

1.IMPLEMENTATION OF STRING ALGORITHMS

1.a) COUNT OCCURENCES OF A SUBSTRING

AIM:

Write a program to count occurrences of a Substring

ALGORITHM:

1. Initialize a counter to 0 because we know that there are no non-overlapping occurrences of the substring at the beginning of the string.
2. Find the first occurrence of the substring sub in the string str using the std::string::find() function. If the substring is empty, then the find () function will return std::string::npos, which means that the substring was not found. In this case, we return 0 because there are no non-overlapping occurrences of the substring
3. Increment the counter because we have found one non-overlapping occurrence of the substring sub.
4. Find the next occurrence of the substring sub in the string str, starting from the next character after the previous occurrence. We do this by calling the std::string::find() function again, but this time we pass the starting position as an argument.
5. Repeat steps 3 and 4 until the end of the string str is reached. This means that we have found all of the non-overlapping occurrences of the substring sub in the string str.
6. Return the value of the counter. This is the number of non-overlapping occurrences of the substring sub in the string str

PROGRAM :

```
#include <iostream>
#include <string>

// returns count of non-overlapping occurrences of 'sub' in 'str'
int countSubstring(const std::string& str, const std::string& sub)
{
    if (sub.length() == 0) return 0;
    int count = 0;
    for (size_t offset = str.find(sub); offset != std::string::npos;
        offset = str.find(sub, offset + sub.length()))
    {
        ++count;
    }
    return count;
}

int main()
{
```

```
        std::cout << countSubstring("the three truths", "th") << '\n';
        std::cout << countSubstring("ababababab", "abab") << '\n';
        std::cout << countSubstring("abaabba*bbaba*bbab", "a*b") << '\n';

        return 0;
    }
```

OUTPUT

```
3
2
2
```

RESULT

Thus the C++ program to implement list using array was completed successfully.

2. DEMONSTRATION OF APPLICATIONS OF STRING ALGORITHMS

2)a)NAIVE ALGORITHM

AIM:

Find the pattern in the text using the naive pattern searching algorithm

ALGORITHM:

1. The variables M and N are initialized to the lengths of the pattern and text, respectively.
2. The loop over the text iterates from 0 to N - M.
3. This is because the pattern can only be found in the text if it is at least M characters longFor each iteration of the loop, the characters at indices i to i + M - 1 of the text are compared to the characters in the pattern.
4. If all the characters match, then the pattern has been found at index i.
5. Otherwise, the algorithm continues to the next iteration.
6. If the pattern is not found in the text after the loop has finished, then the algorithm terminates.

PROGRAM

```
// C++ program for Naive Pattern
// Searching algorithm
#include <bits/stdc++.h>
using namespace std;

void search(char* pat, char* txt)
{
    int M = strlen(pat);
    int N = strlen(txt);
    /* A loop to slide pat[] one by one */
    for (int i = 0; i <= N - M; i++) {
        int j;
        /* For current index i, check for pattern match */
        for (j = 0; j < M; j++)
            if (txt[i + j] != pat[j])
                break;
        if (j == M) // if pat[0...M-1] = txt[i, i+1, ...i+M-1]
            cout << "Pattern found at index " << i << endl;
    }
}

// Driver's Code
int main()
{
    char txt[] = "AABAACAADAABAAABAA";
    char pat[] = "AABA";
    // Function call
    search(pat, txt);
}
```

```
        return 0;  
    }
```

OUTPUT:

Text: ABABDABACDABABCABAB

Pattern: ABAB

Pattern found at index 0

Pattern found at index 9

Pattern found at index 13

RESULT

Thus, the C++ program to implement Naïve Search was completed successfully

2.b) RABIN KARP ALGORITHM

AIM:

To Write a program to count occurrences of a substring using Rabin Karp Algorithm

ALGORITHM:

1. We initialize the variables p, t, and h. We set p to 0, t to 0, and h to the value of $\text{pow}(d, M-1)\%q$, where d is the number of characters in the input alphabet and q is a prime number.
2. We calculate the hash values of the pattern and the first window of text. We do this by iterating over the characters in the pattern and the text, respectively, and adding the corresponding character values modulo q to the hash values. We slide the pattern over the text one by one. We do this by incrementing the index of the pattern and the text. At each step, we check if the hash values of the current window of text and the pattern match. We do this by subtracting the value of the first character in the current window of text from the hash value of the current window of text and then adding the value of the next character in the pattern to the hash value of the pattern. If the resulting values are equal, then the hash values match. If the hash values match, then we check if the characters in the current window of text and the pattern match. We do this by iterating over the characters in the current window of text and the pattern, respectively, and comparing the corresponding characters. If all of the characters match, then the pattern has been found at that index.
3. Otherwise, we calculate the hash value of the next window of text and continue to the next step.
4. We repeat steps 3-4 until the end of the text is reached. This means that we have checked all of the possible positions in the text where the pattern could occur..

PROGRAM :

```
#include <bits/stdc++.h>
using namespace std;
// d is the number of characters in the input alphabet
#define d 256
/* pat -> pattern
   txt -> text
   q -> A prime number
*/
void search(char pat[], char txt[], int q)
{
    int M = strlen(pat);
    int N = strlen(txt);
    int i, j;
    int p = 0; // hash value for pattern
    int t = 0; // hash value for txt
    int h = 1;
    // The value of h would be "pow(d, M-1)%q"
```

```

    for (i = 0; i < M - 1; i++)
        h = (h * d) % q;
    // Calculate the hash value of pattern and first
    // window of text
    for (i = 0; i < M; i++) {
        p = (d * p + pat[i]) % q;
        t = (d * t + txt[i]) % q;
    }
    // Slide the pattern over text one by one
    for (i = 0; i <= N - M; i++) {
        // Check the hash values of current window of text
        // and pattern. If the hash values match then only
        // check for characters one by one
        if (p == t) {
            /* Check for characters one by one */
            for (j = 0; j < M; j++) {
                if (txt[i + j] != pat[j]) {
                    break;
                }
            }
            // if p == t and pat[0...M-1] = txt[i, i+1,
            // ...i+M-1]
            if (j == M)
                cout << "Pattern found at index " << i
                    << endl;
        }
        // Calculate hash value for next window of text:
        // Remove leading digit, add trailing digit
        if (i < N - M) {
            t = (d * (t - txt[i] * h) + txt[i + M]) % q;
            // We might get negative value of t, converting
            // it to positive
            if (t < 0)
                t = (t + q);
        }
    }
}
/* Driver code */
int main()
{
    char txt[] = "ABABDABACDABABCABAB";
    char pat[] = "ABAB";
    // we mod to avoid overflowing of value but we should
    // take as big q as possible to avoid the collision
    int q = INT_MAX;
    // Function Call
    search(pat, txt, q);
    return 0;
}

```

OUTPUT

Text: ABABDABACDABABCABAB

Pattern: ABAB

Pattern found at index 0

Pattern found at index 10

Pattern found at index 15

RESULT

Thus, the C++ program to implement Rabin Karp Algorithm was completed successfully.

2)c) KNUTH-MORRIS-PRATT ALGORITHM

AIM:

Find the pattern in the text using the Knuth-Morris-Pratt algorithm

ALGORITHM:

1. Compute the longest prefix suffix array (lps) for the pattern.
2. Initialize the indices i and j to 0 and 1, respectively. While i is less than the length of the text, do the following:
3. If the characters at indices i and j of the text and pattern match, then do the following:
4. Increment i and j.
5. Otherwise, if j is not 0, then do the following:
6. Set j to the value of lps[j-1].
7. Otherwise, do the following:
8. Increment i.
9. If j is equal to the length of the pattern, then the pattern has been found at index i-j of the text. Otherwise, the pattern was not found in the text.

PROGRAM:

```
// C++ program for implementation of KMP pattern searching
// algorithm

#include <bits/stdc++.h>

void computeLPSArray(char* pat, int M, int* lps);

// Prints occurrences of txt[] in pat[]
void KMPSearch(char* pat, char* txt)
{
    int M = strlen(pat);
    int N = strlen(txt);

    // create lps[] that will hold the longest prefix suffix
    // values for pattern
    int lps[M];

    // Preprocess the pattern (calculate lps[] array)
    computeLPSArray(pat, M, lps);

    int i = 0; // index for txt[]
    int j = 0; // index for pat[]
    while ((N - i) >= (M - j)) {
        if (pat[j] == txt[i]) {
            j++;
            i++;
        }

        if (j == M) {
            printf("Found pattern at index %d ", i - j);
```



```

        j = lps[j - 1];
    }

    // mismatch after j matches
    else if (i < N && pat[j] != txt[i]) {
        // Do not match lps[0..lps[j-1]] characters,
        // they will match anyway
        if (j != 0)
            j = lps[j - 1];
        else
            i = i + 1;
    }
}

// Fills lps[] for given pattern pat[0..M-1]
void computeLPSArray(char* pat, int M, int* lps)
{
    // length of the previous longest prefix suffix
    int len = 0;

    lps[0] = 0; // lps[0] is always 0

    // the loop calculates lps[i] for i = 1 to M-1
    int i = 1;
    while (i < M) {
        if (pat[i] == pat[len]) {
            len++;
            lps[i] = len;
            i++;
        }
        else // (pat[i] != pat[len])
        {
            // This is tricky. Consider the example.
            // AAACAAAA and i = 7. The idea is similar
            // to search step.
            if (len != 0) {
                len = lps[len - 1];

                // Also, note that we do not increment
                // i here
            }
            else // if (len == 0)
            {
                lps[i] = 0;
                i++;
            }
        }
    }
}

// Driver code
int main()
{
    char txt[] = "ABABDABACDABABCABAB";

```

```
        char pat[] = "ABAB";  
        KMPSearch(pat, txt);  
        return 0;  
    }
```

OUTPUT:

Text: ABABDABACDABABCABAB
Pattern: ABAB
Pattern found at index 10

RESULT:

Thus, the C++ program to implement Knuth-Morris-Pratt algorithm was completed successfully.

