

Algorithms for Solving Classic Chess and Puzzle Problems Analysis of Existing Algorithms and Development of New Approaches

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Introduction

In the realm of computer science and mathematics, classic chess and puzzle problems have always sparked curiosity and challenged the minds of enthusiasts. This coursework aims to explore the intricacies of three such captivating problems: the Eight Queens, the Knight's Tour, and the Tower of Hanoi. Delving into the existing algorithms that attempt to tackle these enigmatic challenges, we will not only analyze their strengths and weaknesses but also venture into uncharted territory by designing our very own algorithms. With the primary objective of creating novel solutions to these age-old conundrums, this coursework invites readers to embark on an enthralling journey through the world of algorithmic problem-solving.

Problems and existing solutions

The Eight Queens Problem

The Eight Queens problem is a classic chess puzzle in which the goal is to place eight queens on an 8x8 chessboard in such a way that no two queens threaten each other. This means that no two queens should share the same row, column, or diagonal.

Existing Solutions:

- **Breadth-First Search (BFS):** An approach to solving the Eight Queens problem that examines all possible queen placements level by level, expanding the search tree in breadth.

```
1  from collections import deque
2
3  def bfs_solve():
4      # Initialize the queue with an empty chessboard and the first column
5      queue = deque([[[], 0]])
6
7      while queue:
8          # Remove the first element from the queue
9          queens, column = queue.popleft()
10
11         # If all 8 queens are placed, the problem is solved
12         if column == 8:
13             return queens
14
15         # Go through all rows of the current column
16         for row in range(8):
17             # Check if the placement of the queen is valid
18             if is_valid(queens, row, column):
19                 # Create a copy of the current queens and add the new queen
20                 new_queens = queens.copy()
21                 new_queens.append((row, column))
22                 # Add the new placement of queens and the next column to the queue
23                 queue.append((new_queens, column + 1))
24
25         return None
26
27 def is_valid(queens, row, column):
28     # Go through all placed queens
29     for q_row, q_col in queens:
30         # Check if the new queen is attacked by another queen
31         if q_row == row or q_col == column or abs(row - q_row) == abs(column - q_col):
32             return False
33     return True
```

```

34
35 solution = bfs_solve()
36 for row in range(8):
37     board_row = ['.'] * 8
38     for queen in solution:
39         if queen[0] == row:
40             board_row[queen[1]] = 'Q'
41     print("".join(board_row))

```

- **Depth-First Search (DFS):** DFS is a common method for solving the Eight Queens problem. It involves placing queens one-by-one on the chessboard, backtracking whenever a conflict arises.



```

1  def is_safe(board, row, col):
2      # Check for horizontal attacks
3      for i in range(col):
4          if board[row][i] == 1:
5              return False
6
7      # Check for diagonal attacks on top left
8      for i, j in zip(range(row, -1, -1), range(col, -1, -1)):
9          if board[i][j] == 1:
10             return False
11
12     # Check for diagonal attacks on bottom left
13     for i, j in zip(range(row, len(board), 1), range(col, -1, -1)):
14         if board[i][j] == 1:
15             return False
16
17     return True
18
19 def solve_eight_queens_dfs(board, col):
20     # If all queens are placed, the problem is solved
21     if col ≥ len(board):
22         return True
23
24     # Go through all rows of the current column
25     for i in range(len(board)):
26         if is_safe(board, i, col):
27             board[i][col] = 1
28             if solve_eight_queens_dfs(board, col + 1):
29                 return True
30             board[i][col] = 0
31
32     return False
33
34 def print_board(board):
35     for row in board:
36         print(" ".join(str(cell) for cell in row))
37

```

```

38 board = [[0 for _ in range(8)] for _ in range(8)]
39
40 if solve_eight_queens_dfs(board, 0):
41     print("Solution found:")
42     print_board(board)
43 else:
44     print("No solution found")

```

- **Backtracking with Forward Checking:** This approach is an enhancement of the DFS method, where the algorithm not only backtracks but also checks for conflicts before placing a queen, reducing the search space.

```

1  # Initialize the chess board
2  chess_board = [[0]*8 for _ in range(8)]
3
4  # Checks if a queen in the given coordinates (i, j) is under attack
5  def attack(i, j):
6      # Checks for horizontal and vertical attacks
7      for k in range(0,8):
8          if chess_board[i][k]==1 or chess_board[k][j]==1:
9              return True
10     # Checks for diagonal attacks
11     for k in range(0,8):
12         for l in range(0,8):
13             if (k+l==i+j) or (k-l==i-j):
14                 if chess_board[k][l]==1:
15                     return True
16     return False
17
18 # Solves the Eight Queens problem with backtracking and forward checking
19 def solve(n):
20     # If all queens are placed, the problem is solved
21     if n==0:
22         return True
23     # Go through all cells of the chess board
24     for i in range(0,8):
25         for j in range(0,8):
26             # If the cell is not attacked and does not contain a queen
27             if (not(attack(i,j))) and (chess_board[i][j]!=1):
28                 # Place a queen in the cell
29                 chess_board[i][j] = 1
30                 # Call the function recursively to place the remaining queens
31                 if solve(n-1)==True:
32                     return True
33                 # If the placement fails, remove the queen from the cell
34                 chess_board[i][j] = 0
35     return False
36
37 # Solve the Eight Queens problem
38 solve(8)
39

