

## Practical No. 01

Aim: Write a Lex program to recognize valid strings using regular expression:  $(ab)^*abb$

Test Cases :

Valid strings: abababb, abb, bbbabb

Invalid strings: baba, bba, bbb, ababba

Tools : S.N.	Software/Hardware	Specification	Quantity
1	FLEX Compiler	Win10, 64bit, .exe	-
2	Computer System	Win10 or above, 500GB HDD/250GB SSD i3 processor	1

Theory: Regular Expression :

Every regular set can be expressed by expression called as Regular expression.

Operator	Priority	Meaning
?	1	At most one occurrences of preceding R.E
*	2	Zero or more occurrences of preceding R.E
.	3	Concatenation
+ OR	4	Union

Rules: If  $\Sigma$  is a regular expression, &  $L(\epsilon)$  is  $\{\epsilon\}$ , i.e language whose sole member is empty string.

If  $a$  is symbol in  $\Sigma$ , then  $a$  is regular expression, &  $L(a) = \{a\}$

If  $m$  &  $s$  are two R.E then  $L(m) \cup L(s)$ ,  $L(m)s$  is R.E

$(m)^*$  is regular-expression  $L(m)(s)$  is R.E

Results:

```
abababb  
Valid string: abababb  
abb  
Valid string: abb  
bbbabb  
Valid string: bbbabb  
baba  
bba  
bbb  
ababba
```

Program:

```
% {  
    #include <stdio.h>
```

```
% %
```

```
%%
```

```
(abb)*(abb) { printf("String is Valid %s", yytext); }
```

```
%%
```

```
int main()
```

```
{
```

```
    yylex();  
    return 0;
```

```
}
```



Error & Remedial action:

error: could not locate lex.bak

action: reinstalled the flex software

Result:

|

,

Conclusion: LEX program to recognize given valid string implemented successfully & all test cases verified

Course Outcome Attained & Mapping with Program Outcome:

Course Outcome

Program Outcome

C0-L: Generate scanner &

2. Problem analysis

parser from formal  
specifications

3. Design Development of solution

5. Modern tool usage

## Practical No. 02

Aim: Write a LEX program to recognize tokens 1) Identifier

2) Integer constant (Decimal, Octal, Hexadecimal)

3) Real constants (Floating point and Exponential Format)

Test Cases:

a. Integer Constant Ex +15, -5, 0123, 0x1A

b. Floating Point Constant Ex. +0.5, 230.0, .567, -28.89

c. Exponential format Ex. 1.6e-19, -0.5E+2, +1.7E4, 2.6e-7

Tools:	Sn.	Tools	Specification	Quantity
	1.	LEX Compiler	- Windows, 64 bit	-
	2.	Computer System	- Windows 10+, 500/1GB/4GB - 01	

Theory: a) Regular E for decimal constant

Rules: 1. must have at least one digit

2. May be true or -ve  
3. No spaces, comma, period(.) are allowed.

$$R.E = (+|-)?(digit)^+$$

b) R.E for Real const. (floating point notation.)

1. Must have at least one digit 2. may be true or -ve

3. Must contain decimal point 4. at least one digit after decimal point

5. No (-)(,)(.)

$$R.E: (+|-)?(digit)^*.(digit)^+$$

c) Regular Expression for Real constant (Exponential Notation)

1. Mantissa & exponent separated by 'E' or 'e'

2. Mantissa may be real or integer. 3. true or -ve

4. Integer Exp. 5) Exponent may true or -ve

6. No (-)(,)(.)

$$R.E: (+|-)?(digit)^*.(digit)^+(E|e)(+|-)?(digit)^+$$

Program :

```
%{ #include <stdio.h> %}
Digit [0-9] +
%.
(+|-)?(Digit)+ { printf("Integer/Decimal"); }
(+|-)?(Digit)*.(Digit)+ { printf(" Floating point"); }
(+|-)?(Digit)*.(Digit)+(E|e)(+|-)?(Digit) { printf("Exponential real"); }
(+|-)?(0)(0-7)+ { printf(" Octal Number"); }
(+|-)?(0x10x)(0-9|A-Fa-f)+ { printf("Hexadecimal"); }

%.
int yscanf()
{
    return 1;
}

int main()
{
    yscanf();
    return 0;
}
```

Output / Result

+15	Integer/Decimal
-5	
X	Floating point real const
0123	
X	Octal Number
0x1A	
X	Exponential real constant
+0.5	
X	Floating point real const
236.0	
X	Floating point real const
567	
X	Floating point real const
-26.89	
X	Floating point real const
1.6e-19	
X	Exponential real constant
-0.5E+2	
X	Exponential real constant

## Error & remedial action :

error: Integer underflow: occurs when input program contains -ve integer constant i.e. too large in absolute value to be represented as decimal integer.

