# Design and Analysis of Algorithm

Lecture-3:

# **Contents**



1 Master Method

#### **Master Theorem**

The master method provides a "cookbook" method for solving recurrences of the form

T(n) = a T(n/b) + f(n), where,  $a \ge 1$  and b > 1 and f(n) should always be asymptotically positive function

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1. if f(n) = O(n^{\log_b a - \epsilon}) where \epsilon > 0

then T(n) = \theta(n^{\log_b a})

2. if f(n) = \theta(n^{\log_b a})

then T(n) = \theta(n^{\log_b a})

3. if f(n) = \Omega(n^{\log_b a + \epsilon})

if af(\frac{n}{b}) \leq cf(n) for some c < 1

then T(n) = \theta(f(n))
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# Numerical on Master Theorem

$$T(n) = 9T(n/3) + n$$

$$T(n) = T(2n/3) + 1$$

Step 1: 
$$a = 9$$
 and  $b = 3$  and  $f(n) = n$   
Step 2: compute  $\log_b a \Rightarrow \log_3 9$ 

Step 3: compute 
$$n^{\log_b a} \Rightarrow n^{\log_3 9} \Rightarrow n^2$$
  
Step 4: compare  $f(n)$  with  $n^{\log_b a} \Rightarrow n < n^2$   
Step 5: Choose the case: Case 1

Step 4. Compare 
$$f(n)$$
 with  $n \to n \to n$   
Step 5: Choose the case: Case 1  
Step 6:  $T(n) = \theta(n^2)$ 

$$\begin{array}{l}
^{9} \Rightarrow n^{2} \\
^{a} \Rightarrow n < n^{2}
\end{array}$$

$$n < n^2$$
 Ste

Step 5: Choose the case Step 6: 
$$T(n) = \theta(\log(n))$$

$$T(n) = 3T(n/4) + nlog(n)$$
  
Step 1:  $a = 3$  and  $b = 4$  and  $f(n) = n log(n)$ 

Step 2: compute 
$$\log_b a \Rightarrow \log_4 3$$
  
Step 3: compute  $n^{\log_b a} \Rightarrow n^{\log_4 3} \Rightarrow n^{0.793}$ 

Step 3: compute 
$$n^{\log_b a} \Rightarrow n^{\log_4 3} \Rightarrow n^{0.793}$$
  
Step 4: compare  $f(n)$  with  $n^{\log_b a} \Rightarrow n \log(n) > n^{.793}$   
Step 5: Choose the case: Case 3

Step 6:  $af\left(\frac{n}{h}\right) = 3\left(\frac{n}{4}\right)\lg\left(\frac{n}{4}\right) \le \frac{3}{4}nlog(n) = cf(n); c = 3/4$ 

Step 4: compare 
$$f(n)$$
 with  $n^{\log_b a} \Rightarrow 1 = 1$   
Step 5: Choose the case: Case 2

Step 3: compute 
$$n^{\log_b a} \Rightarrow n^{\log_{3/2} 1} \Rightarrow 1$$
  
Step 4: compare  $f(n)$  with  $n^{\log_b a} \Rightarrow 1 = 1$   
Step 5: Choose the case: Case 2

Step 1: a = 1 and b = 3/2 and f(n) = 1

 $T(n) = \theta(nlog(n))$ 

Step 2: compute  $\log_b a \Rightarrow \log_{3/2} 1$ 

$$b^{a} \Rightarrow 1 = 0$$
se 2

$$T(n) = 4T(n/2) + n^2$$

$$a = 4,$$
  $b = 2,$   $f(n) = n^2$ 

$$f(n) = n^2$$

$$\log_b a = \log_2 4 = 2$$

$$n^{\log_b a} = n^2$$

Compare f(n) with  $n^{\log_b a}$ 

$$f(n) = n^{\log_b a}$$

$$T(n) = \theta(n^2 \log n)$$

$$T(n) = 16T\left(\frac{n}{4}\right) + n!$$

$$a = 16, \qquad b = 4, \qquad f(n) = n!$$

$$\log_b a = \log_4 16 = 2$$

$$n^{\log_b a} = n^2$$
Comp

Compare 
$$f(n)$$
 with  $n^{\log_b a}$ 

$$f(n) > n^{\log_b a}$$
 Case 3  $a\left(f\left(\frac{n}{b}\right)\right) = cf(n) \Rightarrow 16\left(\frac{n}{4}\right)! \leq 0.5 \, n!$  this imples  $c = 0.5$  which is less than 1

$$T(n) = \theta(n!)$$

$$T(n) = 3T(n/2) + n^2$$

$$T(n) = 3T(\frac{n}{2}) + n^{2}$$

$$a = 3 \qquad b = 2 \qquad f(n) = n^{2}$$

$$\log_{b} a = \log_{2} 3$$

$$f(n) > n^{\log_{b} 9}$$

$$n^{\log_{b} 9} = n^{\log_{2} 3}$$

$$Case - 3$$

$$a f(\frac{h}{b}) = 3(\frac{h}{2})^{2} = \frac{3}{4}h^{2}$$

$$c f(h) = ch^{2}$$

$$\frac{3}{4}h^{2} = ch^{2}$$

$$c = \frac{3}{4}h^{3}$$

$$T(n) = \theta(f(n)) = \theta(n^2)$$

a) 
$$T(n) = 2^n T\left(\frac{n}{2}\right) + n^n$$

Master method does not apply as  $\alpha$  is not constant

b) 
$$T(n) = T(n/2) + n$$
  
 $T(n) = \theta(n)$ 

$$=\theta(n)$$

$$c) T(n) = 2T(n/2) + 2$$
$$T(n) = \theta(n)$$

$$d) T(n) = 4T(n/2) + n$$
$$T(n) = \theta(n^2)$$

e) 
$$T(n) = 0.5T(\frac{n}{2}) + 1/n$$

Master method does not apply as  $\alpha$  is less than

f) 
$$T(n) = 8T(n/2) + n^2$$

$$T(n) = \theta(n^3)$$

(a) 
$$T(n) = \sqrt{2} T\left(\frac{n}{2}\right) + \log n$$

(b) 
$$T(n) = 3T\left(\frac{n}{2}\right) + n$$

(c) 
$$T(n) = 3T\left(\frac{n}{3}\right) + \sqrt{n}$$

(d) 
$$T(n) = 7 T\left(\frac{n}{3}\right) + n^2$$

(e) 
$$T(n) = 2T(\sqrt{n}) + c$$

$$(f) \ T(n) = T(\sqrt{n}) + \log(n)$$

(g) 
$$T(n) = T(\sqrt{n}) + c$$