# Theory of Automata and Formal Language Lecture-2

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#### **Definition of Grammar**

A grammar G is defined as a quadruple  $G = (V, \Sigma, S, P)$ , where

 $V \to \mathsf{Finite}$  set of variables or non-terminal symbols

 $\Sigma \to \mathsf{Finite}$  set of terminal symbols

 $S \in V \to Starting$  symbol of the Grammar

 $P \rightarrow Finite set of production rules$ 

And P is defined as following:-

$$P\subseteq (V\cup\Sigma)^*X(V\cup\Sigma)^*$$

 $(u,v) \in P$  is denoted by

 $u \rightarrow v$ , where  $u,v \in (V \cup \Sigma)^*$ 

u always contains at least one variable.

# Direct derivation or derivation in one step

Let  $\alpha, \beta \in (V \cup \Sigma)^*$ .  $\alpha$  directly derives  $\beta$  if  $\alpha \to \beta \in P$ . It is denoted by  $\alpha \Rightarrow \beta$ .

# **Derivation in many steps**

Let  $\alpha, \beta \in (V \cup \Sigma)^*$ .  $\alpha$  derives  $\beta$  if there exists production rules  $\alpha \to A_1$ ,  $A_1 \to A_2, A_2 \to A_3, \ldots, A_n \to \beta$ , such that  $\alpha \Rightarrow A_1 \Rightarrow A_2 \Rightarrow A_3 \Rightarrow \ldots \Rightarrow A_n \Rightarrow \beta$ . It is denoted by  $\alpha \stackrel{*}{\Rightarrow} \beta$ .

#### Sentential Form

Let  $\alpha \in (V \cup \Sigma)^*$ . The string  $\alpha$  is said to be in the sentential form if

 $S \stackrel{*}{\Rightarrow} \alpha$ , where S is the starting symbol.

# Language generated by a Grammar

The set of all the sentental forms consisting of only terminal symbols is said to be language generated Grammar. That is,

$$L(G) = \{ x \in \Sigma^* ! S \stackrel{*}{\Rightarrow} x \}$$

# **Equivalent Grammars**

Two grammars  $G_1$  and  $G_2$  are said to be equivalent grammar if the languages generated by both grammars are the same. That is,

$$\mathsf{L}(\mathit{G}_{1})=\mathsf{L}(\mathit{G}_{2}).$$

# **Examples**

Determine the languages generated by the following grammars:-

- 1.  $S \rightarrow aS/a$
- 2.  $S \rightarrow 0S1/\epsilon$
- 3. S 
  ightarrow aCa , C 
  ightarrow aCa/b
- 4.  $S \rightarrow aS/bS/a/b$
- 5. S 
  ightarrow 0 S A2 , S 
  ightarrow 012 , 2 A 
  ightarrow A2 , 1 A 
  ightarrow 11