Remedial: we can check length of input string before converting it to an integer & raise an error if length for given data type

## Conclusion :

Lex program to recognize tokens  
 1] Identifier  
 2] Integer const. → Recd const. implemented successfully

## Course Outcome Attained & mapping with Program O/c !

Course Outcome	Program Outcome
CO-1: Generate & parser from formal specification	PO-1, PO-2, PO-3, PO-4, PO-5

✓  
 15/09/23  
 02/09/23

## Practical No. 03

Aim: Source code in a text file . constants used in the program are strings. write a program to read integer constants as a string and convert it to decimal Number. Integer constants are represented in Decimal, Octal and Hexadecimal. It may be +ve or -ve.

Test Cases: a. Decimal const: +15, -3, 123.

b. octal const. +0123, 0117, -0777

c. Exponential Format ex: 1.6e-9, -0.5E+2, +1.7E4

Tools:

Sr.	Tools	Specification	Quantity
1	Lex Compiler (FLEX)	- FLEX, wind10,	- --
2	Computer System	- Orbit Windows 10 OS, i3 processor, 4GB RAM	- 01

Theory: In this experiment we are defining rules to match integer constants that are represented in decimal, octal & hexadecimal formats. program reads i/p program character by character & matches it against defined rules

- When match is found the program performs an action, which is in case is to convert matched integer constant to a decimal number and print it to the output
- we define h-3 helper funn to convert decimal, octal, and hexadecimal strings to decimal integers

- 1) First fun will convert any octal no. to decimal by applying applicable method to convert
- 2) Second fun will convert any Hexadecimal, floating into decimal.

## Program:

```
% {
```

```
#include <stdio.h>
```

```
#include <string.h>
```

```
int strToIntDec(const char *str) { int result = 0;
```

```
    int len = strlen(str); for (int i = 0; i < len; i++) { result = result * 10 + (str[i] - '0');
```

```
    return result;
```

```
}
```

```
int octToIntDec(const char *str) { int result = 0, i; int len = strlen(str);
```

```
    for (int i = 0; i < len; i++) { result = result * 8 + (str[i] - '0');
```

```
}
```

```
return result;
```

```
}
```

```
int hexToIntDec(const char *str) { int result = 0, i; int len = strlen(str);
```

```
    for (int i = 0; i < len; i++) {
```

```
        if (str[i] >= '0' & str[i] <= '9') { result = result * 16 + (str[i] - '0');
```

```
        } else if (str[i] >= 'a' & str[i] <= 'f') {
```

```
            result = result * 16 + (str[i] - 'a' + 10);
```

```
        } else if (str[i] >= 'A' & str[i] <= 'F') {
```

```
            result = result * 16 + (str[i] - 'A' + 10);
```

```
}
```

```
return result;
```

```
}
```

```
%
```

```
[-] ? [0-9] + {printf("%d\n", strToIntDec(44text));}
```

```
[-] ? [0-7] + {printf("%d\n", octToIntDec(44text+1));}
```

```
[-] ? [0-7][0-9][0-9][0-9] + {printf("%d\n", HexToIntDec(44text+1));}
```

```
%
```

```
int main()
```

```
44lex();
```

```
return 0;
```

```
}
```

Error or remedial action :

error : lexical error : occurs when ip program contains characters that do not match any of defined rules. lexer will not able to identify input as valid integer const.

remedial : added `printf("Invalid input\n");` statement.

Result :

+15
15
-5
-5
123
123
+0123
83
0117
79

-0777
-511
1.6e-19
1.6e-19
-0.5E+2
-50.0
+1.7E4
17000.0
.26e-7
2.6e-08

Conclusion : Lex program to convert input string constant that could be string, octal or hexadecimal implemented successfully

Course Outcome and mapping with program outcome :

Course Outcome	Program Outcome
CO-1: generate scanner and parser from formal specs	PO1, PO2, PO3, PO5, PO4

## Practical No. 04

Aim: Write a program to eliminate left recursion from the grammar. Test Cases:

1.  $B \rightarrow E \mid E - E \mid id$
2.  $S \rightarrow (L) \mid a$   
 $L \rightarrow L, S \mid S$

Tools :

Sn.	Tools	Specification	Quantity
1	- Lexcompiler, Compiler - any GDB compiler -	-	--
2	- Computer System	- min 150GB HDD/GARRAM - 01	

### Theory:

→ Steps to eliminate left recursion from grammar.

