

Agenda

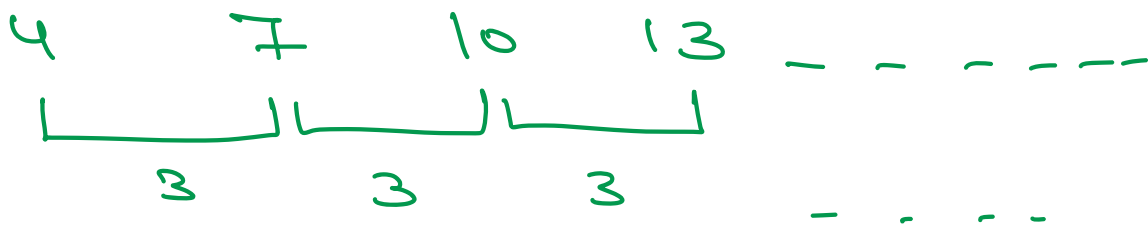
- AP and GP
- Time Complexity
 - Big O notation
 - Comparison of Orders of TC
- Space Complexity

Arithmetic Progression

1 2 3 4 : ... 10

$$\frac{n \times (n+1) / 2}{\cancel{10 \times (10+1)}}$$

$$\frac{n^2 + n}{2}$$



$$a \quad a+d \quad a+2d \quad a+3d \quad \dots$$

$$a=4$$

$$d=3$$

A.P

$$S_n = \left[2a + (n-1)d \right]$$

Ex-2

$$5 \quad 10 \quad 20 \quad 40 \quad \dots$$

x2 x2 x2

$$a \quad ar \quad ar^2 \quad ar^3 \quad \dots$$

$$5 \quad 5 \times 2 \quad 5 \times 2^2$$

$$a=5$$

$$r=2$$

$$\frac{a \times (r^n - 1)}{r - 1}$$

Concept - 1

$$N \rightarrow N/2 \rightarrow N/4 \dots \dots \dots 1$$

② $\log_2 N \Rightarrow$ No. of Steps

$$N \quad N/2 \quad N/4 \quad \dots \dots \dots N/2^k$$

Elemental \rightarrow

$$N/2^k = 1$$

$$N = 2^k$$

$$\log_2 N = \log_2 2^k$$

$$k \Rightarrow \log_2 N$$

$$2^3$$

$$\log_2 (2^3)$$

$$3 \log_2 2$$

$$3 \times 1$$

$$\log_a a \rightarrow 1$$

$$a \rightarrow 1$$

but at Every step, reduce

the number to its half ✓

$\log_2 N$ → Binary Search
Binary Tree
Merge Sort
Heaps

Q11

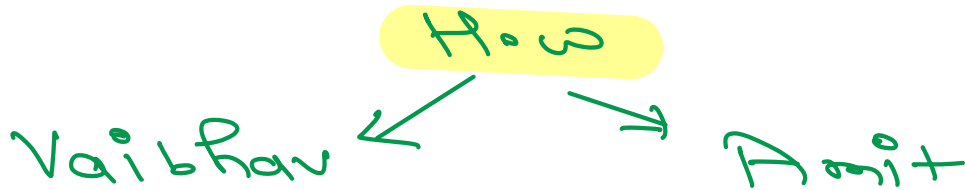
$$1 \rightarrow N$$

but Every time we multiply
curr-num with 2
Total number of Step to
reach to N?

$$\begin{matrix} [3, & 10] \Rightarrow (6-2+1) \\ a & b & \Rightarrow (10-3)+1=8 \end{matrix}$$

