

Lecture 1: GCD and Euclidean Algorithm

Def: Given integers a, b the greatest common divisor or $\gcd(a, b)$ is the largest integer d such that $d \mid a$ and $d \mid b$

Ex.

$$(1) \gcd(180, 72) = 36 \quad (180 = 5 \cdot 36, 72 = 2 \cdot 36)$$

One way of computing \gcd (but very slow) is to look at the divisors of the numbers. Note that `divisors()` is a Sage command

`divisors(180)`

1, 2, 3, 4, 5, 6, 9, 10, 12, 15, 18, 20, 30, 36, 45, 60, 90, 180

`divisors(72)`

1, 2, 3, 4, 6, 8, 9, 12, 18, 24, 36, 72

The largest number appearing in both lists is 36

(2) Using the Sage command `gcd`

$$\gcd(31571776, 443085702) = 1346$$

$$\text{Note: } 31571776 = 1346 \cdot 23456, \quad 443085702 = 1346 \cdot 329187$$

Def. Given positive integers a, b . There exists integers q, r such that $a = qr + b$ where $0 \leq r < b$

Euclidean algorithm

Let a, b be positive integers with $b \leq a$. The following algorithm computes $\gcd(a, b)$. Compute the sequence r_0, r_1, r_2, \dots as follows:

Step 1: let $r_0 = a, r_1 = b$

Step 2: Let $i = 1$

Step 3: Compute r_{i+1} by dividing r_{i-1} by r_i as follows

$$r_{i-1} = q \cdot r_i + r_{i+1} \quad \text{where } 0 \leq r_{i+1} < r_i$$

Step 4: If $r_{i+1} = 0$

then $\gcd(a, b) = r_i$

else let $i = i + 1$ and repeat Step 3

The running time of the algorithm is $O(\log b)$

Ex:

(1) $\gcd(101,97)$

$$101 = 1*97 + 4 \quad \text{Now divide 97 by 4}$$

$$97 = 24*4 + 1 \quad \text{Now divide 4 by 1}$$

$$4 = 4*1 + 0 \quad \text{Done, } \mathbf{\gcd(101,97) = 1} \text{ (the remainder at the previous step)}$$

(2) $\gcd(42823,6409)$

$$42823 = 6*6409 + 4369 \quad \text{Now divide 6409 by 4369}$$

$$6409 = 1*4369 + 2040 \quad \text{Now divide 4369 by 2040}$$

$$4369 = 2*2040 + 289 \quad \text{Now divide 2040 by 289}$$

$$2040 = 7*289 + 17 \quad \text{Now divide 289 by 17}$$

$$289 = 17*17 + 0 \quad \text{Done } \mathbf{\gcd(42823,6409) = 17} \text{ (the remainder at the previous stage)}$$

Extended Euclidean Algorithm

Given positive integers a, b . There exists integers u, v such that

$$ua + vb = \gcd(a, b)$$

Start with the next to last step of the Euclidean algorithm. Solve for the $\gcd(a, b)$. Then keep substituting from the previous equations until reaching a, b .

Ex:

(1) 101,97

$$101 = 1*97 + 4$$

$$97 = 24*4 + 1$$

$$4 = 4*1 + 0$$

$$1 = 97 - 24*4 = 97 - 24*(101 - 97) = 25*97 - 24*101 \quad (\text{so } u = 25, v = -24)$$

(2) 42823,6409

$$42823 = 6*6409 + 4369$$

$$6409 = 1*4369 + 2040$$

$$4369 = 2*2040 + 289$$

$$2040 = 7*289 + 17$$

$$289 = 17*17 + 0$$

$$\begin{aligned} 17 &= 2040 - 7*289 = 2040 - 7*(4369 - 2*2040) = 15*2040 - 7*4369 = 15*(6409 - 4369) - 7*4369 \\ &= 15*6409 - 22*4369 = 15*6409 - 22*(42823 - 6*6409) = 147*6409 - 22*42823 \\ &(\text{so } u = 147, v = -22) \end{aligned}$$