

# Annotated ML DSA Signature Algorithm

March-2026

## The Annotated ML DSA Algorithms

### Conversion Between Data Types

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**Algorithm 14**  $\text{CoeffFromThreeBytes}(b_0 : \mathbb{B}^1, b_1 : \mathbb{B}^1, b_2 : \mathbb{B}^1) \rightarrow z : \mathbb{Z}_q \cup \perp$ 

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Generates an element of  $\{0, 1, 2, \dots, q-1\} \cup \perp$ .

```
1:  $b'_2 \leftarrow b_2$ 
2: if  $b'_2 > 127$  then
3:    $b'_2 \leftarrow b'_2 - 128$  ▷ set the top bit of  $b'_2$  to zero
4: end if
5:  $z \leftarrow 2^{16} \cdot b'_2 + 2^8 \cdot b_1 + b_0$  ▷  $0 \leq z \leq 2^{23} - 1$ 
    $\triangleright \max(z) = 2^{16} \cdot (2^7 - 1) + 2^8 \cdot (2^8 - 1) + (2^8 - 1)$ 
    $\triangleright \quad \quad \quad = 2^{23} - 1$ 
6: if  $z < q$  then return  $z$  ▷ rejection sampling
7: else return  $\perp$ 
8: end if
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**Algorithm 15**  $\text{CoeffFromHalfByte}(b : [0, 15]) \rightarrow z : [-\eta, \eta] \cup \perp$ 

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Generates an element of  $\{-\eta, -\eta + 1, \dots, \eta\} \cup \{\perp\}$ .

$\eta \in \{2, 4\}$ .

```
1: if  $\eta = 2$  and  $b < 15$  then return  $2 - (b \bmod 5)$  ▷ rejection sampling from  $\{-2, \dots, 2\}$ 
    $\triangleright$  case 1: ML-DSA-44 and ML-DSA-87
2: else
    $\triangleright$  case 2: ML-DSA-65
3:   if  $\eta = 4$  and  $b < 9$  then return  $4 - b$  ▷ rejection sampling from  $\{-4, \dots, 4\}$ 
4:   else return  $\perp$ 
5:   end if
6: end if
```

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## Pseudorandom Sampling

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**Algorithm 29**  $\text{SampleInBall}(\rho : \mathbb{B}^{\lambda/4}) \rightarrow c : R_{[-1,0,1]}$

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Samples a polynomial  $c \in R$  with coefficients from  $\{-1, 0, 1\}$  and *Hamming weight*  $\tau \leq 64$ .

```

1:  $c \leftarrow 0$ 
2:  $\text{ctx} \leftarrow \text{H.Init}()$   $\triangleright \text{H} \doteq \text{SHAKE256}$ 
3:  $\text{ctx} \leftarrow \text{H.Absorb}(\text{ctx}, \rho)$ 
    $\triangleright \text{len}(\rho)$  is 32 bytes in ML-DSA-44, 48 in ML-DSA-65, and 64 in ML-DSA-87.
4:  $(\text{ctx}, s : \mathbb{B}^8) \leftarrow \text{H.Squeeze}(\text{ctx}, 8)$ 
5:  $h : \{0, 1\}^{64} \leftarrow \text{BytesToBits}(s)$   $\triangleright h$  is a bit string of length 64
6: for  $i$  from  $256 - \tau$  to  $255$  do
    $\triangleright \tau = 39, 49, 60$  in ML-DSA-44, ML-DSA-65, and ML-DSA-87, respectively.
7:    $(\text{ctx}, j : \mathbb{B}^1) \leftarrow \text{H.Squeeze}(\text{ctx}, 1)$ 
8:   while  $j > i$  do  $\triangleright$  rejection sampling in  $\{0, \dots, i\}$ 
9:      $(\text{ctx}, j : \mathbb{B}^1) \leftarrow \text{H.Squeeze}(\text{ctx}, 1)$ 
10:  end while  $\triangleright j$  is a pseudorandom byte that is  $\leq i$ 
11:   $c_i \leftarrow c_j$ 
    $\triangleright c_j$  is a smaller-degree coefficient, pseudorandomly selected.
    $\triangleright c_i$  is a larger degree coefficient.  $c_i$  receives the value of  $C_j$ .
12:   $c_j \leftarrow (-1)^{h[i+\tau-256]}$ 
    $\triangleright$  access pattern:  $h[0], h[1], \dots, h[\tau]$  where  $39 \leq \tau \leq 60$ , and  $h : \{0, 1\}^{64}$ .
    $\triangleright$  This pseudorandom shuffling is performed  $\tau$  times in total.
13: end for
14: return  $c$ 
```

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**Algorithm 30**  $\text{RejNTTPoly}(\rho : \mathbb{B}^{34}) \rightarrow \hat{a} : T_q$

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Samples a polynomial  $\hat{a} \in T_q$ .

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1:  $j \leftarrow 0$ 
2:  $\text{ctx} \leftarrow \text{G.Init}()$   $\triangleright \text{G} \doteq \text{SHAKE128}$ 
3:  $\text{ctx} \leftarrow \text{G.Absorb}(\text{ctx}, \rho)$ 
4: while  $j < 256$  do
5:    $(\text{ctx}, s : \mathbb{B}^3) \leftarrow \text{G.Squeeze}(\text{ctx}, 3)$ 
6:    $\hat{a}[j] \leftarrow \text{CoeffFromThreeBytes}(s[0], s[1], s[2])$ 
7:   if  $\hat{a}[j] \neq \perp$  then
8:      $j \leftarrow j + 1$ 
9:   end if
10: end while
11: return  $\hat{a}$ 
```

