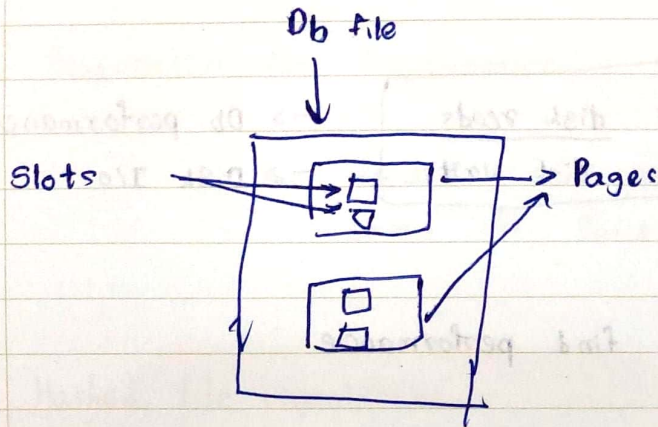


to optimize db we need file organization.

## File organization & indexes

lec 7



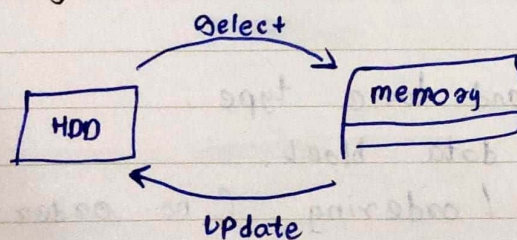
- Db file is a collection of pages, page is a collection of slots.

- in slots store data, each slot have a single data.
- Page is a collection of slots.

if there are 100 records there will be 100 slots.

- actually data stored in HDD. if we run a query to select data it will read data from HDD and send to memory.

- in update store to memory as select and update is in the memory. and write back in HDD.



1 disk block = 1 db file page.

at a time bring one page at a time to memory  
( HDD  $\xrightarrow{\text{read}}$  memory )



db performance Calculated based on disk read and disk write.  $\epsilon$

$$\text{Performance} = \frac{\text{disk reads}}{\text{disk writes}}$$

$\rightarrow$  Db performance

$\rightarrow$  Disk I/O

1000  $\rightarrow$  records

10 records in 1 page find performance

$$= \frac{1000}{10}$$

$$= 100 \text{ (read : write)}$$

3 type of file organization

1. Heap file

2. Sequential file

3. Hashing file

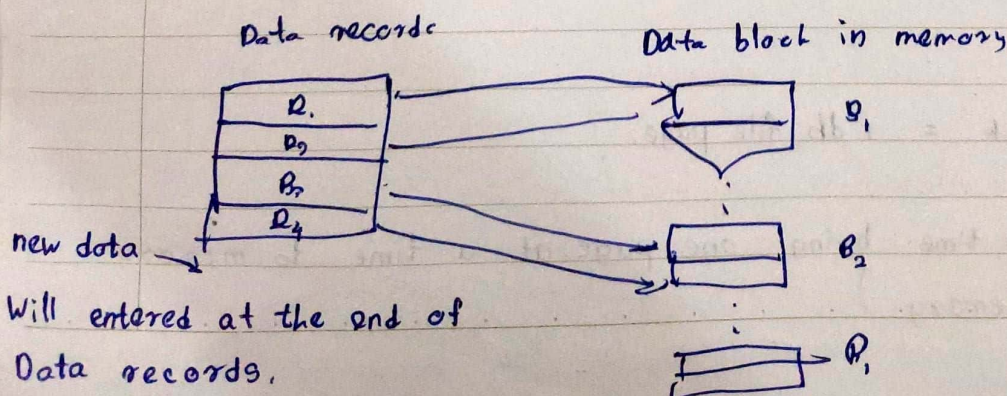
} to optimize the db.

a) Heap file organization

• Simplest and basic type.

• Work with data block.

• No sorting / ordering. (no order to store data)



ProMate



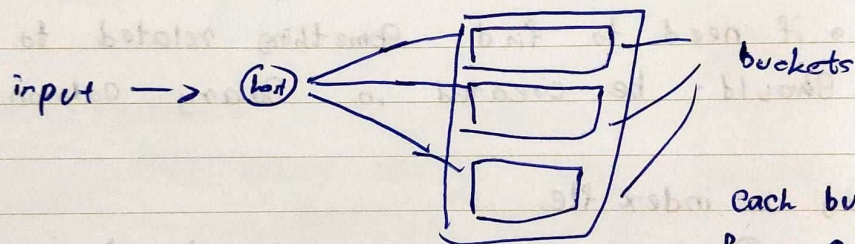
- Update / delete time consuming (need to search)
- not good for large files.

