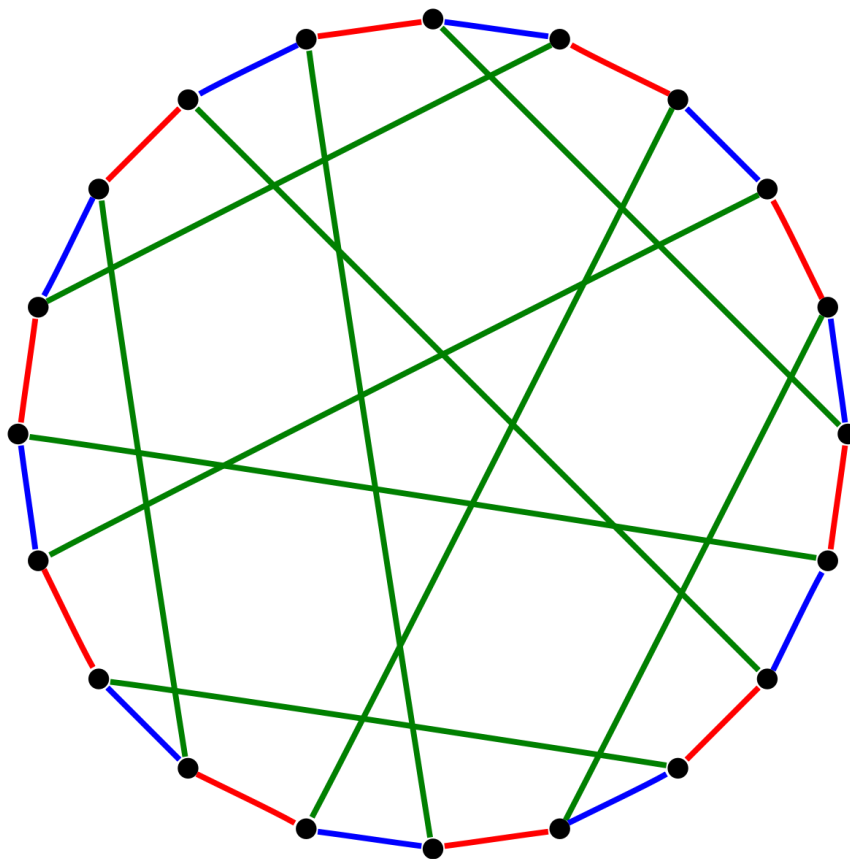


# CS 2150 GRAPH THEORY

## RESEARCH CYCLE-1

### PROJECT REPORT

#### GROUP C6



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**Problem 1- Group C6**

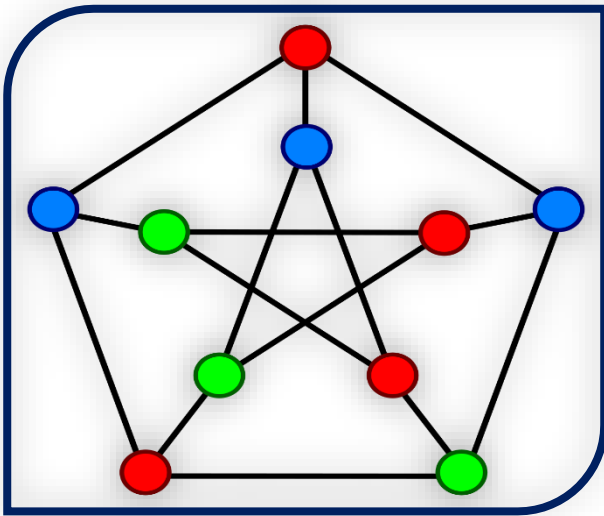
**Problem 2- Group C6**

**References**

**Questions and Solutions**

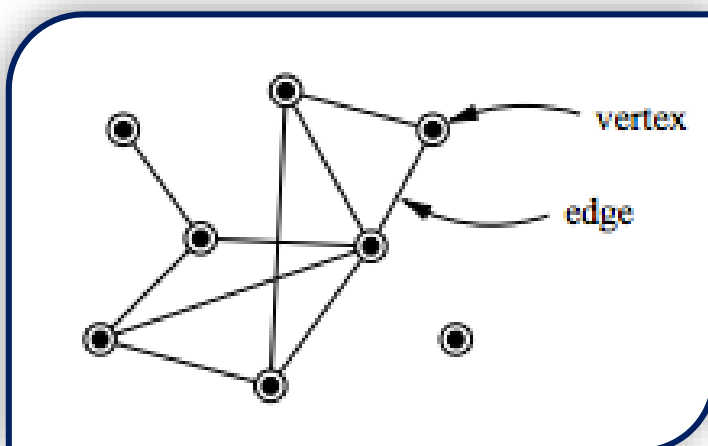
# Definitions

## 1. Graph



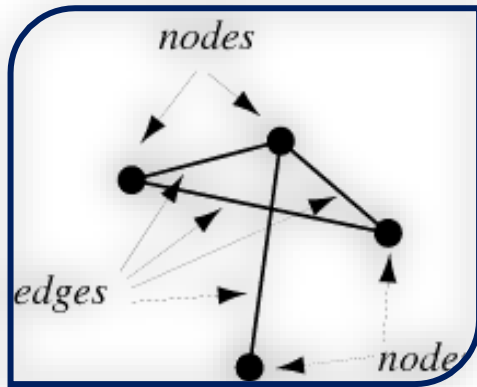
A *graph*  $G = (V, E)$  is an ordered pair of finite sets. Elements of  $V$  are called vertices or nodes, and elements of  $E$  are called edges or arcs. We refer to  $V$  as the vertex set of  $G$ , with  $E$  being the edge set.

