



DHARMSINH DESAI UNIVERSITY, NADIAD
FACULTY OF TECHNOLOGY
B.TECH. SEMESTER VI [CE\IT]

SUBJECT: (CT614) THEORY OF AUTOMATA AND FORMAL LANGUAGES

Examination	: Third Sessional	Seat No.	: _____
Date	: 02/04/2013	Day	: Tuesday
Time	: 12 to 1.15	Max. Marks	: 36

INSTRUCTIONS:

1. Figures to the right indicate maximum marks for that question.
2. The symbols used carry their usual meanings.
3. Assume suitable data, if required & mention them clearly.
4. Draw neat sketches wherever necessary.

Q.1 Do as directed.

- (a) Define: Acceptance by PDA. [2]
- (b) Eliminate left recursion problem from the given grammar. [2]
 $S \rightarrow S_1 \$$ $S_1 \rightarrow S_1 A \mid \Lambda$ $A \rightarrow Aa \mid b$
- (c) Describe Chomsky hierarchy for languages. [2]
- (d) State true or false and Justify. [2]
 1. Turing machine can accept Regular language.
 2. 'C' language is not a CFL.
- (e) Does multi-tapes increase the power of a TM compared to single tape TM? Justify your answer. [2]
- (f) Define Context Sensitive Grammar and Phrase Structure Grammar. [2]

Q.2 Attempt *Any Two* from the following questions. [12]

- (a) Construct a Top down PDA for $\{x \text{ belongs to } \{a,b\}^* \mid x \text{ is a non Palindrome}\}$
- (b) Design a PDA for the language $L = \{a^i b^j c^k \mid i, j, k \geq 1 \text{ \& } i+j = k\}$.
- (c) Using Pumping Lemma prove that $\{a^i b^j c^k \mid i < j < k\}$ is not a context free language.

- Q.3**
- (a) Define Encoding function of UTM. Construct TM that accept $L = \{a^n b^n \mid n \geq 0\}$ and encode TM. [6]
 - (b) Construct **Insert(a)** TM which changes the tape content 'xy' to xay. Where **x** belongs to $(\sum \cup \{\Delta\})^*$, **y** belongs to \sum^* and **a** belongs to $\sum \cup \{\Delta\}$. [6]

OR

- Q.3**
- (a) Construct a TM that computes Reverse function of the input string [6]
 - (b) Discuss Non Deterministic Turing Machine. Construct and Explain a single tape / Multi tape version of 'Execute' 'Sub Turing Machine' with suitable example. [6]

Let $L = \{a^i b^j c^k \mid i, j, k \geq 1 \text{ and } i + j = k\}$

- Find a PDA (which accepts via final state) that recognizes L .
- Find a PDA (which accepts via empty stack) that recognizes L .

Solution

- Hint: Push a 's and b 's into the stack and match them with each other and clear.

The PDA $M = (\{q_0, q_1, q_2, q_3\}, \{a, b, c\}, \{a, b, \$\}, \delta, q_0, \$, \{q_3\})$

Transitions are:

$\delta(q_0, a, \$) \text{ contains } (q_0, a \$)$

$\delta(q_0, a, a) \text{ contains } (q_0, aa)$

$\delta(q_0, b, a) \text{ contains } (q_1, ba)$

$\delta(q_1, b, b) \text{ contains } (q_1, bb)$

$\delta(q_1, c, b) \text{ contains } (q_2, \epsilon)$

$\delta(q_2, c, a) \text{ contains } (q_2, \epsilon)$

$\delta(q_2, c, b) \text{ contains } (q_2, \epsilon)$

$\delta(q_2, \epsilon, \$) \text{ contains } (q_3, \epsilon)$

- The machine M above is the required PDA, q_3 is the final state and all the transitions remain the same as for (i). This machine accepts L via both by empty stack and final state. Note that this is a DPDA.