

Generating Functions

Transforming problems about sequences into problems about functions is known as generating functions.

$$\langle \underline{a_0} \ \underline{a_1} \ \underline{a_2} \ \underline{a_3} \ \dots \ \underline{a_i} \ \dots \rangle \leftrightarrow \underline{a_0}x^0 + \underline{a_1}x^1 + \underline{a_2}x^2 + \dots + a_ix^i + \dots + a_nx^n$$

Rule is a_i is acting as co-efficient of x^i

(where indices are 0, 1, 2, 3,)

$$\langle \underline{1} \ \underline{2} \ \underline{3} \ \underline{4} \ \dots \rangle \mapsto 1 \cdot x^0 + 2 \cdot x^1 + 3 \cdot x^2 + 4 \cdot x^3$$



$$\begin{aligned} * <1111\dots> &\mapsto 1 \cdot x^0 + 1 \cdot x^1 + 1 \cdot x^2 + 1 \cdot x^3 + \dots \\ &= \underline{1 + x + x^2 + x^3 + \dots} \quad \text{C.P.} = S_{\infty} = \frac{a}{1-x} \\ &= \frac{1}{1-x} \\ &= (1-x)^{-1} \end{aligned}$$
$$<1-11-1\dots> \mapsto 1 - x + x^2 - x^3 + \dots$$
$$<1010\dots> \mapsto 1 + x^2 + x^4 + \dots$$
$$<2020\dots> \mapsto$$



<u>n</u>	<u>Elements</u>	<u>No. of ways</u>	<u>$\{a, b\}$</u>
$n=0$	—	1 ✓	
$n=1$	a	1 ✓	
$n=2$	Not possible	0 ✓	
$n=3$	a	0 ✓	

\therefore Sequence = $\langle 1 \ 1 \ 0 \ 0 \ \dots \rangle \longmapsto 1 \cdot x^0 + 1 \cdot x^1 + 0 \cdot x^2 + 0 \cdot x^3 - \dots$
 $= (1+x)$

Q. The generating function for choosing n-elements from

$$\mathcal{G}_{\{a_1\}} = (1+x)$$

$$\mathcal{G}_{\{a_2\}} = (1+x)$$

$$\mathcal{G}_{\{a_1, a_2\}} = (1+x)(1+x)$$

$$\mathcal{G}_{\{a_1, a_2, a_3\}} = (1+x)(1+x)(1+x)$$



<u>n</u>	<u>Elements</u>	<u>No. of ways</u>
n=0	—	1
n=1	$\left. \begin{matrix} a_1 \\ a_2 \end{matrix} \right\} \longrightarrow 2C_1 = 2$	
n=2	$a_1 a_2 \longrightarrow 2C_2 = 1$	
n=3	NOT possible.	0
n=4	"	0
		\vdots

Sequence is

$$\langle 1 \ 2 \ 1 \ 0 \ 0 \dots \rangle \longmapsto 1 \cdot x^0 + 2 \cdot x^1 + 1 \cdot x^2 + 0 \cdot x^3 + \dots$$

$$= 1 + 2x + x^2$$

$$= (1+x)^2$$

$$= \underbrace{(1+x)}_{\text{Seq choosing } n\text{-element from } \{a_1\}} \underbrace{(1+x)}_{\text{Seq choosing } n\text{-element from } \{a_2\}}$$

= Generating function for
Seq choosing n -element
from $\{a_1\} * \{a_2\}$



$\{a_1, a_2, a_3\}$

$(1+x)^3$

<u>n</u>	<u>Elements</u>	<u>No. of ways</u>
$n=0$	—	1
$n=1$	$\begin{bmatrix} a_1 \\ a_2 \\ a_3 \end{bmatrix} \rightarrow 3C_1$	$= 3$
$n=2$	$\begin{bmatrix} a_1 a_2 \\ a_1 a_3 \\ a_2 a_3 \end{bmatrix} \rightarrow 3C_2$	$= 3$
$n=3$	$a_1 a_2 a_3 \rightarrow 3C_3$	$= 1$
$n=4$	Not possible	0

