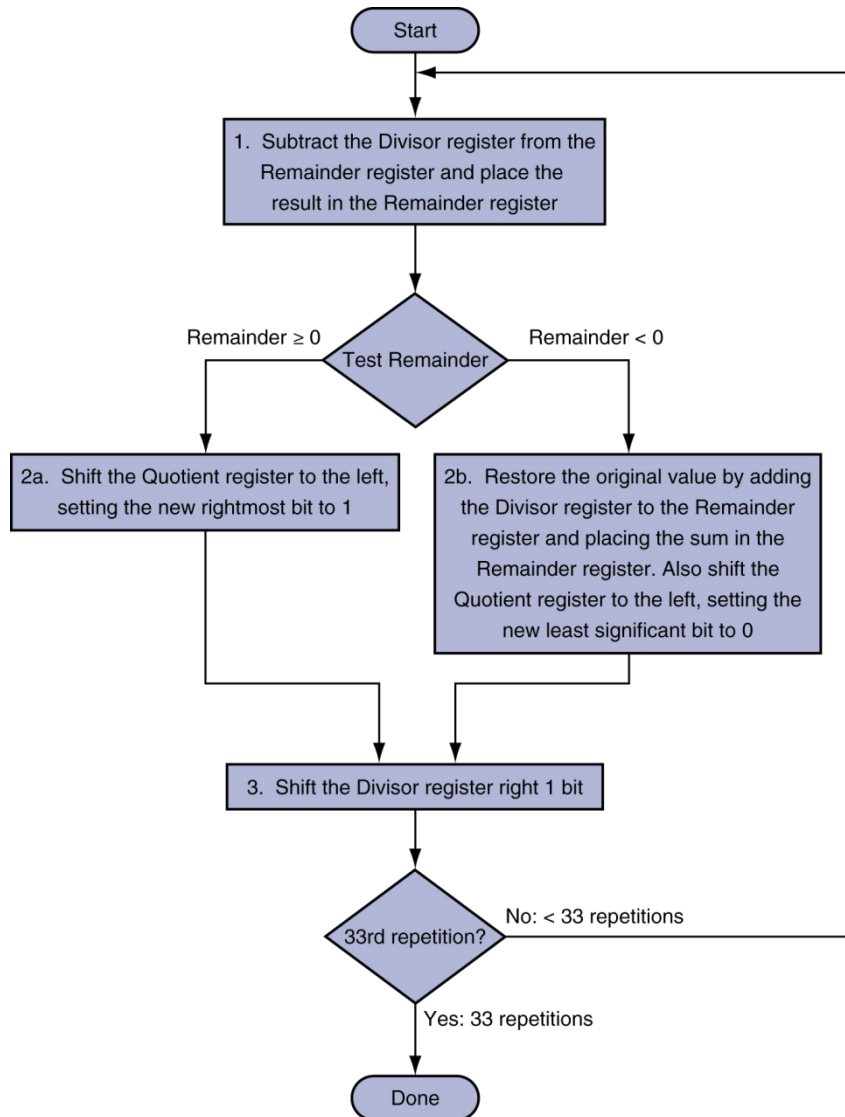
Iter	Step	Quot	Divisor	Remainder
0	Initial values			
1				
2				
3				
4				
5				

Divide Example

- Divide 7_{ten} ($0000\ 0111_{\text{two}}$) by 2_{ten} (0010_{two})

Iter	Step	Quot	Divisor	Remainder
0	Initial values	0000	0010 0000	0000 0111
1	Rem = Rem - Div	0000	0010 0000	1110 0111
	Rem < 0  +Div, shift 0 into Q	0000	0010 0000	0000 0111
	Shift Div right	0000	0001 0000	0000 0111
2	Same steps as 1	0000	0001 0000	1111 0111
		0000	0001 0000	0000 0111
		0000	0000 1000	0000 0111
3	Same steps as 1	0000	0000 0100	0000 0111
4	Rem = Rem - Div	0000	0000 0100	0000 0011
	Rem >= 0  shift 1 into Q	0001	0000 0100	0000 0011
	Shift Div right	0001	0000 0010	0000 0011
5	Same steps as 4	0011	0000 0001	0000 0001

Division Hardware



MIPS Division

- Use HI/LO registers for result
 - HI: 32-bit remainder
 - LO: 32-bit quotient
- Instructions
 - `div rs, rt` / `divu rs, rt`
 - No overflow or divide-by-0 checking
 - Software must perform checks if required
 - Use `mfhi`, `mflo` to access result

Thank you!