Time - Complexity

- measure of Efficiency
- Relationship between number of Ops vs Size of input



2 sec

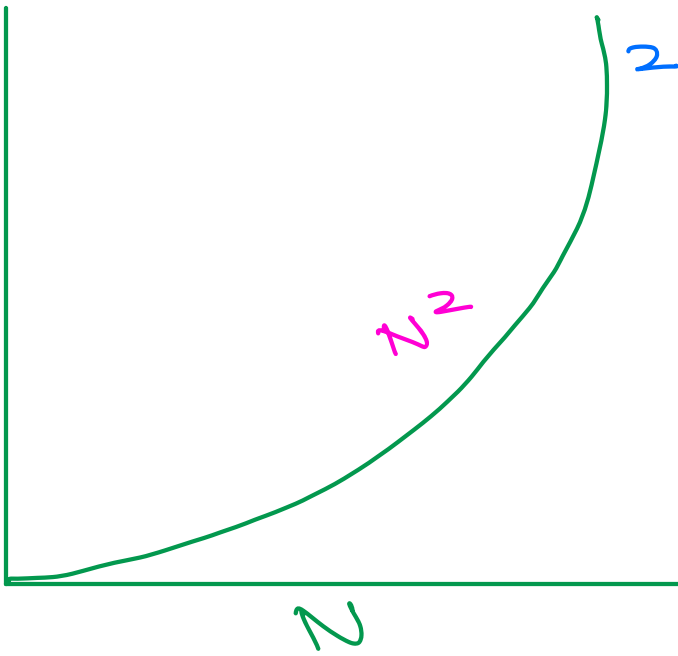
3 sec

3 GR2

1 GR2

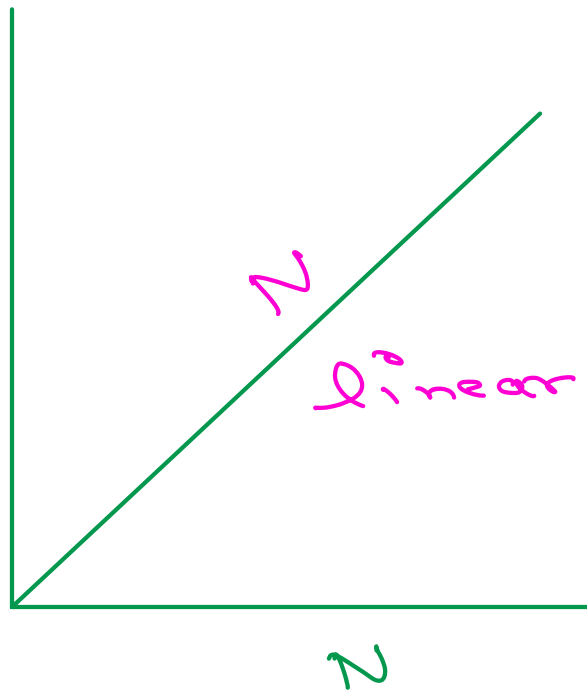
we cannot say?