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**Algorithm 31**  $\text{RejBoundedPoly}(\rho : \mathbb{B}^{66}) \rightarrow a : R_{[-\eta, \eta]}$ 

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Samples an element  $a \in R$  with coefficients in  $[-\eta, \eta]$  computed via rejection sampling from  $\rho$ .

```
1:  $j \leftarrow 0$ 
2:  $\text{ctx} \leftarrow \text{H.Init}()$   $\triangleright \text{H} \doteq \text{SHAKE256}$ 
3:  $\text{ctx} \leftarrow \text{H.Absorb}(\text{ctx}, \rho)$ 
4: while  $j < 256$  do
5:    $z : \mathbb{B}^1 \leftarrow \text{H.Squeeze}(\text{ctx}, 1)$ 
6:    $z_0 \leftarrow \text{CoeffFromHalfByte}(z \bmod 16)$ 
7:    $z_1 \leftarrow \text{CoeffFromHalfByte}(z/16)$ 
8:   if  $z_0 \neq \perp$  then
9:      $a_j \leftarrow z_0$ 
10:     $j \leftarrow j + 1$ 
11:   end if
12:   if  $z_1 \neq \perp$  and  $j < 256$  then
13:      $a_j \leftarrow z_1$ 
14:      $j \leftarrow j + 1$ 
15:   end if
16: end while
17: return  $a$ 
```

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## The Key Generation Algorithm

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**Algorithm 6**  $\text{ML-DSA.KeyGen\_internal}(\xi)$ 

