

## L20: P, NP, and Polynomial-Time Reduction

Problem 7.9 from the Book

7.9) A triangle in an undirected graph is a 3-clique. Show that TRIANGLE  $\in$  P where  $\text{TRIANGLE} = \{ \langle G \rangle \mid G \text{ contains a triangle} \}$

Class P: P is a class of languages that are decidable in polynomial time on a deterministic single tape Turing-machine.

Specified language

A triangle in an undirected graph is a 3-clique.

$$\text{TRIANGLE} = \{ \langle G \rangle \mid G \text{ contains a triangle} \}$$

Now we have to show that  $\text{TRIANGLE} \in P$

Let A be the Turing Machine that decides TRIANGLE in Polynomial time

A can be described as follows:

A = "on input  $G = \langle V, E \rangle$  :

V denotes set of vertices of the graph G.  
E denotes set of edges of the graph G.

For  $u, v, w \in V$  and  $u < v < w$ , we enumerate all triples  $(u, v, w)$ .

Check whether all three edges  $(u, v)$ ,  $(v, w)$ ,  $(w, u)$  exist in E or not. If exist then accept.

Otherwise reject

Enumeration of all triple require  $O(|V|^3)$  time

Checking whether all three edges belong to E take  $O(|E|)$  time

Overall time is  $O(|V|^3 |E|)$  which is polynomial in the length of the input

Therefore  $\text{TRIANGLE} \in P$ .