

Discrete Structures and Theory of Logic

Lecture-36

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Theory of inference for statement calculus

Let A and B be two statement formulas. We say that "B logically follows from A" or "B is a valid conclusion of the premise A" iff $A \rightarrow B$ is a tautology. That is, $A \Rightarrow B$.

A conclusion C follows from a set of premises $\{H_1, H_2, \dots, H_m\}$ iff

$$H_1 \wedge H_2 \wedge \dots \wedge H_m \Rightarrow C$$

Exercise

1. Determine whether the conclusion C follows logically from the premises H_1 and H_2 .

- | | | |
|-----------------------------|-----------------------------|------------------------|
| (a) $H_1 : P \rightarrow Q$ | $H_2 : P$ | $C : Q$ |
| (b) $H_1 : P \rightarrow Q$ | $H_2 : \neg P$ | $C : Q$ |
| (c) $H_1 : P \rightarrow Q$ | $H_2 : \neg(P \wedge Q)$ | $C : \neg P$ |
| (d) $H_1 : \neg P$ | $H_2 : P \leftrightarrow Q$ | $C : \neg(P \wedge Q)$ |
| (e) $H_1 : P \rightarrow Q$ | $H_2 : Q$ | $C : P$ |

Exercise

2. Show that the conclusion C follows from the premises H_1, H_2, \dots in the following cases:-

(a) $H_1 : P \rightarrow Q$

$C : P \rightarrow (P \wedge Q)$

(b) $H_1 : \neg P \vee Q$

$H_2 : \neg(Q \wedge \neg R)$

$H_3 : \neg R$

$C : \neg P$

(c) $H_1 : \neg P$

$H_2 : P \vee Q$

$C : Q$

(d) $H_1 : \neg Q$

$H_2 : P \rightarrow Q$

$C : \neg P$

(e) $H_1 : P \rightarrow Q$

$H_2 : Q \rightarrow R$

$C : P \rightarrow R$

(f) $H_1 : R$

$H_2 : P \vee \neg P$

$C : R$

Exercise

3. Determine whether the conclusion C is valid in the following, when H_1, H_2, \dots are the premises.

- (a) $H_1 : P \rightarrow Q$ $H_2 : \neg Q$ $C : P$
- (b) $H_1 : P \vee Q$ $H_2 : P \rightarrow R$ $H_3 : Q \rightarrow R$ $C : R$
- (c) $H_1 : P \rightarrow (Q \rightarrow R)$ $H_2 : P \wedge Q$ $C : R$
- (d) $H_1 : P \rightarrow (Q \rightarrow R)$ $H_2 : R$ $C : P$
- (e) $H_1 : \neg P$ $H_2 : P \vee Q$ $C : P \wedge Q$

4. Without constructing a truth table, show that $A \wedge E$ is not a valid consequence of

$$A \leftrightarrow B, B \leftrightarrow (C \wedge D), C \leftrightarrow (A \vee E) \text{ and } A \vee E$$

Also show that $A \vee C$ is not a valid consequence of

$$A \leftrightarrow (B \rightarrow C), B \leftrightarrow (\neg A \vee \neg C), C \leftrightarrow (A \vee \neg B) \text{ and } B$$

Implication rules

1. Simplification

$$P \wedge Q \Rightarrow P, P \wedge Q \Rightarrow Q$$

2. Addition

$$P \Rightarrow P \vee Q, Q \Rightarrow P \vee Q$$

3. Disjunctive Syllogism

$$\neg P, P \vee Q \Rightarrow Q$$

4. Modus Ponens

$$P, P \rightarrow Q \Rightarrow Q$$

5. Modus Tollens

$$\neg Q, P \rightarrow Q \Rightarrow \neg P$$

6. Hypothetical Syllogism

$$P \rightarrow Q, Q \rightarrow R \Rightarrow P \rightarrow R$$

7. Dilemma

$$P \vee Q, P \rightarrow R, Q \rightarrow R \Rightarrow R$$

8. Conjunction

$$P, Q \Rightarrow P \wedge Q$$

Rules of Inference

Rule P: A premise may be introduced at any point in the derivation.

Rule T: A formula S may be introduced in a derivation if S is tautology implied by any one or more of the preceding formulas in the derivation.

Example: Demonstrate that R is a valid inference from the premises $P \rightarrow Q$, $Q \rightarrow R$ and P .

Solution:

- | | |
|-----------------------|--|
| (1) $P \rightarrow Q$ | By rule P |
| (2) $Q \rightarrow R$ | By rule P |
| (3) $P \rightarrow R$ | By rule T, (1), (2) and hypothetical syllogism |
| (4) P | By rule P |
| (5) R | By rule T, (3), (4) and modus ponens |

Example: Show that $R \vee S$ follows logically from the premises $C \vee D$, $(C \vee D) \rightarrow \neg H$, $\neg H \rightarrow (A \wedge \neg B)$ and $(A \wedge \neg B) \rightarrow (R \vee S)$.

Solution:

- | | |
|--|--|
| (1) $(C \vee D) \rightarrow \neg H$ | By rule P |
| (2) $\neg H \rightarrow (A \wedge \neg B)$ | By rule P |
| (3) $C \vee D \rightarrow (A \wedge \neg B)$ | By rule T, (1), (2) and hypothetical syllogism |
| (4) $(A \wedge \neg B) \rightarrow (R \vee S)$ | By rule P |
| (5) $C \vee D \rightarrow (R \vee S)$ | By rule T, (3), (4) and hypothetical syllogism |
| (6) $C \vee D$ | By rule P |
| (7) $R \vee S$ | By rule T, (5), (6) and modus ponens |

Example: Show that $R \vee S$ is tautologically implied by $(P \vee Q) \wedge (P \rightarrow R) \wedge (Q \rightarrow S)$.

Solution:

- | | |
|----------------------------|--|
| (1) $P \vee Q$ | By rule P |
| (2) $\neg P \rightarrow Q$ | By rule T, (1) and $(P \rightarrow Q \Leftrightarrow \neg P \vee Q)$ |
| (3) $Q \rightarrow S$ | By rule P |
| (4) $\neg P \rightarrow S$ | By rule T, (2), (3) and hypothetical syllogism |
| (5) $\neg S \rightarrow P$ | By rule T, (4) and $(\neg P \rightarrow S \Leftrightarrow \neg S \rightarrow P)$ |
| (6) $P \rightarrow R$ | By rule P |
| (7) $\neg S \rightarrow R$ | By rule T, (5), (6) and hypothetical syllogism |
| (8) $R \vee S$ | By rule T, (7) and $(P \rightarrow Q \Leftrightarrow \neg P \vee Q)$ |