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```
1:  $(\rho, \rho', K) \leftarrow \text{H}(\xi \parallel \text{Integer2Bytes}(k, 1) \parallel \text{Integer2Bytes}(\ell, 1), 128)$ 
2:  $\triangleright$  expand seed
3:  $\hat{\mathbf{A}} \leftarrow \text{ExpandA}(\rho)$ 
4:  $(\mathbf{s}_1, \mathbf{s}_2) \leftarrow \text{ExpandS}(\rho')$ 
5:  $\mathbf{t} \leftarrow \text{NTT}^{-1}(\hat{\mathbf{A}} \circ \text{NTT}(\mathbf{s}_1)) + \mathbf{s}_2$ 
6:  $(\mathbf{t}_1, \mathbf{t}_0) \leftarrow \text{Power2Round}(\mathbf{t})$   $\triangleright$  compress  $\mathbf{t}$ 
7:  $\triangleright$  Power2Round is applied componentwise
8:  $pk \leftarrow \text{pkEncode}(\rho, \mathbf{t}_1)$ 
9:  $tr \leftarrow \text{H}(pk, 64)$ 
10:  $sk \leftarrow \text{skEncode}(\rho, K, tr, \mathbf{s}_1, \mathbf{s}_2, \mathbf{t}_0)$ 
11: return  $(pk, sk)$ 
```

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## The Annotated Key Generation Algorithm

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### Annotated Algorithm 6 ML-DSA.KeyGen\_internal( $\xi : \mathbb{B}^{32}$ )

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- 1:  $(\rho : \mathbb{B}^{32}, \rho' : \mathbb{B}^{64}, K : \mathbb{B}^{32}) \leftarrow H(\xi \parallel \text{Integer2Bytes}(k, 1) \parallel \text{Integer2Bytes}(\ell, 1), 128)$
  - 2:  $\triangleright$  expand seed
  - 3:  $\widehat{\mathbf{A}} : \mathbb{T}_q^{k \times \ell} \leftarrow \text{ExpandA}(\rho)$
  - 4:  $(\mathbf{s}_1 : R_m^\ell, \mathbf{s}_2 : R_m^k) \leftarrow \text{ExpandS}(\rho')$   
 $m \in [-2, 2]$  if ML-DSA-44 or ML-DSA-87  $\triangleright \eta = 2$   
 $m \in [-4, 4]$  otherwise  $\triangleright \eta = 4$
  - 5:  $\mathbf{t} : R_q^k \leftarrow \text{NTT}^{-1}(\widehat{\mathbf{A}} \circ \text{NTT}(\mathbf{s}_1)) + \mathbf{s}_2$
  - 6:  $(\mathbf{t}_1 : R_{q_1}^k, \mathbf{t}_0 : R_{q_0}^k) \leftarrow \text{Power2Round}(\mathbf{t})$   $\triangleright$  compress  $\mathbf{t}$   
 $\mathbf{t}_1 \in [0, 1023]$   $\triangleright$  10-bit value  
 $\mathbf{t}_0 \in [-4095, 4096]$   $\triangleright \bmod^{\pm} 2^d, d = 13, ([-2^{12} + 1, 2^{12}])$
  - 7:  $\triangleright$  Power2Round is applied componentwise
  - 8:  $pk \leftarrow \text{pkEncode}(\rho, \mathbf{t}_1)$
  - 9:  $tr : \mathbb{B}^{64} \leftarrow H(pk, 64)$
  - 10:  $sk \leftarrow \text{skEncode}(\rho, K, tr, \mathbf{s}_1, \mathbf{s}_2, \mathbf{t}_0)$
  - 11: **return**  $(pk, sk)$
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## The Annotated Signature Algorithm

The ML-DSA Signing (Internal) algorithm named **ML-DSA.Sign\_internal** in FIPS 204 standard is reproduced below for reference. The line numbers and the pseudocode matches exactly with the original version.

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**Algorithm 7** ML-DSA.Sign\_internal( $sk, M', rnd$ )

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```
1:  $(\rho, K, tr, \mathbf{s}_1, \mathbf{s}_2, \mathbf{t}_0) \leftarrow \text{skDecode}(sk)$ 
2:  $\widehat{\mathbf{s}}_1 \leftarrow \text{NTT}(\mathbf{s}_1)$ 
3:  $\widehat{\mathbf{s}}_2 \leftarrow \text{NTT}(\mathbf{s}_2)$ 
4:  $\widehat{\mathbf{t}}_0 \leftarrow \text{NTT}(\mathbf{t}_0)$ 
5:  $\widehat{\mathbf{A}} \leftarrow \text{ExpandA}(\rho)$ 
6:  $\mu \leftarrow \text{H}(\text{BytesToBits}(tr \parallel M', 64))$ 
