Design and Analysis of Algorithms

Lecture-28

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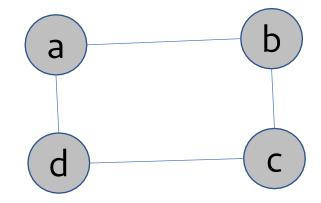
Spanning tree:

Spanning tree is a non-cyclic sub-graph of a connected and undirected graph G that connects all the vertices together.

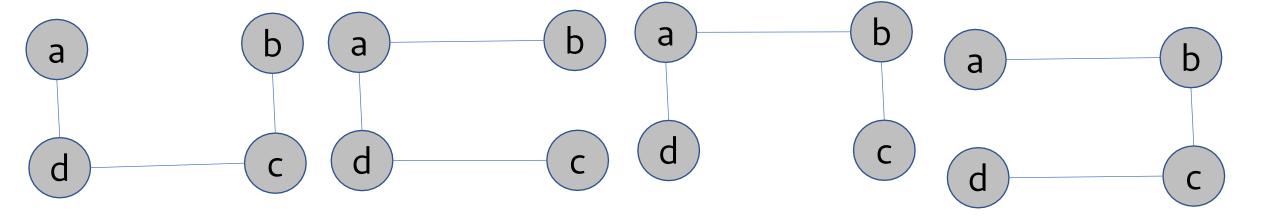
General Properties of Spanning Tree

- A connected graph G can have more than one spanning tree.
- All possible spanning trees of graph G, have the same number of edges and vertices.
- The spanning tree does not have any cycle (loops).
- Removing one edge from the spanning tree will make the graph disconnected, i.e. the spanning tree is **minimally connected**.
- Adding one edge to the spanning tree will create a circuit or loop.
- Spanning tree has **n-1** edges, where **n** is the number of nodes (vertices).
- A complete graph can have maximum n^{n-2} number of spanning trees.

Example: Consider the following graph:-



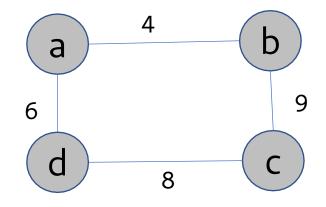
All the spanning tree of this graph are the following:-



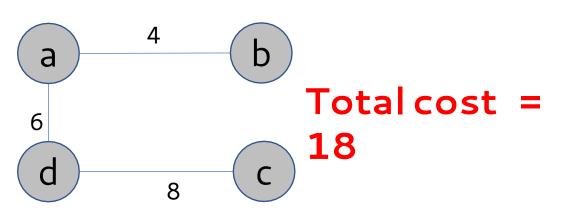
If the given graph is weighted graph, then we define the minimum spanning tree.

<u>Definition</u>: A spanning tree is said to be minimum spanning tree if sum of weights of all the edges in the tree is smallest.

Example: Consider the following graph:-



Graph



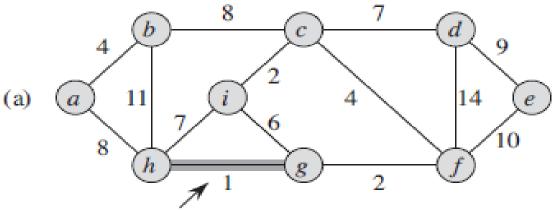
Minimum spanning tree

- In this chapter, we will study two algorithms to find the minimum spanning tree.
- (1) Kruskal's Algorithm
- (2) Prim's Algorithm

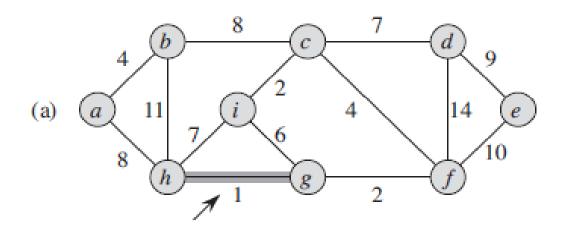
Both algorithms are based on Greedy approach.

