Zach Seymour CS 575 Theory Assignment #1 September 10, 2013

Part I

1 $\sum_{i=0}^{2} 3 = 9$

2
$$\prod_{i=0}^{2} = 27$$

3 $\log_3 1 = 1$

4 $\lg \lg 32 = \lg 5$

$$5 lg^2 8 = (lg8)^2 = 9$$

 $\mathbf{6} \quad \overline{\lim_{n \to \infty} (\frac{1}{2})^n = 0}$

$$7 \quad \lim_{n \to \infty} 2^n = \infty$$

$$8 \quad \frac{d}{dx} \ln x = \frac{1}{x}$$

$$9 \quad \frac{d}{dx}e^x = e^x$$

10 $4^{\log_3 n} = n^x$. What is x?

$$x \ln n = \ln \left(4^{\log_3 n} \right)$$
$$x = \frac{\ln \left(4^{\log_3 n} \right)}{\ln n}$$

11 Prove by induction that $1+3+5+7+\cdots+(2(n-1)+1)=n^2$. (Your proof should include the following: proof of base case, the induction hypothesis, what you are trying to prove, the proof. See document in Blackboard on induction proofs)

Proposition 1.1. $\sum_{i=0}^{n-1} (2i+1) = n^2$.

Proof. For n=1, we have $\sum_{i=0}^{0}(2i+1)=2\cdot 0+1=1$. So the base case is satisfied. Now, $\forall k>1$ we assume $\sum_{i=0}^{k-1}(2i+1)=k^2$. Therefore, we aim to show $\sum_{i=0}^{k}(2i+1)=(k+1)^2$. We have

$$\sum_{i=0}^{k} (2i+1) = k^2 + \sum_{i=k}^{k} (2i+1)$$
$$= k^2 + 2k + 1$$
$$= (k+1)^2$$

12 Use strong induction and the identity $F_i = F_{i-1} + F_{i-2}$ to prove that the *i*th Fibonacci number satisfies

$$F_i = \frac{\phi^i - \hat{\phi}^i}{\sqrt{5}}$$

where $\phi = \frac{1+\sqrt{5}}{2}$ and $\hat{\phi} = \frac{1-\sqrt{5}}{2}$. (Hint: $\phi^2 = \frac{3+\sqrt{5}}{2}$ and $\hat{\phi}^2 = \frac{3-\sqrt{5}}{2}$)

Proof. For i=1, we have

$$F_1 = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^1 - \left(\frac{1-\sqrt{5}}{2}\right)^1}{\sqrt{5}}$$
$$= \frac{\sqrt{5}}{\sqrt{5}}$$
$$= 1$$

And, for i=1, we have

$$F_2 = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^2 - \left(\frac{1-\sqrt{5}}{2}\right)^2}{\sqrt{5}}$$
$$= \frac{\sqrt{5}}{\sqrt{5}}$$
$$= 1$$

Now, let $i \geq 2$. Assume $\forall k, 1 \leq k \leq i, F_k = \frac{\phi^k - \hat{\phi}^k}{\sqrt{5}}$. We will show $F_{k+1} = \frac{\phi^{k+1} - \hat{\phi}^{k+1}}{\sqrt{5}}$. We know $F_{k+1} = F_k + F_{k-1}$. So, we have

$$\begin{split} F_{k+1} &= F_k + F_{k-1} \\ &= \frac{\phi^k - \hat{\phi}^k}{\sqrt{5}} + \frac{\phi^{k-1} - \hat{\phi}^{k-1}}{\sqrt{5}} \\ &= \frac{\left(\phi^k - \hat{\phi}^k\right) + \left(\phi^{k-1} - \hat{\phi}^{k-1}\right)}{\sqrt{5}} \\ &= \frac{\phi^k + \phi^{k-1} - \hat{\phi}^k - \hat{\phi}^{k-1}}{\sqrt{5}} \\ &= \frac{\phi^{k-1}(\phi + 1) - \hat{\phi}^{k-1}(\hat{\phi} + 1)}{\sqrt{5}} \\ &= \frac{\phi^{k-1}\phi^2 - \hat{\phi}^{k-1}\hat{\phi}^2}{\sqrt{5}} \\ &= \frac{\phi^{k+1} - \hat{\phi}^{k+1}}{\sqrt{5}} \end{split}$$

13
$$\sum_{i=1}^{n} n - 1 = n \cdot (n-1)$$

14
$$\sum_{i=1}^{n} 3^i = 3^1 + 3^2 + \dots + 3^n = \frac{1 - 3^{n+1}}{-2}$$

15 In how many different orders can x students sit around a round table? Two orders are identical if one can be rotated to form the other.

There are n! different arrangements, each with n rotations, so (n-1)! different orders.

Part II

- 1 CountMe has count n^4 .
- **2** CountMe has count $\left\lceil \frac{n}{2} \right\rceil$ for n even and $\left\lfloor \frac{n}{2} \right\rfloor$ for n odd.
- **3** CountMe has count $\lceil \frac{n}{2} \rceil$ for n odd and $\lceil \frac{n}{2} \rceil$ for n even.
- 4 Given the pseudo code below for bubble sort...Let length[A] = n. Derive the count of the number of times that the comparison (A[j] < A[j-1]) is executed by the algorithm. Show how you derived your answer.

If length[A] = n, the outer loop executes n-1 times. The inner loop executes $\sum_{i=2}^{n} i = \frac{n^2+n-2}{2}$ times. So, the overall instruction count is $(n-1)\left(\frac{n^2+n-2}{2}\right)$.

5 Suppose we are comparing implementations of bubble sort and merge sort on the same machine. For inputs of size n, bubble sort runs in $15n^2$ steps, while merge sort runs in $150n \lg n$ steps. For which values of n does bubble sort beat merge sort? (Hint: use Excel write a program, or try different values). What is the number of steps done by each for n = 1024?

Bubble sort beats merge sort for all values n, 1 < n < 59, using the following Mathematica code.

$$N[Solve[15 n^2 == 150*n*Log[2, n], n]]$$

For n = 1024, bubble sort takes 15,728,640 steps, while merge sort takes 1,536,000.

6 What is the smallest value of n such that an algorithm whose running time is $2^{20}n$ runs faster than an algorithm whose running time is 2^n on the same machine?

Once n is larger than 24, 2^n exceeds $2^{20}n$. So the smallest such n is 25.

Extra Credit

Prove by induction that the following program computes n! correctly for all integers $n \ge 1$.

Proof. For initialization, we need to show the program is correct when n = 1. Indeed, given the condition if (n==1) return 1, we correctly have 1! = 1.

For maintenance, we observe that for each call where n > 1, the program returns n multiplied by a call to the next lowest factorial, which is correct with regards to the formula for the factorial function, $n! = \prod_{i=1}^{n} i$.

Finally, for termination, we notice that, when repeatedly subtracting one from n, where n > 1, we will eventually have n = 1, whereby the program will terminate.