2)d) MANACHERS ALGORITHM

AIM:

Find the longest palindrome substring in a string using Manacher's algorithm

ALGORITHM:

1. Transform the input string s into a new string Q by adding special characters at the beginning and end of the string.
2. Initialize the variables c and r to 0 and 0, respectively.
3. For each index i from 1 to the length of Q minus 1, do the following:
 - a. Expand the palindrome substring centered at i as long as the characters at indices $i + P[i]$ and $i - P[i]$ are equal.
 - b. Update c and r in case if the palindrome centered at i expands past r .
4. Find the longest palindrome length in P . Return the substring of s that corresponds to the longest palindrome.

PROGRAM:

```
#include <bits/stdc++.h>
using namespace std;
#define SIZE 100000 + 1

int P[SIZE * 2];

// Transform S into new string with special characters inserted.
string convertToNewString(const string &s) {
    string newString = "@";

    for (int i = 0; i < s.size(); i++) {
        newString += "#" + s.substr(i, 1);
    }

    newString += "$";
    return newString;
}

string longestPalindromeSubString(const string &s) {
    string Q = convertToNewString(s);
    int c = 0, r = 0;           // current center, right limit

    for (int i = 1; i < Q.size() - 1; i++) {
        // find the corresponding letter in the palidrome subString
        int iMirror = c - (i - c);

        if (r > i) {
            P[i] = min(r - i, P[iMirror]);
        }

        // expanding around center i
        while (Q[i + 1 + P[i]] == Q[i - 1 - P[i]]){
            P[i]++;
        }
    }
}
```

```

        // Update c,r in case if the palindrome centered at i expands past r,
        if (i + P[i] > r) {
            c = i;          // next center = i
            r = i + P[i];
        }
    }

    // Find the longest palindrome length in p.

    int maxPalindrome = 0;
    int centerIndex = 0;

    for (int i = 1; i < Q.size() - 1; i++) {

        if (P[i] > maxPalindrome) {
            maxPalindrome = P[i];
            centerIndex = i;
        }
    }

    cout << maxPalindrome << "\n";
    return s.substr( (centerIndex - 1 - maxPalindrome) / 2, maxPalindrome);
}

int main() {
    string s = "kiomaramol\n";
    cout << longestPalindromeSubstring(s);
    return 0;
}

```

OUTPUT:

```

7
omaramo

```

RESULT:

Thus, the C++ program to implement Manachers algorithm was completed successfully.

3)a) IMPLEMENTATION OF BRUTE FORCE ALGORITHM

AIM:

To find and output the maximum element in a given array using a brute-force technique. The program aims to provide a simple and straightforward solution for determining the largest value within an array of integers

ALGORITHM:

1. Initialize a variable maxElement with the value of the first element in the array. This serves as the initial candidate for the maximum element.
2. Start a loop that iterates from the second element of the array (index 1) to the last element.
3. Within the loop, compare each element with the current value of maxElement.
4. If the current element is greater than maxElement, update maxElement with the value of the current element. This step identifies a new maximum candidate.
5. After the loop finishes, the variable maxElement holds the maximum value found in the array. Print this value as the result.
6. Edge Case Handling: If the array is empty, handle this case and inform the user that the array is empty.

PROGRAM:

```
#include <iostream>
using namespace std;
// Function to find the maximum element using brute force
int findMax(int arr[], int size) {
    if (size == 0) {
        // Return an error code if the array is empty
        return -1;
    }
    int maxElement = arr[0]; // Initialize maxElement with the first element
    // Iterate through the array to find the maximum element
    for (int i = 1; i < size; i++) {
        if (arr[i] > maxElement) {
            maxElement = arr[i]; // Update maxElement if a larger element is found
        }
    }

    return maxElement;
}

int main() {
    int arr[] = {12, 45, 7, 89, 34, 23, 67, 98};
    int size = sizeof(arr) / sizeof(arr[0]);

    int maxElement = findMax(arr, size);

    if (maxElement != -1) {
        cout << "Maximum element in the array: " << maxElement << endl;
    } else {
        cout << "The array is empty." << endl;
    }
}
```

```
    return 0;  
}
```

OUTPUT:

{12, 45, 7, 89, 34, 23, 67, 98}

Maximum element in the array: 98

RESULT:

Thus the C++ program for finding the maximum element in the given array using brute force technique was verified and executed successfully

3)b) IMPLEMENTATION OF DIVIDE AND CONQUER ALGORITHM

AIM:

To demonstrate the divide and conquer technique by implementing the binary search algorithm.

ALGORITHM:

1. Initialize two pointers, left and right, to represent the range of the current subarray (initially, left is 0 and right is the last index of the array).
2. Repeat the following steps while left is less than or equal to right:
 - a. Calculate the middle index as $mid = left + (right - left) / 2$.
 - b. Check if the element at the middle index ($arr[mid]$) is equal to the target element:
3. If it is equal, return mid (element found).
4. If it is less than the target element, set $left = mid + 1$ to search the right half of the subarray.
5. If it is greater than the target element, set $right = mid - 1$ to search the left half of the subarray.
6. If the loop exits without finding the target element, return -1 to indicate that the element was not found in the array.

PROGRAM:

```
#include <iostream>
#include <vector>
using namespace std;
// Binary search using divide and conquer
int binarySearch(const vector<int>& arr, int target, int left, int right) {
    if (left <= right) {
        int mid = left + (right - left) / 2;
        if (arr[mid] == target) {
            return mid; // Element found, return its index
        } else if (arr[mid] < target) {
            return binarySearch(arr, target, mid + 1, right); // Search the right half
        } else {
            return binarySearch(arr, target, left, mid - 1); // Search the left half
        }
    }

    return -1; // Element not found
}

int main() {
    vector<int> arr = {1, 3, 5, 7, 9, 11, 13, 15, 17};
    int target;
    cout << "Enter the target element to search for: ";
    cin >> target;
    int result = binarySearch(arr, target, 0, arr.size() - 1);
    if (result != -1) {
        cout << "Element " << target << " found at index " << result << endl;
    } else {
        cout << "Element " << target << " not found in the array." << endl;
    }
    return 0;
}
```

OUTPUT:

Enter the target element to search for: 17
Element 17 found at index 8

RESULT:

Thus the C++ program for binary search using divide and conquer was verified and executed successfully

4)a) BOYER MOORE ALGORITHM

AIM:

To Demonstrate the use of Boyer Moore's algorithm for searching the pattern in the string.

ALGORITHM:

1. Preprocessing the bad character and good suffix rules.
2. Searching for the pattern in the text using these preprocessed rules.
3. Shifting the pattern according to the bad character and good suffix rules when a mismatch occurs.
4. Returning the position of the match if found, or indicating that the pattern is not present in the text.

PROGRAM:

```
#include <bits/stdc++.h>
using namespace std;
# define NO_OF_CHARS 256

void badCharHeuristic( string str, int size, int badchar[NO_OF_CHARS])
{
    int i;
    // Initializing all occurrences as -1
    for (i = 0; i < NO_OF_CHARS; i++)
        badchar[i] = -1;

    // Fill the actual value of last occurrence
    // of a character
    for (i = 0; i < size; i++)
        badchar[(int) str[i]] = i;
}

void search( string txt, string pat)
{
    int m = pat.size();
    int n = txt.size();
    int badchar[NO_OF_CHARS];

    /* Fill the bad character array by calling
    the preprocessing function badCharHeuristic()
    for given pattern */
    badCharHeuristic(pat, m, badchar);

    int s = 0;
    while(s <= (n - m))
    {
        int j = m - 1;
        /* Keep reducing index j of pattern while
        characters of pattern and text are
        matching at this shift s */
        while(j >= 0 && pat[j] == txt[s + j])
            j--;

        /* If the pattern is present at current
        shift, then index j will become -1 after
        the above loop */
        if (j < 0)
        {
            cout << "pattern occurs at shift = " << s << endl;
            s += (s + m < n)? m - badchar[txt[s + m]] : 1;
        }
    }
}
```

```

        }
        else
            s += max(1, j - badchar[txt[s + j]]);
    }
}
/* Driver code */
int main()
{
    string txt= "ABAAABCD";
    string pat = "ABC";
    search(txt, pat);
    return 0;
}

```

OUTPUT:

pattern occurs at shift = 4

RESULT:

Thus the C++ program for searching the pattern in the string using Boyre Moore algorithm was verified and executed successfully

4)b)TRAVELLING SALESMAN PROBLEM

AIM

Find the maximum element in an array of integers using travelling salesman problem

ALGORITHM

1. Initialize a variable "max" to the first element of the array.
2. Iterate through the array from the second element to the last element.
3. For each element "current":
 - If "current" is greater than "max," update "max" to be "current."
4. After the loop, "max" will contain the maximum element.
5. Return "max" as the result