Sequence is

$\langle 1 \ 3 \ 3 \ 1 \ 0 \ 0 \dots \rangle$

$$= 1 \cdot x^0 + 3 \cdot x^1 + 3 \cdot x^2 + 1 \cdot x^3 + 0 \cdot x^4 + \dots$$

$$= 1 + 3x + \underline{3x^2} + x^3$$

$$= (1+x)^3$$

$$= \underbrace{(1+x)}_{\text{choosing } a_1} \underbrace{(1+x)}_{\text{choosing } a_2} \underbrace{(1+x)}_{\text{choosing } a_3}$$

generating function for choosing n -elements from $\{a_1\} \times \{a_2\} \times \{a_3\}$

$3 \cdot x^2$

Q. Generating function for choosing n-elements from

$$\{a_1, a_2, a_3, a_4\} = (1+x)(1+x)(1+x)(1+x)$$

$$= (1+x)^4$$

$$= 1 + 4x + 6x^2 + 4x^3 + x^4$$

$$= 1 + \underbrace{4 \cdot x^1} + \underbrace{6 \cdot x^2} + \underbrace{4 \cdot x^3} + \underbrace{1 \cdot x^4}$$

$$\underline{4 \cdot x^1}$$

$6 \cdot x^2 =$ No. of ways choosing 2-elements from $(1+x)^4$
is 6

$$\underline{4 \cdot x^3 = \text{No. of ways choosing 3-elements} = x^3 \text{ co-eff} = 4}$$



Generating function for choosing n-elements of $\{a_1, a_2, a_3, \dots, a_k\}$

$$= \underline{(1+x)(1+x)(1+x) \dots (1+x) [k - \text{times}]}$$

$$= \underline{(1+x)^k}$$

* The number of ways of choosing 3-objects of $\{a_1, a_2, \dots, a_k\}$

$$= \underline{\text{Co-efficient of } x^3 \text{ in } (1+x)^k}$$

* The number of ways of choosing n-objects of $\{a_1, a_2, \dots, a_k\}$

$$= \underline{\text{Co-efficient of } x^n \text{ in } (1+x)^k}$$

$$= {}^kC_n \quad \checkmark$$



* $\langle 1 \ 1 \ 1 \ 1 \ \dots \rangle$ $\leftrightarrow 1 + x + x^2 + x^3 + \dots$

$\leftrightarrow \frac{1}{1-x} = (1-x)^{-1}$

* The generating function for choosing n -objects of $\{a_1\}$ with repetitions

$\rightarrow = 1.x^0 + 1.x^1 + 1.x^2 + 1.x^3 + \dots$

$= 1 + x + x^2 + x^3 + \dots$

$= (1-x)^{-1}$

<u>n</u>	<u>$\{a_1\}$</u> <u>Elements</u>	<u>No. of ways</u>
$n=0$	—	1
$n=1$	a_1	1
$n=2$	a_1, a_1	1
$n=3$	a_1, a_1, a_1	1

$\langle 1 \ 1 \ 1 \ 1 \ \dots \rangle$



* The generating function for choosing n -objects of $\{a_2\}$

with repetition $= (1 - x)^{-1}$ ✓

* The generating function for choosing n-objects of

$$\{\underline{a_1, a_2, a_3, \dots, a_k}\} \text{ with repetitions} = (1-x)^{-1} (1-x)^{-1} \dots (1-x)^{-1} [k \text{ times}]$$

$$= (1-x)^{-k} \checkmark$$

* The number of ways of choosing n-objects of $\{a_1, a_2, a_3, \dots, a_k\}$

With repetitions = Co-efficient of x^n in $(1-x)^{-k}$

$$\sum_{n=0}^{\infty} {}^{n+k-1}C_n \cdot x^n$$

$$= \frac{(n+k-1)!}{(k-1)!n!} = {}^{(n+k-1)}C_n$$

$$\textcircled{{}^{n+k-1}C_n}$$



$$*(1-x)^{-k} = \sum_{n=0}^{\infty} {}^{(n+k-1)}C_n \cdot x^n = \sum_{n=0}^{\infty} {}^{n+k-1}C_n \cdot x^n$$