## 2. Vertex



A *vertex* is a point where multiple lines meet. It is also called a node. Similar to points, a vertex is also denoted by an alphabet.

### 3. Edge



An *edge* is the mathematical term for a line that connects two vertices. Many edges can be formed from a single vertex. Without a vertex, an edge cannot be formed. There must be a starting vertex and an ending vertex for an edge.

### 4. Degree of a vertex

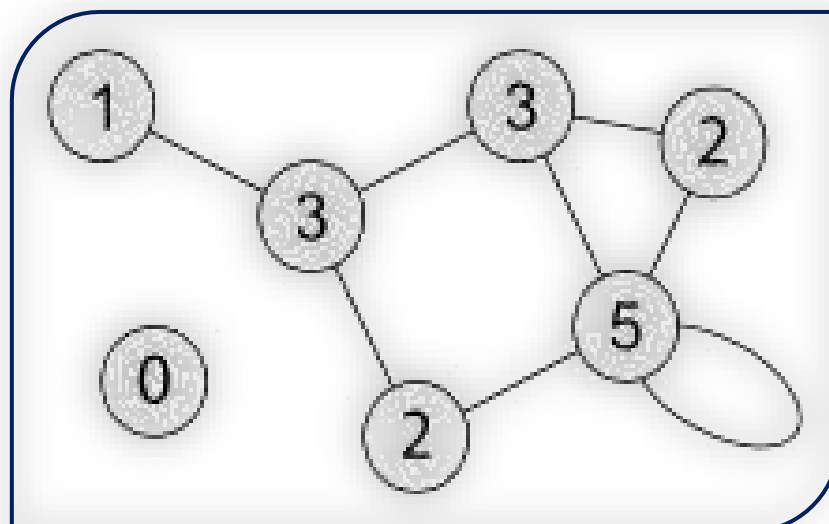
The *degree of vertex* in a graph is number of edges that enter or exit from the vertex. A loop contributes 2 to the degree of its vertex.

Directed graphs have two types of degrees, known as the indegree and the outdegree.