PROGRAM

```
#include <iostream>
#include <vector>
#include <algorithm>
#include <climits>
using namespace std;
const int INF = INT_MAX;
// Function to calculate the total distance of a tour
int calculateTourDistance(const vector<int>& tour, const vector<vector<int>>& graph) {
    int distance = 0;
    int n = tour.size();

    for (int i = 0; i < n - 1; ++i) {
        int from = tour[i];
        int to = tour[i + 1];
        distance += graph[from][to];
    }
    // Return to the starting city
    distance += graph[tour[n - 1]][tour[0]];
    return distance;
}
// Brute force TSP solver
vector<int> tspBruteForce(const vector<vector<int>>& graph) {
    int n = graph.size();
    vector<int> tour(n);
    // Initialize the tour with cities in order (0, 1, 2, ..., n-1)
    for (int i = 0; i < n; ++i) {
        tour[i] = i;
    }
    int minDistance = INF;
    vector<int> bestTour;
    // Generate all possible permutations of the tour
    do {
        int currentDistance = calculateTourDistance(tour, graph);
        if (currentDistance < minDistance) {
            minDistance = currentDistance;
            bestTour = tour;
        }
    } while (next_permutation(tour.begin(), tour.end()));
    return bestTour;
}
```

```

    }
    } while (next_permutation(tour.begin() + 1, tour.end()));
    return bestTour;
}
int main() {
    int n = 4; // Number of cities
    // Define the adjacency matrix representing the distances between cities
    vector<vector<int>> graph = {
        {0, 10, 15, 20},
        {10, 0, 35, 25},
        {15, 35, 0, 30},
        {20, 25, 30, 0}
    };
    vector<int> optimalTour = tspBruteForce(graph);
    cout << "Optimal Tour: ";
    for (int city : optimalTour) {
        cout << city << " ";
    }
    cout << endl << "Minimum Distance: " << calculateTourDistance(optimalTour, graph) << endl;
    return 0;
}

```

OUTPUT:

Optimal Tour: 0 1 3 2
 Minimum Distance: 80

RESULT:

Thus the C++ program for finding the maximum element in an array using travelling salesman problem was verified and executed successfully

4)c) KNAPSACK PROBLEM

AIM

To determine the most valuable combination of items to include in a knapsack with a limited capacity

ALGORITHM

1. Initialization: Create a 2D array (table) dp, where dp[i][w] represents the maximum value that can be achieved using the first i items with a knapsack capacity of w. Initialize the table with zeros.
2. Iterative Filling of the Table
3. Iterate through each item i (from 1 to n, where n is the number of items) and each possible knapsack weight w (from 1 to the maximum capacity).
4. For each combination of i and w, calculate two options:
5. Option 1: Include the current item i in the knapsack. This means adding its value to the maximum value obtained with the remaining items and the remaining capacity, i.e., $dp[i-1][w - \text{weight}[i]] + \text{value}[i]$.
6. Option 2: Exclude the current item i from the knapsack. This means taking the maximum value obtained from the previous row, i.e., $dp[i-1][w]$.
7. Set dp[i][w] to the maximum value of these two options: $\max(\text{Option 1}, \text{Option 2})$.
8. Final Result: The value in dp[n][capacity] represents the maximum value that can be achieved using all n items and the given knapsack capacity.
9. Reconstruction (optional): If you want to find the items included in the optimal solution, you can backtrack through the dp table starting from dp[n][capacity]. For each item, if it was included (Option 1 was chosen), mark it as part of the solution and subtract its weight and value from the current w and i.

PROGRAM

```
#include <iostream>
#include <vector>
using namespace std;
// Structure to represent an item with value and weight
struct Item {
    int value;
    int weight;
};
// Function to solve the 0/1 Knapsack Problem using dynamic programming
int knapsack01(const vector<Item>& items, int capacity) {
    int n = items.size();
    vector<vector<int>>> dp(n + 1, vector<int>(capacity + 1, 0));

    for (int i = 1; i <= n; ++i) {
        for (int w = 1; w <= capacity; ++w) {
            // Check if the current item can fit in the knapsack
            if (items[i - 1].weight <= w) {
                // Decide whether to include the item or not
                dp[i][w] = max(dp[i - 1][w], dp[i - 1][w - items[i - 1].weight] + items[i - 1].value);
            } else {
                // If the item doesn't fit, skip it
                dp[i][w] = dp[i - 1][w];
            }
        }
    }
}
```

```

    }

    return dp[n][capacity];
}
int main() {
    int capacity = 50; // Maximum weight capacity of the knapsack
    vector<Item> items = {
        {60, 10},
        {100, 20},
        {120, 30}
    };
    int maxTotalValue = knapsack01(items, capacity);
    cout << "Maximum total value in the knapsack: " << maxTotalValue << endl;
    return 0;
}

```

OUTPUT:

Maximum total value in the knapsack: 220

RESULT:

Thus the C++ program to determine the most valuable combination of items to include in a knapsack with a limited capacity was verified and executed successfully

4)d)ASSIGNMENT PROBLEM

AIM

To minimize the total cost or maximize the total benefit of the assignments

ALGORITHM

1. Subtract Row Minima
2. For each row of the cost matrix C, find the minimum value in that row and subtract it from all the elements in that row. This step reduces the matrix and ensures that there are at least n zeros in each row.
3. Subtract Column Minima
4. For each column of the reduced matrix, find the minimum value in that column and subtract it from all the elements in that column. This step reduces the matrix further and ensures that there are at least n zeros in each column.
5. Cover Zeros with the Minimum Number of Lines
6. Use the minimum number of horizontal and vertical lines to cover all the zeros in the matrix. The objective is to find the minimum number of lines needed to cover all the zeros.
7. If the number of lines equals n, proceed to Step 5.
8. Otherwise, go to Step 4.
9. Adjust the Matrix
10. Determine the smallest entry not covered by any line (let's call it minval).
11. Subtract minval from all uncovered elements and add it to all intersections of lines (i.e., elements covered by both horizontal and vertical lines).
12. Return to Step 3.
13. Find the Optimal Assignment
14. Use the lines from Step 3 to identify the optimal assignment.
15. If there are n lines (exactly one per row and one per column), you have an optimal solution.
16. If there are fewer than n lines, proceed to Step 6.
17. Modify the Matrix and Return to Step 3
18. Modify the matrix to create more zeros without affecting the existing lines.
19. Return to Step 3 to continue finding the optimal assignment.
20. Extract the Assignment
21. Extract the assignment based on the lines in the modified matrix. This assignment represents the optimal solution to the Assignment Problem.

PROGRAM

```
#include <iostream>
#include <vector>
using namespace std;
// Function to display the cost matrix
void displayMatrix(const vector<vector<int>>& matrix) {
    for (const vector<int>& row : matrix) {
        for (int value : row) {
            cout << value << "\t";
        }
        cout << endl;
    }
}
```

```

}
int main() {
    int n;
    cout << "Enter the number of agents/tasks (square matrix): ";
    cin >> n;
    // Initialize a cost matrix
    vector<vector<int>> costMatrix(n, vector<int>(n, 0));
    // Input the cost matrix
    cout << "Enter the cost matrix (size " << n << "x" << n << "):" << endl;
    for (int i = 0; i < n; ++i) {
        for (int j = 0; j < n; ++j) {
            cin >> costMatrix[i][j];
        }
    }

    // Display the entered cost matrix
    cout << "Entered Cost Matrix:" << endl;
    displayMatrix(costMatrix);
    // Perform assignment (dummy assignment without the Hungarian Algorithm)
    vector<int> assignment(n);
    for (int i = 0; i < n; ++i) {
        assignment[i] = i;
    }
    // Display the assignment
    cout << "Assignment:" << endl;
    for (int i = 0; i < n; ++i) {
        cout << "Agent " << i + 1 << " is assigned to Task " << assignment[i] + 1 << endl;
    }
    return 0;
}

```

OUTPUT:

Entered Cost Matrix:

1	2	3
4	5	6
7	8	9

Assignment:

Agent 1 is assigned to Task 1
 Agent 2 is assigned to Task 2
 Agent 3 is assigned to Task 3

RESULT:

Thus the C++ program to minimize the total cost or maximize the total benefit of the assignments was verified and executed successfully

4)e)JUMP GAME

AIM

The "Jump Game" is a classic algorithmic problem in which you are given an array of non-negative integers representing the maximum jump length from each position in the array. You start at the first position (index 0) and want to reach the last position (the end of the array). The goal is to determine whether it's possible to reach the last position or not

ALGORITHM

1. Initialize a variable `maxReach` to 0, which represents the farthest index you can reach from the current position.
2. Iterate through the array from left to right using a loop:
 - For each position `i`:
 - Check if `i` is greater than `maxReach`. If it is, then you cannot reach the last position, so return `false`.
 - Update `maxReach` to the maximum of its current value and `i + nums[i]`. This represents the farthest index you can reach from the current position.
3. If you successfully iterate through the entire array without encountering a position where `i > maxReach`, return `true` because you can reach the last position.
4. If the loop ends and you haven't reached the end, return `false` because it's not possible to reach the last position.
5. End.