7:  $\rho'' \leftarrow \text{H}(K \parallel rnd \parallel \mu, 64)$ 
8:  $\kappa \leftarrow 0$ 
9:  $(\mathbf{z}, \mathbf{h}) \leftarrow \perp$ 
10: while  $(\mathbf{z}, \mathbf{h}) = \perp$  do
11:    $\mathbf{y} \in R_q^\ell \leftarrow \text{ExpandMask}(\rho'', \kappa)$ 
12:    $\mathbf{w} \leftarrow \text{NTT}^{-1}(\widehat{\mathbf{A}} \circ \text{NTT}(\mathbf{y}))$ 
13:    $\mathbf{w}_1 \leftarrow \text{HighBits}(\mathbf{w})$ 
14:   ▷ HighBits is applied componentwise
15:    $\tilde{c} \leftarrow \text{H}(\mu \parallel \text{w1Encode}(\mathbf{w}_1), \lambda/4)$ 
16:    $\mathbf{c} \in R_q \leftarrow \text{SampleInBall}(\tilde{c})$ 
17:    $\widehat{\mathbf{c}} \leftarrow \text{NTT}(\mathbf{c})$ 
18:    $\langle\langle \mathbf{cs}_1 \rangle\rangle \leftarrow \text{NTT}^{-1}(\widehat{\mathbf{c}} \circ \widehat{\mathbf{s}}_1)$ 
19:    $\langle\langle \mathbf{cs}_2 \rangle\rangle \leftarrow \text{NTT}^{-1}(\widehat{\mathbf{c}} \circ \widehat{\mathbf{s}}_2)$ 
20:    $\mathbf{z} \leftarrow \mathbf{y} + \langle\langle \mathbf{cs}_1 \rangle\rangle$ 
21:    $\mathbf{r}_0 \leftarrow \text{LowBits}(\mathbf{w} - \langle\langle \mathbf{cs}_2 \rangle\rangle)$ 
22:   ▷ LowBits is applied componentwise
23:   if  $\|\mathbf{z}\|_\infty \geq \gamma_1 - \beta$  or  $\|\mathbf{r}_0\|_\infty \geq \gamma_2 - \beta$  then  $(\mathbf{z}, \mathbf{h}) \leftarrow \perp$ 
24:   else
25:      $\langle\langle \mathbf{ct}_0 \rangle\rangle \leftarrow \text{NTT}^{-1}(\widehat{\mathbf{c}} \circ \widehat{\mathbf{t}}_0)$ 
26:      $\mathbf{h} \leftarrow \text{MakeHint}(-\langle\langle \mathbf{ct}_0 \rangle\rangle, \mathbf{w} - \langle\langle \mathbf{cs}_2 \rangle\rangle + \langle\langle \mathbf{ct}_0 \rangle\rangle)$ 
27:     ▷ MakeHint is applied componentwise
28:     if  $\|\langle\langle \mathbf{ct}_0 \rangle\rangle\|_\infty \geq \gamma_2$  or the number of 1's in  $\mathbf{h}$  is greater than  $\omega$  then  $(\mathbf{z}, \mathbf{h}) \leftarrow \perp$ 
29:     end if
30:   end if
31:    $\kappa \leftarrow \kappa + \ell$ 
32: end while
33:  $\sigma \leftarrow \text{sigEncode}(\tilde{c}, \mathbf{z} \bmod^\pm q, \mathbf{h})$ 
34: return  $\sigma$ 
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**Annotated Algorithm 7** ML-DSA.Sign\_internal( $sk, M', rnd$ )

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1: ( $\rho : \mathbb{B}^{32}, K : \mathbb{B}^{32}, tr : \mathbb{B}^{64}, \mathbf{s}_1 : R_m^\ell, \mathbf{s}_2 : R_m^k, \mathbf{t}_0 : R_t^k$ ) = skDecode( $sk$ )
    $\triangleright m \in [-2, 2]$  if ML-DSA-44 or ML-DSA-87  $\triangleright \eta = 2$ 
    $\triangleright m \in [-4, 4]$  otherwise  $\triangleright \eta = 4$ 
    $\triangleright t \in [-2^{12} + 1, 2^{12} - 1]$   $\triangleright d = 13$ 
2:  $\widehat{\mathbf{s}}_1 : \mathbb{T}_q^\ell \leftarrow \text{NTT}(\mathbf{s}_1)$ 
3:  $\widehat{\mathbf{s}}_2 : \mathbb{T}_q^k \leftarrow \text{NTT}(\mathbf{s}_2)$ 
4:  $\widehat{\mathbf{t}}_0 : \mathbb{T}_q^k \leftarrow \text{NTT}(\mathbf{t}_0)$ 
5:  $\widehat{\mathbf{A}} : \mathbb{T}_q^{k \times \ell} \leftarrow \text{ExpandA}(\rho)$ 
6:  $\mu : \mathbb{B}^{64} \leftarrow \text{H}(\text{BytesToBits}(tr \parallel M', 64))$ 
7:  $\rho'' : \mathbb{B}^{64} \leftarrow \text{H}(K \parallel rnd \parallel \mu, 64)$ 
8:  $\kappa : \text{u16} \leftarrow 0$   $\triangleright$  unsigned 16-bit integer
