

Discrete Mathematics

DMS = Discrete Mathematic Structures

DS = Discrete Structures

D.M = D. Maths = Discrete Maths

MFCS = Mathematical Foundations
for Computer Science

DS&T = Discrete Structures & Graph Theory



Discrete Maths	Engg. Maths
<u>I</u> . Mathematical Logic	1. Linear alg
<u>II</u> . Set theory & Algebra	2. calculus
<u>III</u> Combinatorics	3. Probability
<u>IV</u> . Graph Theory	





Discrete Maths (6 marks — 15 marks)

8 Marks

Chapter wise

UNITS

- I. Mathematical Logic
 - 1. propositional logic (10)
 - 2. First-order logic (10)
- II. Set Theory & Algebraic structures
 - 3. sets (4)
 - 4. Relations (4)
 - 5. Partial-ordering & Lattices (4)
 - 6. Functions (4)
 - 7. Algebraic structures (Groups) (6)
- III. Combinatorics
- IV. Graph Theory
 - 8. Combinatorics (20)
 - 9. Graph Theory (20)

80 hrs

Reference Books

1. ✓ Discrete Mathematical Structure with Applications to Computer Science – Tremblay & Manohar
2. ✓ Discrete Mathematics for Computer Scientist & Mathematicians – Mott, Kandell & Baker.
- *** 3. ✓ Discrete Mathematics – Kenneth Rosen
4. ✓ Discrete Mathematics – C.L.Liu

LOGIC



Logic is the basis of all mathematical reasoning, and of all automated reasoning.

It has the practical applications

- To the design of computing machines
- To the specification of systems
- To artificial intelligence
- To programming languages and to other areas of computer science.

Propositions:-

A proposition is a declarative sentence (that is, a sentence that declares a fact) that is either true or false, But not both.

proposition = T (or) F



Eg:-

1. Washington D.C. is the capital of the united states of America. ✓ (T)
2. Bombay is the capital of India. ✓ (F)
3. $1 + 1 = 2$ ✓ (T)
4. $2 + 2 = 3$ ✓ (F)

Eg:-

1. What time is it? ✓ NOT proposition
2. Read this carefully ✓ “
3. $x + 1 = 2$ ✓ “
4. $x + y = z$ “

Propositional Variables : p, q, r, s, \dots

Truth Values : True \rightarrow T ✓
False \rightarrow F ✓

Propositional Calculus : The area of logic that deals with propositions is called as propositional logic (or) propositional calculus.

Compound Propositions : Many Mathematical statements are constructed by combining one or propositions by using logical operators, known as compound propositions.

No. of prop.

No. of poss.

1 \rightarrow 2

2 \rightarrow 4

3 \rightarrow 8

$n \rightarrow 2^n$



Logical Operators:

\wedge (and)	-	<u>Conjunction</u>
\vee (or)	-	<u>Disjunction</u>
\rightarrow (Implies)	-	<u>Implication</u>
\leftrightarrow (Bi-implies)	-	<u>Bi-conditional</u>

Negation:

(p) Mumbai is the capital of Maharashtra ✓

$(\sim p)$: Mumbai is NOT the capital of Maharashtra ✓

Conjunction: when two or more simple propositions are combined with logically connective "and (\wedge)" then the obtained new proposition is known as conjunction

p	q	$p \wedge q$
\checkmark T	\checkmark T	T
T	F	F
F	T	F
F	F	F

Truth Table: Truth table for compound proposition

P	q	$P \wedge q$	$P \vee q$	$P \rightarrow q$	$P \leftrightarrow q$	$q \rightarrow P$	$P \downarrow q$	$P \uparrow q$	$P \oplus q$
T	T	T	T	T	T	T	F	F	F
T	F	F	T	F	F	T	F	T	T
F	T	F	T	T	F	F	F	T	T
F	F	F	F	T	T	T	T	T	F

