



DHARMSINH DESAI UNIVERSITY, NADIAD

FACULTY OF TECHNOLOGY

B.TECH - Semester – VI(CE/IT)

SUBJECT: (CT 614) THEORY OF AUTOMATA & FORMAL LANGUAGE

Examination : 2nd Session

Date : 15/02/2013

Time : 1. Hr. 15 Mins.

Seat No. :

Day :

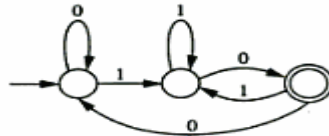
Max. Marks : 36

INSTRUCTIONS:

1. Figures to the right indicate maximum marks for that question.
2. The symbols used carry their usual meanings.
3. Assume suitable data, if required & mention them clearly.
4. Draw neat sketches wherever necessary.

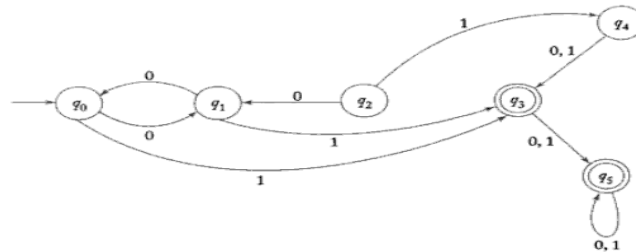
Q.1 Do as Directed [10]

- a. Define Pumping Lemma for regular language 2
- b. Disprove Statement with Counter example 2
If L_1 is regular, L_2 is non-regular then $L_1 \cup L_2$ is non-regular
- c. Identify which of the following languages over $\{0,1\}$ are regular 2
 - i) Set of odd length strings with middle symbol '0'
 - ii) Set of all strings x containing some non-null substring of the form 'ww'
- d. Show that the language $L = \{a^i b^j c^k \mid j > i + k\}$ cannot be written in the form $L = L_1 L_2 L_3$, 2
where L_1, L_2 , and L_3 are subsets of $\{a\}^*$, $\{b\}^*$, and $\{c\}^*$, respectively.
- e. Convert given Finite Automata into Chomsky Normal Form 2



Q.2 Answer the Following (Any Three) [12]

- a. Using pair wise distinguishable strings, show that language of Palindrome can't be accepted by Finite Automata
- b. Generate a Context-Free Grammar for $L = \{a^i b^j c^k \mid i < j + k\}$
- c. Define Ambiguous grammar. Describe "Dangling Else" problem.
- d. Minimize the Finite Automata



Q.3 Answer the Following [14]

- a. State & prove Kleene's theorem Part-I for any two operations of regular language 8
- b. Convert the NFA- \wedge to NFA shown in Figure. 4

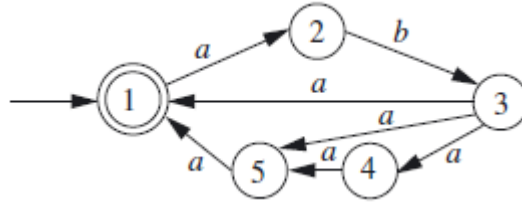
- c. Let $M = \langle Q, \Sigma, q_0, A, \delta \rangle$ be any NFA- \wedge . Let $M_1 = \langle Q, \Sigma, q_0, A_1, \delta_1 \rangle$ be the equivalent NFA. Is it correct to define $\delta_1(q, a) = \delta^*(q, a)$ for some q belongs to Q and a belongs to Σ^* ? 2

OR

Q.3 Answer the Following

[14]

- a. For any NFA $M = (Q, \Sigma, q_0, A, \delta)$ accepting a language L subset of Σ^* , there is a FA $M_1 = (Q_1, \Sigma, q_1, A_1, \delta_1)$ that also accepts L **8**
- b. Convert Non-Deterministic Finite Automata (NFA) to DFA shown in Figure. **4**



- c. What is the relation if any, between Union of \wedge -closures and \wedge -closures of Union? Justify your answer with a suitable example. **2**