1. Identify left-recursive productions: look for prod in grammar that have form  $A \rightarrow A \alpha \beta$ , where  $A$  is nonterminal symbol &  $\alpha \beta$  are strings of terminals & non-terminals

2. Eliminate direct-left recursion: to eliminate, you can replace left-recursive production  $A \rightarrow A \alpha \beta$  with a new production  $A \rightarrow \beta A'$ , where  $A'$  is a new nonterminal symbol. Then add new productions  $A' \rightarrow \alpha A' \mid \epsilon$

3. Eliminate Indirect-left recursion: where  $A \rightarrow B \alpha$  and  $B \rightarrow A \beta$

can use same approach as for direct recursion. Replace  $B$  in  $A \rightarrow B \alpha$  & its new nonterminal symbol say  $B'$ .

4. Repeat step 2 & 3 as necessary: if grammar still contains left recursion repeat process

5. Reorder production: after eliminating left recursion, you may need to reorder productions in grammar to ensure they are in correct order of parsing.

## Program

```
#include <stdio.h>
#include <string.h>
#define MAX_RULES 100
#define MAX_SYMBOLS 100
char rules[MAX_RULES][MAX_SYMBOLS];
int nRules = 0;
void eliminateLr (char nonterm) {
    int i, j, k;
    char newnt = nonterm + 2;
    for (int i = 0; i < nRules; i++) {
        if (rules[i][0] == nonterm) {
            char alpha[MAX_SYMBOLS], beta[MAX_SYMBOLS];
            int alphalen = 0, betalen = 0;
            for (k = 2; rules[i][k] != '\0'; k++) {
                alpha[alphalen++] = rules[i][k];
            }
            alpha[alphalen] = '\0';
            for (k = 2; rules[i][k] != '\0'; k++) {
                betaln[k - alphalen] = rules[i][k];
            }
            betaln[betaln] = '\0';
            strcpy(rules[i], " ");
            strcpy(rules[i], alpha);
            printf("Rules[%d] : %s %s %s\n", i, nonterm, beta, newnt);
            printf("Rules[%d] : %s %s %s | epsilon\n", i, newnonterminal, alpha, newnonterminal);
            rules[i] = '\0';
            else printf("can't be reduced\n");
        }
    }
    return;
}
int main() {
    // string l; char input[100]; l[50], n[50];
    printf("Enter the productions : ");
    scanf("%s -> %s", l, n);
    input[1] = l[0] + " " + n[0];
    eliminateLr(input[1]);
    printf("After eliminating Left Recursion are (%s)\n", eliminateLr(input));
}
```

### Error and remedial action :

error: Buffer overflow : program consumer may length of production rule & max no. of production rules are fixed.

Remedy: program could dynamically allocate memory for rules & symbols.

### Com. Results :

Enter the productions:  $E \rightarrow E+E|T$   
 The productions after eliminating Left  
 Recursion are:  
 $E \rightarrow EE'$   
 $E' \rightarrow TE'$   
 $E \rightarrow S$

Enter the productions:  $S : (L) | a$

$L : L, S | S$   
 The productions after eliminating Left  
 Recursion are:  
 $S \rightarrow (L)$   
 $| a$   
 $L \rightarrow (L) L'$   
 $| a L'$   
 $L' \rightarrow , S L'$   
 $| \epsilon$

Conclusion: C program to eliminate left recursion from grammar implemented successfully

### Course outcomes attained from matching program outcomes :

Course Outcomes	Program Outcomes
CO-1: generate parser from formal specification	PO-1, PO-2, PO-3, PO4, PO5
CO-2: generate top-down and bottom-up parser using predictive parsing SLR & LR passing techniques.	