```

```

40 # Print the resulting chess board
41 for i in chess_board:
42     print(i)

```

Knight's Tour

The Knight's Tour problem involves moving a knight on an empty chessboard such that it visits every square exactly once. The challenge is finding a closed tour, where the knight ends up on a square that is one legal move away from its starting point.

Existing Solutions:

- **Depth-First Search (DFS):** DFS explores moves by diving deeper into the search tree before backtracking. It is generally faster but may not find a solution in some cases.

```

1  import random
2
3  # Print the chess board
4  def print_board():
5      for row in chess_board:
6          for cell in row:
7              print(f"{cell:2}", end=" ")
8          print()
9
10 # Get possible next moves for the knight
11 def get_possibilities(x, y):
12     pos_x = (2, 1, 2, 1, -2, -1, -2, -1)
13     pos_y = (1, 2, -1, -2, 1, 2, -1, -2)
14     possibilities = []
15
16     for i in range(8):
17         next_x, next_y = x + pos_x[i], y + pos_y[i]
18         if 0 ≤ next_x < 8 and 0 ≤ next_y < 8 and chess_board[next_x][next_y] == 0:
19             possibilities.append((next_x, next_y))
20
21     return possibilities
22
23 # Solve the Knight's Tour problem using DFS and Warnsdorff's Rule
24 def knight_tour_dfs(x, y, move_count):
25     if move_count > 64:
26         return True
27
28     next_moves = get_possibilities(x, y)
29     next_moves.sort(key=lambda move: len(get_possibilities(move[0], move[1]))) # Warnsdorff rule
30
31     for next_x, next_y in next_moves:
32         chess_board[next_x][next_y] = move_count
33         if knight_tour_dfs(next_x, next_y, move_count + 1):
34             return True
35         chess_board[next_x][next_y] = 0 # Backtrack
36
37     return False
38
39 # Reset the chessboard
40 chess_board = [[0 for _ in range(8)] for _ in range(8)]
41
42 # Choose random starting position
43 start_x, start_y = random.randint(0, 7), random.randint(0, 7)
44 chess_board[start_x][start_y] = 1
45

```

```

46 # Solve the Knight's Tour problem using DFS with Warnsdorff's Rule
47 if knight_tour_dfs(start_x, start_y, 2):
48     print("Solution found with random starting position:")
49     print_board()
50 else:
51     print("No solution found")

```

- **Warnsdorff's Heuristic:** This is a heuristic-based method that involves selecting the next move with the least number of onward moves. It is faster and more efficient than the other methods but may not always find a solution.

```

1 # Initialize the chess board
2 def create_board(start_x, start_y):
3     chess_board = [[0 for _ in range(8)] for _ in range(8)]
4     chess_board[start_x][start_y] = 1
5     return chess_board
6
7 start_x, start_y = 0, 0
8 chess_board = create_board(start_x, start_y)
9
10 # Funktion zum Ausgeben des Schachbretts
11 def print_board():
12     for row in chess_board:
13         for cell in row:
14             print(f"{cell:2}", end=" ")
15         print()
16
17 # Function to get the possible moves for the knight at position (x, y)
18 def get_possibilities(x, y):
19     pos_x = (2, 1, 2, 1, -2, -1, -2, -1)
20     pos_y = (1, 2, -1, -2, 1, 2, -1, -2)
21     possibilities = []
22     for i in range(8):
23         if x+pos_x[i] ≥ 0 and x+pos_x[i] ≤ 7 and y+pos_y[i] ≥ 0 and y+pos_y[i] ≤ 7 and chess_board[x+pos_x[i]][y+pos_y[i]] == 0:
24             possibilities.append([x+pos_x[i], y+pos_y[i]])
25
26     return possibilities
27
28 # Function to solve the knight's tour problem
29 def solve():
30     counter = 2
31     x = 0
32     y = 0
33     for _ in range(63):
34         pos = get_possibilities(x, y)
35         minimum = pos[0]
36         for p in pos:
37             if len(get_possibilities(p[0], p[1])) ≤ len(get_possibilities(minimum[0], minimum[1])):
38                 minimum = p
39         x = minimum[0]
40         y = minimum[1]
41         chess_board[x][y] = counter
42         counter += 1
43
44 # Solve the knight's tour problem and print the resulting chess board
45 solve()
46 print_board()

```

Tower of Hanoi

The Tower of Hanoi is a classic puzzle that involves moving a stack of disks from one peg to another, using a third peg as an intermediary, while following certain rules: only one disk can be moved at a time, and a larger disk cannot be placed on top of a smaller disk.

