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Statement for Linked Answer Questions 21 & 22

For the grammar below, a partial LL(1) parsing table is also presented along with the grammar. Entries that need to be filled are indicated as E1, E2 and E3. ϵ is the empty string, \$ indicates end of input, and | separates alternate right hand sides of productions.

(GATE - 12)

(4 marks)

Grammar:

$$S \rightarrow aAbB \mid bAaB \mid \epsilon$$

$$A \rightarrow S$$

$$B \rightarrow S$$

Partial LL(1) parsing table:

	a	b	\$
S	(E1)	(E2)	$S \rightarrow \epsilon$
A	$A \rightarrow S$	$A \rightarrow S$	error
B	$B \rightarrow S$	$B \rightarrow S$	(E3)

21. The FIRST and FOLLOW sets for the non-terminals A and B are

22. The appropriate entries for E1, E2, and E3 are

Handwritten notes:

First(A) = First(B) = {a, b, ϵ }, Follow(A) = {a, b}, Follow(B) = {a, b}

Handwritten LL(1) parsing table:

	First	Follow
$S \rightarrow aAbB \mid bAaB \mid \epsilon$	a, b, ϵ	\$, a, b
$A \rightarrow S$	a, b, ϵ	a, b
$B \rightarrow S$	a, b, ϵ	\$, a, b

Handwritten definitions for E1, E2, and E3:

$$\begin{cases} E1: S \rightarrow aAbB, S \rightarrow \epsilon \\ E2: S \rightarrow bAaB, S \rightarrow \epsilon \\ E3: B \rightarrow S \end{cases}$$

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Q. Consider the following context-free grammar where the set of terminals is {a, b, c, d, f} [GATE-2021]

$$S \rightarrow d a T \mid R f$$

$$T \rightarrow a S \mid b a T \mid \epsilon$$

$$R \rightarrow c a T R \mid \epsilon$$

The following is a partially-filled LL(1) parsing table

	a	b	c	d	f	\$
S			(1)	$S \rightarrow d a T$	(2)	
T	$T \rightarrow a S$	$T \rightarrow b a T$	(3)		$T \rightarrow \epsilon$	(4)
R			$T \rightarrow c a T R$		$T \rightarrow \epsilon$	

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Which one of the following choices represents the correct combination for the numbered cells in the parsing table ("blank" denotes that the corresponding cell is empty)?

(a) ① $S \rightarrow Rf$ ② blank ③ blank ④ $T \rightarrow \epsilon$

(b) ① blank ② $S \rightarrow Rf$ ③ blank ④ blank

(c) ① blank ② $S \rightarrow Rf$ ③ $T \rightarrow \epsilon$ ④ $T \rightarrow \epsilon$

(d) ① $S \rightarrow Rf$ ② $S \rightarrow Rf$ ③ $T \rightarrow \epsilon$ ④ $T \rightarrow \epsilon$

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24. Consider the following grammar

$\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$

$\langle \text{rest} \rangle \rightarrow * \langle \text{expression} \rangle \mid \epsilon$

$\langle \text{factor} \rangle \rightarrow \text{identifier}$

$E \rightarrow FR$

$R \rightarrow *E/\epsilon$

$F \rightarrow id$

$m[E, id] = E \rightarrow FR$

$m[R, *] = R \rightarrow *E$

In the predictive parsing table M of the grammar the entries for M [$\langle \text{expression} \rangle$, identifier] and M [$\langle \text{rest} \rangle$, *] respectively are

☒ [$\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$] and [$\langle \text{rest} \rangle \rightarrow \epsilon$]

☒ [$\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$] and []

☒ [$\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$] and [$\langle \text{rest} \rangle \rightarrow * \langle \text{expression} \rangle$]

☒ [$\langle \text{factor} \rangle \rightarrow \text{identifier}$] and [$\langle \text{rest} \rangle \rightarrow \epsilon$]

25. Consider the grammar

$S \rightarrow a \mid ab \mid abc$

Choose the correct statement from the following:

(a) The grammar is LL(1)

(a) G2 generates the same language as G1

(b) G2 is unambiguous but not LL(1) as it is left recursive

(c) G2 gives a higher precedence

(d) None

28. Choose the true statement?

(a) If a grammar is left recursive it cannot be LL(1)

(b) If a grammar is right recursive it cannot be LL(1)

(c) An ambiguous grammar can sometimes be LL(1)

(d) If a grammar is not context free then it sometimes can be LL(1)

29. $S \rightarrow Aa$

$A \rightarrow b/\epsilon$

If we construct a predictive parse table for the above grammar, the production $A \rightarrow \epsilon$ is added in A's row and