$N \cdot y$
ops



Variable

y



Amits

10 \rightarrow 100

1000 \rightarrow 1,000,000

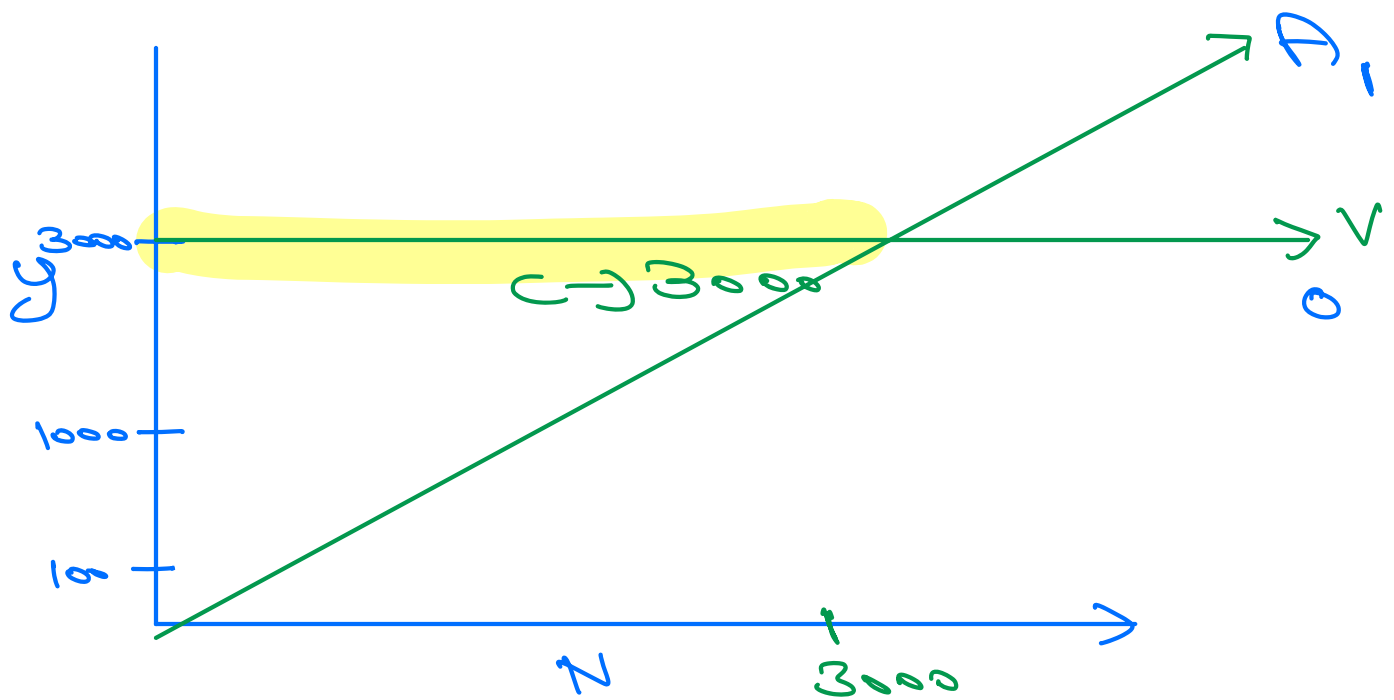
10 \rightarrow 10

1000 \rightarrow 1000

1 million order

million

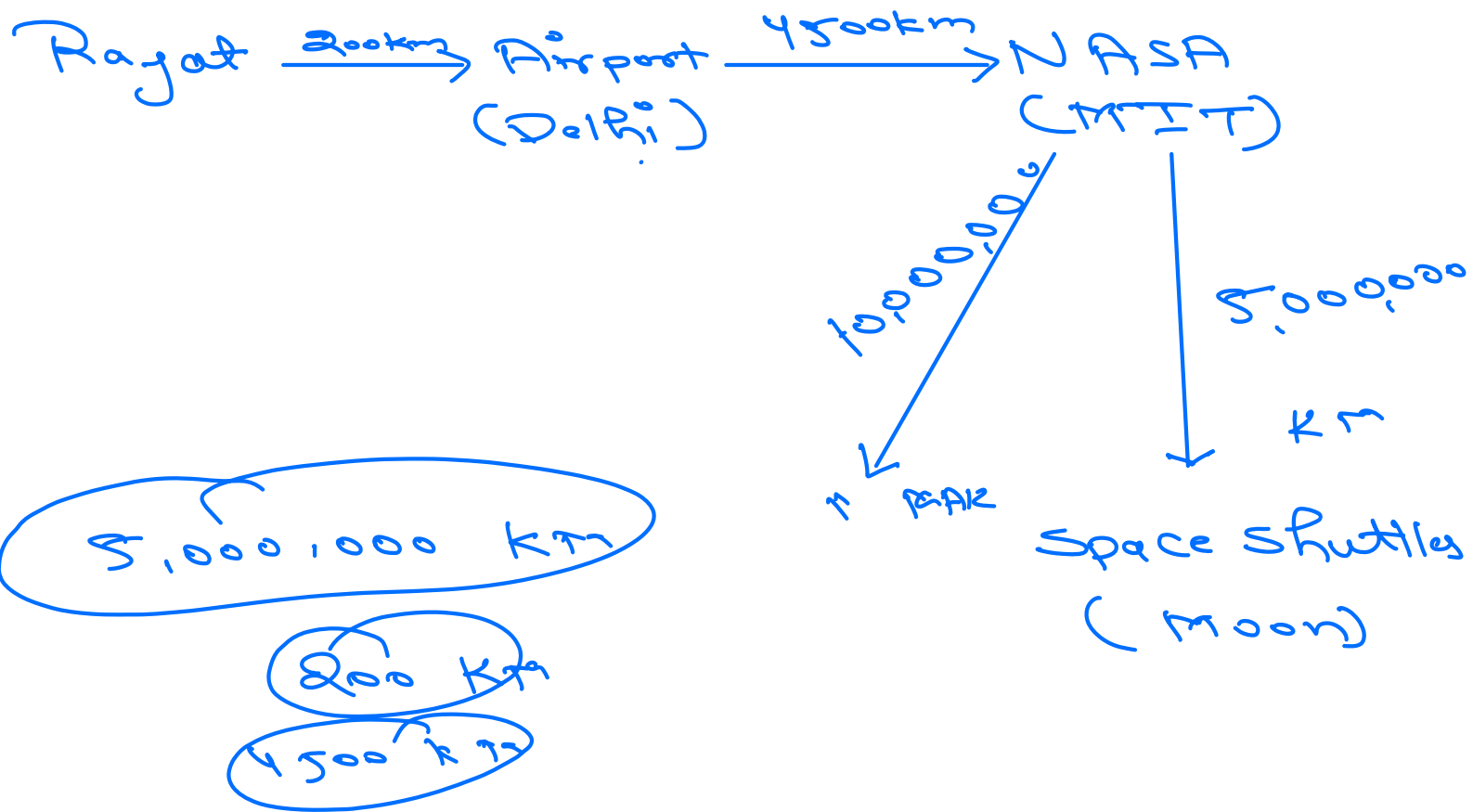
100 million



Long Term $\Rightarrow V$