Program :

```

#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <math.h>

int main() {
    char str[50];
    int base = 10;
    long long int num = 0;
    int i, len, exp;
    scanf("%s", str);
    len = strlen(str);
    if (str[0] == '-') return 0;
    if (len > 1 && str[0] == '0') {
        if (str[1] == '*' || str[1] <= 'X') {
            base = 16; i = 2;
        } else {
            base = 8;
            i = 1;
        }
    } else {
        i = 0;
    }
    while (i < len) {
        if (str[i] >= '0' && str[i] <= '9')
            num = num * base + (str[i] - '0');
        else if (str[i] >= 'a' && str[i] <= 'f') {
            num = num * base + (str[i] - 'a' + 10);
        } else if (str[i] >= 'A' && str[i] <= 'F') {
            num = num * base + (str[i] - 'A' + 10);
        }
    }
}

```

```

        else if (str[i] == 'e' || str[i] == 'E') {
            i++;
            exp = 0;
            while (i < len) {
                exp = exp * 10 + (str[i] - '0');
                i++;
            }
            exp = -exp;
        } else if (str[i] == '+') {
            exp = 0;
            i++;
            while (i < len) {
                exp = exp * 10 + (str[i] - '0');
                i++;
            }
        } else if (str[i] == '-') {
            exp = 0;
            i++;
            while (i < len) {
                exp = exp * 10 + (str[i] - '0');
                i++;
            }
        } else {
            printf("Invalid constant format");
            return 0;
        }
    }
    num *= num pow(10, exp);
    break;
} else {
    printf("Invalid");
    return 0;
}
printf("%d\n", num);
return 0;
}

```

$i \rightarrow (i+1) \rightarrow \dots \rightarrow (i+n)$   
 $88^0 = 88$        $i \rightarrow (i+1) \rightarrow \dots \rightarrow (i+n)$   
 $i(0) \rightarrow (i+1) \rightarrow \dots \rightarrow (i+n) = m$  ?  
 $i(1) \rightarrow (i+1) \rightarrow \dots \rightarrow (i+n) = 10(i+1) + m$  ?  
 $(01)^n \cdot (i+1) + 220d + m = m$   
 $'1' \rightarrow (i+1) \rightarrow \dots \rightarrow (i+n) = 10(i+1) + m$  ?  
 $i(01 + A) \rightarrow (i+1) \rightarrow \dots \rightarrow (i+n) = m$  ?

Program :

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#define MAX_RULES 100
#define MAX_PROD 100
#define MAX_SYMBOLS 20
typedef struct {
    char lhs[MAX_SYMBOLS];
    char rhs[MAX_PROD][MAX_SYMBOLS];
    int num_prods;
} Rule;
```

// nonterminal (uppercase letter)

```
int is_nt(char symbol) {
```

```
    return symbol >= 'A' && symbol <= 'Z'; }
```

```
void eliminate_nt(Rule rules[], int num_rules) {
```

```
    for (int i=0; i<num_rules; i++) {
```

```
        for (int j=0; j<rules[i].num_prods; j++) {
```

```
            char fs = rules[i].rhs[j][0];
```

```
            if (is_nt(fs) && fs == rules[i].lhs[0]) {
```

```
                char new_nt = rules[i].lhs[0] + 1;
```

```
                Rule new_rule = {num_prods = 0};
```

```
                strcpy(new_rule.lhs, &new_nt);
```

```
                for (int k=0; k<rules[i].num_prods; k++) {
```

```
                    char fs = rules[i].rhs[k][0];
```

```
                    if (is_nt(fs) && fs == rules[i].lhs[0]) {
```

```
                        new_rule.num_prods++;
```

```
} else {
```

```
    rules[i].num_prods++;
```

Teacher's Signature.....

```
    sprintf(rules[i].rhs[mrules[i].num-prods], "%c%c",  
           new_nt, '\n');  
    rules[i].num-prods++;
```

{ } { }

```
int main()
```

```
    Rule rules[ ];  
    printf("Enter the productions: ");
```

```
    scanf("%s",
```

```
    while(1){
```

```
        Rule rule;  
        memset(&rule, 0, sizeof(Rule));
```

```
        if (num_rules >= MAX_RULES) {
```

```
            printf("Too many rules");
```

```
            return 1; // Error
```

```
        }
```

```
        printf("y.d: ", num_rules+1);
```

```
        if (scanf("%s", rule.lhs) == EOF) break;
```

```
        strcpy(rule.rhs[rule.num-prods].prod-str,  
               MAX_LEN);
```

