# VeloxDB - A Locally Grown Database

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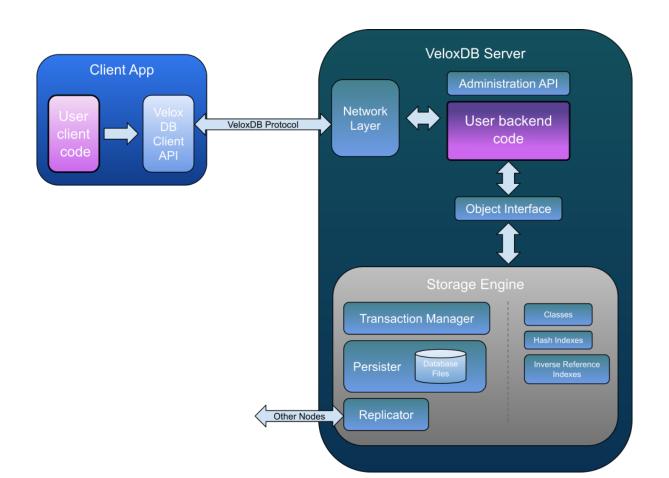


#### •Members:

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- •Fast (2.5 million trans/second on AWS c6a.8xlarge)
- ACID
- Object Oriented (ORM as first class citizen)
- .NET based
- Cross platform
- ◆Open Source (MIT license)
- High Availability
- •Read scale out
- •In Memory



- User provides model as C# abstract classes
- Each class has a pointer to unmanaged data.
- Getters are implemented by directly reading data

VeloxDB implements abstract classes with code generation at runtime

Setters accumulate changes for writing to disk and replication

#### Example model:

```
[DatabaseClass]
public abstract class Post : DatabaseObject
   [DatabaseProperty]
   public abstract string Title { get; set; }
   [DatabaseProperty]
   public abstract string Content { get; set; }
```

### Example database operation:

```
[DbAPIOperation]
public long CreatePost(ObjectModel om, PostDTO
post)
{
   Post newPost = om.CreateObject<Post>();
   newPost.Title = post.Title;
```

- Complexity
- Performance problems
- ●N+1 query problem
- •Incomplete abstraction
- Difficult to debug

management.

•Uses Multi Version Concurrency Control (MVCC) for transaction

Handles updates by treating them as inserts with new versions.

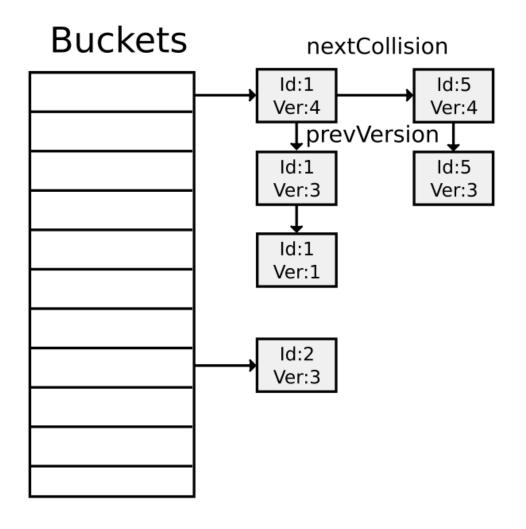
 Stores multiple versions of the same chiest (Multiversioning).

Stores multiple versions of the same object (Multiversioning).

T2: AddBook("To Kill a Mockingbird", "Harper Lee")
T3: AddBook("One Hundred Years of Solitude", "Gabriel Garcia Marquez")
T4: UpdateBookAuthor(1, "F. Scott Fitzegarld")

T1: AddBook("The Great Gatsby", "F. Scott Fitzegarld")

