

Section 1

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Let us define equi-satisfiable relation over the set of wffs as follows. A and B are equi-satisfiable iff (A is satisfiable iff B is satisfiable). Which of the following formulas are true?

- ☒ Equi-satisfiable is a reflexive relation
- ☒ Equi-satisfiable is a symmetric relation
- ☒ Equi-satisfiable is a transitive relation
- ☒ Equi-satisfiable is an equivalence relation
- ☐ None of these



Which of the following statements are true?

- ☒ Every atomic proposition is a literal
- ☐ Every formula is a literal
- ☐ Every literal is an atomic proposition
- ☐ None of these
- ☒ Every literal is a formula

Which of the following statements are true? " \cup " denotes set union and " \setminus " denotes set minus.

- ☒ If $U \models A$ then $U \cup \{B\} \models A$ for any B
- ☒ If $U \models A$ then $U \cup \{B\} \models A$ for any valid formula B
- ☐ If $U \models A$ then $U \setminus \{B\} \models A$ for any B
- ☐ If $U \models A$ then $U \setminus \{B\} \models A$ for any valid B



Let us define a set of good formulas inductively.

- Every atomic proposition p is a good formula.
- If A is a good formula, then $(\sim A)$ is a good formula.
- If A and B are good formulas then $(A/\wedge B)$ is a good formula.

Consider a semantic tableaux T of a good formula. Recall $U(n)$ denotes the set of formulas associated with the node n in T .

For which of the following measures, $W(n) < W(n_1)$ where n is child of n_1 ?

- ☒ $W(n)$ = number of operators in $U(n)$.
- ☐ $W(n)$ = number of conjunct operators ($/\wedge$) in $U(n)$
- ☒ $W(n)$ = number of unary operators in $U(n)$ + 2 * number of conjunct operators in $U(n)$
- ☒ $W(n)$ = number of conjunct operators in $U(n)$ + 2 * number of unary operators in $U(n)$
- ☐ None of these

Which of the following statements are true about box clause?

- ☐ Box clause is valid
- ☒ Box clause is unsatisfiable
- ☒ Box clause is a clause without any literals
- ☒ Any clause set containing box clause is unsatisfiable
- ☐ None of these



Consider the following logic program $P = \{\text{likes}(\text{john}, \text{flowers}), \text{likes}(\text{mary}, \text{food}), \text{likes}(\text{mary}, \text{oj}), \text{likes}(\text{john}, \text{oj}), \text{likes}(\text{john}, \text{mary}), \text{likes}(\text{paul}, \text{mary})\}$. What Prolog will return for the query $\text{likes}(\text{mary}, X)$?

- ☐ None of these
- ☐ X
- ☒ oj
- ☐ Paul
- ☐ John

Let us define a set of good formulas inductively. Two fixed atomic proposition p and q are good formulas. If A and B are good formulas then $(A \leftrightarrow B)$ is a good formula and $(A \wedge B)$ is a good formula. Which of the following statements are true?

- ☐ Every good formula is valid or equivalent to $(p \wedge q)$
- ☐ There are no good formulas which are valid.
- ☐ Every good formula is equivalent to p or q or $(p \rightarrow q)$
- ☒ None of these
- ☐ Every good formula is equivalent to either p or q

Let $C1 = \{p, \sim q, \sim r\}$, $C2 = \{p, q, \sim r\}$ be two clauses. Which of the following statements are true?

- ☒ $\{p, \sim r\}$ is a resolvent of $C1$ and $C2$
- ☐ $\{p, r, \sim r\}$ is a resolvent of $C1$ and $C2$
- ☐ $\{p\}$ is a resolvent of $C1$ and $C2$
- ☐ $\{p, q, \sim q\}$ is a resolvent of $C1$ and $C2$
- ☐ None of these



Which of the following statements are true? Recall, we will use set notation for CNF formulas only.

- ☐ $\{\{p,q,r\}\}$ is a horn formula
- ☒ $\{\{p,\sim q,\sim r\},\{\sim q,\sim r\}\}$ is a horn formula
- ☒ $\{\{p\},\{q\},\{r\}\}$ is a horn formula
- ☒ $\{\{\sim q,\sim r\},\{\sim p,\sim r\}\}$ is a horn formula
- ☐ None of these

Let us define equi-satisfiable relation over the set of wffs as follows. A and B are equi-satisfiable iff (A is satisfiable iff B is satisfiable). Which of the following statements are true?