*Indegree* - The number of inward directed graph edges from a given graph vertex

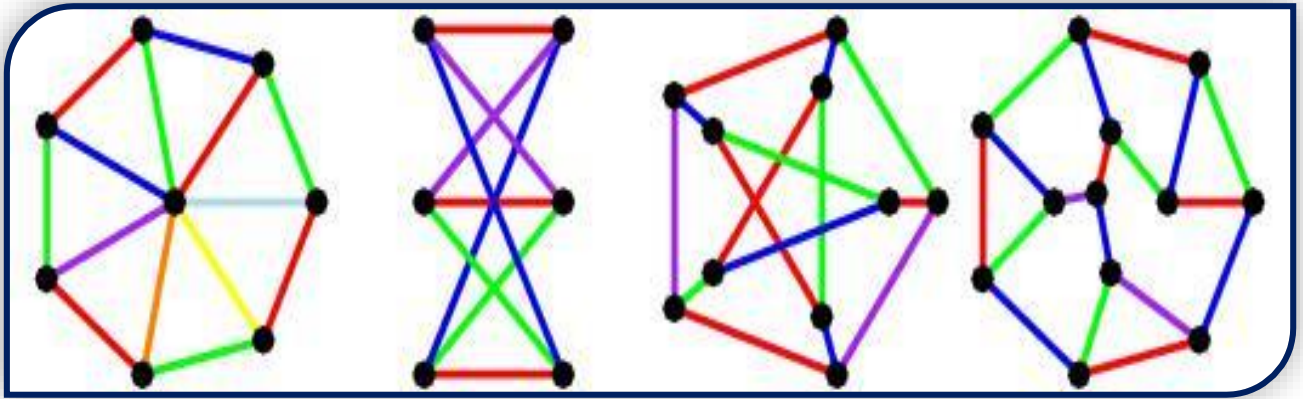
*Outdegree* - The number of outward directed graph edges from a given graph.

In the following figure, degree of vertex is shown inside vertex.



## 5. Proper Edge Colouring

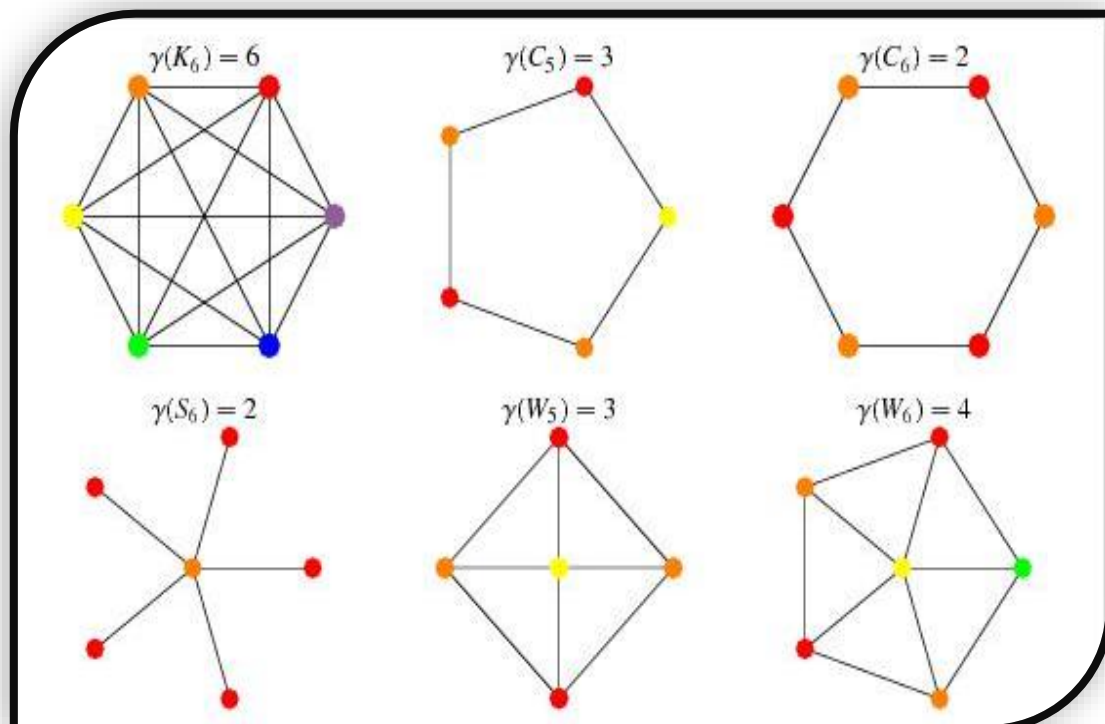
An *edge colouring* of a graph is an assignment of colours to edges of a graph so that no vertex is incident to two edges of the same colour. An edge colouring with  $k$  colours is called a  **$k$ -edge-colouring**.



