

MENG INDIVIDUAL PROJECT

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Cryptocurrency Statistical Arbitrage on Decentralised Exchanges

Author:
Devam Savjani

Supervisor:
Dr. Thomas E. Lancaster

Second Marker:
Mr. Ivan Procaccini

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Abstract

SOME ABSTRACT

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Chapter 1

Introduction

1.1 Motivation

Since the introduction of Bitcoin, a peer-to-peer payment network and cryptocurrency, there have been countless new cryptocurrencies that are used and traded. Along with this, their high volatility has piqued a lot of retail investor's interest with some investors gaining a high return of interest and others losing a lot of money [1]. In addition to this, larger investment institutions have also sought to gain profits from this new type of tradable asset [2]. This has lead to more sophisticated forms of cryptocurrency trading.

However, trading cryptocurrencies was initially very difficult for people without technical know-how, and the first recorded transaction was on 12th October 2009 via a paypal transaction [3]. Since then, other cryptocurrency exchanges have emerged, many of which are centralized, which provide a more traditional trading terminal and support for investors with little technical know-how, and others are decentralised, which operate on blockchain networks and allow users to directly with each other using smart contracts. Centralized exchanges are predominant due to their easy-to use and familiar interface for traders. We can see in Figure 1.1 the proportion of trades on decentralised exchanges compares to the number of trades on centralized exchanges, we can also see that the volume traded in DEXes have been near 0% until the summer of 2020 and has been increasing since.

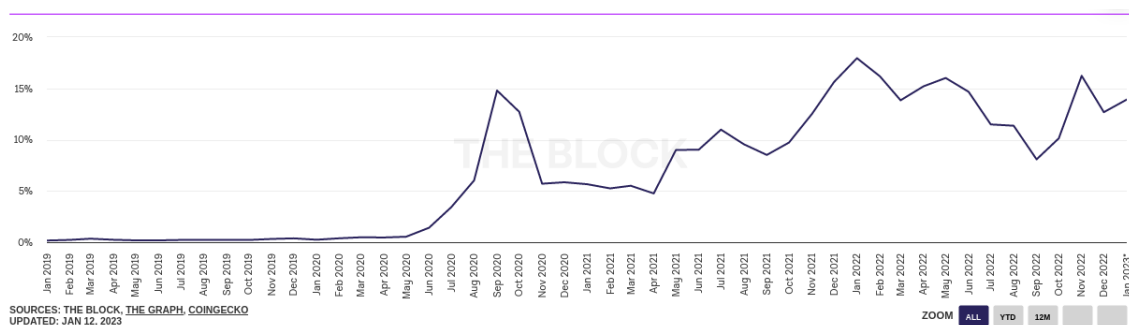


Figure 1.1: DEX to CEX Spot Trade Volume [4]

This poses the question that about which trading strategies can exploit arbitrage opportunities on decentralised exchanges. There has been some research on this topic, mainly focussing on triangular and cyclic arbitrage on DEXes such as Uniswap and SushiSwap, however there has been no research into analysing the performance of statistical arbitrage methods on decentralised exchanges.

1.2 Contributions - TODO

Chapter 2

Background

2.1 Cryptocurrencies

Before delving into the financial side of the project, it is important to understand the underlying assets and the technology that drives them.

2.1.1 Blockchain

The building blocks of cryptocurrencies come from blockchain technology. Blockchain is a distributed ledger that stores data, in blocks, in a chain, comprising the data itself as well as a full transaction history [5]. They can be thought of as a State Transition System where the state is the store of data, i.e. the owners of each token, and the state-transition function is a function of state and a transaction [6]. The transition function defines how the transaction, T , should affect the state, σ .

$$\sigma' = transition_func(\sigma, T)$$

Thus the state, σ , and transition function $transition_func$ are defined by the blockchain implementation, which tends to be more complex.

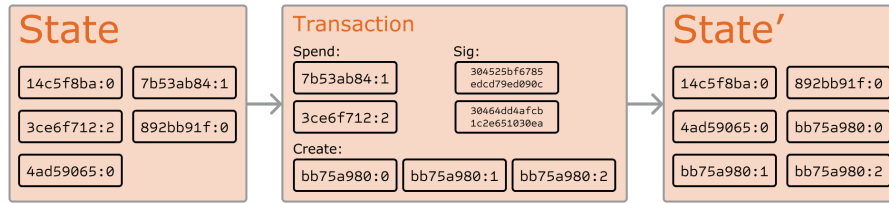


Figure 2.1: State Transition Diagram [6]

In reality, a collection of transactions are batched together into a block which is then committed to the network. The components of a block can be seen below.

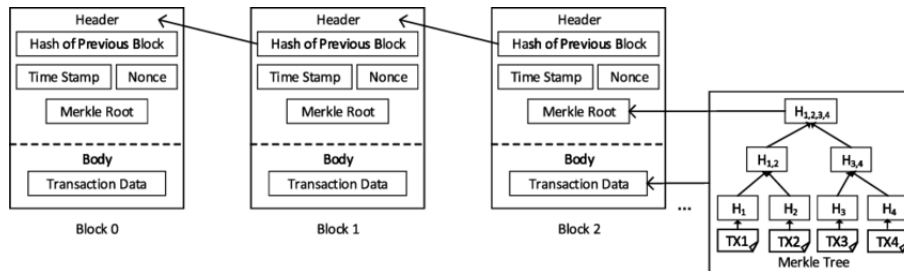


Figure 2.2: Blockchain Diagram [7]

A block can have two significant components, the header and the body. The body of the block contains the list of transactions that are to be applied. Whereas, the components of the header may depend on the blockchain's implementation however they typically have the following key components.

- **Block number** - A unique identifier assigned to each block in the chain. It serves as a chronological ordering mechanism, indicating the position of the block within the blockchain.
- **Timestamp** - Timestamp of the block's creation time.
- **Merkle Root** - A Merkle root is also stored in each block to validate transactions in the body of the block efficiently, in terms of storage and searching [8]. A Merkle tree is a tree of hashes where each leaf node is its data hash and its parent node is the hash of their children's hashes. In storing the Merkle root, we do not need to directly store each transaction in each block, and also allows a quick search for any malicious alterations in differing blocks.
- **Hash of Previous Block** - Acts as a pointer to the predecesing block. Hence, the term "Blockchain" is coined as one would be able to track back to the first block from the latest block in a chain like manner.

In addition to these components, a major feature of blockchain technology is its property of being decentralised. The state of the blockchain is managed by a distributed network of nodes that are maintained by voluntary node operators. Node operators maintain the blockchain by synchronizing with the network, validating, propagating transactions and blocks and finally participating in consensus mechanisms to ensure all of the nodes agree on the current state before updating the blockchain. For this blockchain networks are typically built on a peer-to-peer (P2P) architecture. Each participant in the network, or node, has an equal status and communicates directly with other nodes. Transactions and data are shared among participants without the need for intermediaries or central servers. When transitioning to the next state, the node broadcasts the transaction, T , to the blockchain network and each node applies the transition function to obtain the new and same state, $\sigma' = transition_func(\sigma, T)$. All nodes arrive to the same state as the transition function is deterministic. Once each node has applied the transition function, the network uses a consensus mechanism to the globally agreed σ' and the transactions become actualised.

Consensus mechanisms are used to achieve agreement among participants on the validity of transactions and the order in which they are added to the blockchain. Examples of consensus mechanisms include Proof of Work (PoW), Proof of Stake (PoS), and Delegated Proof of Stake (DPoS), however current blockchains predominantly use Proof of Work or Proof of Stake [9]. In addition to reaching consensus, the consensus algorithms also determine which node should be the creator of a new block, which node should be granted a monetary reward.

The Proof of Work (PoW) consensus algorithm was first used in Bitcoin and is where the nodes compete to solve complex mathematical problems to verify that the transactions from a new block are valid in order to add them to the blockchain. These nodes are called miners. Once the first miner obtains the solution of the mathematical puzzle, it broadcasts the new block on the network so that the other nodes verify the solution and update their local replica. Consensus is then reached when 51% of the nodes agree on the new state of the blockchain. The main drawback of this algorithm is that it requires miners to perform extensive computational calculations repeatedly until they find the solution to the mathematical problem, hence once the solution has been found by a miner, after expensive work, the solution is sent along in the broadcast to the other nodes to easily and cheaply verify the solution of the problem.