- ☐ If A and B are equi-satisfiable then A is satisfiable
- ☒ If A and B are valid then A and B are equi-satisfiable
- ☒ If A and B are satisfiable then A and B are equi-satisfiable
- ☒ Atomic propositions p and q are equi-satisfiable.
- ☐ None of these

Consider the formula $(p \vee \sim q) \wedge (p \wedge \sim q)$.

- ☐ None of these
- ☐ All the semantic tableaux of the above formula are closed
- ☐ There exists an unique semantic tableaux for the above formula
- ☐ Any semantic tableaux of the above formula will have at least 6 nodes
- ☒ All the semantic tableaux of the above formula are open



Which of the following is true about bi-implication operator in a semantic tableau?

- ☐ None of these
- ☒ $(A \leftrightarrow B)$ is an alpha-formula with $\alpha_1 = (A \rightarrow B)$, $\alpha_2 = (B \rightarrow A)$
- ☐ $(A \leftrightarrow B)$ is an alpha-formula with $\alpha_1 = A$, $\alpha_2 = B$
- ☒ $\sim(A \leftrightarrow B)$ is a beta-formula with $\beta_1 = (\sim A) \wedge B$, $\beta_2 = A \wedge (\sim B)$
- ☐ $(A \leftrightarrow B)$ is a beta formula with $\beta_1 = A$, $\beta_2 = B$

Let S be a set of clauses. Each clause in S has at most three literals. Number of clauses in S is k and the number of atomic propositions used in S is at most n . Which of the following statements are true?

- ☐ None of these
- ☒ Number of clauses in $\text{Res}^*(S)$ is at most $n^3 + n$
- ☒ Number of clauses in $\text{Res}^*(S)$ is at least k
- ☐ Each clause in $\text{Res}^*(S)$ has at most four literals
- ☐ Number of clauses in $\text{Res}^*(S)$ is at most $k^3 + k$

Which of the following restrictions of propositional logic formulas lead to polynomial running time algorithm for the satisfiability problem?

- ☐ Restricting to DNF formulas
- ☐ Restricting to CNF formulas
- ☒ Restricting to Horn formulas
- ☒ Restricting to CNF formulas in which if a literal occurs in the formula then its complementary pair will not occur in the formula



Which of the following statements are true about Semantic tableau method?

- ☒ It is a syntactic method
- ☐ It is an efficient method
- ☐ To apply this method, formula should be in DNF
- ☐ None of these
- ☐ To apply this method, formula should be in CNF

Which of the following statements are true about Resolution?

- ☐ Resolution is an efficient method
- ☒ To apply resolution method, formula should be in CNF
- ☐ To apply resolution method, formula should be in DNF
- ☐ None of these
- ☒ Resolution is a syntactic method

Let us define equi-satisfiable relation over the set of wffs as follows. A and B are equi-satisfiable iff (A is satisfiable iff B is satisfiable). Which of the following statements are true?

- ☐ If A and B are equi-satisfiable, then A and B are valid
- ☐ If A and B are equi-satisfiable, then A and B are satisfiable
- ☒ Atomic propositions p and $\sim p$ are equi-satisfiable.
- ☐ $(p \rightarrow \sim p)$ and $(p \wedge \sim p)$ are equi-satisfiable
- ☐ None of these



Which of the following statements are true? Formulas are represented using clause set notation. Here \Box denotes the box clause.

- ☐ \Box is equivalent to the empty clause set $\{\}$
- ☒ $\{p\}, \Box$ is equi-satisfiable to \Box
- ☐ $\{p\}, \Box$ is equi-satisfiable to $\{p\}$
- ☐ None of these
- ☐ $\{p\}, \Box$ is equivalent to the empty clause set $\{\}$

Which of the following statements are true about Res operator? Recall that for a given set of clauses S , $\text{Res}(S) = S \cup \{R : R \text{ is a resolvent of two clauses in } S\}$ and \Box denotes the box clause.

- ☐ For all set of clauses S , $\text{Res}^*(S) = \text{Res}(S)$
- ☐ For all set of clauses S , $\text{Res}(S) = S$
- ☒ For all set of clauses S , $\text{Res}(S)$ is equivalent to S
- ☒ If S is unsatisfiable, then $\text{Res}(S) = \{\Box\}$
- ☐ None of these



Which of the statements are true about the following algorithm for validity problem.

Algorithm simple-validity(A)

If A is an atomic. proposition then return Yes

else return No

- ☐ Algorithm simple-validity is sound for validity problem
- ☐ None of these
- ☐ Algorithm simple-validity is complete for validity problem
- ☒ Algorithm simple-validity is neither complete nor sound for validity problem
- ☒ Algorithm simple-validity terminates on all inputs

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