



Examination : Third Sessional

Seat No. :

Date : 9/10/2017

Day : Monday

Time : : to :00

Max. Marks : 36

**INSTRUCTIONS:**

1. Figures to the right indicate maximum marks for that question.
2. The symbols used carry their usual meanings.
3. Assume suitable data, if required & mention them clearly.
4. Draw neat sketches wherever necessary.

- Q.1 Do as directed. [12]
- (a) Turing machine is more powerful than: [01]  
1) Finite automata 2) Push down automata 3) Both (1) and (2) 4) None of these
- (b) Which of the following statement is wrong? [01]  
1) A Turing machine(TM) cannot solve halting problem  
2) Set of recursively enumerable language is closed under union  
3) A finite state machine with 3 stacks is more powerful than finite state machine with 2 stacks  
4) Context sensitive grammar can be recognized by a linearly bounded memory machine
- (c) Context free language can be recognized by – [01]  
1) Finite State Automaton 2) Linear bounded automaton  
3) Pushdown automaton 4) Both 2 and 3
- (d) Assume statements S1 and S2 defined as : [01]  
S<sub>1</sub> : L<sub>2</sub>-L<sub>1</sub> is recursive enumerable where L<sub>2</sub> and L<sub>1</sub> are recursive and recursive enumerable respectively.  
S<sub>2</sub> : The set of all Turing machines is countable.  
Which of the following is true ?  
1) S<sub>1</sub> is correct and S<sub>2</sub> is not correct. 2) Both S<sub>1</sub> and S<sub>2</sub> are correct.  
3) Both S<sub>1</sub> and S<sub>2</sub> are not correct. 4) S<sub>1</sub> is not correct and S<sub>2</sub> is correct.
- (e) Consider three decision problems P<sub>1</sub>, P<sub>2</sub> and P<sub>3</sub>. It is known that P<sub>1</sub> is decidable and P<sub>2</sub> is undecidable. Which one of the following is TRUE? [02]  
1) P<sub>3</sub> is decidable if P<sub>1</sub> is reducible to P<sub>3</sub> 2) P<sub>3</sub> is undecidable if P<sub>3</sub> is reducible to P<sub>2</sub>  
3) P<sub>3</sub> is undecidable if P<sub>2</sub> is reducible to P<sub>3</sub> 4) P<sub>3</sub> is decidable if P<sub>3</sub> is reducible to P<sub>2</sub>'s complement
- (f) Which of the following statements is/are FALSE? [02]  
1) For every non-deterministic Turing machine, there exists an equivalent deterministic Turing machine.  
2) Turing recognizable languages are closed under union and complementation.  
3) Turing decidable languages are closed under intersection and complementation.  
4) Turing recognizable languages are closed under union and intersection.  
5) The halting problem for Turing machines is undecidable  
6) determining whether a context-free grammar is ambiguous is un-decidable  
7) Given two arbitrary context-free grammar, G<sub>1</sub> and G<sub>2</sub>, it is undecidable if L(G<sub>1</sub>) = L(G<sub>2</sub>)  
8) Given two regular grammars G<sub>1</sub> and G<sub>2</sub>, it is undecidable whether L(G<sub>1</sub>) = L(G<sub>2</sub>)
- (g) L<sub>1</sub> = {a<sup>n+m</sup> b<sup>n</sup> c<sup>m</sup> | n, m ≥ 0}  
L<sub>2</sub> = {a<sup>n+m</sup> b<sup>n+m</sup> c<sup>m</sup> | n, m ≥ 0}  
L<sub>3</sub> = {a<sup>n+m</sup> b<sup>n+m</sup> c<sup>m+n</sup> | n, m ≥ 0} [02]  
Which of these languages are not CF.  
1) L<sub>1</sub> only 2) L<sub>3</sub> only 3) L<sub>1</sub> and L<sub>2</sub> 4) L<sub>2</sub> and L<sub>3</sub>
- (h) Let L<sub>1</sub> be a regular language and L<sub>2</sub> a deterministic CFL. L<sub>3</sub> is recursively enumerable but not recursive. Which one of the following statement is FALSE? [02]  
1) L<sub>1</sub> ∩ L<sub>2</sub> is a DCFL 2) L<sub>3</sub> ∩ L<sub>1</sub> is recursive  
3) L<sub>1</sub> ∪ L<sub>2</sub> is context-free 4) L<sub>1</sub> ∩ L<sub>2</sub> ∩ L<sub>3</sub> is recursively enumerable