## 6. Chromatic Number

The *chromatic number* of a graph  $G$  is the smallest number of colours needed to colour the vertices of  $G$  so that no two adjacent vertices share the same colour.

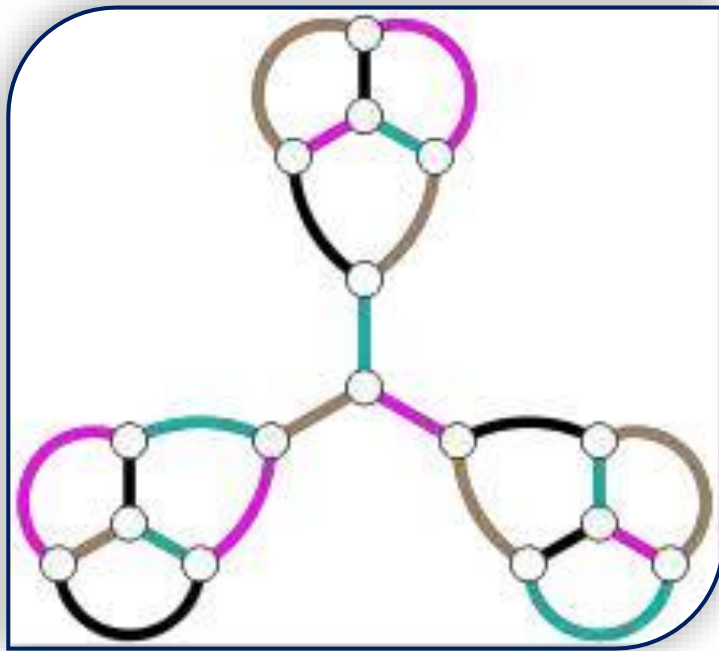
The chromatic of a graph  $G$  is most commonly denoted  $\chi(G)$ .



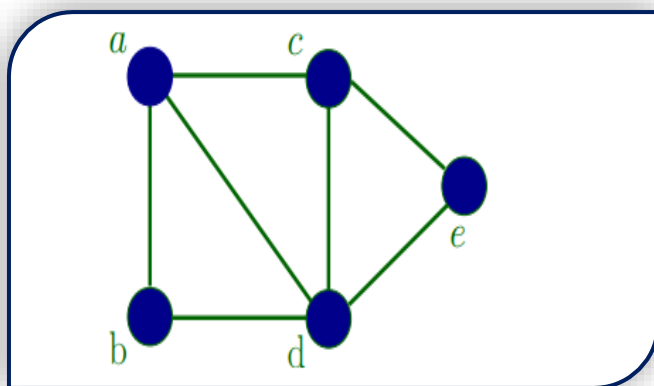
### 7. Edge Chromatic Number

The *edge chromatic number*, sometimes also called the chromatic index, of a graph  $G$  is the fewest number of colours necessary to colour each edge of  $G$  such that no two edges incident on the same vertex have the same colour. In other words, it is the number of distinct colours in a minimum edge colouring.

The edge chromatic number of following figure is 4.