* $(1 - ax)^{-k} = \sum_{n=0}^{\infty} {}^{(n+k-1)}C_n \cdot a^n \cdot x^n$ where 'k' is positive integer

$$(1-ax)^{-k} = \sum_{n=0}^{\infty} {}^{n+k-1}C_n \cdot a^n \cdot x^n$$

Q. $(1 - 3x)^{-1} = ?$

$$a=3, k=1$$

$$(1-ax)^{-k} = \sum_{n=0}^{\infty} {}^{n+k-1}C_n \cdot a^n \cdot x^n$$

$$= \sum_{n=0}^{\infty} {}^{n+1-1}C_n \cdot 3^n \cdot x^n$$

$$= \sum_{n=0}^{\infty} 3^n \cdot x^n$$



Q. The generating function of the sequence

$$\{1, -2, 4, -8, 16, -32, \dots\}$$

$x^0 \quad x^1 \quad x^2 \quad x^3 \quad x^4$

$$\{1, -2, 4, -8, 16, -32, \dots\} \longrightarrow 1 \cdot x^0 - 2 \cdot x^1 + 4 \cdot x^2 - 8 \cdot x^3 + 16 \cdot x^4 - 32 \cdot x^5 + \dots$$

$$= 1 - 2x + 4x^2 - 8x^3 + 16x^4 - 32x^5 + \dots$$

$$= \frac{1}{1+2x}$$

$$= (1+2x)^{-1}$$

$$\begin{cases} a=1, & r=-2x \\ S_{\infty} = \frac{a}{1-r} \end{cases}$$

$$(1-2x)^{-1}$$
$$(1-ax)^{-k}$$

Q. The generating function of the sequence

$$\begin{aligned}\{0, 1, 3, 9, 27, \dots\} &\mapsto 0 \cdot x^0 + 1 \cdot x^1 + 3 \cdot x^2 + 9 \cdot x^3 + 27 \cdot x^4 + \dots \\ &= x + 3x^2 + 9x^3 + 27x^4 + \dots \\ &= \frac{x}{1-3x} \quad a=x, \quad r=3x \\ &= x(1-3x)^{-1}\end{aligned}$$

Q. The Co-efficient of x^{27} in the expansion of

$(x^4 + 2x^5 + 3x^6 + 4x^7 + \dots + \infty)^5$ is

$$= [x^4 (1 + 2x + 3x^2 + 4x^3 + \dots + \infty)]^5$$

$$= x^{20} (1 + 2x + 3x^2 + 4x^3 + \dots + \infty)^5$$

$$= x^{20} [(1-x)^{-2}]^5$$

$$= x^{20} (1-x)^{-10}$$

$16C_7$

$$x^{20} (1-x)^{-10}$$

$$= x^{20} \sum_{n=0}^{\infty} \frac{n+10-1}{n!} C_n \cdot x^n$$

$$= x^{20} \sum_{n=0}^{\infty} \frac{n+9}{n!} C_n \cdot x^n$$

$$= x^{20} \sum_{n=0}^{\infty} \frac{n+9}{n!} C_n \cdot x^n$$

Put $n=7$, $16C_7$

$\langle 1 \ 2 \ 3 \ 4 \ \dots \rangle$

$$(1-x)^{-K} = \sum_{n=0}^{\infty} \frac{n+K-1}{n!} C_n \cdot x^n$$

$n=0, K=0, 2$

$$n+K-1 C_n = 0+2-1 C_0 = 1 C_0 = 1$$

$$1+2-1 C_1 = 2 C_1 = 2$$

$$2+2-1 C_2 = 3 C_2 = 3$$

Q. Consider the polynomial

$$P(x) = a_0 + a_1x + a_2x^2 + a_3x^3, \checkmark$$

Where $a_i \neq 0, \forall i$. The minimum number of multiplications needed to evaluate p on an input x is **(GATE-CS-06)**

a) 3

b) 4

c) 6

d) 9

$$\begin{aligned} P(x) &= a_0 + a_1x + a_2x^2 + a_3x^3 \\ &= a_0 + x[a_1 + a_2x + a_3x^2] \\ &= a_0 + x[a_1 + x(a_2 + a_3x)] \end{aligned}$$