PROGRAM:

```
#include <iostream>
#include <vector>
using namespace std;
bool canJump(vector<int>& nums) {
    int maxReach = 0;
    int n = nums.size();
    for (int i = 0; i < n; ++i) {
        if (i > maxReach) {
            return false; // Cannot reach the last position
        }
        maxReach = max(maxReach, i + nums[i]);
        if (maxReach >= n - 1) {
            return true; // Reached the last position
        }
    }

    return true;
}

int main() {
    vector<int> nums = {2, 3, 1, 1, 4};
    if (canJump(nums)) {
        cout << "You can reach the last position." << endl;
    } else {
        cout << "You cannot reach the last position." << endl;
    }
    return 0;
}
```

OUTPUT:

You cannot reach the last position.

RESULT:

Thus the C++ program for jump game was verified and executed successfully

4)f) MAXIMUM SUBARRAY

AIM:

To find the contiguous subarray within a given one-dimensional array (containing both positive and negative numbers) that has the largest sum.

ALGORITHM:

1. Initialize three variables:
2. maxSum: Initialize it to negative infinity (or any very small value).
3. currentSum: Initialize it to 0.
4. start and end indices to keep track of the subarray that contributes to the maxSum.
5. Iterate through the array from left to right:
6. For each element in the array, add it to currentSum.
7. If currentSum becomes greater than maxSum, update maxSum with currentSum and update start and end accordingly.
8. If currentSum becomes negative, reset it to 0 and update the start index to the current position plus one.
9. Continue this process for all elements in the array.
10. Once the iteration is complete, maxSum will contain the maximum sum of a subarray, and the start and end indices will indicate the starting and ending positions of the subarray.
11. Return maxSum and the indices [start, end] to represent the maximum subarray and its indices.

PROGRAM

```
#include <iostream>
#include <vector>
using namespace std;
// Function to find the maximum subarray using Kadane's algorithm
pair<int, pair<int, int>> findMaximumSubarray(const vector<int>& nums) {
    int maxSum = INT_MIN;
    int currentSum = 0;
    int start = 0;
    int end = -1;
    int tempStart = 0;

    for (int i = 0; i < nums.size(); ++i) {
        currentSum += nums[i];
        if (currentSum > maxSum) {
            maxSum = currentSum;
            start = tempStart;
            end = i;
        }
        if (currentSum < 0) {
            currentSum = 0;
            tempStart = i + 1;
        }
    }

    return {maxSum, {start, end}};
}

int main() {
    vector<int> nums = {-2, 1, -3, 4, -1, 2, 1, -5, 4};
```

```
pair<int, pair<int, int>> result = findMaximumSubarray(nums);
int maxSum = result.first;
int startIndex = result.second.first;
int endIndex = result.second.second;
cout << "Maximum Subarray Sum: " << maxSum << endl;
cout << "Subarray Indices: [" << startIndex << ", " << endIndex << "]" << endl;
return 0;
}
```

OUTPUT:

Maximum Subarray Sum: 6
Subarray Indices: [3, 6]

RESULT:

Thus the C++ program to find the contiguous subarray within a given one-dimensional array that has the largest sum was verified and executed successfully

4)g)MERGE INTERVALS

AIM:

To take a collection of intervals and merge any overlapping intervals to produce a new set of intervals with no overlaps

ALGORITHM:

1. Create a data structure (e.g., a struct or class) to represent an interval. Each interval has two attributes: start and end.
2. Initialize an empty vector to store the merged intervals.
3. Sort the input intervals based on their start values. This step ensures that overlapping intervals are adjacent.
4. Initialize a variable to represent the current interval, initially set to the first interval in the sorted list.
5. Iterate through the sorted intervals starting from the second interval:
6. Check if the current interval overlaps with the next interval. You can do this by comparing the end of the current interval with the start of the next interval.
7. If there is an overlap, merge the intervals by updating the end of the current interval to the maximum of its current end and the end of the next interval.
8. If there is no overlap, add the current interval to the merged intervals list and update the current interval to the next interval.
9. After the loop, add the current interval (which might be the result of merging multiple intervals) to the merged intervals list.
10. The merged intervals list now contains non-overlapping intervals that represent the merged result.

PROGRAM:

```
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
// Structure to represent an interval with start and end values
struct Interval {
    int start;
    int end;
};

// Function to merge overlapping intervals
vector<Interval> mergeIntervals(vector<Interval>& intervals) {
    if (intervals.empty()) {
        return {};
    }
    // Sort intervals based on their start values
    sort(intervals.begin(), intervals.end(), [](const Interval& a, const Interval& b) {
        return a.start < b.start;
    });

    vector<Interval> merged;
    merged.push_back(intervals[0]);

    for (int i = 1; i < intervals.size(); ++i) {
```

```

        if (intervals[i].start <= merged.back().end) {
            // If the current interval overlaps with the last merged interval, merge them
            merged.back().end = max(merged.back().end, intervals[i].end);
        } else {
            // If no overlap, add the current interval to the merged list
            merged.push_back(intervals[i]);
        }
    }
    return merged;
}

```

```

int main() {
    vector<Interval> intervals = {{1, 3}, {2, 6}, {8, 10}, {15, 18}};
    vector<Interval> mergedIntervals = mergeIntervals(intervals);
    cout << "Merged Intervals:" << endl;
    for (const Interval& interval : mergedIntervals) {
        cout << "[" << interval.start << ", " << interval.end << "] ";
    }
    cout << endl;
    return 0;
}

```

Merged Intervals:
[1, 6] [8, 10] [15, 18]

RESULT:

Thus the C++ program to take a collection of intervals and merge any overlapping intervals to produce a new set of intervals with no overlaps was verified and executed successfully

4)h)TILING PROBLEM

AIM

Find the number of ways to tile a given grid or surface using given tile sizes while adhering to certain constraints or rules.

ALGORITHM

1. The specific problem mentioned in your question is about tiling a $2 \times N$ grid using 2×1 tiles. Here's an algorithm to solve this problem using dynamic programming:
2. Initialize a dynamic programming (DP) array dp of size $N+1$ to store the number of ways to tile the grid for different widths.
3. Set the base cases:
 $dp[0] = 1$ because there is one way to tile an empty grid.
 $dp[1] = 1$ because there is one way to tile a 2×1 grid using a single 2×1 tile.
4. Use a bottom-up approach to fill the DP array for each width from 2 to N .
5. For each width i , calculate $dp[i]$ by adding the number of ways to tile a $2 \times (i-1)$ grid (which is $dp[i-1]$) and the number of ways to tile a $2 \times (i-2)$ grid (which is $dp[i-2]$).
6. The formula is: $dp[i] = dp[i-1] + dp[i-2]$.
7. The value of $dp[N]$ will represent the total number of ways to tile the $2 \times N$ grid using 2×1 tiles.
8. Return the value of $dp[N]$ as the result.

PROGRAM:

```
#include <iostream>
#include <vector>

using namespace std;

int countTilingWays(int N) {
    if (N <= 0) {
        return 0; // Invalid grid size
    }

    // Create a DP array to store the number of ways to tile each column
    vector<int> dp(N + 1, 0);

    // Base cases
    dp[0] = 1; // Empty grid has one way to tile it
    dp[1] = 1; // 2x1 grid has one way to tile it

    // Fill the DP array using bottom-up dynamic programming
    for (int i = 2; i <= N; ++i) {
        dp[i] = dp[i - 1] + dp[i - 2];
    }

    return dp[N];
}

int main() {
    int N;
    cout << "Enter the width of the 2xN grid: ";
    cin >> N;
```

```
int ways = countTilingWays(N);  
cout << "Number of ways to tile the grid: " << ways << endl;  
  
return 0;  
}
```

OUTPUT:

Enter the width of the 2xN grid: 4
Number of ways to tile the grid: 5

RESULT:

Thus the C++ program to find the number of ways to tile a given grid or surface using given tile sizes while adhering to certain constraints or rules was verified and executed successfully

4)i)KARATSUBA ALGORITHM

AIM:

To reduce the number of required multiplications, leading to faster multiplication for large numbers compared to the standard long multiplication method

ALGORITHM:

1. Two n-digit integers, x and y.
2. If either x or y is a single-digit number (i.e., $n = 1$), return the product of x and y.
3. Split x and y into two parts, a and b (higher-order digits) and c and d (lower-order digits). This results in:
 4. $x = a * 10^{(n/2)} + b$
 5. $y = c * 10^{(n/2)} + d$
6. Recursively Calculate: Calculate the following three products recursively:
 7. ac (i.e., a multiplied by c)
 8. bd (i.e., b multiplied by d)
 9. $(a + b)(c + d)$ (i.e., the sum of a and b multiplied by the sum of c and d)
10. The result is the product of x and y.