### 02) Sequential file Organization.

- Records are stored in a order. (Age / desc)
- based on some column. Can Store data.

### 03) Hashed file organization.

- insert our input through a hash function. and taking a hash output, and assigned to a buckets.



Each bucket has label from 0 to 9999.

according to the remainder it will allocate for bucket

$$\text{bucket} = \frac{\text{Value}}{\text{bucket count}}$$

$$= \frac{90000}{10}$$

= take the remainder only '0'

= insert to '0' bucket.

- Hashed file organization not sorted well.
- Store in different buckets according to remainders.

ProMate



# INDEX.

index mainly 2 part

- Search key (key)
- data reference (value)

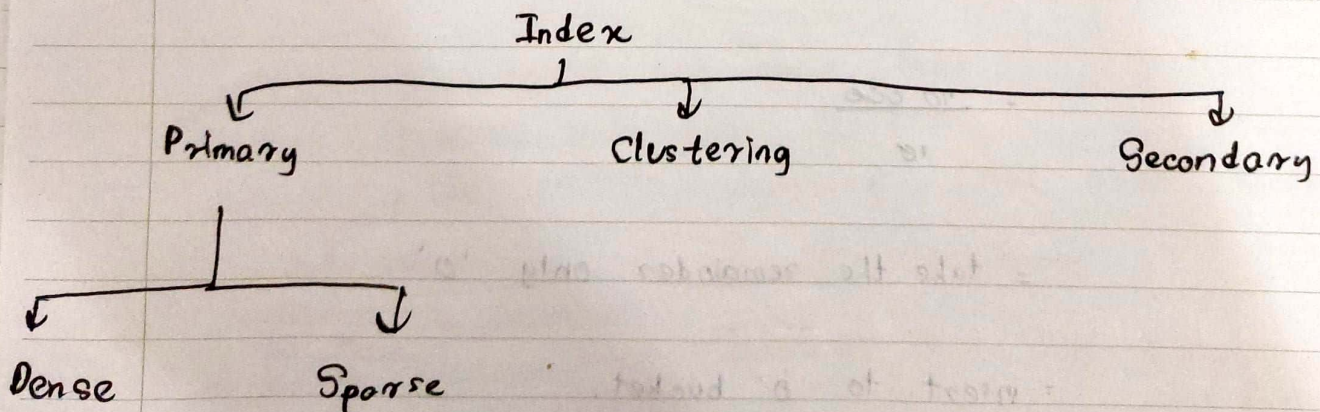
• together called index. Used to optimize the Speed of db.

• it reduce the read and write to the HDD. to create a index we can use more than 1 column.

• if need to find something related to Salary index should be created in Salary column.

Creating an index file.

- ① Create a separate index file for the column(s).
- ② Enter all values to index file.
- ③ Sort the values (ASC / DESC).





## Primary index.

• is a  $\neq$  ordered file which is fixed length size with tow fields. (ordered file)

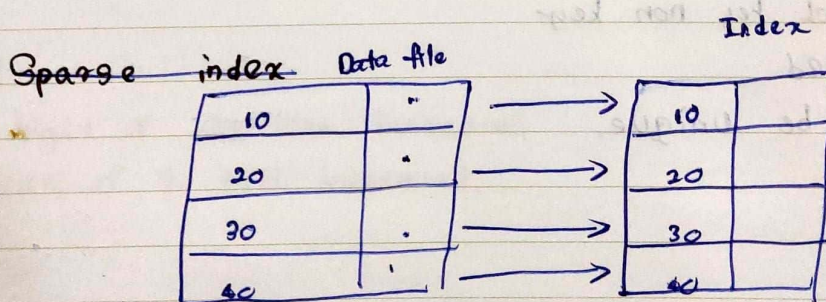
• Created using primary key.

• 1 to 1 relationship between the entries in index table.