Q. $\sum_{x=1}^{99} \frac{1}{x(x+1)} = \underline{\hspace{2cm}}$

(GATE-15-Set1)



$$\sum_{x=1}^{99} \frac{1}{x(x+1)} = \frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \frac{1}{3 \cdot 4} + \frac{1}{4 \cdot 5} + \dots$$

$$= \left(1 - \cancel{\frac{1}{2}}\right) + \left(\cancel{\frac{1}{2}} - \cancel{\frac{1}{3}}\right) + \left(\cancel{\frac{1}{3}} - \cancel{\frac{1}{4}}\right) + \left(\cancel{\frac{1}{4}} - \cancel{\frac{1}{5}}\right) + \dots \left(\cancel{\frac{1}{99}} - \frac{1}{100}\right)$$

$$\Rightarrow 1 - \frac{1}{100}$$

$$= \frac{99}{100} = 0.99$$

Q. Let $G(x) = 1/(1-x)^2 = \sum_{i=0}^{\infty} g(i)x^i$,

Where $|x| < 1$ what is $g(i)$?

(GATE-CS-05)

a) 1

b) $i + 1$

c) $2i$

d) $2i$

$$G(x) = \frac{1}{(1-x)^2} = \sum_{i=0}^{\infty} g(i)x^i$$

Rules for finding (C.F) are given below.

1. Characteristics roots are real and distinct say t_1, t_2, \dots, t_k

$$C.F = c_1 \cdot t_1^n + c_2 \cdot t_2^n + \dots + c_k \cdot t_k^n = c_1 \cdot t_1^n + c_2 \cdot t_2^n + c_3 \cdot t_3^n + \dots$$

2. Roots are real and two roots are equal say $t_1, t_1, t_3, t_4, \dots, t_k$

t_1, t_1, t_3

$$C.F = (c_1 + c_2 \cdot n) t_1^n + c_3 \cdot t_3^n + \dots + c_k \cdot t_k^n =$$

3. Roots are real and 3 roots are equal say $t_1, t_1, t_1, t_4, \dots, t_k$

$$C.F = (c_1 + c_2 \cdot n + c_3 \cdot n^2) t_1^n + c_4 \cdot t_4^n + \dots + c_k \cdot t_k^n$$

4. Suppose if all the roots are equal say $t_1, t_1, t_1, \dots, t_1$

$$C.F = (c_1 + c_2 \cdot n + c_3 \cdot n^2 + \dots + c_k \cdot n^{k-1}) t_1^n$$



5. A pair of roots are complex say $(\alpha \pm i\beta)$

$$C.F = r^n(c_1 \cos(n\theta) + c_2 \sin(n\theta))$$

Where

$$r = \sqrt{\alpha^2 + \beta^2} \text{ and}$$

$$\theta = \tan^{-1}\left(\frac{\beta}{\alpha}\right)$$



Particular Solution (P.S):-

From equation (2),

$$\text{P.S} = \frac{1}{\phi(E)} \{F(n)\} \checkmark$$

The solution of equation (1) is

$$a_n = \text{C.F} + \text{P.S}$$

$$\begin{aligned} \phi(E) a_n &= F(n) \\ a_n &= \frac{F(n)}{\phi(E)} \end{aligned}$$



Rule to find P.S:-

When $F(n) = b^n$ we have

$$P.S = \frac{b^n}{\phi(E)} = \frac{b^n}{\phi(b)} \text{ provided } \phi(b) \neq 0$$

$$P.S = \frac{F(n)}{\phi(E)}$$

$$E^2 - 2E + 2$$

$$F(n) = b^n \cdot n^k$$



Case of failure:-

When $\phi(b) = 0$, use the formula

$$\frac{b^n}{(E-b)^k} = \underline{C(n, k)} \cdot \underline{b^{n-k}} \quad (k = 1, 2, 3, \dots)$$

$$\begin{aligned}\phi(E) &= E^2 - 2E + 1 \\ &= (E-1)^2\end{aligned}$$

Method of undetermined Co-efficient:

$$\phi(E) a_n = F(n) \checkmark$$

$$\underline{F(n)} = \underline{b^n \cdot n^k}$$

$$b \neq t \quad k=2$$

i) If 'b' is not characteristic root then

$$P.S = \underline{b^n(A_0 n^k + A_1 n^{k-1} + \dots + A_k)} = b^n [c n^2 + d n + e]$$