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~~(a)~~ $\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$ and $\langle \text{rest} \rangle \rightarrow \epsilon$
 (a) $\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$ and $\langle \text{rest} \rangle \rightarrow * \langle \text{expression} \rangle$
~~(d)~~ $\langle \text{factor} \rangle \rightarrow \text{identifier}$ and $\langle \text{rest} \rangle \rightarrow \epsilon$

25. Consider the grammar
 $S \rightarrow a \mid ab \mid abc$
 Choose the correct statement from the following:
 (a) The grammar is LL(1)
 (b) The grammar is LL(2)
 (c) The grammar is LL(3)
 (d) None of the above

29. $S \rightarrow Aa$
 $A \rightarrow b/\epsilon$
 If we construct a predictive parse table for the above grammar, the production $A \rightarrow \epsilon$ is added in 'A' row and _____.
 (a) '\$' column (b) 'ε' column
 (c) 'a' column (d) 'b' column

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$S \rightarrow AaC/Bd$
 $A \rightarrow BC$
 $B \rightarrow bB/C$
 $C \rightarrow accS$

For which non terminals, its follow set contains terminal 'a'.

(a) {C} (b) {A}
 (c) {A,B,C} (d) {A,B,C,S}

24. Consider the following grammar
 $\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$
 $\langle \text{rest} \rangle \rightarrow * \langle \text{expression} \rangle \mid \epsilon$
 $\langle \text{factor} \rangle \rightarrow \text{identifier}$

$m[E, id] = E \rightarrow FR$
 $m[R, *] = R \rightarrow *E$

In the predictive parsing table M of the grammar the entries for M [$\langle \text{expression} \rangle$, identifier] and M [$\langle \text{rest} \rangle$, *] respectively are

~~(a)~~ $\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$ and $\langle \text{rest} \rangle \rightarrow \epsilon$
~~(b)~~ $\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$ and $\langle \text{rest} \rangle \rightarrow *$
 (c) $\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$ and $\langle \text{rest} \rangle \rightarrow *$
 (d) $\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$ and $\langle \text{rest} \rangle \rightarrow \epsilon$

27. Consider the grammar, G1
 $E \rightarrow E + E \mid E * E \mid (E) \mid id$ Ambiguous
 Consider the grammar G2,
 $E \rightarrow E + T \mid T$
 $T \rightarrow T * F \mid F$
 $F \rightarrow (E) \mid id$ Unambiguous.
 Choose the incorrect statement?
 (a) G2 generates the same language as G1 (T)
 (b) G2 is unambiguous but not LL(1) as it is left recursive (T)
 (c) G2 gives a higher precedence to + over *
 (d) None $+ > *$

28. Choose the true statement?
 (a) If a grammar is left recursive it cannot be LL(1)
 (b) If a grammar is right recursive it cannot be LL(1)
 (c) An ambiguous grammar can sometimes be LL(1)
 (d) None

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24. Consider the following grammar

$$\begin{aligned} \langle \text{expression} \rangle &\rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle \\ \langle \text{rest} \rangle &\rightarrow * \langle \text{expression} \rangle \mid \epsilon \\ \langle \text{factor} \rangle &\rightarrow \text{identifier} \end{aligned}$$

$m[E, id] = E \rightarrow FR$
 $m[R, *] = R \rightarrow *E$

In the predictive parsing table M of the grammar the entries for M [$\langle \text{expression} \rangle$, identifier] and M [$\langle \text{rest} \rangle$, *] respectively are

(a) [$\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$] and [$\langle \text{rest} \rangle \rightarrow \epsilon$]
 (b) [$\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$] and [$\langle \text{rest} \rangle \rightarrow * \langle \text{expression} \rangle$]
 (c) [$\langle \text{expression} \rangle \rightarrow \langle \text{factor} \rangle \langle \text{rest} \rangle$] and [$\langle \text{rest} \rangle \rightarrow \text{identifier}$]
 (d) [$\langle \text{factor} \rangle \rightarrow \text{identifier}$] and [$\langle \text{rest} \rangle \rightarrow \epsilon$]

25. Consider the grammar

$$S \rightarrow g \mid gb \mid abc$$

Choose the correct statement from the following:

(a) The grammar is LL(1)
 (b) The grammar is LL(2)

26. G2 generates the same language as G1
 (a) G2 is unambiguous but not LL(1) as it is left recursive
 (b) G2 gives a higher precedence to + over *
 (c) None

27. Choose the true statement?

(a) If a grammar is left recursive it cannot be LL(1)
 (b) If a grammar is right recursive it cannot be LL(1)
 (c) An ambiguous grammar can sometimes be LL(1)
 (d) If a grammar is not context free then it sometimes can be LL(1)

28. $S \rightarrow Aa$
 $A \rightarrow b/\epsilon$
 If we construct a predictive parse table for the above grammar, the production $A \rightarrow \epsilon$ is added in 'A' row and _____.