9:  $(\mathbf{z} : R_{\text{mod}^\pm q}^\ell, \mathbf{h} : R_2^k) \leftarrow \perp$ 
10: while  $(\mathbf{z}, \mathbf{h}) = \perp$  do  $\triangleright$  Rejection sampling loop
11:    $\mathbf{y} : R_q^\ell \leftarrow \text{ExpandMask}(\rho'', \kappa)$   $\triangleright$  Sample a fresh mask  $\mathbf{y}$  mixing  $\rho''$  and new  $\kappa$ 
12:    $\mathbf{w} : R_q^k \leftarrow \text{NTT}^{-1}(\widehat{\mathbf{A}} \circ \text{NTT}(\mathbf{y}))$ 
13:    $\mathbf{w}_1 : R_{m_1}^k \leftarrow \text{HighBits}(\mathbf{w})$ 
    $\triangleright m_1 \in [0, 21]$  if ML-DSA-44
    $\triangleright m_1 \in [0, 15]$  if ML-DSA-65 or ML-DSA-87
14:    $\triangleright$  HighBits is applied componentwise
15:    $\tilde{c} : \mathbb{B}^{\lambda/4} \leftarrow \text{H}(\mu \parallel \text{w1Encode}(\mathbf{w}_1), \lambda/4)$   $\triangleright \lambda$  is the collision strength of  $\tilde{c}$ 
    $\triangleright \lambda/4 = 32$  (ML-DSA-44) or 48 (ML-DSA-65) or 64 (ML-DSA-87)
16:    $\mathbf{c} : R_q \leftarrow \text{SampleInBall}(\tilde{c})$   $\triangleright$  polynomial in  $R$  with coefficients from  $\{-1, 0, 1\}$ 
17:    $\widehat{\mathbf{c}} : T_q \leftarrow \text{NTT}(\mathbf{c})$ 
18:    $\langle\langle \mathbf{cs}_1 \rangle\rangle : R_q^\ell \leftarrow \text{NTT}^{-1}(\widehat{\mathbf{c}} \circ \widehat{\mathbf{s}}_1)$   $\triangleright \widehat{\mathbf{c}}$  is multiplied with each of  $\widehat{\mathbf{s}}_1^0, \dots, \widehat{\mathbf{s}}_1^{\ell-1}$ 
19:    $\langle\langle \mathbf{cs}_2 \rangle\rangle : R_q^k \leftarrow \text{NTT}^{-1}(\widehat{\mathbf{c}} \circ \widehat{\mathbf{s}}_2)$   $\triangleright \widehat{\mathbf{c}}$  is multiplied with each of  $\widehat{\mathbf{s}}_2^0, \dots, \widehat{\mathbf{s}}_2^{k-1}$ 
20:    $\mathbf{z} : R_{\text{mod}^\pm q}^\ell \leftarrow \mathbf{y} + \langle\langle \mathbf{cs}_1 \rangle\rangle$   $\triangleright \mathbf{z} : R_{\text{mod}^\pm q}^\ell \leftarrow \mathbf{y} : R_q^\ell + \langle\langle \mathbf{cs}_1 \rangle\rangle : R_q^\ell$ 
21:    $\mathbf{r}_0 : R_q^k \leftarrow \text{LowBits}(\mathbf{w} - \langle\langle \mathbf{cs}_2 \rangle\rangle)$ 
22:    $\triangleright$  LowBits is applied componentwise
23:   if  $\|\mathbf{z}\|_\infty \geq \gamma_1 - \beta$  or  $\|\mathbf{r}_0\|_\infty \geq \gamma_2 - \beta$  then  $(\mathbf{z}, \mathbf{h}) \leftarrow \perp$ 
24:   else
25:      $\langle\langle \mathbf{ct}_0 \rangle\rangle : T_q^k \leftarrow \text{NTT}^{-1}(\widehat{\mathbf{c}} \circ \widehat{\mathbf{t}}_0)$   $\triangleright \widehat{\mathbf{c}}$  is multiplied with each of  $\widehat{\mathbf{t}}_0^0, \dots, \widehat{\mathbf{t}}_0^{k-1}$ 
26:      $\mathbf{h} : R_2^k \leftarrow \text{MakeHint}(-\langle\langle \mathbf{ct}_0 \rangle\rangle, \mathbf{w} - \langle\langle \mathbf{cs}_2 \rangle\rangle + \langle\langle \mathbf{ct}_0 \rangle\rangle)$ 
27:      $\triangleright$  MakeHint is applied componentwise
28:     if  $\|\langle\langle \mathbf{ct}_0 \rangle\rangle\|_\infty \geq \gamma_2$  or the number of 1's in  $\mathbf{h}$  is greater than  $\omega$  then  $(\mathbf{z}, \mathbf{h}) \leftarrow \perp$ 
29:     end if
30:   end if
31:    $\kappa \leftarrow \kappa + \ell$ 
32: end while
33:  $\sigma \leftarrow \text{sigEncode}(\tilde{c}, \mathbf{z} \bmod^\pm q, \mathbf{h})$ 
34: return  $\sigma$ 

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