- Search key is
  - Sorted
  - not null
  - fast

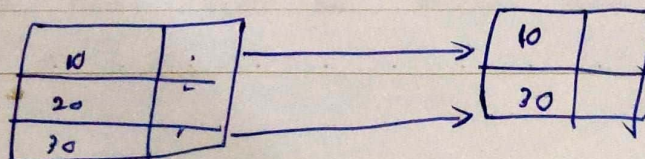
## Dense index

- for every search key record is created
  - fast
  - accurate
  - need more storage.



## Sparse index.

- appeared only for some B. Search key.



ProMate

ProMate



## Secondary indexing

- Can be generated by a candidate key. With a level indexing,

- data is not sorted. non clustered index.

- more time required
- not sorted
- Search key cant be null
- slower than primary and faster than clustered.

## Clustering index.

- used multiple related records found at one place.

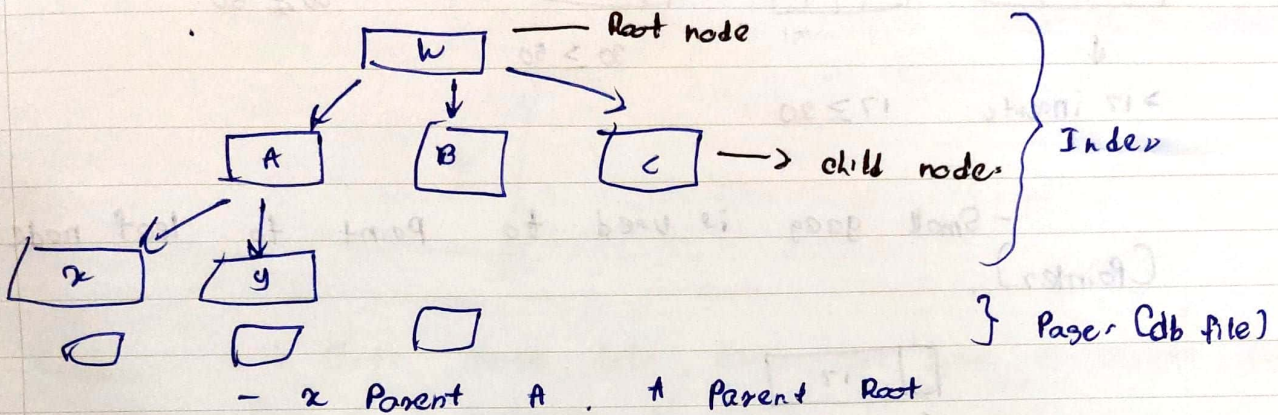
- Sorted data.
- records stored in indexes.
- Created in non - prim columns.
- Can group columns.
- Search key non key
- Sorted
- not be unique.



Using B+ tree.

for range searches.

- Use to implement the db indexes (B+ trees)



- upper layer Parent, lower layer child.

- last nodes are leaf nodes. We actually store db file records in leaf node.

- Can use for insert, update, delete.

Insert.

- height of B+ tree increased.
- width of B+ tree increased.

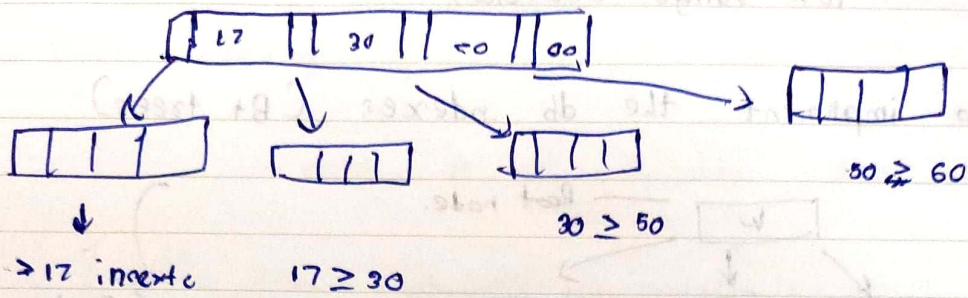
Delete.

- height ↓
- width ↓

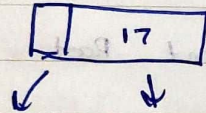
- each node must be filled 50% without root.
- Can't avoid it.



## B+ tree



- Small gaps are used to point to leaf node.  
(Pointer).



Pointer key node