$\Rightarrow TC \Rightarrow$ measures the trend
ignoring constant factors
and Co-efficient

Big O notation \Rightarrow Only the
highest poly
degree



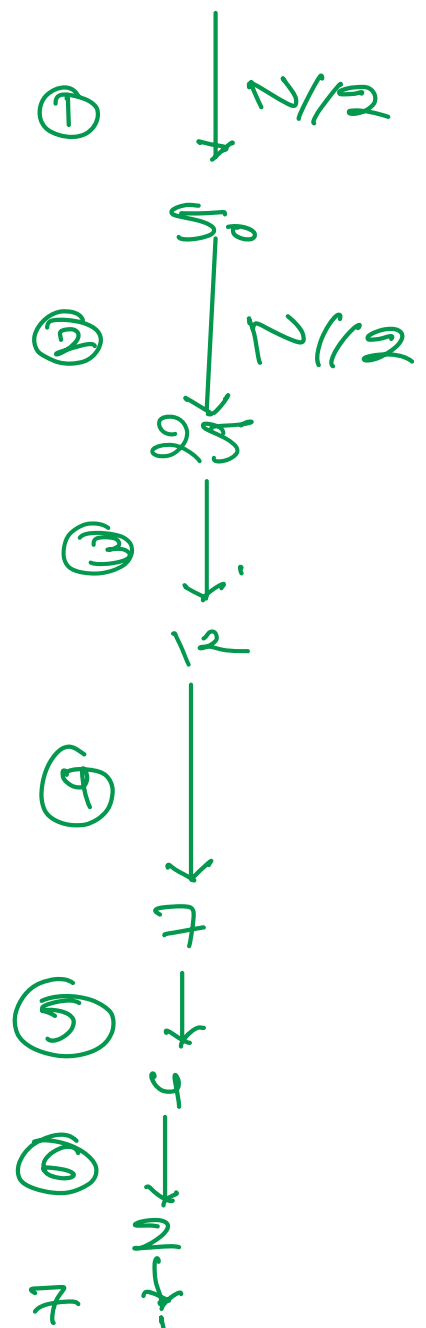
* Big O Notation only focuses on the highest polynomial Degree i.e. Trend of the Equation