To address this issue, some blockchains have opted to use the Proof of Stake (PoS) consensus algorithm as it is more energy efficient and scalable. Unlike Proof of Work, Proof of Stake use validators to create new blocks. Validators stake an amount of cryptocurrency as collateral and the creator of a block is chosen at random from the validators, with the probability of selection based on the amount they have staked. Due to the advantages that PoS has over PoW, Ethereum has transitioned from PoW to PoS in Ethereum 2.0.

2.1.2 Ethereum

One of them was proposed by Vitalik Buterin, the co-founder of Ethereum, in a whitepaper that proposed the idea of using smart contracts to create financial products and services that could operate independently of traditional financial institutions, hence decentralised finance was birthed [10].

Smart Contracts

Smart contracts are programs that are self-executing contracts between buyers and sellers that deploy on a blockchain. The programs are reactive in the sense that they are only executed on the blockchain providing that certain conditions are met [11]. These conditions are set in the lines of code in the smart contract. To possess the reactive property smart contracts are stored and executed on every participating node on the network as part of the state σ . They are also immutable and decentralised which make smart contract advantageous as they can be used to facilitate, verify, and enforce the negotiation or performance of a contract [12, 13].

Ethereum's Architecture

Ethereum's architecture is similar to bitcoin's but has a few differences, one of which is rather than managing a distributed ledger, it uses a distributed state machine. The Ethereum Virtual Machine (EVM) defines the rules of changing states from block to block. Each node on the Ethereum blockchain contains an immutable instance of the EVM [14].

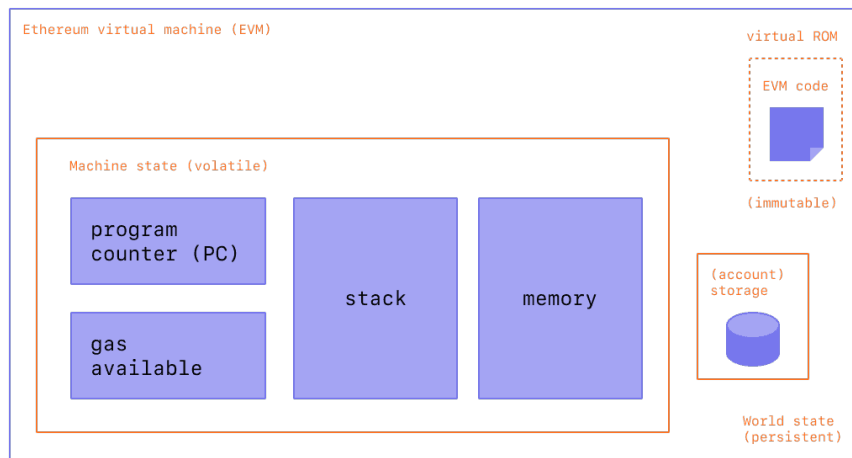


Figure 2.3: EVM components [14]

Ethereum's state σ consists of the states of all accounts where the state of an account is represented by four components: balance of ETH (the native cryptocurrency on Ethereum), a nonce, hash of the smart contract code (if account is a smart contract) and the Merkle root of smart contract storage which maintains the contracts variables (again, if account is a smart contract) [15]. A nonce, 'Number only used once', is a number that is added to a hashed block to make the transaction more secure. It is randomly generated and used

to validate a transaction. A miner first guesses a nonce and appends the guess to the hash of the current header. The miner then rehashes the value and compares this to the target hash. If the guess was correct, the miner is granted the block [16].

Ethereum supports two types of transactions: contract creation and message calls. Message calls are used to transfer ETH from account to account and also for invoking smart contract. The process of how transactions are validated is below:

1. Validate the parent block
2. Validate that the current timestamp is greater than the previous timestamp
3. Check that the Ethereum concepts are valid
4. Perform Proof of Stake on the block
5. Check for errors and gas
6. Validate the final state

Gas Fees

Due to the capability of expressing loops in EVM bytecode, there is a possibility that smart contract execution could continue indefinitely. This presents a challenge as invoking a contract with an infinite loop would cause all Ethereum network nodes to become stuck executing it, resulting in a denial of service. To mitigate this issue, Ethereum has implemented a pricing mechanism. Every computation performed by a smart contract requires the payment of a fee called gas, which is denominated in ether.

To measure how much computational effort is required to execute operations on the Ethereum network, gas is used [17]. Every block has a base fee, derived from the demand for the block space, which is burnt. Therefore, users of the network are expected to set a tip (priority fee) to reimburse miners for adding their transaction in blocks, thus the higher the tip, the greater the incentive for miners to validate the transaction. Using gas means that the Ethereum network is tolerant to spam and also has a maximum gas fee to make Ethereum tolerant to malicious code that would be used to waste resources.

2.1.3 Decentralised Finance

One of the applications of Ethereum and smart contracts is Decentralised Exchanges (DEXes). Before delving into DEXes it is important to understand centralized exchanges.

Centralized Exchanges

Centralized exchanges(CEXes) allow agents to discover and trade assets. CEXes facilitate trading between buyers and sellers by providing an online platform that manages and maintains an order book. An order book aggregates buy and sell orders and execute matching buy and sell orders. The order book and transactions are typically managed on a central database. When trading, exchanges charge trading fees for the maker and the taker to operate the exchange and do not charge any gas fees as there is no interaction with the blockchain.

Decentralized Exchanges

In contrast, DEXes utilize blockchain technology and smart contracts to execute trades thus providing a high level of determinism, by nature of the technology. These trades are executed on the blockchain via smart contracts and on-chain transactions. There are two types of DEX: order book DEXes and Automated Market Makers (AMMs). An order book DEX is less common and is akin to a CEX, an order book is stored on the blockchain

rather than on a central database. This means each order placed requires the order book to be posted on the blockchain at each transaction. Automated Market Makers are more common and provide instant liquidity by using liquidity pools so that users can swap their tokens for a price that is determined by the portions within the liquidity pool [18]. DEXes have multiple pros including lower transaction fees, privacy, diversity and trustless transactions but they also have their drawbacks such as scalability and poor liquidity a lot of the DEXes are quite new [19].

2.2 Arbitrage

Arbitrage is the process in which a trader simultaneously buys and sells an asset to take advantage of a market inefficiency [20]. Arbitrage is also possible in other types of securities by finding price inefficiencies in the prices of options, forward contracts and other exotics.

Sources have shown that the word “*Arbitrage*” has been used as early as the Renaissance era when surviving documents showed a large number of bills being exchanged [21]. There has also been some evidence to suggest that arbitrage was used as early as the Greek and Roman eras. Early forms of arbitrage would likely have been purchasing a commodity then transporting them to a foreign land and selling them at a higher price. This type of arbitrage is called commodity arbitrage and is still applicable today. With the example above, transporting the goods takes a significant amount of to the merchant, or trader, which could cause variations in the price, however, in the modern day this has been reduced and with electronic exchanges, this time to buy and sell is very small. This means inefficiencies in the market, where a trader can profit purely by buying and selling, should not exist. This is called the “Law of One Price”. The “Law of One Price” states that every identical commodity or asset should have the same price regardless of exchange or location, given there are no transaction costs, no transportation costs, no legal restrictions, the exchange rates are the same and no market manipulation occurs [22]. This is because if this were not the case, an arbitrage opportunity would arise and someone would take advantage of the scenario causing the prices on both markets to converge due to the market forces. In the real world arbitrage opportunities are tremendously common, thus allowing a risk-free investment [23, 24].

There are countless types of arbitrage such as spatial arbitrage, which profits off of different prices on exchanges in different locations, temporal arbitrage, which takes advantage of price differences at different times, risk arbitrage, which profits from perceived discrepancies in their risk-return profiles and finally market arbitrage which takes advantages of different prices on different exchanges/markets. Statistical methods include pairs trading, which involves buying and selling assets that are believed to be mispriced relative to one another, momentum trading, which identifies if assets have a strong momentum (either up or down) and profiting off of that, and finally, algorithmic trading which uses algorithms to analyze data and trades based on statistical analysis. This project shows how these opportunities can be exploited both in a pure manner as well as using statistical methods.

