

Theory of Automata and Formal Language

Lecture-18

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Equivalent regular expressions

Two regular expressions \bar{p} and \bar{q} are said to be equivalent if they represent the same set of strings.

Identities for regular expressions

Let p, q, r are all the regular expressions. Then

$$(1) \phi + r = r$$

$$(2) \phi \cdot r = \phi = r \cdot \phi$$

$$(3) \epsilon \cdot r = r \cdot \epsilon = r$$

$$(4) \epsilon^* = \epsilon \text{ and } \phi^* = \epsilon \quad (5) r + r = r$$

$$(6) r^* r^* = r^*$$

$$(7) r r^* = r^* r$$

$$(8) (r^*)^* = r^*$$

$$(9) \epsilon + r r^* = r^* = \epsilon + r^* r$$

$$(10) (p \cdot q)^* p = p \cdot (p \cdot q)^*$$

$$(11) (p + q)^* = (p^* q^*)^* = (p^* + q^*)^*$$

$$(12) (p + q) \cdot r = p \cdot r + q \cdot r \text{ and } r \cdot (p + q) = r \cdot p + r \cdot q$$

ARDEN's Theorem

Statement: Let P and Q be two regular expressions over Σ . If P does not contain ϵ , then the following equation in R,

$$R = Q + RP \dots\dots\dots(1)$$

has a unique solution $R = QP^*$.

Proof:

$$\begin{aligned} Q + RP &= Q + QP^*P \\ &= Q(\epsilon + P^*P) \\ &= QP^* \\ &= R \end{aligned}$$

Therefore equation (1) is satisfied when $R = QP^*$.

Therefore $R = QP^*$ is a solution of equation (1).

Regular Expression

To prove uniqueness of (1), replacing R by $Q+RP$.

$$R = Q + RP = Q + (Q+RP)P$$

$$= Q + QP + RP^2$$

$$= Q + QP + (Q + RP)P^2$$

$$= Q + QP + QP^2 + RP^3$$

.....

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$$= Q + QP + QP^2 + \dots + QP^i + RP^{i+1}$$

$$= Q(\epsilon + P + P^2 + \dots + P^i) + RP^{i+1}$$

$$\text{Therefore, } R = Q(\epsilon + P + P^2 + \dots + P^i) + RP^{i+1} \dots \dots \dots (2)$$

Regular Expression

Let $w \in R$ and $|w| = i$.

From equation (2), w will belong into right hand side of equation (2).

Therefore, $w \in Q(\epsilon + P + P^2 + \dots + P^i) + RP^{i+1}$

Clearly, $w \notin RP^{i+1}$ since $|w| = i$.

Therefore, $w \in Q(\epsilon + P + P^2 + \dots + P^i)$

Therefore, $w \in QP^*$

Therefore, $R \subseteq QP^* \dots \dots \dots (3)$

Regular Expression

Let $w \in QP^*$.

Then $w \in QP^k$ for some $k \geq 0$.

Therefore, $w \in Q(\epsilon + P + P^2 + \dots + P^k)$

Therefore, w belongs into the right hand side of equation (2).

Therefore, $w \in R$.

Therefore, $QP^* \subseteq R \dots\dots\dots(4)$

From (3) and (4),

$R = QP^*$

Example:

(a) Give a regular expression for representing the set L of strings in which every 0 is immediately followed by at least two 1's.

(b) Prove that the regular expression

$$r = \epsilon + 1^*(011)^*(1^*(011)^*)^*$$

also describes the same set of strings.

Example: Prove that

$$(1 + 00^*1) + (1 + 00^*1)(0 + 10^*1)^*(0 + 10^*1) = 0^*1(0 + 10^*1)^*$$