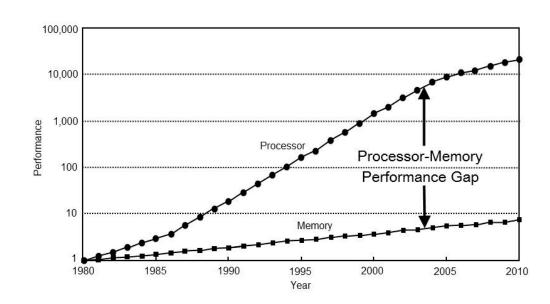
ld	Title	Author	Versio n
1	The Great Gatsby	F. Scott Fitzegarld	1
1	The Great Gatsby	F. Scott Fitzgerald	4
2	To Kill a Mockingbird	Harper Lee	2
3	One Hundred Years of Solitude	Gabriel Garcia Marquez	3

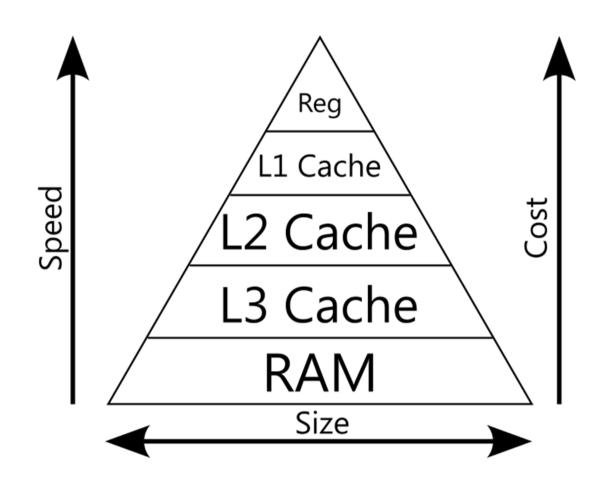


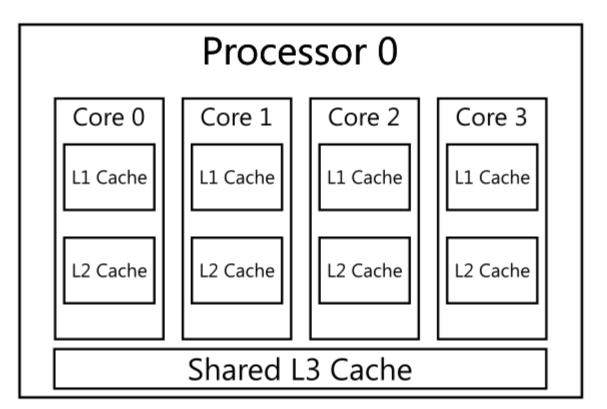
- •We use very fine grained locks employing a significant number of locks
- NET locks are objects, each object has 24-byte overhead
- Spinlock
- ○1 bit needed
- Suitable for rare contentions
- Pure userspace

```
void EnterLock()
   while(true)
       observedState = state;
       if (GetBit(observedState, LOCK_BIT) != LOCKED)
           // Calculate lock state with lock bit set
           lockedState = SetBit(observedState, LOCK_BIT, LOCKED)
           if (CompareAndExchange(ref state, lockedState, observedState) == observedState):
               return // Lock acquired
```

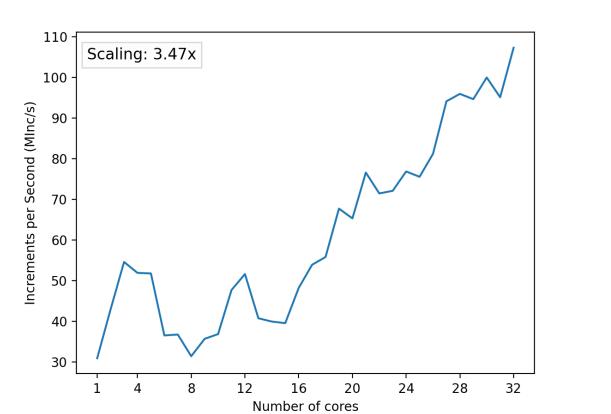
- •CPUs are improving at faster rate than memory
- •Bottlenecks:
- oBandwidth
- oLatency