ii) If 'b' is the characteristic root with multiplicity 'm' then

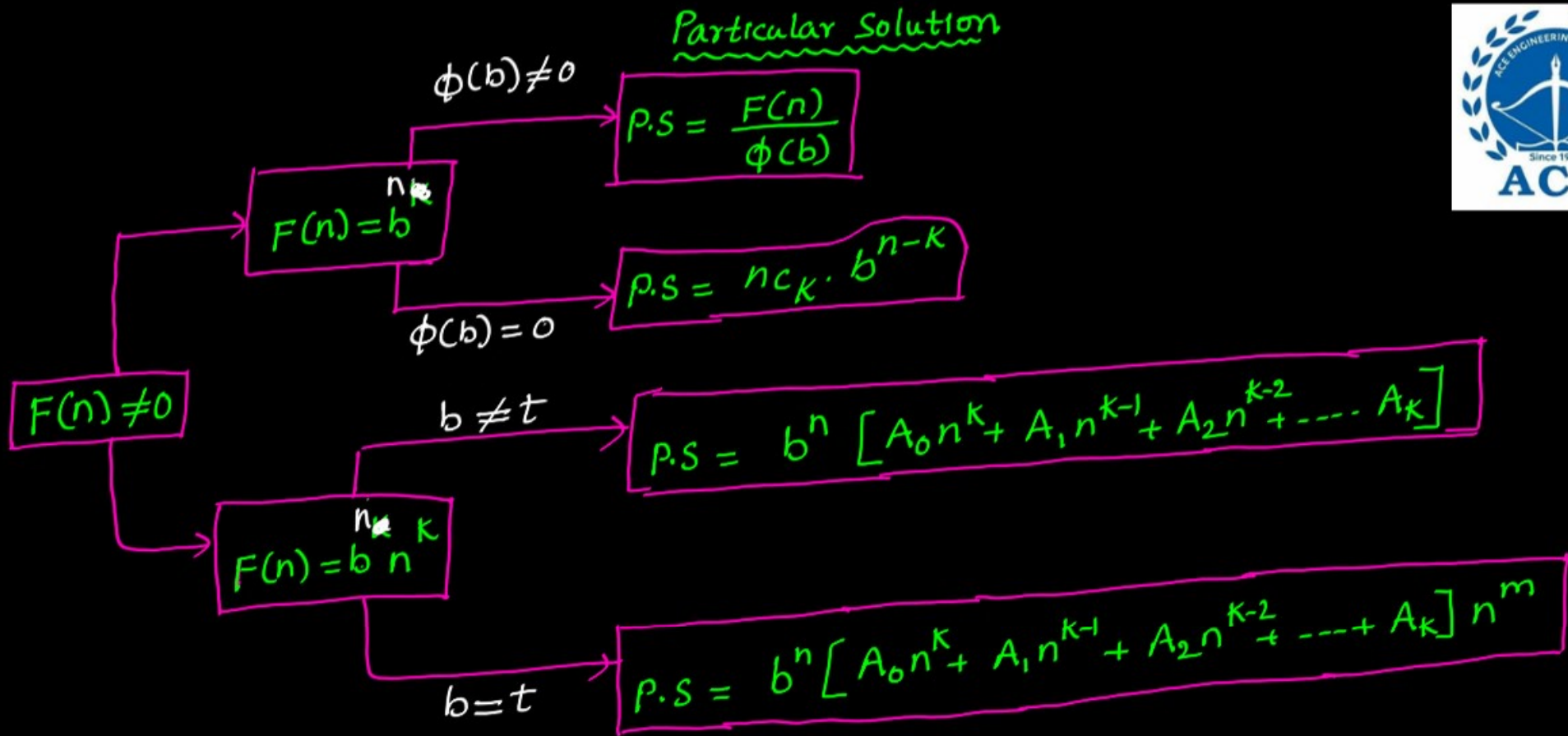
$$P.S = \underline{b^n(A_0 n^k + A_1 n^{k-1} + \dots + A_k) n^m} =$$

$$b = t$$

$$(t - s)^2 = 0$$

$$t = s, s$$

$$m = 2$$



Q. Solve the following recurrence relation (GATE-CS-98)

$$X_n = 2X_{n-1} - 1 \quad n > 1$$

$$X_1 = 2$$

Given Recurrence Relation is

$$x_n = 2x_{n-1} - 1$$

$$x_n - 2x_{n-1} = -1$$

$$\text{put } n = n+1$$

$$x_{n+1} - 2x_n = -1$$

$$E^1(x_n) - 2E^0(x_n) = -1$$

$$(E-2)x_n = -1$$

This is in form $\phi(E) a_n = F(n)$

$$\phi(E) = E-2$$

$$\text{char. eq. } \phi(t) = t-2=0 \Rightarrow t=2$$

$$C.F = C_1 t_1^n = C_1 2^n \quad \checkmark$$

Particular solution:

$$F(n) = -1 = (-1)(1)^n$$

$$P.S = \frac{F(n)}{\phi(b)} = -1 \left[\frac{(1)^n}{b-2} \right] = -1 \left[\frac{1^n}{1-2} \right] = 1 \quad \checkmark$$

Solution to RR is

$$x_n = C.F + P.S$$

$$x_n = C_1 2^n + 1$$

Initial condition $x_1 = 2$ ✓

$$x_1 = C_1 2^1 + 1 = 2$$

$$2C_1 = 1$$

$$C_1 = \frac{1}{2}$$

$$\therefore x_n = \frac{1}{2} \cdot 2^n + 1$$

$$x_n = 2^{n-1} + 1$$



Q. Solve the recurrence equations

(GATE-CS-87)

$$T(n) = T(n-1) + n$$

$$T(1) = 1, a_1 = 1$$

Given Recurrence relation

$$T(n) = T(n-1) + n$$

$$a_n = a_{n-1} + n$$

$$a_n - a_{n-1} = n$$

Put $n = n+1$

$$a_{n+1} - a_n = n+1$$

$$E'(a_n) - E^0(a_n) = n+1$$

$$(E-1)a_n = n+1$$

$$\phi(E)a_n = F(n)$$

$$\phi(E) = E-1$$

$$\text{char. eq. } \phi(t) = (t-1)^1 = 0 \Rightarrow \boxed{t=1} \quad m=1$$

$$\boxed{C.F = C_1 t_1^n = C_1 1^n = C_1}$$

$$F(n) = n+1 = 1^n (n+1) = b^n \cdot n^k$$

$$\boxed{b=1}, k=1$$

Here $b=t$

$$P.S = b^n [A_0 n^k + A_1 n^{k-1} + \dots] n^m = 1^n [cn+d] n^1$$

$$\boxed{P.S = cn^2 + dn} \quad \checkmark$$

\therefore Solution to R.R,

$$a_n = C.F + P.S$$

$$\boxed{a_n = c_1 + cn^2 + dn} \quad \checkmark \checkmark \checkmark$$

* Let us Substitute P.S in R.R.

$$a_n - a_{n-1} = n$$

$$(cn^2 + dn) - [c(n-1)^2 + d(n-1)] = n$$

$$(cn^2 + dn) - [c(n-1)^2 + d(n-1)] = n$$

put $n=0$

$$0 - c + d = 0$$

$$-c + d = 0 \rightarrow \textcircled{1}$$

put $n=1$

$$c + d = 1 \rightarrow \textcircled{2}$$

$$c = \frac{1}{2} ; d = \frac{1}{2}$$

$$\therefore \boxed{a_n = c_1 + \frac{1}{2}n^2 + \frac{1}{2}n} \quad \checkmark$$

initial $a_1 = 1$

$$a_1 = c_1 + \frac{1}{2}(1)^2 + \frac{1}{2}(1) = 1$$

$$\Rightarrow c_1 = 0$$

$$\therefore \boxed{a_n = \frac{1}{2}n^2 + \frac{1}{2}n}$$

$$\boxed{a_n = \frac{n(n+1)}{2}}$$





Method - II :

Given Recurrence Relation is

$$T(n) = T(n-1) + n$$

Initial $T(1) = 1$

$$T(2) = T(1) + 2 = 1 + 2 = 3$$

$$T(3) = T(2) + 3 = (1 + 2) + 3$$

$$T(4) = T(3) + 4 = (1 + 2 + 3) + 4$$

$$T(5) = T(4) + 5 = (1 + 2 + 3 + 4) + 5$$

⋮

$$T(n) = 1 + 2 + 3 + 4 + 5 + \dots + n = \frac{n(n+1)}{2} \quad \checkmark$$

Method of undetermined Co-efficient:

$$\phi(E) a_n = F(n)$$

$$F(n) = b^n \cdot n^k$$

i) If 'b' is not characteristic root then

$$P.S = b^n (\underline{A_0} n^k + \underline{A_1} n^{k-1} + \dots + \underline{A_k}) \quad \checkmark$$

ii) If 'b' is the characteristic root with multiplicity 'm' then

$$P.S = b^n (\underline{A_0} n^k + \underline{A_1} n^{k-1} + \dots + \underline{A_k}) n^m \quad \checkmark$$