2.3 Pure Arbitrage Techniques

Research into cryptocurrency arbitrage is still in its infancy and previous research has mainly focussed on the economics of cryptocurrencies, i.e. miner/trader behaviour and influence of cryptocurrency trading [25, 26, 27, 28, 29, 30, 31]. Furthermore, there has been very limited research comparing statistical strategies and pure methods of arbitrage of cryptocurrencies. Despite this, there has been plentiful research on arbitrage as a whole as it is immensely profitable; as a result of this people/institutions tends to keep their

newly found research secret. Of the published research, I have looked into the arbitrage techniques that are used. As arbitrage can be highly profitable, it can be found in countless types of assets, such as options, stocks, bonds and many other types of products. Research into all types of products exist going into the theory and practical aspects of each [32, 33]. The most similar type of asset class to cryptocurrencies is fiat currencies, traditional currencies that are issued by governments, such as the US Dollar (USD) and the Great British Pound Sterling (GBP). The research in arbitrage in foreign exchanges shows that using a triangular/cyclic arbitrage is highly profitable and effective [34, 35, 36]. This trading strategy takes advantage of price discrepancies between three or more different currencies in the foreign exchange market. It involves executing a series of trades to profit from the imbalance in exchange rates between the currencies involved. An example of this can be seen in Figure 2.4, given the exchange rates are $\$1 = \text{€}0.85$, $\text{€}1 = \text{£}0.75$, $\text{£}1 = \$1.20$, we can make a series of exchanges (trades) such that by starting with \$100, the result of this cyclic arbitrage I am left with \$130.72, hence a \$30.72 risk-free profit.

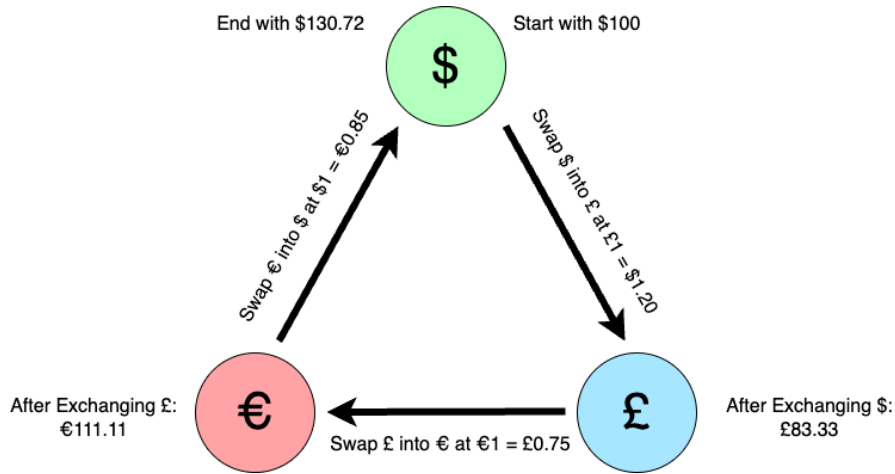


Figure 2.4: Triangular Arbitrage Diagram

Although this project aims to focus on statistical arbitrage techniques, it is still important to look at research into purer forms of arbitrage, i.e. triangular and cyclic arbitrage.

As previously mentioned, research into this topic is still in its infancy thus which means a very thin slice of exploration on the subject matter. The majority of the research has been into the arbitrage on centralized exchanges [37, 38, 39]. Cristian Pauna investigates and implements an arbitrage strategy in [39]. The paper details the technical details of arbitrage trading from the data and the system architecture used. Pauna finds complications such as requesting data from multiple exchanges, converting the data such that it is homogeneous and also managing server load. Pauna presents the architecture such that the servers request data from the necessary exchanges, aggregating prices in a relational database which then triggers a server that is used to generate trading signals.

As previously mentioned triangular and cyclic arbitrage is one of the most used and purest forms of arbitrage to implement and analyse, [40] explores triangular arbitrage on decentralised exchanges. Algorithm 1 is the algorithm used to find the most profitable arbitrage route on a particular platform, once this is calculated, it is compared with other routes on other platforms. Initially, the system converts the base token into another token and converts it back into the base token, using only one token is used as a middle route, then using the algorithm below, increases the number of middle tokens.

On evaluating the performance of the strategy on differing platforms depended on three main features of each exchange:

1. Portion size - Depending on how much the “trader” invested revenues differed and with the larger portion size, the revenue decreases as the token pair prices are adjusted

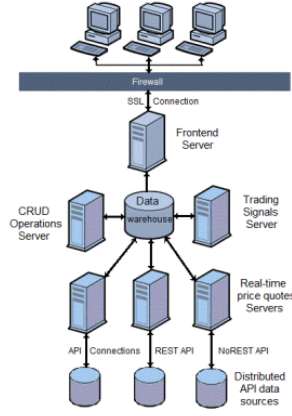


Figure 2.5: Arbitrage system architecture [39]

Algorithm 1 Maximum Profit Route Searching (R)

Input: T (token list), P (price graph), n (current route)

```

for  $i = 1, \dots, T$  do
   $r = \text{get\_profit}(n + i)$ 
  for  $j = 1, \dots, P[i]$  do
     $p = \max(r, R(T, P, n_j))$ 
  end for
end for
return  $p$ 

```

based on supply/demand.

2. Transaction fees - Each exchange has its own transaction fee.
3. Other considerations such as price slippage - Exchanges have different liquidity levels which depend on the usage and liquidity providers that the exchange employs.

Figure 2.6 displays the revenues obtained by same trading token route, $\text{ETH} \rightarrow \text{MKR} \rightarrow \text{OMG} \rightarrow \text{USDT} \rightarrow \text{ETH}$. As we can see upon applying the strategy on multiple exchanges; Uniswap, 1inch, Kyberswap and Bancor, 1 inch was the only exchange that generated a profit whereas the others lose money.

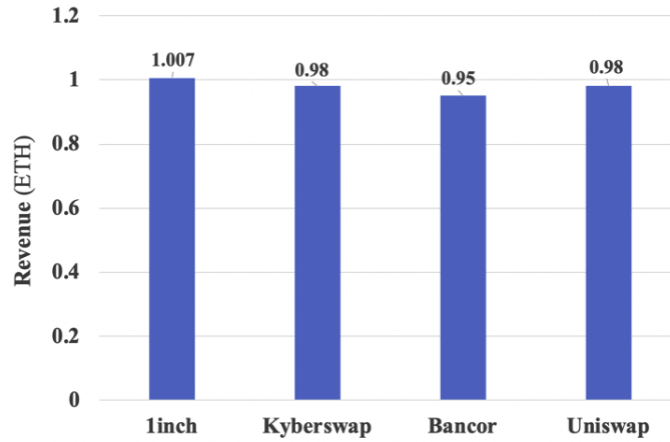


Figure 2.6: Trading profits same token routes within different exchanges [40]

Another paper that implemented and evaluated a cyclic arbitrage opportunity is [41]. The research consists of proposing a theoretical arbitrage model and further evaluation of real

transactional data. The arbitrage model used is simple to understand, as it searches for a cyclic transaction between n tokens, A_1, A_2, \dots, A_n is a sequence of n trades:

Trade 1: Exchange δ_1 of A_1 to δ_2 of A_2

Trade 2: Exchange δ_2 of A_2 to δ_3 of A_3

\dots

Trade n : Exchange δ_n of A_n to δ'_1 of A_1

It is important to note that $\delta_i = \delta_{i+1}$, i.e. the output of trade is equivalent to the input of the next. The revenues within a cycle are defined as $\delta_{i+1} - \delta_i$, and the overall profit is $\delta'_1 - \delta_1$. This is not as simple as the revenues depend on how liquid the exchange is, thus the liquidity pools of each possible trading pair are hugely important. Therefore, the paper proposes a theorem, below:

Theorem 1 *For a given cycle $A_1 \rightarrow A_2 \rightarrow \dots \rightarrow A_n \rightarrow A_1$ with n tokens, there exists an arbitrage opportunity for the cyclic transaction if the product of exchange rates $\frac{a_{2,1}a_{3,2}\dots a_{1,n}}{a_{1,2}a_{2,3}\dots a_{n,1}} > \frac{1}{r_1^n r_2^n}$ where $a_{i,j}$ denotes the liquidity of token A_i in the liquidity pool with token A_j . [41]*

In addition to the theorem, to obtain an optimal strategy we need to compute the optimal trading volume of a cycle, $A_1 \rightarrow A_2 \rightarrow \dots \rightarrow A_n \rightarrow A_1$. The paper proposes the optimal trading volume to be $\delta_a^{op} = \frac{\sqrt{r_1 r_2 a' a} - a}{r_1}$ where $a = \frac{a'_{1,n} a_{n,1}}{a_{n,1} + r_1 r_2 a'_{n,1}}$ and $a' = \frac{r_1 r_2 a'_{1,n} a_{n,1}}{a_{n,1} + r_1 r_2 a'_{n,1}}$. Thus to calculate such arbitrage opportunities knowing the liquidity of tokens in other tokens' liquidity pools, algorithm 2 infers the direction and volumes to trade to get the optimal revenue.