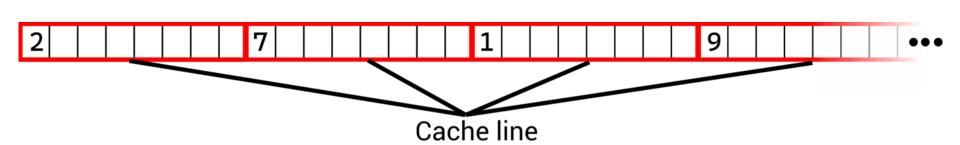


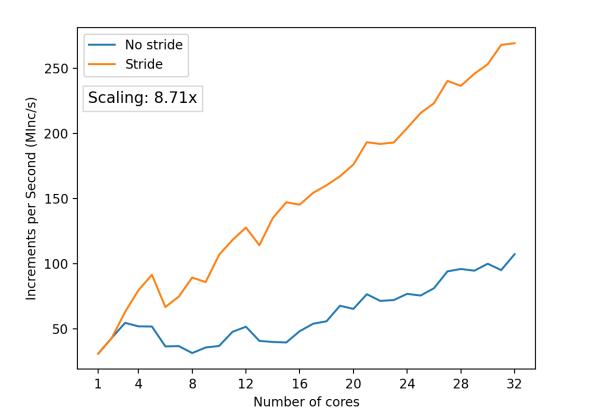


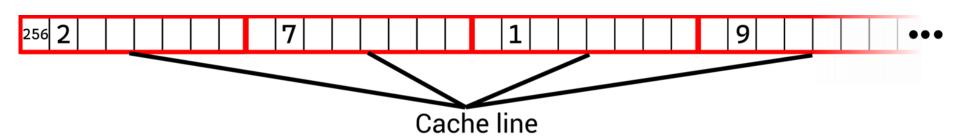


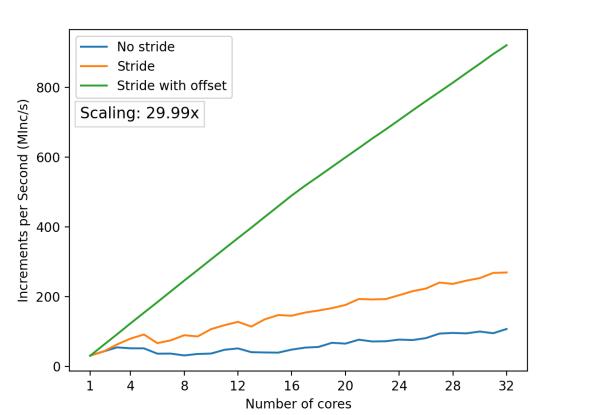


Cache line

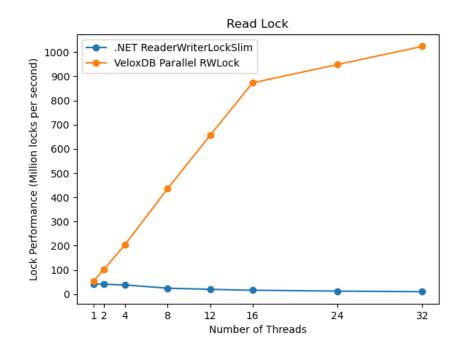


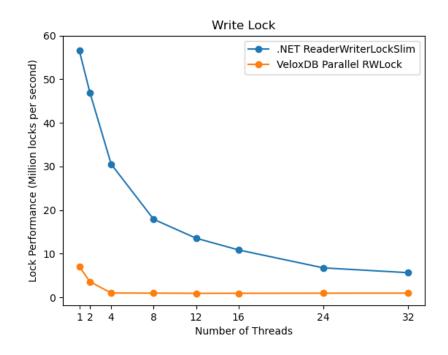






- Very fast read locking
- No contentions when reading
- Slower write
- Each CPU gets its own lock, on separate cache-line
- Read lock acquires only per-cpu lock
- Writes acquire all locks