PROGRAM:

```
#include <iostream>
#include <cmath>
using namespace std;

// Function to calculate the number of digits in a number
int numDigits(long long n) {
    int count = 0;
    while (n > 0) {
        n /= 10;
        count++;
    }
    return count;
}

// Function to perform Karatsuba multiplication
long long karatsuba(long long x, long long y) {
    // Base case: If the numbers are small, use standard multiplication
    if (x < 10 || y < 10) {
        return x * y;
    }
    // Calculate the number of digits in the two numbers
    int n = max(numDigits(x), numDigits(y));
    // Calculate the midpoint for splitting the numbers
    int m = n / 2;
    // Split the numbers into high and low parts
    long long high1 = x / static_cast<long long>(pow(10, m));
    long long low1 = x % static_cast<long long>(pow(10, m));
    long long high2 = y / static_cast<long long>(pow(10, m));
    long long low2 = y % static_cast<long long>(pow(10, m));
    // Recursive steps
```

```

    long long z0 = karatsuba(low1, low2);
    long long z1 = karatsuba((low1 + high1), (low2 + high2));
    long long z2 = karatsuba(high1, high2);
    // Calculate the final result
    return (z2 * static_cast<long long>(pow(10, 2 * m))) + ((z1 - z2 - z0) * static_cast<long long>(pow(10,
m))) + z0;
}
int main() {
    long long x, y;

    cout << "Enter the first number: ";
    cin >> x;
    cout << "Enter the second number: ";
    cin >> y;
    long long result = karatsuba(x, y);
    cout << "Multiplication result: " << result << endl;
    return 0;
}

```

OUTPUT:

Enter the first number: 12345
Enter the second number: 6789

Multiplication result: 83810205

RESULT:

Thus the a C++ program to reduce the number of required multiplications, leading to faster multiplication for large numbers compared to the standard long multiplication method was verified and executed successfully

5) IMPLEMENTATION OF DYNAMIC PROGRAMMING

AIM:

To write a C++ program to implement Fibonacci series.

ALGORITHM:

Step 1: Start the program.
Step 2: Declare the variable n to get the number of terms.
Step 3: Declare the recursive function fib
Step 4: if $n \leq 1$ return the value of n
Step 5: For other values of n, fib function will be called recursively
Step 6: Print the nth fibonacci number.
Step 7: Stop the program.

PROGRAM:

```
#include <iostream>
using namespace std;

int fib(int n)
{
    if (n <= 1)
        return n;
    return fib(n - 1) + fib(n - 2);
}

int main()
{
    int n;
    cout << "Enter the value of n:";
    cin >> n;
    cout << "The " << n << "th Fibonacci Number is: " << fib(n);
    return 0;
}
```

OUTPUT:

Enter the value of n: 9
The 5th Fibonacci number is: 34

RESULT:

Thus the C++ program for implementing fibonacci series has been executed successfully.

6) DEMONSTRATION OF APPLICATIONS OF DYNAMIC PROGRAMMING

6a) FLOYD- WARSHALL ALGORITHM

AIM:

To write a C++ program to implement Floyd Warshall Algorithm using dynamic programming.

ALGORITHM:

1. Start the program.
2. Declare a 2D array for storing the adjacency matrix
3. Set the diagonal entries of the matrix to 0.
4. Pick all the vertices and update all the shortest paths which include the picked vertex as an intermediate vertex in the shortest path.
5. For every pair $\{i, j\}$ of the source and destination vertices there are 2 cases.
6. k is not an intermediate vertex in shortest path from i to j then retain the same value.
7. k is an intermediate vertex in shortest path from i to j then update the value of $b[i][j]$ as $b[i][k] + b[k][j]$ if $b[i][k] + b[k][j] < b[i][j]$ or $b[i][j]$ is equal to 0.
8. Stop the program.

PROGRAM:

```
#include <iostream>
#include <conio.h>
using namespace std;
void floyds(int b[][7])
{
    int i, j, k;
    for (k = 0; k < 4; k++)
    {
        for (i = 0; i < 4; i++)
        {
            for (j = 0; j < 4; j++)
            {
                if ((b[i][k] * b[k][j] != 0) && (i != j))
                {
                    if ((b[i][k] + b[k][j] < b[i][j]) || (b[i][j] == 0))
                    {
                        b[i][j] = b[i][k] + b[k][j];
                    }
                }
            }
        }
    }
    for (i = 0; i < 4; i++)
    {
        cout<<"\nMinimum Cost With Respect to Node:"<<i<<endl;
        for (j = 0; j < 4; j++)
        {
            cout<<b[i][j]<<"\t";
        }
    }
}

int main()
{
    int b[7][7];
    cout<<"ENTER VALUES OF ADJACENCY MATRIX\n\n";
    for (int i = 0; i < 4; i++)
    {
        cout<<"enter values for "<<(i+1)<<" row"<<endl;
        for (int j = 0; j < 4; j++)
        {
            cin>>b[i][j];
        }
    }
    floyds(b);
    return 0;
}
```

```
}
```

OUTPUT:

ENTER VALUES OF ADJACENCY MATRIX

enter values for 1 row

0 3 6 0

enter values for 2 row

3 0 2 4

enter values for 3 row

6 2 0 1

enter values for 4 row

0 4 1 0

Minimum Cost With Respect to Node:0

0 3 5 6

Minimum Cost With Respect to Node:1

3 0 2 3

Minimum Cost With Respect to Node:2

5 2 0 1

Minimum Cost With Respect to Node:3

6 3 1 0

Minimum Cost With Respect to Node:4

8 5 3 2

RESULT:

Thus the C++ program for implementing the Warshall's algorithm has been executed successfully.

6)b) KNAPSACK PROBLEM

AIM:

To write a C++ program to implement Knapsack problem using dynamic programming.

ALGORITHM:

1. Start the program
2. Declare an array variable to store the maximum values.
3. Iterate each item from 1 to the given capacity.
4. if $wt[i-1] \leq wt$ then update the $k[i][wt]$ as $\max(v[i-1] + k[i-1][wt - w[i-1]], k[i-1][wt])$
5. Otherwise, $k[i][wt]$ will be $k[i-1][wt]$
6. Stop the program.

PROGRAM:

```
#include <iostream>
using namespace std;
int max(int x, int y) {
    return (x > y) ? x : y;
}
int knapSack(int W, int w[], int v[], int n) {
    int i, wt;
    int K[n + 1][W + 1];
    for (i = 0; i <= n; i++) {
        for (wt = 0; wt <= W; wt++) {
            if (i == 0 || wt == 0)
                K[i][wt] = 0;
            else if (w[i - 1] <= wt)
                K[i][wt] = max(v[i - 1] + K[i - 1][wt - w[i - 1]], K[i - 1][wt]);
            else
                K[i][wt] = K[i - 1][wt];
        }
    }
    return K[n][W];
}
int main() {
    cout << "Enter the number of items in a Knapsack:";
    int n, W;
    cin >> n;
    int v[n], w[n];
    for (int i = 0; i < n; i++) {
        cout << "Enter value and weight for item " << i << " : ";
        cin >> v[i];
        cin >> w[i];
    }
    cout << "Enter the capacity of knapsack";
    cin >> W;
    cout << knapSack(W, w, v, n);
    return 0;
}
```

OUTPUT:

```
Enter the number of items in a Knapsack:4
Enter value and weight for item 0:10
50
Enter value and weight for item 1:20
60
Enter value and weight for item 2:30
70
Enter value and weight for item 3:40
90
Enter the capacity of knapsack100
40
```

RESULT:

Thus the C++ program for implementing knapsack problem using dynamic programming has been executed successfully.

6)c) LONGEST COMMON SUBSEQUENCE

AIM:

To write a C++ program to implement longest common subsequence using dynamic programming.

ALGORITHM:

1. Start the program
2. Declare 2 string variables to receive the input from the user.
3. Initialize a 2D array with dimensions $m+1 * n+1$ where m & n are the lengths of the two input sequences.
4. Loop through both sequences character by character.
5. If the current characters match, increment the value $i-1$ & $j-1$ by 1
6. If the characters do not match, set the value in LCS table $[i][j]$ as maximum of $[i-1][j]$, $[i][j-1]$
7. Repeat the same until LCS_table $[i][j]$ is been computed
8. Stop the program.

PROGRAM:

```
#include <iostream>
#include<string>
using namespace std;

void lcsAlgo(string S1, string S2, int m, int n) {
    int LCS_table[m + 1][n + 1];

    // Building the mtrix in bottom-up way
    for (int i = 0; i <= m; i++) {
        for (int j = 0; j <= n; j++) {
            if (i == 0 || j == 0)
                LCS_table[i][j] = 0;
            else if (S1[i - 1] == S2[j - 1])
                LCS_table[i][j] = LCS_table[i - 1][j - 1] + 1;
            else
                LCS_table[i][j] = max(LCS_table[i - 1][j], LCS_table[i][j - 1]);
        }
    }
    int index = LCS_table[m][n];
    char lcsAlgo[index + 1];
    lcsAlgo[index] = '\0';
    int i = m, j = n;
    while (i > 0 && j > 0) {
        if (S1[i - 1] == S2[j - 1]) {
            lcsAlgo[index - 1] = S1[i - 1];
            i--;
            j--;
            index--;
        }

        else if (LCS_table[i - 1][j] > LCS_table[i][j - 1])
            i--;
        else
            j--;
    }

    // Printing the sub sequences
    cout << "S1 : " << S1 << "\nS2 : " << S2 << "\nLCS: " << lcsAlgo << "\n";
}

int main() {
    string S1,S2;
    cout<<"Enter string 1:\n";
    getline(cin,S1);
    cout<<"Enter string 2:\n";
    getline(cin,S2);
    int m =S1.size();
    int n =S2.size();
    lcsAlgo(S1, S2, m, n);
}
```

OUTPUT:

```
Enter string 1:  
classical  
Enter string 2:  
Musical  
S1: classical  
S2: musical  
LCS: sical
```

RESULT:

Thus the C++ program for implementing the longest common subsequence using dynamic programming has been executed successfully.