```
        rule.num-prods++;
```

```
    } // while loop
```

```
    eliminate_LH(rules, num_rules);
```

```
    printf("\n Rules after eliminating left recursion are: ");
```

```
    for (int i = 0; i < num_rules; i++) {
```

```
        printf("%s: ", rules[i].lhs);
```

```
        for (int j = 0; j < rules[i].num-prods; j++) {
```

```
            printf(" %s ", rules[i].rhs[j]);
```

```
            if (j < rules[i].num-prods - 1) {
```

```
                printf(" | ");
```

```
            }
```

```
        }
```

```
    }
```

```
    printf("\n");
```

```
}
```

Date.....

Page. No..... 01 .....

Expt. No. 05 Experiment No. 05

Aim: Design Recursive descent parser for the grammar.

$$E \rightarrow +EE | -EE | a | b$$
Tools:

Sn.	Tools	Specification	Caty.
1	C Compiler	Turbo, GCC, or any	-
2	Computer System	Windows, 4GB RAM & 1TB HDD	1

Theory:

A recursive descent parser is a top-down parsing technique where a parse tree is built from top & constructed down to leaves. This type of parser starts with highest level of grammar & recursively expands non-terminals until a terminal symbol is reached.

- To implement this grammar using recursive descent parser, we need to follow steps to implement parser:
  1. Analyze grammar, identify terminals, non-terminals production rules
  2. Write function to parse each non-terminal
  3. Use tokenizer to convert input string into list of tokens
  4. Pass list tokens as input, function should return whether string can be derived from grammar or not.

Teacher's Signature.....

Date.....

Expt. No.....

Date.....

Page. No.....

Program :

```

#include <stdio.h>
#include <stdbool.h>
#include <string.h>

bool parse_E(char *input, int *index);
bool parse_EE(char *input, int *index);
bool parse_EF(char *input, int *index) {
    if (input[*index] == 'a' || input[*index] == 'b') {
        (*index)++;
        return true;
    } else if (input[*index] == '+' || input[*index] == '-')
        return parse_EE(input, index) &&
parse_EF(input, index);
    }
    else { return false; } // end
}

bool parse_EE(char *input, int *index) {
    if (input[*index] == '+' || index == '-' ) {
        (*index)++;
        return parse_EF(input, index); }
    else { return false; }
}

int main() {
    char input[100];
    int EIndex = 0;
    printf("Enter an input string: \n");
    scanf("%s", input);
}

```

Teacher's Signature.....

```
bool success = parse_E(input, &index);
if (success && index == strlen(input))
{
    printf("The input string can be accepted
           using grammar\n");
}
else
{
    printf("The input string cannot be accepted
           from grammar.\n");
}
return 0;
}
```

Output / Result :

```
ab (ab) legal
; (ab, ab) legal
```

```
; (ab) legal
; (ab, ab) legal
; ; (ab) legal
; ; ; (ab) legal
```

```
; () legal
; (ab) legal
; ; (ab) legal
; ; ; (ab) legal
; ; ; ; (ab) legal
```

Date.....

Expt. No..... Date.....

Page. No.....

Error & remedial action :

1. Program assumes input string contains only symbols 'a' 'b' 't' '(-)', if contain any other symbol, program may produce error.

Remedial Action: program can be manipulated to check each symbol that is terminals.

Conclusion: The program to implement recursive descent parser for grammar designed successfully.

Course Outcome attained Q mapping with program outcomes:

Course Outcomes	Program Outcomes
CO-1	PO1, PO2, PO3
CO-2	PO4, PO5

Teacher's Signature.....

Date.....

Page. No.....

Expt. No. 06 Experiment No. 06

Aim: Design a program for basic calculator using YACC or BISON.

Tools :

Sn No.	Tools	Specs	Qty
01	YACC compiler	YACC or BISON	1
02	Comp. System	windows, LIBRARY 2TB HDD	1

Theory : YACC (Yet Another Compiler Compiler) & BISON are tools that generate parsers (programs analyze structure of text or code) based on formal desc of language grammar.  
 - commonly used in development of compilers, interpreters, & other language-based s/w.

Syntax :