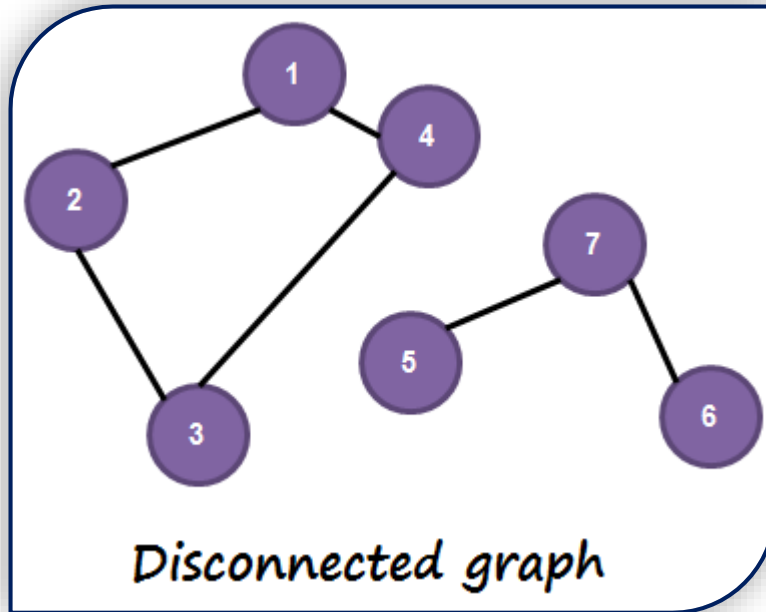


### 8. Connected Graph



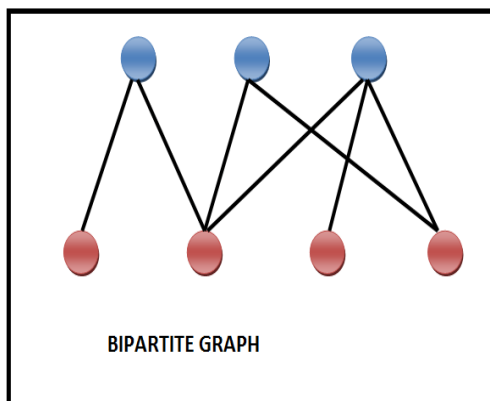
A graph  $G$  is said to be *connected* if there exists a path between every pair of vertices. There should be at least one edge for every vertex in the graph. So that we can say that it is connected to some other vertex at the other side of the edge.

### 9. Disconnected Graph



A graph  $G$  is *disconnected*, if it does not contain at least two connected vertices.

### 10. Bipartite Graph

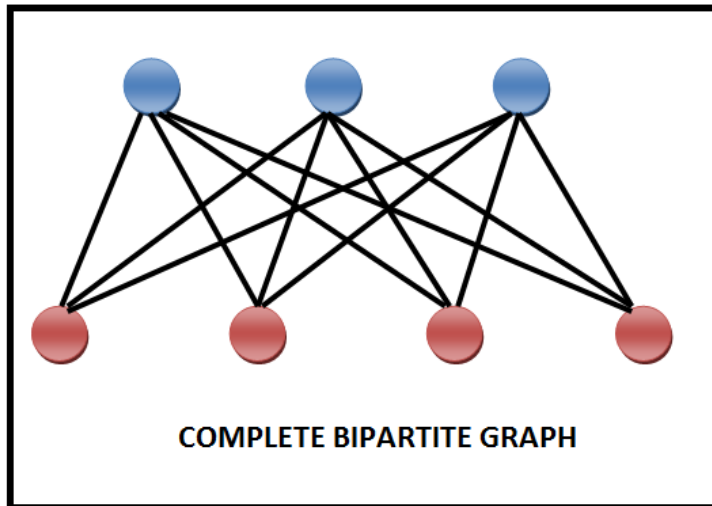


A simple graph  $G = (V, E)$  with vertex partition  $V = \{V_1, V_2\}$  is called a *bipartite graph* if every edge of  $E$  joins a vertex in  $V_1$  to a vertex in  $V_2$ .

In general, a Bipartite graph has two sets of vertices, let us say,  $V_1$  and  $V_2$ , and if an edge is drawn, it should connect any vertex in set  $V_1$  to any vertex in set  $V_2$ .