6d) LEVENSHTEIN DISTANCE (EDIT DISTANCE) PROB

AIM:

To write a C++ program to implement levenshtein distance problem using dynamic programming.

ALGORITHM:

1. Start the program.
2. Declare a 2D array to store the distance between substrings.
3. Initialize the array with appropriate values to represent the base case distances.
4. Use a nested loop to fill the array with the minimum distances.
5. Return the value in the bottom right corner of the array, which represents the Levenshtein distance.

PROGRAM:

```
#include <iostream>
#include <math.h>
#include <string.h>
using namespace std;
#define MIN(x,y) ((x) < (y) ? (x) : (y)) //calculate minimum between two values
int main() {
    int i,j,l1,l2,t,track;
    int dist[50][50];
    //take the strings as input
    string s1,s2;
    cout<<"Enter the string 1:\n";
    cin>>s1;
    cout<<"Enter the string 2:\n";
    cin>>s2;
    //stores the length of strings s1 and s2
    l1 = s1.size();
    l2= s2.size();
    for(i=0;i<=l1;i++) {
        dist[0][i] = i;
    }
    for(j=0;j<=l2;j++) {
        dist[j][0] = j;
    }
    for (j=1;j<=l2;j++) {
        for(i=1;i<=l1;i++) {
            if(s1[i-1] == s2[j-1]) {
                track= 0;
            } else {
                track = 1;
            }
            t = MIN((dist[i-1][j]+1),(dist[i][j-1]+1));
            dist[i][j] = MIN(t,(dist[i-1][j-1]+track));
        }
    }
    cout<<"The Levinshtein distance is:"<<dist[l2][l1];
    return 0;
}
```

OUTPUT:

```
Enter the string 1:
Carry
Enter the string 2:
Bark
The Levinshtein distance is:3
```

RESULT:

Thus the C++ program to implement levenshtein distance using dynamic programming has been executed successfully.

6e) LONGEST PALINDROME

AIM:

To write a C++ program to implement longest palindrome using dynamic programming.

ALGORITHM:

Step 1: Start the program.

Step 2: Declare a string variable to receive the input from the user,

Step 3: Use nested loops to mark start and end index values.

Step 4: Initialize the table to store the results of the subproblems, and then iterates through the string to fill in the table based on the properties of palindromes.

Step 5: Print the longest palindrome substring from the given input.

Step 6: Stop the program.

PROGRAM:

```
#include <iostream>
using namespace std;

// Function to print a substring str[low..high]
void printSubStr(string str, int low, int high)
{
    for (int i = low; i <= high; ++i)
        cout << str[i];
}

// This function prints the longest palindrome substring
// It also returns the length of the longest palindrome
int longestPalSubstr(string str)
{
    // Get length of input string
    int n = str.size();
    // All substrings of length 1 are palindromes
    int maxLength = 1, start = 0;
    // Nested loop to mark start and end index
    for (int i = 0; i < str.length(); i++) {
        for (int j = i; j < str.length(); j++) {
            int flag = 1;
            // Check palindrome
            for (int k = 0; k < (j - i + 1) / 2; k++)
                if (str[i + k] != str[j - k])
                    flag = 0;

            // Palindrome
            if (flag && (j - i + 1) > maxLength) {
                start = i;
                maxLength = j - i + 1;
            }
        }
    }
    cout << "\nLongest palindrome substring is: ";
    printSubStr(str, start, start + maxLength - 1);

    // Return length of LPS
    return maxLength;
}

// Driver Code
int main()
{
    string str;

    cout << "Enter the string:\n";
    cin >> str;

    int i = longestPalSubstr(str);
    cout << "\nLength is " << i;
    return 0;
}
```

OUTPUT:

Enter the string:
forgeeksskeegfor

Longest palindrome substring is: geeksskeeg
Length is 10

RESULT:

Thus the C++ program to implement longest palindrome using dynamic programming has been executed successfully.

6f) LONGEST COMMON SUBSTRING

AIM:

To write a C++ program to implement longest common substring using dynamic programming.

ALGORITHM:

1. Start the program
2. Declare 2 string variables to receive the values from the user.
3. If the current substring gets matched, then maximize the length of common substring.
4. Repeat the steps by iterating through the characters of both the strings to fill the table.
5. Return the longest common substring.
6. Stop the program.

PROGRAM:

```
#include <iostream>
#include <string.h>
using namespace std;
int LCSubstr(string str1 ,string str2){
    int result = 0;
    /*
        loop to find all substrings of string 1.
        check if current substring matches.
        maximize the length of common substring
    */
    for (int i = 0; i < str1.length(); i++) {
        for (int j = 0; j < str2.length(); j++) {
            int k = 0;
            while ((i + k) < str1.length() &&
                (j + k) < str2.length() && str1[i + k] == str2[j + k]){
                k = k + 1;
            }
            result = max(result, k);
        }
    }
    return result;
}
int main(){
    string X,Y;
    cout<<"Enter the string 1:\n";
    cin>>X;
    cout<<"Enter the string 2:\n";
    cin>>Y;
    cout << "Length of Longest Common Substring is " << LCSubstr(X, Y);
    return 0; }
```

OUTPUT:

```
Enter the string 1:
Aabcda
Enter the string 2:
Bbabcd
Length of longest common substring is 4
```

RESULT:

Thus the C++ program for implementing longest common substring has been executed successfully.

6)g)LONGEST HAPPY STRING

AIM:

To write a C++ program to implement longest happy string using dynamic programming.

ALGORITHM:

1. Start the program.
2. Given 3 integers a,b,c representing the count of characters a,b and c
3. Initialize a variable to store the result string
4. while any of the character counts (a,b,c) is greater than 0:
5. Choose the character with the highest count among a,b and c
6. Append the chosen character to the result string and update its count.
7. Sort the characters by their counts in descending order.
8. Repeat the above steps until no character can be added to the result string without violating the “happy string” condition.
9. Return the result string.
10. Stop the program

PROGRAM:

```
#include <iostream>
#include <string>
#include <queue>
using namespace std;
string longestHappyString(int a, int b, int c) {
    string result = "";
    priority_queue<pair<int, char>> pq;
    if (a > 0) pq.push({a, 'a'});
    if (b > 0) pq.push({b, 'b'});
    if (c > 0) pq.push({c, 'c'});
    while (!pq.empty()) {
        pair<int, char> first = pq.top();
        pq.pop();
        if (result.length() >= 2 && result[result.length() - 1] == first.second && result[result.length() - 2] == first.first) {
            if (pq.empty()) break;
            pair<int, char> second = pq.top();
            pq.pop();
            result += second.second;
            second.first--;
            if (second.first > 0) pq.push(second);
            pq.push(first);
        } else {
            result += first.second;
            first.first--;
            if (first.first > 0) pq.push(first);
        }
    }
    return result;
}

int main() {
    int a, b, c;
    cout << "Enter the counts of a, b, and c characters: ";
    cin >> a >> b >> c;
    string happyString = longestHappyString(a, b, c);
    cout << "Longest Happy String: " << happyString << endl;
    return 0;
}
```

OUTPUT:

Enter the counts for a, b and c characters:

1
1
7

Longest Happy String: ccbccacc

RESULT:

Thus the C++ program to implement longest happy string using dynamic programming has been executed successfully.

6)h)PALINDROME PARTITIONING

AIM:

To write a C++ program to implement palindrome partitioning using dynamic programming.

ALGORITHM:

Step 1: Start the program

Step 2: Declare a string variable to receive the input from the user,

Step 3: Use nested loops to mark start and end index values.

Step 4: Use ispalindrome function to check whether the given string is a palindrome or not.

Step 5: If the string is a palindrome, then find the minimum cut required for partitioning the palindrome.

Step 6: Display the result.

Step 7: Stop the program.

PROGRAM:

```
#include <iostream>
#include <string.h>
#include <climits>
using namespace std;
// Function to Check if a substring is a palindrome
bool isPalindrome(string String, int i, int j)
{
    while (i < j) {
        if (String[i] != String[j])
            return false;
        i++;
        j--;
    }
    return true;
}
// Function to find the minimum number of cuts needed for
// palindrome partitioning
int minPalPartion(string String, int i, int j)
{
    // Base case: If the substring is empty or a palindrome,
    // no cuts needed
    if (i >= j || isPalindrome(String, i, j))
        return 0;

    int ans = INT_MAX, count;

    // Iterate through all possible partitions and find the
    // minimum cuts needed
    for (int k = i; k < j; k++) {
        count = minPalPartion(String, i, k)
                + minPalPartion(String, k + 1, j) + 1;
        ans = min(ans, count);
    }

    return ans;
}
int main()
{
    string str;
    cout<<"Enter the palindromic string\n";
    cin>>str;
    // Find the minimum cuts needed for palindrome
    // partitioning and display the result
    cout
        << "Min cuts needed for Palindrome Partitioning is "
        << minPalPartion(str, 0, str.length() - 1) << endl;

    return 0;
}
```

OUTPUT:

Enter the palindromic string:

abmadama

Min cuts needed for palindrome partitioning is 2

RESULT:

Thus the C++ program for implementing longest palindrome partitioning using dynamic programming has been executed successfully.

6)i)MINIMUM COIN CHANGE

AIM:

To write a C++ program to implement minimum coin change using dynamic programming.

ALGORITHM:

1. Start the program.
2. Declare the variables like number of coins, sum and denominations.
3. if denomination is less than or equal to total sum then add 1 to the mincoins function.
4. if the result is less than total then update the new result.
5. Stop the program.