$$TC \Rightarrow 10n^2 + 3n + 50$$

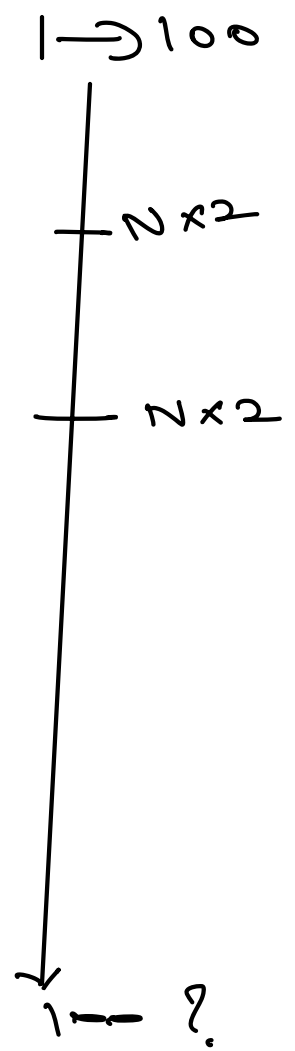


$$O(n) \Rightarrow \cancel{n^2} \Rightarrow n^2$$

$$N \Rightarrow 100$$



$$\log_2 100$$



```
def fun(N):
    s = 0
    for i in range(1, N+1):
        s += i
    return s
```

Annotations: $\rightarrow ①$ (for `s = 0`), $\rightarrow ①$ (for `for i in range(1, N+1):`), $\rightarrow N-1+1$ (for `s += i`), $\rightarrow ①$ (for `return s`).
 Formula: $b - a + 1$

$$N - \cancel{1} + 1 + 1 + 1$$

$$TC \Rightarrow N + \cancel{2}$$

$$O \Rightarrow N$$

$$[a \rightarrow b] \\ \rightarrow \\ b - a + 1$$

```
def fun(N):
    sum = 0
    for i in range(1, N+1, N//2):
        sum += i
```

$$\Rightarrow C \Rightarrow O(1)$$

for i in range(1, 100, 50)

$$\downarrow \\ 2$$

$$1, 100, 500 \\ \downarrow \\ 2$$

```

def fun(N):
    i = N
    while i >= 1:
        i = i // 2

```

	Before i	After i
①	N	$N//2$
②	$N//2$	$N//2^2$
③	$N/2^2$	$N/2^3$
⋮		
④		$N/2^k$

↓
1

$$O(\log_2 N)$$

$$O(\infty)$$

```
def fun(N):
    s = 0
    for i in range(N):
        for j in range(N):
            s += j
    return s
```

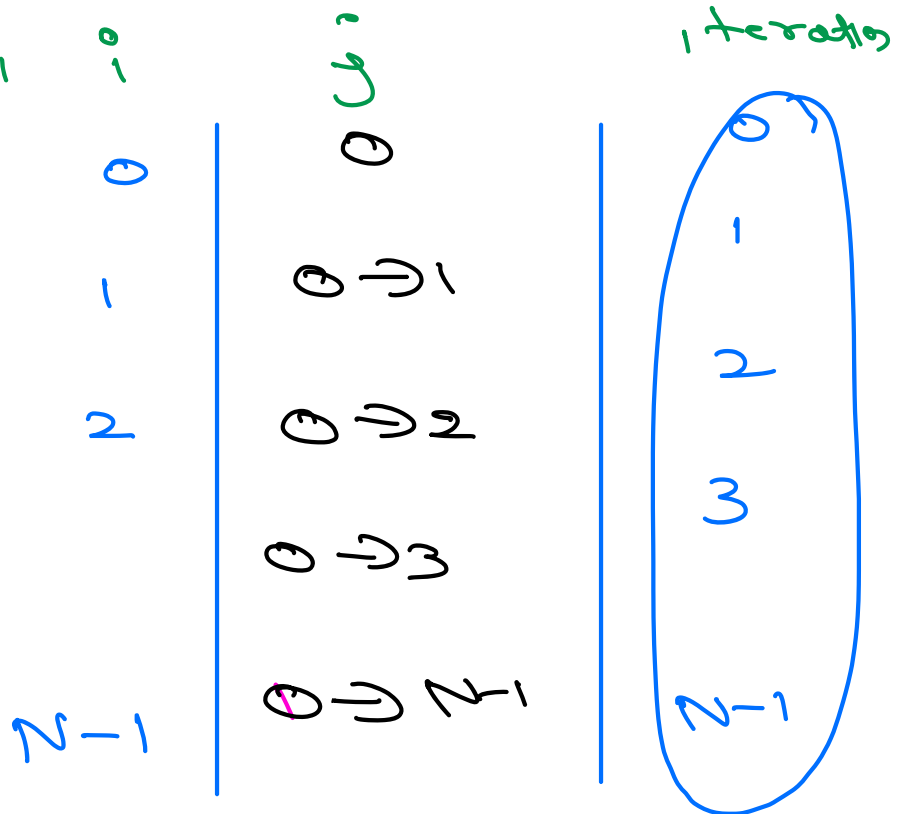
i	j	iteration
①	1 → N	1
②	1 → N	2
③	1 → N	3
⋮	⋮	⋮
N	1 → N	N