Algorithm 2 Computing the equivalent liquidity of the cycle

```

 $a'_{1,n} \leftarrow a_{1,2}$ 
 $a'_{n,1} \leftarrow a_{2,1}$ 
for  $i$  from 2 to  $n - 1$  do
   $a'_{1,n} \leftarrow \frac{a'_{1,n} a_{i,i+1}}{a_{i,i+1} + r_1 r_2 a'_{n,1}}$ 
   $a'_{n,1} \leftarrow \frac{r_1 r_2 a'_{1,n} a_{i+1,i}}{a_{i,i+1} + r_1 r_2 a'_{n,1}}$ 
end for

```

After analyzing Ethereum block data and applying this strategy to identify the number of arbitrage opportunities, it was found that between May 4, 2020, and April 15, 2021, there were numerous exploitable and profitable arbitrage opportunities. These opportunities grew consistently to reach 1,750 in 11 months, as depicted in Figure 2.7. Only cycles with length 3 were experimented with and only cycles including ETH as 80% of the liquidity pools on Uniswap include ETH and another cryptocurrency [42]. Furthermore, it is found that 287,241 of the 292,606 arbitrages executed started with ETH, and 85% of the arbitrages used a cycle of length 3. The total revenue of the cyclic arbitrage was 34,429 ETH. However, gas fees account for 24.6% of the total revenue leaving an approximate 25,971 ETH profit.

The paper then delves into the implementation of the smart contract, and explores how both *sequential* and *atomic* implementations would affect the revenue and execution of the contracts. It was found that 52.3% of the arbitrages that were executed sequentially generated a loss, likely due to the fact that, when one submits n orders, the n blockchain transactions are executed sequentially, meaning some external transactions can be inserted between these transactions. Thus using atomic transactions avoids this issue of external transactions does not affect the market price that may affect the outcome of the arbitrage.

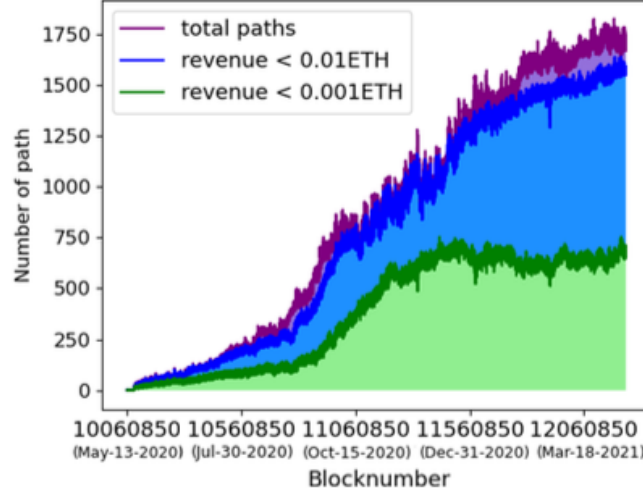


Figure 2.7: Number of exploitable opportunities in Uniswap V2 over time. The purple line represents the number of cycles that provide revenue higher than 0.0001 ETH. The green represents the number of cycles whose revenue is under 0.001 ETH. The blue line represents the number of cycles whose revenue is under 0.01 ETH. [41]

Furthermore, the authors of the paper also investigated the performance differences between using private smart contracts and public contracts. Deploying a smart contract that calls Uniswap functions, i.e. a private smart contract, is intuitively better and achieves a higher success rate of a lower bound of 52% and a higher bound of 90% in comparison to calling a public Uniswap smart contract which has a success rate of 27.3%. Overall the paper provides an insightful look into cyclic arbitrage in DEXes and highlights important decisions made such as liquidity calculations and smart contracts while comparing the performance of different options available.

2.4 Statistical Arbitrage Techniques

As mentioned previously mentioned and the subject of the project is to optimize statistical arbitrage methods to be able to compete with a purer form of arbitrage, i.e. cyclic arbitrage. As previously mentioned there are many methods of stat arb, pairs trading, momentum trading and algorithmic trading. Within these methods there are countless strategies to adopt and profit from, thus to limit the scope, this project I will be investigating strategies within pair trading. Research within Pair trading has been vast with many streams of approaches emerging; distance approach, cointegration approach, time-series approach, stochastic approach and some others, including using machine learning [43]. However, for this project, we will only look at cointegration/co-correlation approaches.

2.4.1 Mean Reversion

The cointegration approach follows three key steps. The first is the selection of pairs based on similarity measures, the next is assessing the tradability and finally, thresholds are set for trading. The spread is defined as

$$\varepsilon_{ij,t} = P_{i,t} + \gamma P_{j,t}$$

where $P_{i,t}$ and $P_{j,t}$ denote the $I(1)$ non-stationary price processes of the assets i and j , γ is the cointegration coefficient, also referred to in literature as the hedge ratio. $\varepsilon_{ij,t}$ is the linear combination of the non-stationary prices and is $I(0)$ stationary and hence mean-reverting, note that stationary processes are those of which have a constant mean. Rad's

implementation of this approach on stocks results in a 0.83% return before considering transaction costs [44]. Another paper, [45], looked into setting the thresholds and setting a minimum profit, MP_{ij,t_c} :

$$MP_{ij,t_c} = \frac{n(\varepsilon_{ij,t_0} - \varepsilon_{ij,t_c})}{|\gamma|}$$

Where t_0 and t_c are the opening and closing times, n is the volume longed of asset j .

2.4.2 Optimal Portfolio Design for Mean Reversion

There has been further research into optimizing mean reversion, one of which was to use the successive convex approximation method on the mean reverting portfolio design [46]. The paper initially proposes the mean reversion portfolio:

- For each asset, the price at time t is denoted as p_t and its corresponding log-price $y_t \triangleq \log(p_t)$, its vector form of M assets $\mathbf{y}_t \triangleq [y_{1,t}, \dots, y_{M,t}]^T$.
- The log-price spread is given by $y_t \triangleq \beta^T \mathbf{y}_t$, where $\beta \triangleq [\beta_1, \dots, \beta_M]^T$ denotes the hedge ratios.
- The cointegration space with N relations is defined by $\mathbf{B} \triangleq [\beta_1, \dots, \beta_N]$, thus the N spreads are $s_t \triangleq \mathbf{B}^T \mathbf{y}_t$.
- For these N spreads, the portfolio weight matrix is denoted as $\mathbf{w} \triangleq [w_1, \dots, w_N]^T$.
- The auto-covariance matrix for the spreads s_t is defined as $M_i \triangleq \text{Cov}(s_t, s_{t+i}) = \mathbb{E}[(s_t - \mathbb{E}[s_t])(s_{t+i} - \mathbb{E}[s_{t+i}])^T]$

Now that we have defined everything required, we can now formalize the problem. The general problem of mean reversion portfolio design problem is formalized by:

$$\begin{aligned} & \underset{\mathbf{w}}{\text{minimize}} \quad F(\mathbf{w}) \triangleq U(\mathbf{w}) + \mu V(\mathbf{w}) + \gamma S(\mathbf{w}) \\ & \text{subject to} \quad \mathbf{w} \in \left\{ \mathbf{w} \mid \|\mathbf{B}\mathbf{w}\|_0 \leq L \right\}, \quad \text{where } L \text{ is the total leveraged investment} \end{aligned}$$

- μ defines the trade-off between the mean reversion measure and the variance preference.
- γ defines the regularization parameter of how sparse we would like the cointegration space to be.

Where the Mean Reversion term:

$$U(\mathbf{w}) \triangleq \xi \frac{\mathbf{w}^T \mathbf{H} \mathbf{w}}{\mathbf{w}^T \mathbf{M}_0 \mathbf{w}} + \zeta \left(\frac{\mathbf{w}^T \mathbf{M}_1 \mathbf{w}}{\mathbf{w}^T \mathbf{M}_0 \mathbf{w}} \right)^2 + \eta \sum_{i=2}^p \left(\frac{\mathbf{w}^T \mathbf{M}_i \mathbf{w}}{\mathbf{w}^T \mathbf{M}_0 \mathbf{w}} \right)^2$$

And the variance term:

$$V(\mathbf{w}) \triangleq \begin{cases} 1/\mathbf{w}^T \mathbf{M}_0 \mathbf{w} & \text{VarInv}(\mathbf{w}) \\ 1/\sqrt{\mathbf{w}^T \mathbf{M}_0 \mathbf{w}} & \text{StdInv}(\mathbf{w}) \\ -\mathbf{w}^T \mathbf{M}_0 \mathbf{w} & \text{VarNeg}(\mathbf{w}) \\ -\sqrt{\mathbf{w}^T \mathbf{M}_0 \mathbf{w}} & \text{StdNeg}(\mathbf{w}) \end{cases}$$

The variance term can be represented in any of the four forms.