```
%{ // declaration section
```

```
%}
```

```
%%
```

```
% token definition
```

```
grammar declaration
```

```
%%
```

```
int main() {
```

```
yyparse();
```

```
return 1;
```

```
}
```

Teacher's Signature.....

## Result?

• Cytogenetic analysis of patients  
showing skin rash that are  
suspected to have  
chromosomal abnormalities  
revealed normal karyotype in  
skin cells from patients  
with constitutional chromosomal

abnormalities

• Mitochondrial DNA  
abnormalities common

• Other factors  
• Age  
• Gender

Date.....

Expt. No.....

Page. No.....

Programs:

% \$ if include <stdio.h>

#include <lex.yy.c> int yylex();  
int yyerror();

% %

%def '+' '-' ;

%def '\*' '/' ;

expression :

1 expression '+' 2 expression { \$\$ = \$1 \* \$2 ;

printf ("\* : %d\n", \$\$); }

1 expression '+' 2 expression { \$\$ = \$1 + \$2 ;

printf ("+ : %d\n", \$\$); }

1 expression '-' 2 expression { \$\$ = \$1 - \$2 ;

printf ("- : %d\n", \$\$); }

1 expression '/' 2 expression { \$\$ = \$1 / \$2 ;

printf ("/ : %d\n", \$\$); }

% %

int main () {

yyparse();

return 1;

}

Error & Remedial Action:

1. Parsing error for invalid input expression  
yyerror() will generate error

Remedial action: print appropriate error message .

Teacher's Signature.....

Date.....

Expt. No..... 06

Page. No.....

Conclusion: The program for basic calculation using yacc or BISON designed successfully

Course Outcomes Attained & mapping with program outcomes:

Course Outcomes

CO1

CO3

Program Outcomes

PO1 PO2 PO3 PO4

PO5 , PO7

Teacher's Signature.....

Date.....

Expt. No. 07

Experiment No. 07

Page. No.....

Aim: Design a program to evaluate postfix expression.

Tools :

	Specs.	Qty.
1 - C compiler	- Turbo, Acc, any	- -
2 - Computer System	- Windows, 4GB RAM	- 01

Theory :

- Postfix Notation : way of writing arithmetic expressions in which each operator follows its Operands. Ex. "3 + 4"  $\Rightarrow$  "3 4 +"

- Steps to evaluate postfix expression: Ex. "3 4 +"

1. Scan Expression from left to right:

- Push 3 onto stack (stack: 3)
- Push 4 onto stack (stack: 3 4)
- pop 3 & 4 from stack, add them together ( $4 + 3 = 7$ ) & push result (7) back onto stack

2. return expression is in top element of stack, which is 7.

Teacher's Signature.....

Program :-

```

#include < stdio.h >
#include < ctype.h >
#define MAXSTACK 100
#define POSTFIXSIZE 100
int stack[MAXSTACK];
int top = -1;
void push(int item)
{
    if (top >= MAXSTACK - 1)
        printf("stack overflow");
    return;
    else
        top = top + 1;
        stack[top] = item;
}

int pop()
{
    int item;
    if (top < 0)
        printf("stack under flow");
    else
        item = stack[top];
        top = top - 1;
        return item;
}

void EvalPostfix(char postfix[])
{
    int i;
    char ch;
    int val;
    int A, B;
    for (i = 0; postfix[i] != '\0'; i++)
    {
        ch = postfix[i];
        if (isdigit(ch))
            push(ch - '0');
    }
}

```

```

else if(ch == '+' || ch == '-' || ch == '*' || ch == '/')
{
    A = pop();
    B = pop();
    switch(ch)
    {
        case '*': val = B * A; break;
        case '/': val = B / A; break;
        case '+': val = B + A; break;
        case '-': val = B - A; break;
    }
    push(val);
}
printf("\n Result of expression : %d \n", pop());
}

int main()
{
    int i;
    char postfix[POSTFIXSIZE];
    printf("Enter postfix expression, press right\n"
           "Parenthesis ')' for end expression : ");
    for(i=0; i < POSTFIXSIZE-1; i++)
    {
        scanf("%c", &postfix[i]);
        if(postfix[i] == ')') { break; }
    }
    EvalPostfix(postfix);
    return 0;
}

```

Result:

Date.....

Expt. No....07.....

Page. No.....09.....

Error & Remedial Action :

Error : Divide by zero

Remedial : need to apply proper error message  
for divide by zero.

Conclusion : The program to evaluate postfix expression designed successfully.

Course Outcomes attained & mapping with Program O/C.

CO's

CO-1

PO's

PO-1, PO-2, PO-3, PO-4

Teacher's Signature.....