Q. Solve the recurrence equations:

$$T(n) = T\left(\frac{n}{2}\right) + 1$$

$$T(1) = 1$$

Given Recurrence Relation

$$T(n) = T\left(\frac{n}{2}\right) + 1$$

$$T(n) = aT\left(\frac{n}{b}\right) + \theta(n^k \log_b^p n)$$

$$a=1, b=2, k=0, p=0$$

$$a=1, b^k = 2^0 = 1$$

clearly $a = b^k$

and also $p > -1$

$$T(n) = \theta\left[n^{\log_b a}, \log_b^{p+1} n\right]$$

(GATE-CS-88)

$$T(n) = \theta\left[n^{\log_2 1}, \log_2^{0+1} n\right]$$

$$= \theta(\log n) \checkmark$$

Method-II:

$$T(n) = T\left(\frac{n}{2}\right) + 1$$

$$\text{let } n = 2^x$$

$$\therefore T(2^x) = T\left(\frac{2^x}{2}\right) + 1$$

$$T(2^x) = T(2^{x-1}) + 1$$

$$a_x = a_{x-1} + 1$$

$$a_x - a_{x-1} = \textcircled{1}$$

put $x = x+1$

$$a_{x+1} - a_x = 1$$

$$E^1(a_x) - E^0(a_x) = 1$$

$$(E-1)a_x = 1$$

$$\Phi(E)a_n = F(n)$$

$$\boxed{\Phi(E) = E-1}$$

$$= (E-1)^1 = (E-1)^{1=K}$$

$$\text{char. eq. } \Phi(t) = t-1=0 \Rightarrow t=1$$

$$\boxed{C.F = C_1 \cdot t_1^n = C_1 \cdot 1^n = C_1} \checkmark$$



$$F(n) = 1 = 1^n = b^n$$

$$\phi(b) = 0$$

$$P.S = \frac{b^n}{(E-b)^k} = n c_k \cdot b^{n-k}$$

$$= x c_k \cdot b^{x-k}$$

$$= x c_1 \cdot (1)^{x-1}$$

$$= x$$

$$n = 2^x$$

$$\log_2 n = x$$

Solution to R.R.

$$a_x = c \cdot F + P.S$$

$$\boxed{a_x = c_1 + x} \quad \checkmark$$

initial condition $T(1) = 1$

$$T(1) = T(n) = T(2^x) = T(2^0)$$

$$= a_0 = c_1 + 0 = 1$$

$$c_1 = 1$$

$$\therefore a_x = c_1 + x$$

$$\boxed{a_x = 1 + x}$$



$$T(n) = T\left(\frac{n}{2}\right) + 1$$

$$n = 2^x \neq 2^0$$

$$T(1)$$

$$a_x = T(2^x) = 1 + x$$

$$= \boxed{T(n) = 1 + \log_2 n}$$

$$\Theta(1 + \log n)$$

Q. Find the particular solution of

$$a_n - 2a_{n-1} + a_{n-2} = n^2 + n + 1$$

$$a_n - 2a_{n-1} + a_{n-2} = n^2 + n + 1$$

$$\boxed{p.s = cn^4 + dn^3 + en^2}$$



Q. The recurrence equation

(GATE-CS-04)

$$T(1) = 1$$

$$T(n) = 2T(n-1) + n, n \geq 2 \text{ evaluates to}$$

a) $2^{n+1} - n - 2$

b) $2^n - n$

c) $2^{n+1} - 2n - 2$

d) $2^{n+1} + n$



Q. The solution to the recurrence equation $T(2^k) = 3T(2^{k-1}) + 1$, $T(1) = 1$ is

(GATE-CS-02)

a) 2^k

b) $(3^{k+1} - 1)/2$

c) $3^{\log 2^k}$

d) $2^{\log 3^k}$