And the asset selection term:

$$S(\mathbf{w}) \triangleq \|\mathbf{B}\mathbf{w}\|_0 = \sum_{m=1}^M \text{sgn}(|[\mathbf{B}\mathbf{w}]_m|)$$

This asset selection criterion is not necessary however as trading incurs a cost, selecting all of the assets is costly, thus selecting a subset of assets to trade is more profitable. To formalize this goal, we would like to minimize the cointegration space thus we use the ℓ_0 norm.

The paper then goes on to solve the optimization problem using the successive convex approximation (SCA) method [47]. The SCA method takes an optimization problem in the form of:

$$\begin{aligned} & \underset{\mathbf{x}}{\text{minimize}} && f(\mathbf{x}) \\ & \text{subject to} && \mathbf{x} \in \mathcal{X} \end{aligned}$$

Where $\mathcal{X} \subseteq \mathbb{R}^N$ is convex and $f(\mathbf{x})$ is non-convex. The SCA method involves starting at an initial point $\mathbf{x}^{(0)}$ and solving a series of subproblems of surrogate functions $\tilde{f}(\mathbf{x}; \mathbf{x}^{(k)})$ over the set \mathcal{X} . The sequence $\{\mathbf{x}^{(k)}\}$ is generated by:

$$\begin{cases} \hat{\mathbf{x}}^{(k+1)} = \underset{\mathbf{x} \in \mathcal{X}}{\text{argmin}} \tilde{f}(\mathbf{x}; \mathbf{x}^{(k)}) \\ \mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \gamma^{(k)}(\hat{\mathbf{x}}^{(k+1)} - \mathbf{x}^{(k)}) \end{cases}$$

The first step is to generate a descent direction and then update the variable with a step size of $\gamma^{(k)}$. After applying this method to the MRP problem and further analysis of the paper, the following algorithm is proposed and used to solve the MRP design problem:

Algorithm 3 SCA-Based Algorithm for The Optimal MRP Design Problem

Require: $\mathbf{H}, \mathbf{M}_i, \mu, \gamma, \mathbf{B}, L$ and τ

- 1: Set $k = 0, \gamma^{(0)}$ and $\mathbf{w}^{(0)}$
 - 2: **repeat**
 - 3: Compute $\mathbf{A}^{(k)}$ and $\mathbf{b}^{(k)}$
 - 4: $\hat{\mathbf{w}}^{(k+1)} = \underset{\mathbf{w} \in \mathcal{W}}{\text{argmin}} \mathbf{w}^T \mathbf{A}^{(k)} \mathbf{w} + \mathbf{b}^{(k)T} \mathbf{w}$
 - 5: $\mathbf{w}^{(k+1)} = \mathbf{w}^{(k)} + \gamma^{(k)}(\hat{\mathbf{w}}^{(k+1)} - \mathbf{w}^{(k)})$
 - 6: $k \leftarrow k + 1$
 - 7: **until** convergence
-

However, 4 line is a convex problem and has no closed-form solution thus to solve this subproblem using the ADMM method, this is done by introducing an auxiliary variable $\mathbf{z} = \mathbf{B}\mathbf{w}$.

$$\begin{aligned} & \underset{\mathbf{x}, \mathbf{z}}{\text{minimize}} && \mathbf{w}^T \mathbf{A} \mathbf{w} + \mathbf{b}^T \mathbf{w} \\ & \text{subject to} && \|\mathbf{z}\|_1 \leq B, \mathbf{B}\mathbf{w} - \mathbf{z} = \mathbf{0} \end{aligned}$$

This is then summarized into Algorithm 4:

Algorithm 4 An ADMM-Based Algorithm for Problem on line 4 in Algorithm 3

Require: $\mathbf{A}, \mathbf{b}, \mathbf{B}, B, \rho$

- 1: Set $\mathbf{w}^{(0)}, \mathbf{z}^{(0)}, \mathbf{u}^{(0)}$ and $k = 0$
 - 2: **repeat**
 - 3: $\mathbf{w}^{(k+1)} = -(2\mathbf{A} + \rho\mathbf{B}^T\mathbf{B})^{-1}(\mathbf{b} + \rho\mathbf{B}^T(\mathbf{u}^{(k)} - \mathbf{z}^{(k)}))$
 - 4: $\mathbf{h}^{(k)} = \mathbf{B}\mathbf{w}^{(k+1)} + \mathbf{u}^{(k)}$
 - 5: $\mathbf{z}^{(k+1)} = \Pi_{\mathcal{C}}(\mathbf{h}^{(k)})$
 - 6: $\mathbf{u}^{(k+1)} = \mathbf{u}^{(k)} + \mathbf{B}\mathbf{w}^{(k+1)} - \mathbf{z}^{(k+1)}$
 - 7: $k \leftarrow k + 1$
 - 8: **until** convergence
-

After all of this analysis, the authors of the paper, [48, 46], ran simulations on real data comparing underlying spread. It found that it resulted in consistent profits as shown in figure 2.8:

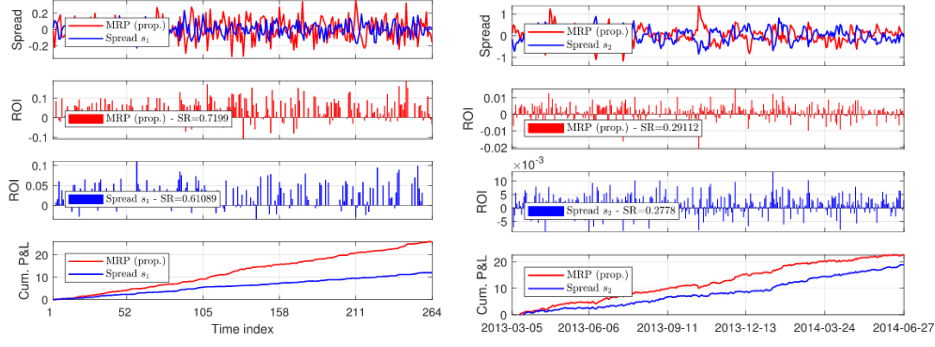


Figure 2.8: A mean-reversion trading based on real data [46]

Overall, this research successfully formulizes, solves the optimization problem mathematically, and goes further to implement the algorithms to solve the problem programmatically. In addition, the author compares the implementation with other benchmark algorithms, showing that it results in a greater P&L and Sharpe ratio.

2.4.3 Statistical Arbitrage using the Kalman Filter

Another method that is used in statistical arbitrage is using the Kalman Filter. Recall the equation for spread:

$$\varepsilon_{ij,t} = P_{i,t} + \gamma P_{j,t}$$

The Kalman filter is a recursive algorithm for estimating the state of noisy data and that needs to be filtered to be able to estimate the state of a system based on a sequence of observations, taking into account both current measurements and the system dynamics [49]. This makes it very useful to estimate the hedge ratio γ . Initially, a book by Vidyamurthy discusses best practices for choosing cointegrated equities and found that the Kalman filter was found optimal when the state-space and observation equations are linear and the noise is Gaussian [50]. Since then there have been many extensions of the filter such as the Extended Kalman Filter (EKF) and Unscented KF aimed to handle when the state-space and observation equations are non-linear and the noise is not Gaussian.