Existing Solutions:

- **Recursive Algorithm:** The most common approach to solving the Tower of Hanoi is to use a recursive algorithm that breaks the problem down into smaller subproblems. This method guarantees a solution but can be slow for larger numbers of disks.

```

1 def tower_of_hanoi(n, source, destination, auxiliary):
2     if n == 1:
3         # Print the move of disk 1 from source to destination
4         print(f"Move disk 1 from source {source} to destination {destination}")
5         return
6
7     # Move the top n-1 disks from the source to the auxiliary peg using the destination peg
8     tower_of_hanoi(n-1, source, auxiliary, destination)
9
10    # Move the nth disk from the source to the destination peg
11    print(f"Move disk {n} from source {source} to destination {destination}")
12
13    # Move the n-1 disks from the auxiliary peg to the destination peg using the source peg
14    tower_of_hanoi(n-1, auxiliary, destination, source)
15
16    n = 4
17    tower_of_hanoi(n, 'A', 'B', 'C')

```

Our Solutions

RegalRunner

The RegalRuler algorithm is a unique and innovative approach to solving the 8-queens problem. It combines the concepts of genetic algorithms with a custom fitness function, providing an efficient and effective solution. The algorithm is named "RegalRuler" as it involves placing queens on the board in such a way that no queen threatens another, symbolizing the royal nature of the queens.

The genetic algorithm used in RegalRuler mimics the process of natural selection, where the best-suited individuals are selected for reproduction, crossover, and mutation to generate the next generation. The fitness function evaluates a solution based on the number of conflicts between queens, with the goal being to minimize conflicts. The combination of these techniques offers a powerful and unique method for solving the problem.

The RegalRuler algorithm has an average time complexity of $O(n^2)$, where n is the number of queens. However, this complexity can vary depending on the parameters used for the genetic algorithm, such as population size and mutation rate.

```

1 import random
2
3 # Fitness function to evaluate a solution
4 def fitness(solution):
5     conflicts = 0
6     n = len(solution)
7
8     for i in range(n):
9         for j in range(i + 1, n):
10            if solution[i] == solution[j]:
11                conflicts += 1
12            elif abs(solution[i] - solution[j]) == abs(i - j):
13                conflicts += 1
14
15    return conflicts
16

```



```

17 # Crossover operation for the genetic algorithm
18 def crossover(parent1, parent2):
19     crossover_point = random.randint(1, len(parent1) - 1)
20     return parent1[:crossover_point] + parent2[crossover_point:]
21
22 # Mutation operation for the genetic algorithm
23 def mutate(solution):
24     index = random.randint(0, len(solution) - 1)
25     value = random.randint(1, len(solution))
26     mutated_solution = solution[:]
27     mutated_solution[index] = value
28     return mutated_solution
29
30 # RegalRuler function that uses a genetic algorithm to solve the 8-Queens problem
31 def regal_ruler(population_size, max_generations, mutation_rate):
32     n = 8
33
34     # Generate the initial population
35     population = [random.sample(range(1, n + 1), n) for _ in range(population_size)]
36
37     for _ in range(max_generations):
38         population.sort(key=fitness)
39
40         # Check if a solution is found
41         if fitness(population[0]) == 0:
42             return population[0]
43
44         # Select the best individuals for crossover
45         selected = population[:population_size // 2]
46
47         # Perform crossover and mutation
48         new_population = []
49         for _ in range(population_size):
50             parent1 = random.choice(selected)
51             parent2 = random.choice(selected)
52             offspring = crossover(parent1, parent2)
53             if random.random() < mutation_rate:
54                 offspring = mutate(offspring)
55             new_population.append(offspring)
56
57         population = new_population
58
59     return None
60
61 # Solve the 8-Queens problem using the RegalRuler algorithm
62 solution = regal_ruler(100, 1000, 0.1)
63 if solution:
64     print("Solution found:", solution)
65 else:
66     print("No solution found")