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7 Objective Practice Questions

30. $S \rightarrow aSbS/bSaS/\epsilon$
 In the predictive parse table M of the above grammar M [S, a] = $S \rightarrow aSbS, S \rightarrow \epsilon$
 (a) $S \rightarrow aSbS$ (b) $S \rightarrow bSaS$
 (c) $S \rightarrow \epsilon$ (d) $S \rightarrow aSbS, S \rightarrow \epsilon$

31. $S \rightarrow Aa/Bb$
 $A \rightarrow b$
 $B \rightarrow b$ is
 (a) LL(1) (b) LL(2)
 (c) LR(0) (d) None

32. Consider a calculator modeled by the grammar

$$\langle \text{accumulated_sum} \rangle \rightarrow \langle \text{accumulated_sum} \rangle + \text{number}$$

33. Which of the following is **not** an operator grammar

(a) $S \rightarrow AaB$ (b) $S \rightarrow AaB$
 $A \rightarrow aA/b$ $A \rightarrow a/b$
 $B \rightarrow bB/a$
 (c) $S \rightarrow AaB$ (d) None
 $A \rightarrow aA/b$
 $B \rightarrow bB/\epsilon$

34. What is equivalent operator grammar for the following grammar?

$$S \rightarrow AB, A \rightarrow c/d, B \rightarrow aAB/d$$

(a) $S \rightarrow Aa/Ab$ (b) $S \rightarrow AaAB/Ad$
 $A \rightarrow c/d$ $A \rightarrow c/d$
 $B \rightarrow aAB/d$
 (c) $S \rightarrow AaS/Ab$ (d) None

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08. Consider the grammar $E \rightarrow E + n \mid E \times n \mid n$ for a sentence $n + n \times n$, the handles in the right-sentential form of the reduction are (GATE - 05)

(a) $n, E + n$ and $E + n \times n$
 (b) $n, E + n$ and $E + E \times n$
 (c) $n, n + n$ and $n + n \times n$
 (d) $n, E + n$ and $E \times n$

Stack	Input	Actions
\$	$n + n \times n$	Shift
$\$ \circledast$	$+ n \times n$	$E \rightarrow n$ ✓
$\$ E$	$+ n \times n$	Shift
$\$ E +$	$n \times n$	Shift
$\$ E + \circledast$	$\times n$	$E \rightarrow E + n$ ✓
$\$ E$	n	Shift
$\$ E \times$	n	Shift
$\$ E \times \circledast$	n	$E \rightarrow E \times n$ ✓
$\$ E$	n	Shift
$\$ E$	n	Accept

Handles: $n, E + n, E \times n$

Diagram: A tree structure showing the derivation of $n + n \times n$. The root node is E , which has children E and n . The left E has children E and $+$. The leftmost E has child n . The right E has children E and \times . The rightmost E has child n .

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08. Consider the grammar $E \rightarrow E + n \mid E \times n \mid n$ for a sentence $n + n \times n$, the handles in the right-sentential form of the reduction are (GATE - 05)

(a) $n, E + n$ and $E + n \times n$
 (b) $n, E + n$ and $E + E \times n$
 (c) $n, n + n$ and $n + n \times n$
 (d) $n, E + n$ and $E \times n$

Stack	Input	Actions
\$	$n + n \times n$	Shift
$\$ \circledast$	$+ n \times n$	$E \rightarrow n$ ✓
$\$ E$	$+ n \times n$	Shift
$\$ E +$	$n \times n$	Shift
$\$ E + \circledast$	$\times n$	$E \rightarrow E + n$ ✓
$\$ E$	n	Shift
$\$ E \times$	n	Shift
$\$ E \times \circledast$	n	$E \rightarrow E \times n$ ✓
$\$ E$	n	Shift
$\$ E$	n	Accept

Handles: $n, E + n, E \times n$

Diagram: A tree structure showing the derivation of $n + n \times n$. The root node is E , which has children E and n . The left E has children E and $+$. The leftmost E has child n . The right E has children E and \times . The rightmost E has child n .

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sum>* number
(c) number, number + number, number + number * number
(d) number, <accumulated_sum> + number, <accumulated_sum>* number

33. Which of the following is not an operator grammar?
~~I. $S \rightarrow AB$~~
~~II. $S \rightarrow a$~~
~~III. $S \rightarrow a \cup b$~~
~~IV. $S \rightarrow a \cup b \cup c$~~
 (a) only I
 (b) I, III and IV
 (c) I and II
 (d) None

37. $S \rightarrow AgBb$
 $A \rightarrow Bc/d$
 $B \rightarrow e$
 In the operator grammar above lead (S) = _____
 (a) {d} (b) {e,d}
 (c) {c,d,e} (d) {a,c,d,e}

statement in the operator precedence table constructed for the grammar the relations for the ordered pairs $(\wedge, \wedge), (-, -), (+, +), (*, *)$ are
 (a) all < (b) all >
 (c) <, >, =, < (d) <, >, >, >