The Kalman Filter works in 2 phases, prediction and update. The prediction phase is as follows

$$\begin{aligned}\hat{\mathbf{x}}_k &= \mathbf{F}_k \hat{\mathbf{x}}_{k-1} + \mathbf{B}_k \vec{\mathbf{u}}_k + \mathbf{w}_k \\ \mathbf{P}_k &= \mathbf{F}_k \mathbf{P}_{k-1} \mathbf{F}_k^T + \mathbf{Q}_k\end{aligned}$$

Where $\hat{\mathbf{x}}_k$ is the new best estimate (prediction) that is derived from $\hat{\mathbf{x}}_{k-1}$, the previous estimate and the prediction function \mathbf{F}_k . $\vec{\mathbf{u}}_k$ is the correction term, called the control vector, that is used when it is known that there are external influences in combination with \mathbf{B}_k which is called the control matrix. In addition to this, the new uncertainty (covariance matrix), \mathbf{P}_k , is calculated using the previous uncertainty and additional uncertainty from the environment, \mathbf{Q}_k , \mathbf{w}_k is called the state noise. The update is as follows

$$\begin{aligned}\hat{\mathbf{x}}'_k &= \hat{\mathbf{x}}_k + \mathbf{K}'(\vec{\mathbf{z}}_k - \mathbf{H}_k \hat{\mathbf{x}}_k) \\ \mathbf{P}'_k &= \mathbf{P}_k - \mathbf{K}' \mathbf{H}_k \mathbf{P}_k \\ \mathbf{K}' &= \mathbf{P}_k \mathbf{H}_k^T (\mathbf{H}_k \mathbf{P}_k \mathbf{H}_k^T + \mathbf{R}_k)^{-1}\end{aligned}$$

Where \mathbf{K}' is defined as the Kalman gain, \mathbf{H}_k is the measurement matrix, $\vec{\mathbf{z}}_k$ is mean of the observed values, which is also calculated by $\vec{\mathbf{z}}_k = \mathbf{H}_k \hat{\mathbf{x}}_k + \mathbf{v}_k$ where \mathbf{v}_k is the measurement noise, and \mathbf{R}_k is the covariance of the uncertainty of the observed values [51].

A paper that investigated the use of the Kalman Filter on ETFs found that the strategy it employed worked well for in-sample data points and worse, but still profitable, results for out-of-sample data [52]. The paper adapted the Kalman Filter to be able to use it for pairs trading to the following:

$$\mathbf{y}_t = \mathbf{x}_t \beta_t + \epsilon_t$$

$$\beta_t = \mathbf{I} \beta_{t-1} + \omega_t$$

Then calculating the Kalman Gain:

$$\text{Kalman Gain} = \frac{\text{Error in the estimate}}{\text{Error in the estimate} + \text{Error in the measurement}}$$

Then to calculate the estimate:

$$\text{Estimate}_t = \text{Estimate}_{t-1} + \text{Kalman Gain} \times (\text{Measurement} - \text{Estimate}_{t-1})$$

And finally, calculating the new error:

$$E_{\text{estimate}_t} = \frac{E_{\text{measurement}} \times E_{\text{estimate}_{t-1}}}{E_{\text{measurement}} + E_{\text{estimate}_{t-1}}}$$

$$E_{\text{estimate}_t} = E_{\text{estimate}_{t-1}} \times (1 - \text{Kalman Gain})$$

The author states “pairs trading strategies have gained widespread acceptance thus making profitability much more elusive” to justify the disappointing results, however, the author fails to find evidence or provide sufficient evidence to justify the claim [52].

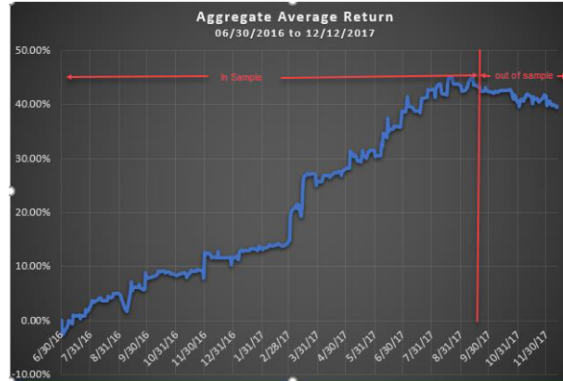


Figure 2.9: Aggregate average return of using the Kalman filter for pairs trading on ETFs [52]

Another paper used the combination of the Kalman Filter and Machine Learning, more specifically Extreme Learning Machine and Support Vector Regression (SVR) to build a statistical arbitrage strategy on the Brazilian Stock Exchange [53]. The strategies can simply be explained as using SVR and ELM to forecast returns and using the Kalman Filter to improve the forecast.

The paper also compares methods, such as LASSO, BMA, and GRR, to benchmark the performance of the Kalman Filter. The research found that using simply ELM and SVR forecasts results in a return of 20.19% and 21.32% respectively for out-of-sample data points and using a combination with the Kalman Filter gives a return of 26.13% for out-of-sample data points. The full results can be seen below in Figure 2.11. In addition to this, it can be seen that the volatility of the return also decreases which is ideal for investment managers.

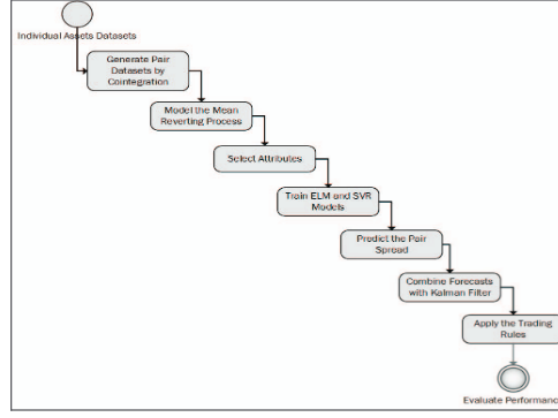


Figure 2.10: Visualisation of the trading strategy used in [53]

TABLE IV. ECONOMETRIC PERFORMANCE – ELM AND SVR – IN-SAMPLE

MODEL	MAX. DD	SHARPE	VOLATILITY	RETURN
ELM	-2.31%	1.80	3.05%	5.83%
SVR	-2.18%	1.73	3.20%	5.39%

TABLE V. ECONOMETRIC PERFORMANCE ELM AND SVR – OUT-OF-SAMPLE

MODEL	MAX. DD	SHARPE	VOLATILITY	RETURN
ELM	-2.80%	3.83	5.64%	20.18%
SVR	-2.72%	4.31	5.29%	21.32%

TABLE VI. ECONOMETRIC PERFORMANCE – COMBINATION MODELS– IN-SAMPLE

MODEL	MAX. DD	SHARPE	VOLATILITY	RETURN
BMA	-2.24%	1.38	3.05%	4.94%
GRR	-2.34%	2.07	3.20%	6.37%
KALMAN	-2.30%	1.97	3.57%	6.89%
LASSO	-2.43%	2.06	3.39%	6.30%

TABLE VII. ECONOMETRIC PERFORMANCE – COMBINATION MODELS – OUT-OF-SAMPLE

MODEL	MAX. DD	SHARPE	VOLATILITY	RETURN
BMA	-2.97%	3.83	5.12%	19.33%
GRR	-2.53%	4.76	5.49%	23.69%
KALMAN	-2.64%	5.29	5.17%	26.13%
LASSO	-2.64%	4.76	5.49%	23.79%

Figure 2.11: Econometric results [53]

Other papers/articles such as [54, 55, 56] have designed, compared and analysed other statistical arbitrage techniques using Machine Learning algorithms and revealed that some algorithms are profitable. The majority of research on machine learning trading strategies has been on assets such as stocks on centralized exchanges. The little research that has been done on statistical arbitrage on cryptocurrencies has all been on analysing arbitrage on centralized exchanges and not decentralised exchanges. One of the research projects that analysed machine learning methods of statistical arbitrage on cryptocurrencies on a centralized exchange, compared a logistic regression approach with a random forest approach [56].

2.4.4 Analysis on Cryptocurrency Arbitrage on Centralized Exchanges

Although the research in the papers previously mentioned does not investigate the cointegration approach on cryptocurrencies, the takeaways are the mathematical fundamentals that are used in statistical arbitrage. Kristoufek and Bouri researched the sources of stat. arb. of bitcoin in multiple centralized exchanges. The Grey correlation is built on top of the Grey system theory [57], and can capture non-linear correlations without assuming a Gaussian distribution, thus using the Grey correlation provides a more robust metric to understand correlations between both series. The Grey correlation $\gamma(X_0, X_i)$ is defined with two steps:

1. $\gamma(x_0(k), x_i(k)) = \frac{\min_i \min_k |x_0(k) - x_i(k)| + \varepsilon \max_i \max_k |x_0(k) - x_i(k)|}{|x_0(k) - x_i(k)| + \varepsilon \max_i \max_k |x_0(k) - x_i(k)|}$
2. $\gamma(X_0, X_i) = \frac{1}{n} \sum_{i=1}^n \gamma(x_0(k), x_i(k))$

With $\varepsilon \in [0, 1]$, the standard is set to $\varepsilon = 0.5$.

The DCC-GARCH(1,1), [58], model is also used to obtain conditional correlations for Bitcoin exchanges. The model was designed to use a combination of parameters such as the standard deviation of Bitcoin returns, traded volume, the volume of on-chain transactions, fees paid to miners, the ratio of current price and recent price history and internet hype/trends.