Q.2 Attempt any two from following

- (a) Explain in detail all components of Chomsky hierarchy. [06]  
 (b) Write short note on : Universal Turing machine and other variations of TM. [06]  
 (c)

If  $L$  is a CFL, then  $\exists p$  (pumping length) such that  $\forall z \in L$ , if  $|z| \geq p$  then  $\exists u, v, w, x, y$  such

[01  
+  
05]

that  $z = uvwxy$  and

1.  $|vwx| \leq p$  2.  $|vx| > 0$  3.  $\forall i \geq 0, uv^iwx^iy \in L$ .

What is the above theorem popularly known as ?

Use it to prove the following language is non context free.

$L = \{a^i b^j c^k d^l \mid i, j \geq 0\}$ . Assume  $z = a^p b^p c^p d^p$  belongs to  $L$ , to prove your claim.

- Q-3 (a) Consider TM depicted in Figure1 and answer following questions with reference to it :-  
 In figure1, "B" indicates Blank symbol on tape.

- 1) Give the informal working logic of Turing machine. [02]
- 2) Is given TM Decider? Justify. [02]
- 3) What is the corresponding language used by TM? [02]
- 4) Give trace of the machine on appropriate input. [02]

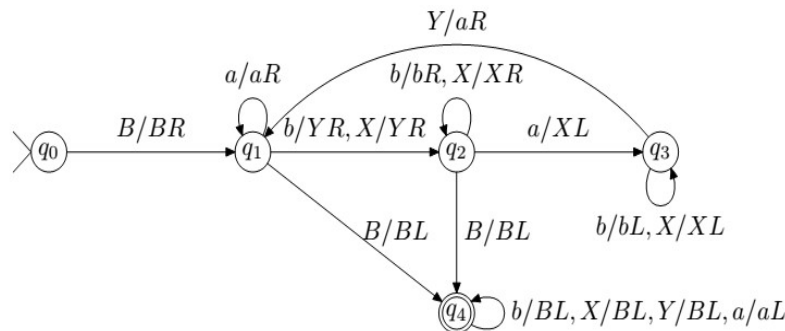


Figure 1

- (b) Explain with example any one Undecidable problem. [04]

OR

Q.3

- (a) How to find time complexity of Turing machine? Explain with example. [04]

(b)

Consider TM designed as follows-

$Q = \{q_0, q_1, q_2, q_3, q_f\}$  where  $q_0$  is the start state and  $q_f$  the only final state.

$\Sigma = \{1\}$   $\Gamma = \{1, i, \square\}$  where  $\square$  is the blank tape symbol, and  $\delta$  is defined as follows:

$$\begin{aligned}
\delta(q_0, 1) &= (q_0, 1, R) & \delta(q_1, 1) &= (q_2, 1, R) \\
\delta(q_0, \square) &= (q_1, \square, L) & \delta(q_2, 1) &= (q_2, 1, R) \\
\delta(q_1, 1) &= (q_1, 1, L), & \delta(q_2, \square) &= (q_3, 1, R) \\
\delta(q_3, \square) &= (q_1, 1, L) & \delta(q_1, \square) &= (q_f, \square, R)
\end{aligned}$$

- 1) Give the pictorial representation of above TM.
- 2) Explain the informal logic of above Turing machine.
- 3) Show trace on some appropriate input string.

[02]

[03]

[03]