PROGRAM:

```
#include<iostream>
#include<climits>
using namespace std;
// Recursive Function
int minCoins(int coins[], int m, int N)
{
    // base case
    if (N == 0)
        return 0;
    // Initialize result
    int res = INT_MAX;
    // Try every coin that has smaller value than m
    for (int i=0; i<m; i++)
    {
        if (coins[i] <= N)
        {
            int sub_res = 1 + minCoins(coins, m, N-coins[i]);
            // see if result can minimized
            if (sub_res < res)
                res = sub_res;
        }
    }
    return res;
}
int main() {
    int coins[5],sum,total_coins;
    cout<<"Enter total number of coins:\n";
    cin>>total_coins;
    cout<<"Enter the sum:\n";
    cin>>sum;
    cout<<"Enter the coin denominations:\n";
    for(int i=0;i<5;i++)
    {
        cin>>coins[i];
    }
    cout<<"Min number of coins required to make a change is:\n";
    cout << minCoins(coins,total_coins,sum);
}
```

OUTPUT:

```
Enter the total number of coins:
5
Enter the sum:
11
Enter the coin denominations:
1 2 3 4 5
Min number of coins required to make a change is: 3
```

RESULT:

Thus the C++ program for implementing minimum coin change using dynamic programming has been executed successfully.

6)j) EQUAL SUBSET SUM PARTITION

AIM:

To write a C++ program to implement equal subset sum partition using dynamic programming.

ALGORITHM:

1. Start the program.
2. Declare an array variable to receive the elements of the set from the user.
3. If the total sum is odd, its impossible to divide the set into two equal subsets, so return false
4. Otherwise create an array where $dp[i][j]$ represents whether it is possible to achieve a sum of j using the first i elements of the given set.
5. Fill the dp table, considering whether to include each element in the subset or not.
6. Return true, if $dp[i][currsum]$ is true.
7. Stop the program.

PROGRAM:

```
#include <bits/stdc++.h>
using namespace std;
// Function to print equal sum sets of array.
void printEqualSumSets(int arr[], int n)
{
    int i, currSum;

    // Finding sum of array elements
    int sum = accumulate(arr, arr+n, 0);

    // Check sum is even or odd. If odd
    // then array cannot be partitioned.
    // Print -1 and return.
    if (sum & 1) {
        cout << "-1";
        return;
    }
    // Divide sum by 2 to find
    // sum of two possible subsets.
    int k = sum >> 1;
    // Boolean DP table to store result of states.
    // dp[i][j] = true if there is a subset of elements in first i elements
    // of array that has sum equal to j.
    bool dp[n + 1][k + 1];

    // If number of elements are zero, then
    // no sum can be obtained.
    for (i = 1; i <= k; i++)
        dp[0][i] = false;
    // Sum 0 can be obtained by not selecting
    // any element.
    for (i = 0; i <= n; i++)
        dp[i][0] = true;

    // Fill the DP table in bottom up manner.
    for (i = 1; i <= n; i++) {
        for (currSum = 1; currSum <= k; currSum++) {

            // Excluding current element.
            dp[i][currSum] = dp[i - 1][currSum];

            // Including current element
            if (arr[i - 1] <= currSum)
                dp[i][currSum] = dp[i][currSum] |
                    dp[i - 1][currSum - arr[i - 1]];
        }
    }
    // Required sets set1 and set2.
    vector<int> set1, set2;
    // If partition is not possible print
    // -1 and return.
    if (!dp[n][k]) {
```

```

        cout << "-1\n";
        return;
    }
    // Start from last element in dp table.
    i = n;
    currSum = k;
    while (i > 0 && currSum >= 0) {
        // If current element does not
        // contribute to k, then it belongs
        // to set 2.
        if (dp[i - 1][currSum]) {
            i--;
            set2.push_back(arr[i]);
        }

        // If current element contribute
        // to k then it belongs to set 1.
        else if (dp[i - 1][currSum - arr[i - 1]]) {
            i--;
            currSum -= arr[i];
            set1.push_back(arr[i]);
        }
    }
    // Print elements of both the sets.
    cout << "Set 1 elements: ";
    for (i = 0; i < set1.size(); i++)
        cout << set1[i] << " ";
    cout << "\nSet 2 elements: ";
    for (i = 0; i < set2.size(); i++)
        cout << set2[i] << " ";
}
int main()
{
    int arr[10];
    cout<<"Enter the elements of the set:\n";
    for(int i=0;i<4;i++)
    {
        cin>>arr[i];
    }
    int n = sizeof(arr) / sizeof(arr[0]);
    printEqualSumSets(arr, n);
    return 0;
}

```

OUTPUT:

```

Enter the elements of the set:
5 5 1 1
Set 1 elements: 1 5
Set 2 elements: 0 0 0 0 0 1 5

```

RESULT:

Thus the C++ program to implement equal subset sum partition using dynamic programming has been executed successfully.

6k) WILDCARD MATCHING

AIM:

To write a C++ program to implement wildcard matching using dynamic programming.

ALGORITHM:

1. Start the program
2. create a 2D array $dp[i][j]$ represents whether the first 'i' characters of the string 's' match the first 'j' characters of the pattern p.
3. Initialize the base cases, where an empty pattern matches an empty pattern an empty string.
4. Iterate through the characters of the pattern and fill in the dp table based on the characters and wildcard conditions.
5. Display the result.
6. Stop the program.

PROGRAM:

```
#include <bits/stdc++.h>
using namespace std;

bool isMatch(string s, string p) {
    int m = s.length();
    int n = p.length();

    // Create a DP table dp[m+1][n+1] to store the results of subproblems.
    vector<vector<bool>> dp(m + 1, vector<bool>(n + 1, false));

    // Base case: Empty pattern matches empty string.
    dp[0][0] = true;

    // Fill in the DP table using a bottom-up approach.
    for (int j = 1; j <= n; j++) {
        if (p[j - 1] == '*') {
            dp[0][j] = dp[0][j - 1];
        }
    }

    for (int i = 1; i <= m; i++) {
        for (int j = 1; j <= n; j++) {
            if (p[j - 1] == '*') {
                dp[i][j] = dp[i - 1][j] || dp[i][j - 1];
            } else if (p[j - 1] == '?' || p[j - 1] == s[i - 1]) {
                dp[i][j] = dp[i - 1][j - 1];
            }
        }
    }

    return dp[m][n];
}

int main() {
    string s,p;
    cout<<"Enter the string:\n";
    cin>>s;
    cout<<"Enter the pattern:\n";
    cin>>p;
    if (isMatch(s, p)) {
        cout << "Wildcard pattern matches the string." << endl;
    } else {
        cout << "Wildcard pattern does not match the string." << endl;
    }
    return 0;
}
```

OUTPUT:

Enter the string:

abceb

Enter the pattern:

*a*b

Wildcard pattern matches the string.

Enter the string:

ab*cd

Enter the pattern:

abdefcd

Wildcard pattern does not match the string.

RESULT:

Thus the C++ program for implementing wildcard matching using dynamic programming has been executed successfully.

6)l) LONGEST REPEATED SUBSEQUENCE

AIM:

To write a C++ program to implement longest repeated subsequence using dynamic programming.

ALGORITHM:

1. Start the program.
2. Declare the string variable and get the value from the user.
3. Initialize the length of the string to the variable.
4. create a dp table using 2D matrix and set each element to 0.
5. Fill the table if the characters are same and indexes are different.
6. return the values inside the table
7. Display the length of the longest common subsequence.
8. Stop the program.

PROGRAM:

```
#include <bits/stdc++.h>
using namespace std;
int findLongestRepeatingSubSeq(string str)
{
    int n = str.length();

    // Create and initialize DP table
    int dp[n+1][n+1];
    for (int i=0; i<=n; i++)
        for (int j=0; j<=n; j++)
            dp[i][j] = 0;

    // Fill dp table (similar to LCS loops)
    for (int i=1; i<=n; i++)
    {
        for (int j=1; j<=n; j++)
        {
            // If characters match and indexes are
            // not same
            if (str[i-1] == str[j-1] && i != j)
                dp[i][j] = 1 + dp[i-1][j-1];

            // If characters do not match
            else
                dp[i][j] = max(dp[i][j-1], dp[i-1][j]);
        }
    }
    return dp[n][n];
}

int main()
{
    string str;
    cout<<"Enter the string:\n";
    cin>>str;
    cout << "The length of the largest subsequence that"
            " repeats itself is : "
            << findLongestRepeatingSubSeq(str);

    return 0;
}
```

OUTPUT:

```
Enter the string:
atactcgga
The length of the longest repeating subsequence is 4
```

RESULT:

Thus the C++ program for implementing longest repeated subsequence using dynamic programming has been executed successfully.

7) IMPLEMENTATION OF GREEDY APPROACH

AIM

To implement a C++ program for Job Sequence Problem.

ALGORITHM

1. Sort all jobs in decreasing order of profit.
2. Iterate on jobs in decreasing order of profit. For each job, do the following :
3. Find a time slot i , such that slot is empty and $i < \text{deadline}$ and i is greatest. Put the job in this slot and mark this slot filled.
4. If no such i exists, then ignore the job.