Upon analysis of Grey and DCC-GARCH(1,1) correlations, it is found that the DCC correlations show a little variability whereas the Grey correlations are a lot more variable ranging from 0.29 to 1. In addition, the paper then further investigates these sources and finds that these opportunities are introduced when there is a large number of inter-exchange transfer requests, i.e. the network is congested, and high price volatility. In contrast, the high volume of exchanges and on-chain activity cause the arbitrage opportunities to decrease. This paper finds and explains these sources of statistical arbitrage however does not implement or devise an algorithm that uses statistical arbitrage to generate a profit from price discrepancies of Bitcoin on different exchanges.

A paper that investigates statistical arbitrage on multiple cryptocurrencies is [59]. The authors of this paper analysed co-movements and cointegration of different cryptocurrencies on a centralized exchange using Augmented Dickey-Fuller (ADF) and Kwiatkowski-Phillips-Schmidt-Shin (KPSS), Ljung-Box autocorrelation tests on both stationary forms ($I(0)$) and the original form ($I(1)$). The paper then develops a dynamic factor model based on the assumption that the price dynamics of cryptocurrencies are driven by Bitcoin [60], this is then evidenced by similar paths found in cryptocurrencies shown in Figure 2.12.

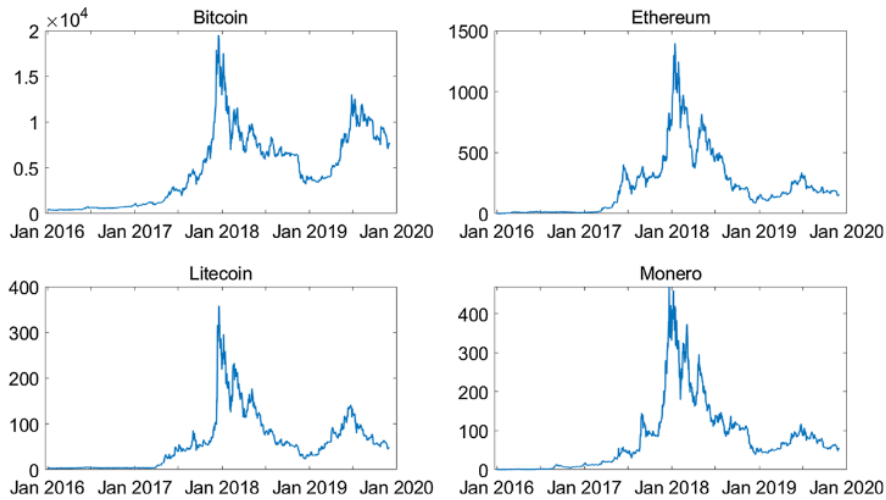


Figure 2.12: Price behaviour of Bitcoin, Ethereum, Litecoin, Monero [59]

For simplicity the authors set the number of hidden factors to 2 and upon analysis f_1 is a $I(1)$ process and the second factor f_2 is a stationary process that is independent of f_1 . It is also found after overlaying f_1 with the price of Bitcoin, that the first factor strongly correlates with the price of Bitcoin.

The paper then uses this model to build an investment strategy, using forecasting using

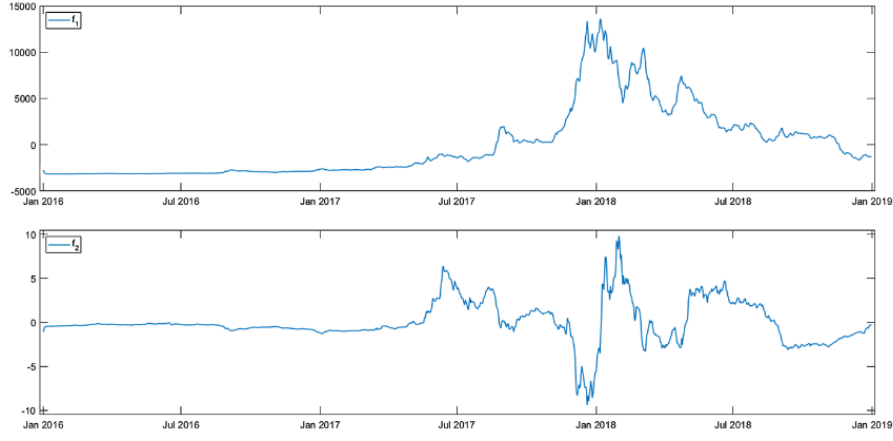


Figure 2.13: Hidden factors f_1 and f_2 from Jan 2016 to Dec 2018 [59]

the estimated parameters:

$$\hat{p}_{i,\tau+1} = \mathbb{E}_\tau(p_{i,\tau+1}) = \hat{\alpha}_i + \hat{\beta}_{i1}\mathbb{E}_\tau(f_{1,\tau+1}) + \hat{\beta}_{i2}\mathbb{E}_\tau(f_{2,\tau+1})$$

Where

$$f_{1,t} = \lambda_1 f_{1,t-1} + \eta_{1,t}$$

$$f_{2,t} = \lambda_2 f_{2,t-1} + \eta_{2,t}$$

The expected gains one day ahead are given by:

$$g_{\tau+1} = \mathbb{E}_\tau[v_{\tau+1}] = \sum_{i=1}^{\lfloor I/2 \rfloor} \hat{p}_{\tau+1}^{(i)} - \sum_{i=\lfloor I/2 \rfloor + 1}^I \hat{p}_{\tau+1}^{(i)}$$

Using this and a threshold which is calculated by the combination of the current price and standard deviation of the trading position value:

- if $g_{\tau+1} > v_\tau + c\sigma_\tau^v$, go long
- if $g_{\tau+1} < v_\tau - c\sigma_\tau^v$, go short
- if $v_\tau - c\sigma_\tau^v \leq g_{\tau+1} \leq v_\tau + c\sigma_\tau^v$, no trade

The researchers of the paper evaluated their trading strategy for 334 days and a moving window of 3 years, 1096 observations, every day to estimate the parameters for the dynamic factor model. We can see in Figure 2.14 that the strategy was able to consistently generate a profit even when considering transaction costs.

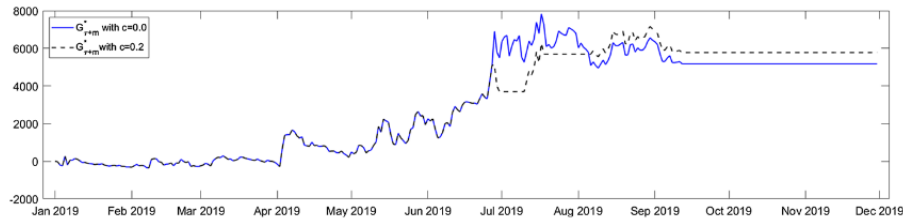


Figure 2.14: Net gains taking transaction fees into account [59]

Chapter 3

Uniswap

3.1 Overview

Uniswap is a decentralized exchange protocol built on the Ethereum blockchain. It allows users to trade ERC-20 tokens directly from their wallets without the need for intermediaries or traditional order books. Uniswap utilizes automated market-making, where liquidity providers contribute funds to liquidity pools, earning fees on trades made in the pool. The protocol employs a mathematical formula called the constant product formula to maintain balanced token ratios in the pool. When users want to make a trade, Uniswap calculates the conversion based on the pool's token ratios and executes the trade through a smart contract.

3.2 How it works

3.2.1 Automated Market Maker (AMM) Model and Liquidity Pools

The AMM model is a system that replaces traditional order books with liquidity pools to facilitate trading between different tokens. Automated Market Makers do this with the aid of liquidity pools. They are pools of tokens contributed by liquidity providers (LPs) to facilitate trading between different tokens within the exchange. One paper describes pools as “a smart contract that holds at least two cryptoassets and allows trading through depositing a token of one type and thereby withdrawing tokens of the other type” [61]. These liquidity pools are maintained by smart contracts on Ethereum and Liquidity Providers (LPs). Liquidity providers are individuals that voluntarily contribute an equal amount of cryptocurrencies liquidity pools, for example, in a pool for trading ETH and DAI, LPs would contribute an equal value of ETH and DAI tokens. LPs are incentivised to do provide liquidity in exchange for a share of the trading fees that occur in the liquidity pool. LPs can later withdraw their shares along with the accumulated fees. On Uniswap V3, the fee is 0.3% on Ethereum, however they can be any of 0.01%, 0.05%, 0.3%, or 1% depending on blockchain network.