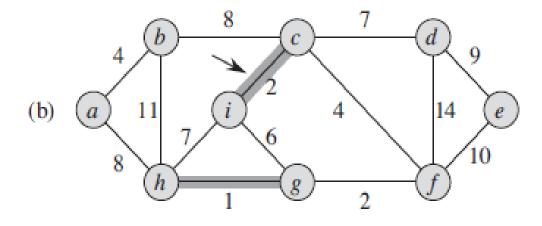
Example: Find the minimum spanning of the following graph using Kruskal

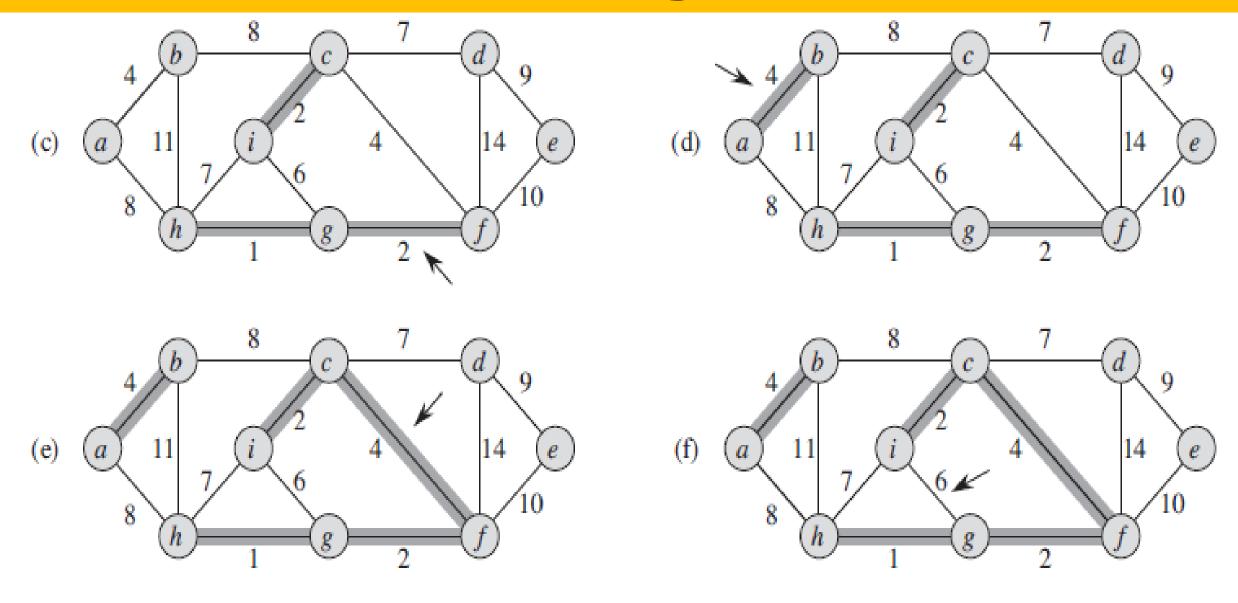
algorithm.

