

Theory of Automata and Formal Language

Lecture-11

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Finite Automata (FA)

Theorem: For a given NFA M , there exists a DFA M' such that $L(M) = L(M')$.

Proof:

Suppose the given NFA M is

$$M = (Q, \Sigma, \delta, q_0, F)$$

Now, we construct a DFA M' as following

$$M' = (Q', \Sigma, \delta', q'_0, F')$$

Where $Q' = P(Q)$, $q'_0 = \{q_0\}$

$$F' = \{P \subseteq Q \mid P \cap F \neq \emptyset\}$$

and $\delta' \equiv \hat{\delta}$

$$\text{i.e. } \delta'(P, a) = \hat{\delta}(P, a) = \bigcup_{p \in P} \delta(p, a) = \bigcup_{p \in P} \delta(p, a)$$

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Now, we have to show that

$$L(M) = L(M') \dots\dots\dots (1)$$

To prove the equation (1), first we will prove that

$$\hat{\delta}'(q'_0, w) = \hat{\delta}(q_0, w), \forall w \in \Sigma^* \dots\dots\dots (2)$$

We will prove this by using induction method on the length of string w .

For $|w| = 0$, i.e. $w = \epsilon$

$$\hat{\delta}'(q'_0, \epsilon) = q'_0 = \{q_0\} = \hat{\delta}(q_0, \epsilon)$$

$$\text{Therefore, } \hat{\delta}'(q'_0, \epsilon) = \hat{\delta}(q_0, \epsilon)$$

Therefore, it is proved for $w = \epsilon$.

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For $|w| = 1$, i.e. $w = a$, where $a \in \Sigma$

$$\begin{aligned}\hat{\delta}'(q'_0, a) &= \delta'(q'_0, a) \\ &= \hat{\hat{\delta}}(\{q_0\}, a) \\ &= \hat{\delta}(q_0, a)\end{aligned}$$

Therefore, $\hat{\delta}'(q'_0, a) = \hat{\delta}(q_0, a)$

Therefore, it is proved for $w = a$, i.e. $|w| = 1$.

Suppose equation (2) is true for string $w = x$.

Therefore, $\hat{\delta}'(q'_0, x) = \hat{\delta}(q_0, x)$,(3)

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Now, we will prove for $w = xa$.

$$\begin{aligned}\hat{\delta}'(q'_0, xa) &= \delta'(\hat{\delta}'(q'_0, x), a) \\ &= \hat{\hat{\delta}}(\hat{\delta}(q_0, x), a) \\ &= \bigcup_{p \in \hat{\delta}(q_0, x)} \hat{\delta}(p, a) \\ &= \bigcup_{p \in \hat{\delta}(q_0, x)} \delta(p, a) \\ &= \hat{\delta}(q_0, xa)\end{aligned}$$

Therefore, $\hat{\delta}'(q'_0, xa) = \hat{\delta}(q_0, xa)$

Therefore, it is proved for $w = xa$.

Therefore, equation (2) is proved.

Now, we prove the equation (1) by using equation (2).

$$\begin{aligned}\text{Let } w \in L(M') &\Leftrightarrow \hat{\delta}'(q'_0, w) \in F' \\ &\Leftrightarrow \hat{\delta}'(q'_0, w) \cap F \neq \phi \\ &\Leftrightarrow \hat{\delta}(q_0, w) \cap F \neq \phi \\ &\Leftrightarrow w \in L(M)\end{aligned}$$

Therefore, $L(M) = L(M')$

Now, it is proved.

Some Examples

Example: Find DFA equivalent to the following NFA:-

$$M = (\{q_0, q_1, q_2\}, \{a, b\}, \delta, q_0, \{q_2\}),$$

where δ is given by

| δ | a | b |
|-------------------|------------|------------|
| $\rightarrow q_0$ | q_0, q_1 | q_2 |
| q_1 | q_0 | q_1 |
| q_2 | | q_0, q_1 |

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Example: Find DFA equivalent to the following NFA:-

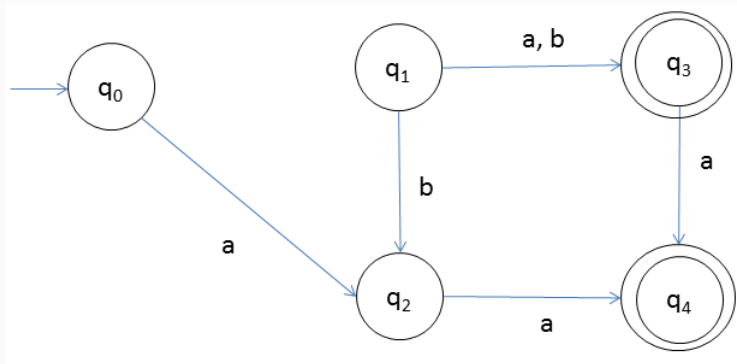
$$M = (\{q_0, q_1, q_2, q_3\}, \{a, b\}, \delta, q_0, \{q_3\}),$$

where δ is given by

| δ | a | b |
|-------------------|------------|-------|
| $\rightarrow q_0$ | q_0, q_1 | q_0 |
| q_1 | q_2 | q_1 |
| q_2 | q_3 | q_3 |
| q_3 | | q_2 |

Finite Automata (FA)

Example: Construct DFA equivalent to the following NFA:-



Finite Automata (FA)

Example: Construct DFA equivalent to the following NFA:-

| δ | a | b |
|-------------------|------------|------------|
| $\rightarrow q_0$ | q_1, q_3 | q_2, q_3 |
| q_1 | q_1 | q_3 |
| q_2 | q_3 | q_2 |
| q_3 | | |