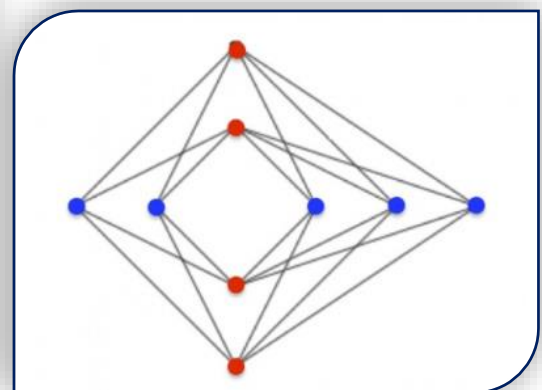
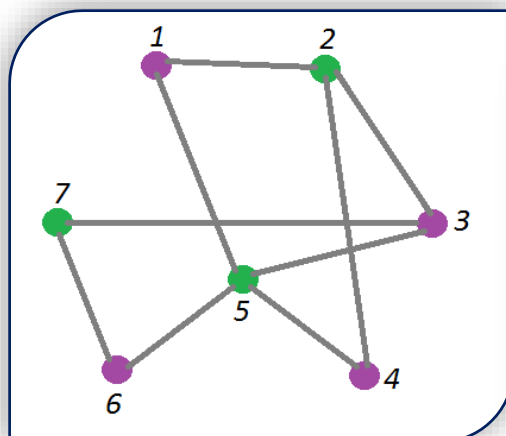
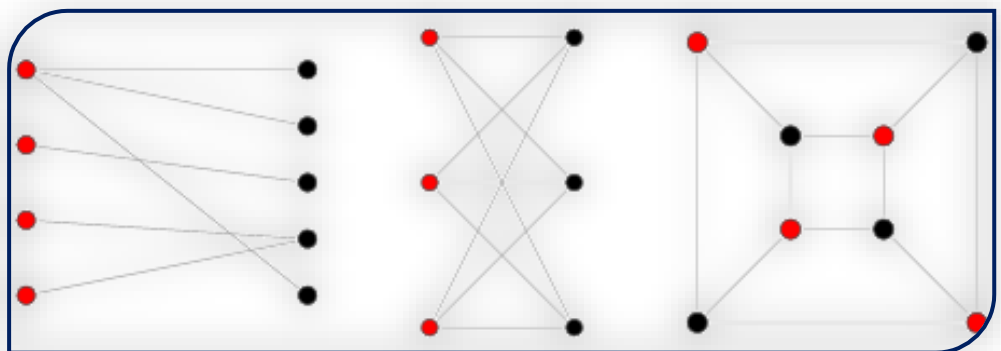
## 11. Complete Bipartite Graph

A bipartite graph 'G',  $G = (V, E)$  with partition  $V = \{V_1, V_2\}$  is said to be a *complete bipartite graph* if every vertex in  $V_1$  is connected to every vertex of  $V_2$ .



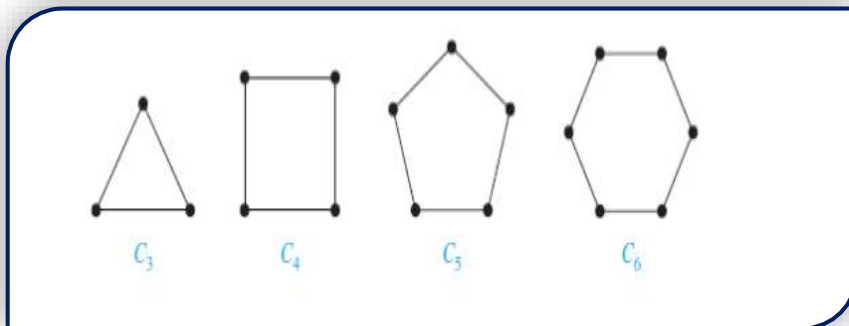
In general, a complete bipartite graph connects each vertex from set  $V_1$  to each vertex from set  $V_2$ .

Some bipartite graphs examples:





## 12. Cycle Graph



A simple graph with 'n' vertices ( $n \geq 3$ ) and 'n' edges is called a cycle graph if all its edges form a cycle of length 'n'.

If the **degree of each vertex in the graph is two**, then it is called a *Cycle Graph*.

**Notation** –  $C_n$

## 13. Odd Cycle

An *odd cycle* is a cycle with odd length, that is, with an odd number of edges.