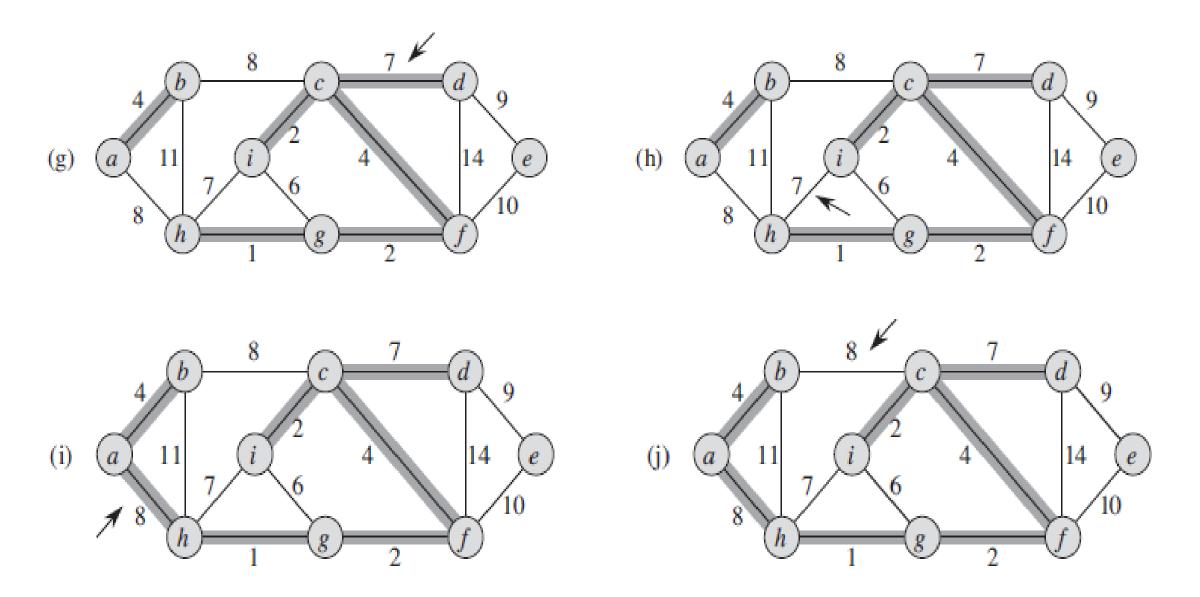


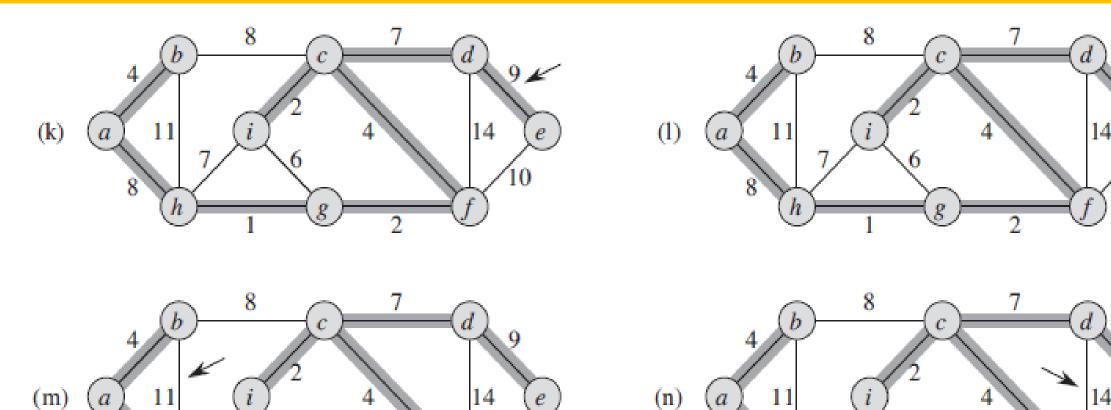
Solution:











Total cost = 37

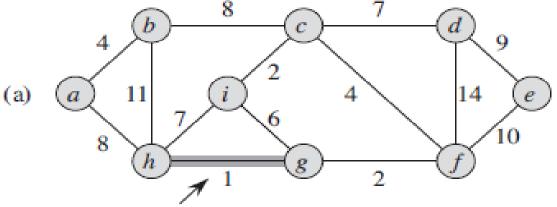
Final minimum spanning tree

```
MST-KRUSKAL(G, w)
1 \quad A = \emptyset
   for each vertex v \in G.V
       MAKE-SET(\nu)
   sort the edges of G.E into nondecreasing order by weight w
   for each edge (u, v) \in G.E, taken in nondecreasing order by weight
       if FIND-SET(u) \neq FIND-SET(v)
            A = A \cup \{(u, v)\}
            Union(u, v)
   return A
```

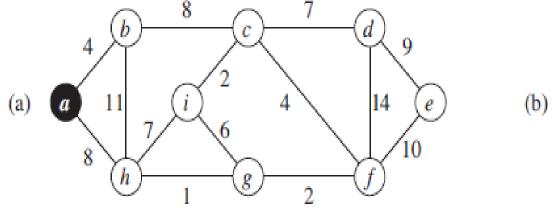
Time complexity of this algorithm is O(E lg E) or O(E lgV).