NOR

NAND

Boolean Sum

Construct Truth Tables of the following compound proposition:

$$\textcircled{1} \sim (p \wedge q) \vee (\sim p \rightarrow \sim q) = TTTT \quad \checkmark \text{ satisfiable}$$

$$\begin{aligned} \textcircled{2} \sim (p \vee q) &= FFFT \quad \checkmark \\ \textcircled{3} \sim p \wedge \sim q &= FFFT \quad \checkmark \end{aligned} \quad \left. \begin{array}{l} \\ \end{array} \right\} \begin{array}{l} \text{contingency} \\ \text{satisfiable.} \end{array}$$

$$\begin{aligned} \textcircled{4} \sim p \vee q &= TF TT \quad \checkmark \\ \textcircled{5} \sim (\sim p) &= TF \quad \checkmark \end{aligned} \quad \left. \begin{array}{l} \\ \end{array} \right\} \text{Contingency}$$

$$\textcircled{6} (\sim p \wedge \sim q) \vee (p \wedge q) = TF FT \quad \checkmark$$

$$\textcircled{7} (\sim p \vee q) \wedge (p \vee \sim q) = TF FT \quad \checkmark$$

$$\textcircled{8} p \vee (q \rightarrow r) = TTTTTF TT \quad \checkmark$$

$$\textcircled{9} p \vee \sim p = TT = \text{always True}$$

$$\textcircled{10} p \wedge \sim p = FF = \text{always False}$$

$$\sim (p \vee q) \equiv \sim p \wedge \sim q$$

$$p \rightarrow q \equiv \sim p \vee q$$

$$\sim (\sim p) \equiv p$$

$$p \leftrightarrow q \equiv (\sim p \wedge \sim q) \vee (p \wedge q) \equiv (\sim p \vee q) \wedge (p \vee \sim q)$$

⑧ $p \vee (q \rightarrow r)$

p	q	r	$q \rightarrow r$	$p \vee (q \rightarrow r)$
T	T	T		
T	T	F		
T	F	T		
T	F	F		
F	T	T		
F	T	F		
F	F	T		
F	F	F		



$$\sim q \rightarrow \sim p$$

						$X \leftrightarrow Y$
P	q	$P \rightarrow q$	$\sim q$	$\sim p$	$\sim q \rightarrow \sim p$	$(P \rightarrow q) \leftrightarrow (\sim q \rightarrow \sim p)$
T	T	T ✓	F	F	T ✓	T ✓
T	F	F ✓	T ✓	F ✓	F ✓	F ✓
F	T	T ✓	F	T	T ✓	T ✓
F	F	T ✓	T	T	T ✓	T

$$(P \rightarrow q) \equiv (\sim q \rightarrow \sim p)$$

Logical Equivalence: Compound propositions that have the same truth values in all possible cases are called logically equivalent.

The compound propositions \underline{p} and \underline{q} are called logically equivalent if $p \leftrightarrow q$ is a tautology

$X \cong Y$ if and only if $\overset{X}{p} \leftrightarrow \overset{Y}{q}$ Tautology $p \cong q$

$p \leftrightarrow q$ is a Tautology
 $p \cong q$

GATE: The binary equation \square is defined as follows

p	q	$p \square q$
T	T	T
T	F	T
F	T	F
F	F	T

Then compound proposition $(p \vee q) \Leftrightarrow$ \simeq

a) $\sim q \square \sim p$

b) $p \square \sim q$

c) $\sim q \square p$

d) $p \square q$

$p \vee q \simeq$

- a)
- b)
- c)
- d)

Δ $p \odot q$ $p \ominus q$
 $p * q$

* P	q	$P \sqcap q$	* $P \vee q$	* $\sim q$	$\sim P$	$\sim q \sqcap \sim P$	$P \sqcup \sim q$
T✓	T	T	T✓	F	F	T✓	T
T✓	F	T	T✓	T	F	T✓	T
(F)✓	(T)	(F)	T✓	(F)	(T)	F✓	T
(F)✓	F	T	F✓	(T)	T	T✓	F

$$P \vee q \equiv P \sqcup \sim q$$

Tautology, Contradiction, Contingency, Satisfiable and Unsatisfiable:

Always TRUE = Tautology

Always False = Contradiction

Neither Tautology
Nor contradiction } = Contingency

at least one TRUE = Satisfiable

NOT Satisfiable = unsatisfiable.

* Every contingency is satisfiable,
But converse need not be TRUE



Eg: Which of the following is NOT a Tautology

a) $\sim (p \rightarrow q) \rightarrow p$

c) $[\sim p \wedge (p \vee q)] \rightarrow q$

b) $\sim (p \rightarrow q) \rightarrow \sim q$

d) $(p \rightarrow q) \wedge q \rightarrow p$

I. Truth table

II. Logical Approach

III. Properties (or) laws