## 14. Even Cycle

An *even cycle* is a cycle with even length, that is, with an even number of edges.

### Problem 1– Group C6

**Prove that the number of colours required for a proper edge colouring of a graph  $G$  is  $\geq$  the maximum degree of any vertex of  $G$ .**

Let's prove the statement by Contra positive method.

#### **Statement:**

The number of colours required for a proper edge colouring of a graph  $G$  is  $\geq$  The maximum degree of any vertex of  $G$

Let,

- The number of colours required for proper edge colouring of  $G$  (Edge chromatic number) be ' $n$ '
- The maximum degree of any vertex of  $G$  be ' $x$ '.

#### **Prepositions in the Statement:**

$P$  : Graph  $G$  is properly coloured

$\neg P$  : Graph  $G$  is not properly coloured

$Q$  :  $n \geq x$  (number of colours required for proper edge colouring is greater than or equal to the Maximum degree of any vertex of  $G$ )

$\neg Q$  :  $n < x$

#### **To prove:**

$P \Rightarrow Q$

By contra positive method, we have to prove that,  $\neg Q \Rightarrow \neg P$  That is, if  $n < x$ , then graph  $G$  is not properly coloured.

#### **Proof:**

Let's take the following graph

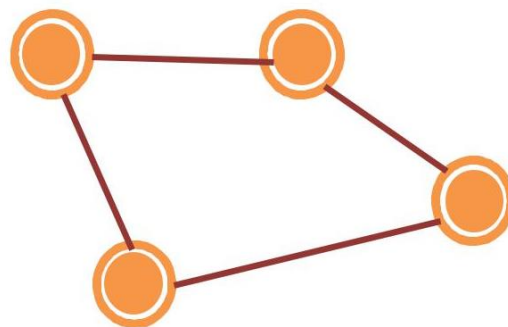
Here maximum degree of any vertex = 2

So  $x = 2$

For  $n < x$  to be true, we must be able to colour the edges using at most 1 colour

This violates the definition of proper edge colouring because two edges incident on the same vertex have the same colour.

That is the graph  $G$  is not properly coloured. ( $\neg P$ )



Taking  $\neg Q$  as true, we have proved that  $\neg P$ . ( $\neg Q \Rightarrow \neg P$ ).

By the contrapositive method,  $P \Rightarrow Q$

**If Graph G is properly coloured, then the number of colours required is  $\geq$  maximum degree of any vertex of G.**

### Problem 2– Group C6

**Prove that a graph G is bipartite if and only if it does not have a cycle of odd length.**

$\Rightarrow$  **G is bipartite  $\Rightarrow$  G does not have a cycle of odd length**

Let G is bipartite then

$$V(G) = A \cup B \quad A \cap B = \emptyset$$

For every  $e \in E(G)$   $e = ab$  where  $a \in A$ ,  $b \in B$

(This is definition of bipartite graph)

Suppose G has (at least) one odd length cycle.

Let the length of C be n

$$\text{Let } C = (v_1, v_2, \dots, v_n)$$

Let  $v_1 \in A$  then  $v_2 \in B$  then  $v_3 \in A$  and so on...

Hence we see that for all  $k = \{1, 2, \dots, n\}$

$$v_k \in A : k \text{ odd}$$

$$v_k \in B : k \text{ even}$$

But as n is odd,  $v_n \in A$

$$v_1, v_n \in A$$

So  $v_1 v_n \in E(G)$  which contradicts the assumption that G is bipartite

Therefore, G is bipartite  $\Rightarrow$  G does not have a cycle of odd length

$\Leftarrow$  **G does not have a cycle of odd length  $\Rightarrow$  G is bipartite**

If G is disconnected then any cycle in G is contained in one of the connected components.

It suffices to show the claim for connected graphs. Assume G is connected graph.

Suppose G has no odd length cycle.