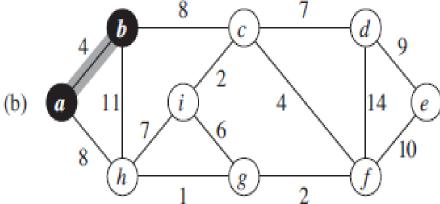
Example: Find the minimum spanning of the following graph using Prim's

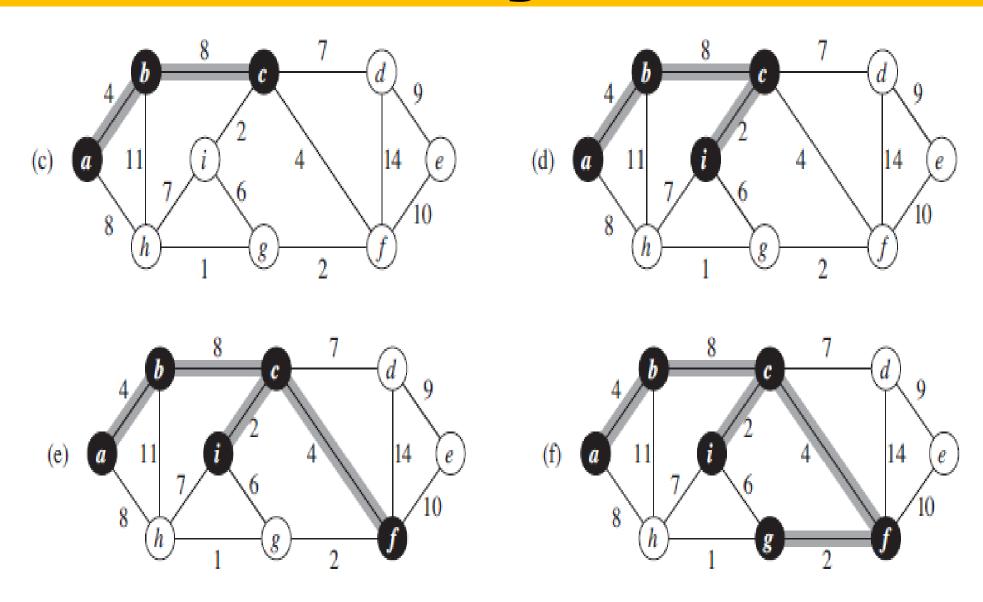
algorithm.

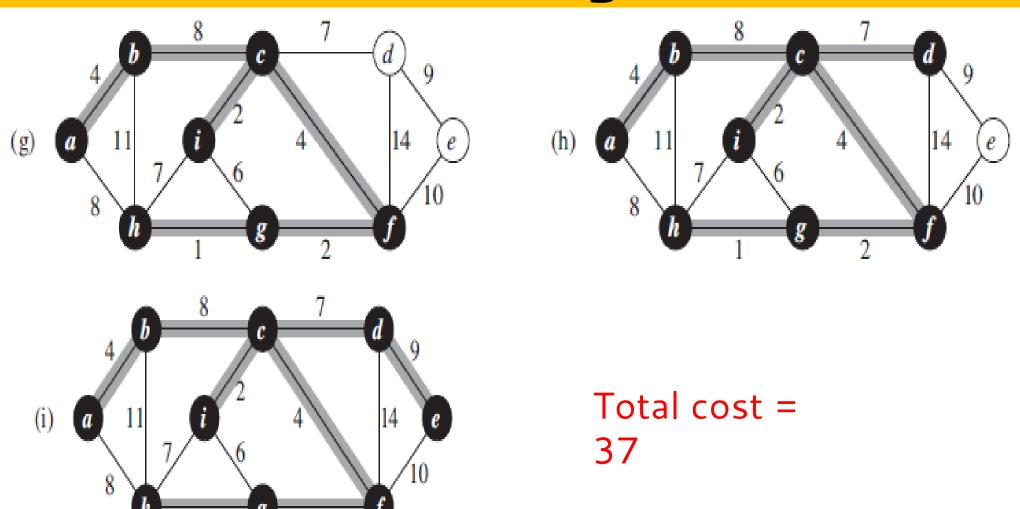


Solution:









Final minimum spanning

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```
MST-PRIM(G, w, r)
    for each u \in G, V
    u.key = \infty
      u.\pi = NIL
   r.key = 0
 5 \quad Q = G.V
    while Q \neq \emptyset
         u = \text{EXTRACT-MIN}(Q)
         for each v \in G. Adj[u]
              if v \in Q and w(u, v) < v.key
10
                  \nu.\pi = u
                  v.key = w(u, v)
11
```

Time complexity of this algorithm is O(E lgV + V lgV) = O(E lgV).