$N + N + \dots + N$
 $\underbrace{\hspace{10em}}_{N \times N \Rightarrow N \times 2}$

i → 1, 2	①	⑤	$\left. \begin{matrix} 1 \rightarrow 2 \\ 2 \rightarrow 2 \end{matrix} \right\} 2 + 2$ 4
j → 1, 2	①	②	

$2 + 2 + 2 + \dots + N \Rightarrow N \times 2$

```
def fun(N):
    s = 0
    for i in range(N):
        for j in range(0, i):
            s += j
    return s
```



0 1 2 3 ... N-1

$$\frac{n \times (n+1)}{2} \Rightarrow \frac{n^2}{2} + \frac{n}{2}$$

$O(n^2)$

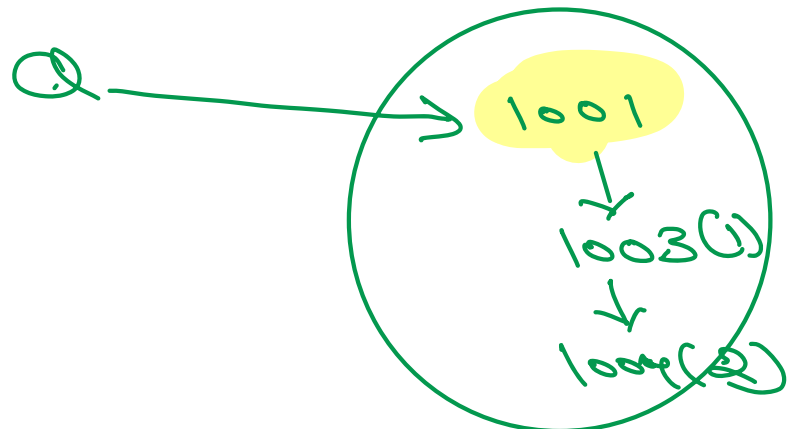
Space Complexity

- ⑤ Measure Extra Memory a program/Algo needs to complete Execution
- ⑤ It's Analysis of memory usage w.r.t Input Size

Q = []

Q.append(1)

Q.append(2)



T.C

⑥	Constant	C
⑦	Logarithmic	$\log_2 n$
⑧	Linear	n
⑨	Linear Logarithmic	$n \log_2 n$
⑩	Quadratic	n^2
⑪	Exponential	2^n or n^n

Higher Order Function

H O F

① A function that returns a function

```
def pow(x):  
    def func(n):  
        return x ** n  
    return func
```

func2 = pow(2)

pow(2) →

def func(n)
 return 2^n

return func

func2 → func2(10) → 2^{10}

anotherFunc → pow(100)
↓

anotherFunc(50) → 100^{50}