3.2.2 Constant Product Formula

To maintain a balanced ratio between the tokens in the pool and price any swaps, the AMM model relies on a mathematical formula called the constant product formula ($xy = k$). As trades occur, the product of the token balances remains constant. When one token's value increases, its proportion in the pool decreases, ensuring an automatic adjustment in prices. When a trade is executed, the change in prices can be described in this formula $(x + \Delta x)(y + \Delta y) = k$. Hence, after rearranging: $\Delta y = \frac{k}{x + \Delta x} - y$. Under this model, the balance of the tokens in a liquidity pool can never be depleted as the token will get infinitely more expensive as the reserves approach 0.

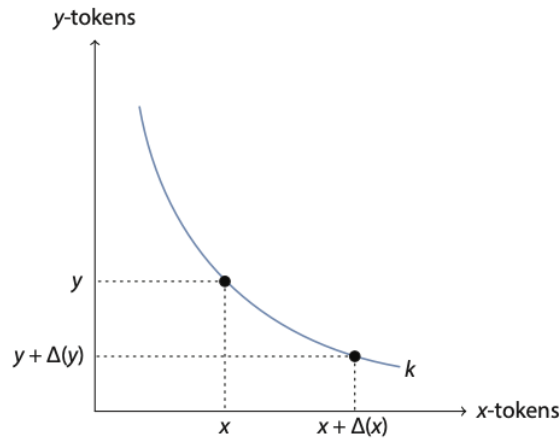


Figure 3.1: Constant Product Formula [61]

3.2.3 Slippage

As seen in Figure 3.2, we can see how the constant product formula is used in Uniswap. Uniswap also allow users to set slippage tolerance levels, which determine the maximum acceptable difference between the expected and executed prices. Slippage refers to the difference between the expected price of a trade and the executed price due to market volatility and liquidity conditions. Slippage happens because the constant product formula adjusts prices based on the ratio of tokens in the pool. As trades are executed, the token balances change, and the prices change accordingly. Thus, larger trades can cause more significant price impact, resulting in slippage.

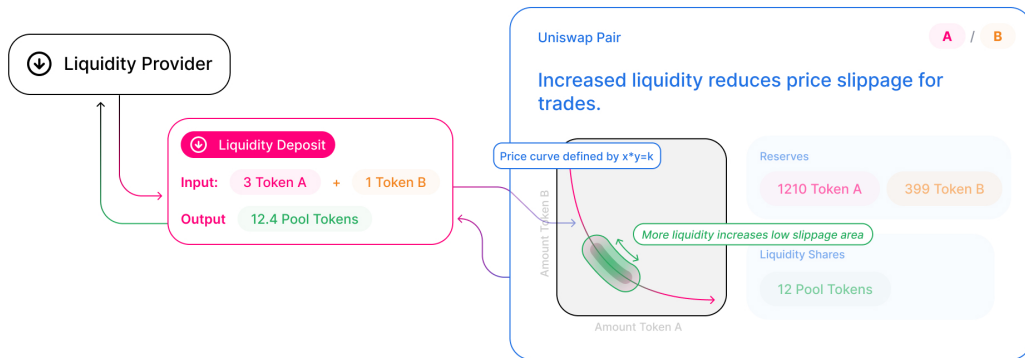


Figure 3.2: How Uniswap works [62]

Chapter 4

Aave

4.1 Overview

Aave is a decentralized lending and borrowing protocol built on the Ethereum blockchain. It enables users to lend and borrow a wide range of cryptocurrencies directly, without the need for intermediaries such as banks. Aave operates through liquidity pools and smart contracts, providing a secure, transparent, and efficient platform for decentralized finance (DeFi) activities.

Users can deposit their cryptocurrency assets into Aave's liquidity pools and earn interest on their deposits. These funds contribute to the available liquidity for borrowers. On the other hand, borrowers can use their deposited assets as collateral to borrow other assets from the pool. The amount they can borrow is determined by the value of their collateral and specific borrowing parameters set by the protocol.

4.2 How it works

4.2.1 Liquidity Pools

The fundamental mechanism that enables Aave's functionality of lending and borrowing is liquidity pools. Users deposit their cryptocurrency assets into Aave's liquidity pools. These assets serve as collateral and contribute to the overall liquidity of the protocol. The deposited assets in the liquidity pools create reserves of available liquidity. These reserves are utilized to fulfill borrowing demands from other users within the Aave ecosystem.

4.2.2 Lending and Borrowing

Lending works by lenders depositing their cryptocurrency assets into Aave's liquidity pools. These assets act as collateral and are represented by interest-bearing tokens called aTokens. The aTokens represent the user's share of the deposited assets. Interest is earned immediately and is accrued in real-time and reflected in an increase in the quantity of aTokens held by the depositor.

Borrowing works by borrowers using their deposited assets as collateral to borrow other assets from the liquidity pools. The value of the collateral determines the borrowing capacity. The borrowing capacity is calculated by parameters set by Aave, one of which is the maximum loan-to-value (LTV) ratio. Once the borrower requests a loan, if the borrower's collateral meets the necessary requirements, they can proceed with the loan and the borrowed funds are transferred to the borrower's wallet. In addition to this, borrowers pay interest on the borrowed amount based on the prevailing interest rates. Aave offers both variable and stable interest rates for borrowers. Variable rates fluctuate based on market dynamics, while stable rates remain fixed. This flexibility allows users to

choose the borrowing option that best suits their needs. The interest on a loan is accrued in real-time, second by second, and the borrower decides when to repay it, as long as the loan is safe from liquidation. Liquidation is what happens if the value of a borrower's collateral falls below the liquidation threshold due to market volatility or other factors, the collateral may be liquidated. Liquidators can purchase the collateral at a discounted price to ensure the solvency of the liquidity pool.

Chapter 5

Key Decisions and Outline of Trading Strategies

5.1 Liquidity Pools

5.1.1 Cointegrated Liquidity Pools

5.2 Strategy

5.2.1 Hedge Ratio

Mean Reversion Strategy

Kalman Filter

5.2.2 Volume of Trades

Chapter 6

Implementation

6.1 Backtesting System

6.2 Live Trading

Chapter 7

Evaluation

7.1 Number Arbitrage opportunities found

One of the ways I will be evaluating the strategies I implemented will be finding the number of opportunities that each strategy finds as well as highlighting opportunities that resulted in both profits and loss. Furthermore, I will show how the P&L is effected by transaction/gas fees by plotting both the P&L including and excluding any fees.

7.2 Return on theoretical trading

One major problem with this project is the timeline, from background research, I can see that most papers evaluate their strategies over at least a 6 month period. As this limits development time and evaluation time, the majority of the evaluation will be run by a back-testing system and evaluating each transaction cost at each time for any transaction that I theoretically make to simulate the trade environment. Using this trading environment, I will run each strategy and plot it's revenue and return.

7.3 Return on actual trading

In addition to the trading on the in-sample data points, I plan to trade and evaluate the performance and returns of each strategy in real time.

7.4 Sharpe Ratio

Another metric I will use is the Sharpe Ratio, the reason for this is that it is important to evaluate how risky the strategies are compared to the risk free rate. The metric is given by the formula below:

$$\text{Sharpe Ratio} = \frac{R_p - R_f}{\sigma_p}$$

Where R_p is the return of a portfolio, R_f is the risk-free rate and σ_p is the standard deviation of the portfolio. I will calculate the Sharpe Ratio for both in-sample and out-of-sample results.

7.5 Performance of Signal Generation

I also plan to evaluate how long it takes for each strategy to generate a signal when prices change. In analysing this metric, we may be able to find times where once the signal is created, the opportunity could be lost hence by looking at this metric, we may be able to find processes that may decrease the signal creation time.

Chapter 8

Ethical Issues

The ethics of cryptocurrencies are widely debated for reasons such as anonymity, leading it to be the choice of currency used by criminals and illegal institutions, volatility and lack of regulation. The high volatility makes cryptocurrencies and decentralised finance very risky for retail investors that don't have the technical or financial know-how making investing in cryptocurrencies.

Another aspect of cryptocurrencies that has raised ethical questions is the energy consumption and carbon dioxide emission from the mining of cryptocurrencies. Formal research about this has also been completed and found that 'approximately 69 million metric tons of CO₂ (Carbon dioxide) emission as a result of bitcoin mining' [63]. Thus, this is an ethical concern that I have thought about when designing the strategies so that the number of transactions that don't result in a profit, i.e. do not add value to the project, is limited.

In addition to the concerns above, although this project aims to find riskless profits, *'free lunches'*, it is not, in any form, of financial advice, and those who use the research or software that used in the development and research process to attempt to get favourable results, are liable for the losses or gains.

Chapter 9

Conclusion

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