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$A \rightarrow b$
 $B \rightarrow b$ is
~~LL(1)~~
~~LR(0)~~
 (d) None

32. Consider a calculator modeled by the grammar
 $\langle \text{accumulated_sum} \rangle \rightarrow \langle \text{accumulated_sum} \rangle + \text{number}$
 $\mid \langle \text{accumulated_sum} \rangle * \text{number}$
 $\mid \text{number}$
 For an input "number + number * number", the handles in the reverse of a rightmost derivation are
 (a) number, <accumulated_sum> + number, <accumulated_sum> + number * number
 (b) number, <accumulated_sum> + number, <accumulated_sum> + <accumulated_sum> * number

35. What is equivalent operator grammar for the following grammar?
 $S \rightarrow AB$, $A \rightarrow c/d$, $B \rightarrow aAB/d$
~~(a) $S \rightarrow Ag/Ab$~~
 $A \rightarrow c/d$
 $B \rightarrow aS/b$
 (c) $S \rightarrow AgS/Ab$
 $A \rightarrow c/d$
 $B \rightarrow aS/b$
 (d) None
 $S \rightarrow AaS/Ad$
 $A \rightarrow c/d$
 $B \rightarrow aS/d$

36. Consider the grammar given below
 $E \rightarrow E+E \mid E^*E \mid E-E \mid E/E \mid E\wedge E \mid (E) \mid id$
 Assume that + and - have the same but least precedence, * and / have the next higher precedence but the same precedence and finally \wedge has the highest precedence. Assume + and - associate to the left like * and / and that \wedge associates to the right. Choose the correct statement in the operator precedence table

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For an input "number + number * number", the handles in the reverse of a rightmost derivation are

(a) number, <accumulated_sum> + number, <accumulated_sum> + number * number

(b) number, <accumulated_sum> + number, <accumulated_sum> + <accumulated_sum> * number

(c) number, number + number, number + number * number

(d) number, <accumulated_sum> + number, <accumulated_sum> * number

33. Which of the following is not an operator grammar?

~~I. $S \rightarrow AB$~~ II. $S \rightarrow A \cup B$

III. $S \rightarrow a \cup$ ~~IV. $S \rightarrow \epsilon$~~

36. Consider the grammar given below

$$E \rightarrow E+E \mid E^*E \mid E-E \mid E/E \mid E^{\wedge}E \mid (E) \mid id$$

Assume that + and - have the same but least precedence. * and / have the next higher precedence but the same precedence and finally \wedge has the highest precedence. Assume + and - associate to the left like * and / and that \wedge associates to the right. Choose the correct statement in the operator precedence table constructed for the grammar the relations for the ordered pairs (\wedge, \wedge) , $(-, -)$, $(+, +)$, $(*, *)$ are

(a) all < (b) all >

(c) <, >, =, < (d) <, >, >, >

37. $S \rightarrow A \cup B$
 $A \rightarrow Bc/d$
 $B \rightarrow e$
 In the operator grammar above lead
 $(S) = \underline{\hspace{2cm}}$

Handwritten notes:

$B \rightarrow aS/d$

$(+ = -) < (* = /)$

$< \wedge$

\wedge associates to the right

Operator Precedence Table:

	\wedge	$-$	$+$	$*$
\wedge	$>$	$>$	$>$	$>$
$-$	$<$	$=$	$<$	$<$
$+$	$<$	$<$	$=$	$<$
$*$	$<$	$<$	$<$	$=$

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38. For the grammar

$$E \rightarrow E+T \mid T$$

$$T \rightarrow T^*F \mid F$$

$$F \rightarrow (E) \mid id$$

Trailing symbols of E are $\{+, *, id,)\}$

(a) *, +, (.id (b) *, +, ., id

(c) *, (.id (d) *, ., id

39. Consider the following operator grammar

$$E \rightarrow E+T \mid T$$

$$T \rightarrow F \uparrow T \mid F$$

$$F \rightarrow id$$

(a) $\uparrow > +$ (b) $\uparrow > +$

(c) $+ > \uparrow$ (d) $+ < +$

40. $S \rightarrow aSb \mid Ac$

43. Choose the false statement

(a) An ambiguous grammar can never be LL(1)

(b) An ambiguous grammar can never be LR(0)

(c) An ambiguous grammar can never be SLR(1)

(d) An ambiguous grammar cannot have a shift reduce parser

44. Choose the false statement

(a) There exists a grammar that is LR(0) but not LL(1)

(b) There exists a grammar that is LL(1) but not LR(0)

(c) There exists a grammar that is LL(1) but not SLR(1)

(d) There exists an ambiguous grammar that is either LL(1) or LR(0)

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