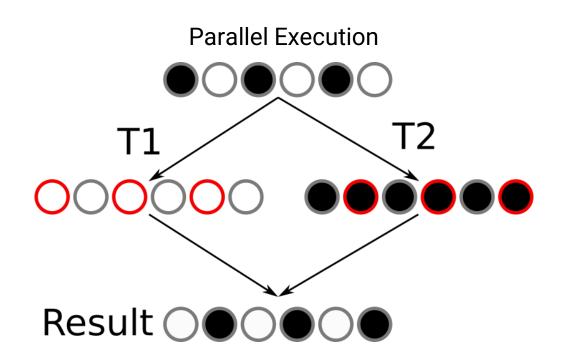




- •Strict Serializability Transactions view the database as if they are the only ones executing
- Transactions are executed in parallel
- Optimistic Concurrency Control
- Conflict When transactions cannot be executed concurrently.
- OWrite-Write conflict Two transactions trying to modify the same object
- oRead-Write conflict A transaction trying to write data that has been read by another transaction

- •Table with black/white balls
- •T1 Scans the table and flips black to white
- •T2 Scans the table and flips white to black





- Objects must keep track of every transaction that reads them
- Lists are expensive

Multiversion Concurrency Control (MVCC)

- •Small 8-byte structure that can hold 3 transaction IDs and a count
- •If more than 3 transactions read the same object, a dedicated list is created

- Better understanding of data lifecycle
- oData stored in database is going to be long lived
- Fewer distinct types
- oEasier to detect what data can be discarded
- •Reduces pressure on the .NET GC

•Each memory manager is composed of smaller per-cpu memory managers

•Freed objects are linked in free object list for reuse

Each table has its own memory manager

Fixed width allocators

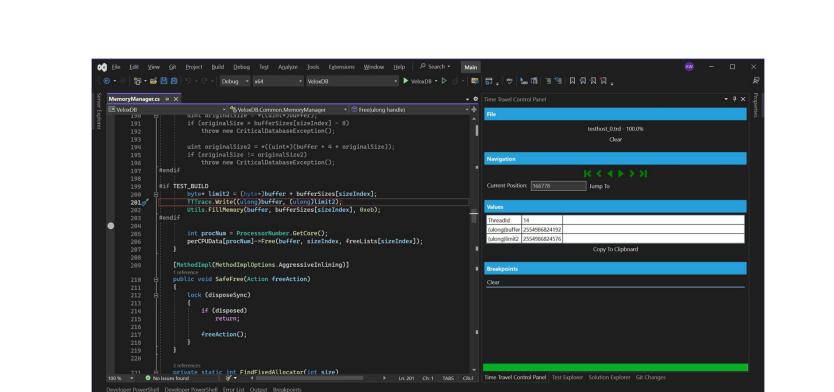
- •It's important to verify that database behaves correctly
- Large body of tests (unit, component, end to end)
- Most useful "stress" tests
- Concurrent
- Semantic Fuzzing
- Verified using Reference database

Challenging bugs (concurrent and memory corruption)

Hard to reproduce

•We built our own time travel solution

```
public void CreateObjectDiff(ClassObject* obj, ulong commonVersion, ChangesetWriter writer,
   IdSet partnerIds, List<long> otherIds, GenerateAlignDelegate alignDelegate)
   TTTrace.Write(TraceId, ClassDesc.Id, obj→id, obj→version, obj→IsDeleted, commonVersion);
    if (obj→IsDeleted)
       return;
   bool sbyMissing = partnerIds ≠ null && !partnerIds.Contains(obj→id);
    if (obj→version > commonVersion || sbyMissing)
       // If the standby deleted the object (due to split brain scenario) we need to pack an entire object which
       // will be achieved with the commonVersion being set to zero (this forces strings and blobs to be packed as well).
       TTTrace.Write(TraceId, ClassDesc.Id, obj→id, obj→version, commonVersion, sbyMissing);
       ulong version = sbyMissing ? 0 : commonVersion;
       alignDelegate(obj, stringStorage, blobStorage, writer, version);
       writer.LastValueWritten();
    else
       TTTrace.Write(TraceId, ClassDesc.Id, obj→id, obj→version, commonVersion);
       otherIds ??= new List<long>();
       otherIds.Add(obj→id);
```



## Questions