



Master's method

We can solve recurrence relation using a formula denoted by master's method.

$$T(n) = a T(n/b) + F(n) \quad \text{where } n \geq d \text{ \& } d \text{ is some constant}$$

Then the master theorem can be stated for efficiency analysis as -

if $F(n)$ is $\Theta(n^d)$ where $d \geq 0$ in the recurrence relation then

- 1) $T(n) = \Theta(n^d)$ if $a < b^d$
- 2) $T(n) = \Theta(n^d \log n)$ if $a = b^d$
- 3) $T(n) = \Theta(n^{\log_b a})$ if $a > b^d$

- 1) Solve the following recurrence relation
 $T(n) = 4 T(n/2) + n$

Solⁿ - $T(n) = a T(n/b) + f(n)$

Now

$$f(n) = n \quad \text{ie. } n^1 \quad \text{Hence } d = 1$$

$$a = 4 \quad b = 2 \quad \& \quad a > b^d$$

$$4 > 2^1$$

$$\therefore T(n) = \Theta(n^{\log_b a})$$

$$= \Theta(n^{\log_2 4})$$

$$= \Theta(n^2) \quad \because \log_2 4 = 2$$



Another variation of Master theorem is for

$$T(n) = a T(n/b) + f(n) \quad \text{if } n \gg d.$$

1) If $f(n) = O(n^{\log_b a - \epsilon})$ then

$$T(n) = \Theta(n^{\log_b a})$$

2) If $f(n) = \Theta(n^{\log_b a} \log^k n)$ then

$$T(n) = \Theta(n^{\log_b a} \log^{k+1} n)$$

3) If $f(n) = \Omega(n^{\log_b a + \epsilon})$ then

$$T(n) = \Theta(f(n))$$

1) Solve the following recurrence relation.

$$T(n) = 2 T(n/2) + n \log n.$$

Solⁿ

$$f(n) = n \log n$$

$$a = 2 \quad b = 2$$

$$\log_b a = \log_2 2 = 1.$$

According to case 2 given in above Master theorem.

$$f(n) = \Theta(n^{\log_2 2} \log^1 n) \quad \text{i.e. } k=1.$$

$$\text{then } T(n) = \Theta(\log_b a \log^{k+1} n)$$

$$= \Theta(\log_2 2 \log^2 n)$$



$$T(n) = \Theta(n^1 \log^2 n)$$

$$\therefore T(n) = \Theta(n \log^2 n).$$

$$2) T(n) = 8T(n/2) + n^2.$$

Solⁿ

Here $f(n) = n^2$

$a=8$ $b=2$

$$a > b^d$$

$$\therefore \log_b a = \log_2 8 = 3.$$

Then according to case 1 of above given Master theorem -

$$\begin{aligned} f(n) &= O(n^{\log_b a - \epsilon}) \\ &= O(n^{\log_2 8 - \epsilon}) \\ &= O(n^{3 - \epsilon}) \end{aligned}$$

if we put $\epsilon = 1$ then

$$O(n^{3-1}) = O(n^2) = f(n).$$

then

$$\begin{aligned} T(n) &= \Theta(n^{\log_b a}) \\ &= \Theta(n^{\log_2 8}) \end{aligned}$$

$$\therefore \underline{T(n) = \Theta(n^3)}$$



$$3) T(n) = 9 T(n/3) + n^3.$$

solⁿ Here $a = 9$ $b = 3$ & $f(n) = n^3$

$$\log_b^a = \log_3^9 = 2. \quad \begin{array}{cc} a & 3^3 \\ 9 & = 9 \end{array}$$

According to case 3 in above master theorem -

$$\text{As } f(n) = \Omega(n^{\log_3^9 + \epsilon})$$

$$\text{i.e. } \Omega(n^{2+\epsilon}) \text{ \& we have } f(n) = n^3.$$

then to have

$$f(n) = \Omega(n), \text{ we must put } \epsilon = 1.$$

$$\therefore \text{ Then } T(n) = \Theta(f(n))$$

$$T(n) = \Theta(n^3).$$