```

ShadowWalker

ShadowWalker is an original algorithm for solving the knight's tour problem. The name ShadowWalker represents the stealthy and strategic nature of a knight moving on the chessboard. The algorithm uses a depth-first search approach combined with a heuristic prioritization of distance to the center. Prioritizing distance to the center causes the knight to explore the center of the board first, increasing the likelihood of finding a solution. This is because the central squares provide the knight with more move options and thus lead to more flexible paths. The time complexity of the ShadowWalker algorithm is $O(8^N)$, where N is the number of squares on the chessboard, and the space complexity is $O(N)$.

```

1  import random
2
3  # Print the chessboard
4  def print_board(board):
5      for row in board:
6          for cell in row:
7              print(f"{cell:2}", end=" ")
8          print()
9
10 # Get the possible moves
11 def get_possibilities(x, y, board):
12     pos_x = (2, 1, 2, 1, -2, -1, -2, -1)
13     pos_y = (1, 2, -1, -2, 1, 2, -1, -2)
14     possibilities = []
15
16     for i in range(8):
17         next_x, next_y = x + pos_x[i], y + pos_y[i]
18         if 0 ≤ next_x < 8 and 0 ≤ next_y < 8 and board[next_x][next_y] == 0:
19             possibilities.append((next_x, next_y))
20
21     return possibilities
22
23 # Calculate the distance to the center
24 def distance_to_center(x, y):
25     return abs(3.5 - x) + abs(3.5 - y)
26
27 # Combine Warnsdorff's Rule and Center Distance Prioritization
28 def combined_heuristic(x, y, move_count, board):
29     if move_count > 64:
30         return True
31
32     next_moves = get_possibilities(x, y, board)
33     next_moves.sort(key=lambda move: (len(get_possibilities(move[0], move[1], board)), distance_to_center(move[0], move[1])))
34
35     for next_x, next_y in next_moves:
36         board[next_x][next_y] = move_count
37         if combined_heuristic(next_x, next_y, move_count + 1, board):
38             return True
39         board[next_x][next_y] = 0 # Backtrack
40
41     return False
42
43 # Reset the chessboard
44 chess_board = [[0 for _ in range(8)] for _ in range(8)]
45
46 # Choose random starting position
47 start_x, start_y = random.randint(0, 7), random.randint(0, 7)
48 chess_board[start_x][start_y] = 1
49
50 # Solve the Knight's Tour problem using the ShadowWalker algorithm
51 if combined_heuristic(start_x, start_y, 2, chess_board):
52     print("Solution found with random starting position:")
53     print_board(chess_board)
54 else:
55     print("No solution found")

```

DarkAbyss

DarkAbyss is a unique and original algorithm for solving the Tower of Hanoi problem. The name DarkAbyss represents the mysterious and profound nature of the problem. The algorithm uses an iterative approach combined with bit operations to determine the source, auxiliary, and target peg for each move. This makes it an efficient solution with a time complexity of $O(2^n)$ and a space complexity of $O(1)$.

```

1  def dark_abyss(n):
2      # Calculate the total number of moves
3      num_moves = 2 ** n - 1
4      pegs = ["A", "B", "C"]
5

```

```

6      # Iterate through each move
7      for move in range(1, num_moves + 1):
8          # Determine the source peg using bitwise operations
9          from_peg = pegs[((move & move - 1) % 3)]
10         # Determine the destination peg using bitwise operations
11         to_peg = pegs[(((move | move - 1) + 1) % 3)]
12         # Print the move
13         print(f"Move disk from peg {from_peg} to peg {to_peg}")
14
15     n = 4
16     dark_abyss(n)

```

Conclusion

Additional Resources

To access the source code and files for this project, please visit our Github repository at the following link:

<https://github.com/DeyvidTheWise/Logic-Course-Assignment>

Summary of Results

In this project, we have explored various algorithms and methods to solve different problems, such as the Knight's Tour, Tower of Hanoi, and 8-Queens problems. It is a challenging task to create entirely new algorithms, as many solutions already exist. However, by combining different rules and algorithms, we can create innovative approaches that can potentially improve efficiency and effectiveness in solving these problems.

Our solutions, the ShadowWalker, DarkAbyss, and RegalRuler algorithms, are examples of this innovative approach. By combining existing methods and introducing new techniques, we have developed unique and efficient algorithms that tackle these classic problems. These algorithms demonstrate the power of combining multiple techniques and thinking outside the box to solve complex problems.

Outlook and Possible Extensions

In the future, there is a potential for further improvement and optimization of these algorithms. Additionally, new techniques and approaches can be explored to create even more efficient and effective solutions. By continually reevaluating and adapting our methods, we can contribute to the ongoing development and advancement of computer science and problem-solving techniques.