Let  $v \in V(G)$

$A = \{a \mid d(v, a) \text{ is even} \}$

$B = \{b \mid d(v, b) \text{ is odd} \}$

$d$ - shortest path

Since  $G$  is connected,

$A \cup B = V(G) \quad A \cap B = \emptyset$

We need to show that any  $e \in E(G)$  is at the form  $e = ab \mid a \in A, b \in B$

Need to show that  $a_1, a_2 \in E(G) \mid a_1, a_2 \in A$  OR  $b_1, b_2 \in E(G) \mid b_1, b_2 \in B$  it is impossible.

Suppose  $a_1, a_2 \in A$  are adjacent,  $a_1 a_2 \in E(G)$

Then there would be a closed walk at odd length  $(v, \dots, a_1, a_2, \dots, v)$

But from graph containing closed walk of odd length also contains odd cycle it follows that  $G$  would then contain an odd cycle.

This contradicts our assumption  $G$  has no odd length cycles,

Therefore  $a_1 a_2 \notin E(G)$

By same argument if  $b_1, b_2 \in B \Rightarrow b_1 b_2 \notin E(G)$

Therefore, any  $e \in E(G)$  is the form  $e = ab \mid a \in A, b \in B$

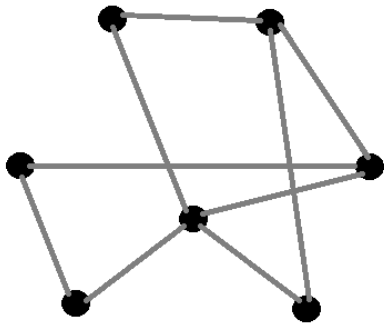
Therefore,  $G$  is bipartite.

## References

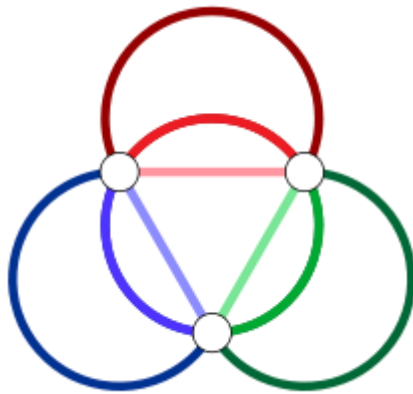
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- [https://proofwiki.org/wiki/Graph\\_is\\_Bipartite\\_iff\\_No\\_Odd\\_Cycles](https://proofwiki.org/wiki/Graph_is_Bipartite_iff_No_Odd_Cycles)
- <https://www.youtube.com/watch?v=YiGFhWxtHjQ>

## Questions

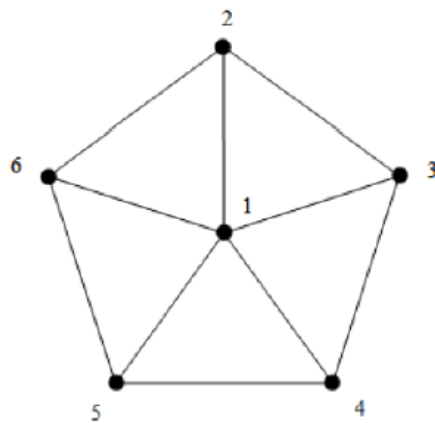
a) The following graph is a bipartite graph. (True/False)



b) The following graph edge chromatic number is 9. (True/False)

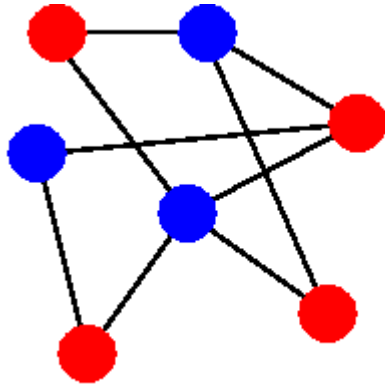


c) What is the chromatic number of following wheel graph?



## Solutions

a) True



b) False

Edge chromatic number is 3.

c) Chromatic number is 4.

$W_n$  ( $n \geq 3$ ) the wheel graph with  $n+1$  vertices then

$\chi(W_n) = 3$  ; if  $n$  is even

$\chi(W_n) = 4$  ; if  $n$  is odd