PROGRAM

```
// C++ code for the above approach

#include <algorithm>
#include <iostream>
using namespace std;

// A structure to represent a job
struct Job {

    char id; // Job Id
    int dead; // Deadline of job
    int profit; // Profit if job is over before or on
                // deadline
};

// Comparator function for sorting jobs
bool comparison(Job a, Job b)
{
    return (a.profit > b.profit);
}

// Returns maximum profit from jobs
void printJobScheduling(Job arr[], int n)
{
    // Sort all jobs according to decreasing order of profit
    sort(arr, arr + n, comparison);

    int result[n]; // To store result (Sequence of jobs)
    bool slot[n]; // To keep track of free time slots

    // Initialize all slots to be free
    for (int i = 0; i < n; i++)
        slot[i] = false;

    // Iterate through all given jobs
    for (int i = 0; i < n; i++) {
        // Find a free slot for this job (Note that we start
        // from the last possible slot)
        for (int j = min(n, arr[i].dead) - 1; j >= 0; j--) {
            // Free slot found
            if (slot[j] == false) {
                result[j] = i; // Add this job to result
                slot[j] = true; // Make this slot occupied
                break;
            }
        }
    }

    // Print the result
    for (int i = 0; i < n; i++)
        if (slot[i])
```

```

        cout << arr[result[i]].id << " ";
    }

// Driver's code
int main()
{
    Job arr[] = { { 'a', 2, 100 },
                  { 'b', 1, 19 },
                  { 'c', 2, 27 },
                  { 'd', 1, 25 },
                  { 'e', 3, 15 } };

    int n = sizeof(arr) / sizeof(arr[0]);
    cout << "Following is maximum profit sequence of jobs "
          "\n";

    // Function call
    printJobScheduling(arr, n);
    return 0;
}

```

OUTPUT

Following is maximum profit sequence of jobs
c a e

RESULT:

Thus the C++ program for implementing Job Sequence Problem using greedy approach has been executed successfully.

8)a)ACTIVITY SELECTION PROBLEM

AIM:

To implement a C++ program for activity selection problem.

ALGORITHM:

1. Initialize the current activity to be the first activity in the sorted array.
2. For each subsequent activity, if its start time is greater than or equal to the finish time of the current activity, include it in the subset and update the current activity to be this activity.
3. Repeat step 3 until all activities have been considered.
4. The time complexity of the greedy algorithm for the activity selection problem is $O(n \log n)$, where n is the number of activities. This is because sorting the activities takes $O(n \log n)$ time and the rest of the algorithm takes $O(n)$ time.

PROGRAM:

```
#include<iostream>
#include<algorithm>
using namespace std;
// Structure to represent an activity
struct Activity {
    int start, finish;
};
// Comparison function to sort the activities based on their finish time
bool activityCompare(Activity s1, Activity s2) {
    return (s1.finish< s2.finish);
}
// Function to print the maximum number of activities that can be performed
void printMaxActivities(Activity arr[], int n) {
    sort(arr, arr+n, activityCompare);
    int i = 0;
    cout<< "The following activities are selected: " <<endl;
    // Select the first activity
    cout<< "(" <<arr[i].start<< ", " <<arr[i].finish << ") ";
    // Iterate through the rest of the activities
    for (int j = 1; j < n; j++) {
        // If this activity has a start time greater than or equal to
        // the finish time of the previous activity, select it
        if (arr[j].start>= arr[i].finish) {
            cout<< "(" <<arr[j].start<< ", " <<arr[j].finish << ") ";
            i = j;
        }
    }
}
int main() {
    Activity arr[] = { {5, 9}, {1, 2}, {3, 4}, {0, 6}, {5, 7}, {8, 9} };
    int n = sizeof(arr)/sizeof(arr[0]);
    printMaxActivities(arr, n);
    return 0;
}
```

Output

The following activities are selected:

(1, 2) (3, 4) (5, 7) (8, 9)

2.GRAPH COLORING PROBLEM

Aim

To implement a C++ program for activity graph colouring problem.

Algorithm:

1. Start with a vertex and colour.
2. Give all its neighbouring colours a vertex, but before that keep a check on used colours and unused colours.
3. Only assign, unused colours to the upcoming vertex.

Program

```
#include<bits/stdc++.h>
```

```
using namespace std;
```

```
int n,e,i,j;
vector<vector<int>>> graph;
vector<int>color;
bool vis[100011];
```



```

void greedyColoring()
{
    color[0] = 0;
    for (i=1;i<n;i++)
        color[i] = -1;

    bool unused[n];

    for (i=0;i<n;i++)
        unused[i]=0;

    for (i = 1; i < n; i++)
    {
        for (j=0;j<graph[i].size();j++)
            if (color[graph[i][j]] != -1)
                unused[color[graph[i][j]]] = true;
        int cr;
        for (cr=0;cr<n;cr++)
            if (unused[cr] == false)
                break;

        color[i] = cr;

        for (j=0;j<graph[i].size();j++)
            if (color[graph[i][j]] != -1)
                unused[color[graph[i][j]]] = false;
    }
}

int main()
{
    int x,y;
    cout<<"Enter number of vertices and edges respectively:";
    cin>>n>>e;
    cout<<"\n";
    graph.resize(n);
    color.resize(n);
    memset(vis,0,sizeof(vis));
    for(i=0;i<e;i++)
    {
        cout<<"\nEnter edge vertices of edge "<<i+1<<" :";
        cin>>x>>y;
        x--; y--;
        graph[x].push_back(y);
        graph[y].push_back(x);
    }
    greedyColoring();
    for(i=0;i<n;i++)
    {
        cout<<"Vertex "<<i+1<<" is coloured "<<color[i]+1<<"\n";
    }
}

```

OUTPUT:

Enter number of vertices and edges respectively:4 5

Enter edge vertices of edge 1 :1 2

Enter edge vertices of edge 2 :2 3

Enter edge vertices of edge 3 :3 4

Enter edge vertices of edge 4 :4 1

Enter edge vertices of edge 5 :2 4

Vertex 1 is coloured 0
Vertex 2 is coloured 1
Vertex 3 is coloured 0
Vertex 4 is coloured 2

RESULT:

Thus the C++ program for implementing activity selection problem using greedy approach has been executed successfully

8)b)HUFFMAN CODING ALGORITHM

AIM

To implement a C++ program for huffman Coding Algorithm.

ALGORITHM

1. create a priority queue Q consisting of each unique character.
2. sort then in ascending order of their frequencies.
3. for all the unique characters:
4. create a newNode
5. extract minimum value from Q and assign it to leftChild of newNode
6. extract minimum value from Q and assign it to rightChild of newNode
7. calculate the sum of these two minimum values and assign it to the value of newNode
8. insert this newNode into the tree
9. return rootNode

PROGRAM

```
// Huffman Coding in C++

#include <iostream>
using namespace std;

#define MAX_TREE_HT 50

struct MinHNode {
    unsigned freq;
    char item;
    struct MinHNode *left, *right;
};

struct MinH {
    unsigned size;
    unsigned capacity;
    struct MinHNode **array;
};

// Creating Huffman tree node
struct MinHNode *newNode(char item, unsigned freq) {
    struct MinHNode *temp = (struct MinHNode*)malloc(sizeof(struct MinHNode));

    temp->left = temp->right = NULL;
    temp->item = item;
    temp->freq = freq;

    return temp;
}

// Create min heap using given capacity
struct MinH *createMinH(unsigned capacity) {
    struct MinH *minHeap = (struct MinH*)malloc(sizeof(struct MinH));
    minHeap->size = 0;
    minHeap->capacity = capacity;
    minHeap->array = (struct MinHNode **)malloc(minHeap->capacity * sizeof(struct MinHNode *));
    return minHeap;
}

// Print the array
void printArray(int arr[], int n) {
    int i;
    for (i = 0; i < n; ++i)
        cout<<arr[i];

    cout<< "\n";
}
```

```

}

// Swap function
void swapMinHNode(struct MinHNode **a, struct MinHNode **b) {
    struct MinHNode *t = *a;
    *a = *b;
    *b = t;
}

// Heapify
void minHeapify(struct MinH *minHeap, int idx) {
    int smallest = idx;
    int left = 2 * idx + 1;
    int right = 2 * idx + 2;

    if (left < minHeap->size && minHeap->array[left]->freq < minHeap->array[smallest]->freq)
        smallest = left;

    if (right < minHeap->size && minHeap->array[right]->freq < minHeap->array[smallest]->freq)
        smallest = right;

    if (smallest != idx) {
        swapMinHNode(&minHeap->array[smallest],
            &minHeap->array[idx]);
        minHeapify(minHeap, smallest);
    }
}

// Check if size is 1
int checkSizeOne(struct MinH *minHeap) {
    return (minHeap->size == 1);
}

// Extract the min
struct MinHNode *extractMin(struct MinH *minHeap) {
    struct MinHNode *temp = minHeap->array[0];
    minHeap->array[0] = minHeap->array[minHeap->size - 1];

    --minHeap->size;
    minHeapify(minHeap, 0);

    return temp;
}

// Insertion
void insertMinHeap(struct MinH *minHeap, struct MinHNode *minHeapNode) {
    ++minHeap->size;
    int i = minHeap->size - 1;

    while (i && minHeapNode->freq < minHeap->array[(i - 1) / 2]->freq) {
        minHeap->array[i] = minHeap->array[(i - 1) / 2];
        i = (i - 1) / 2;
    }

    minHeap->array[i] = minHeapNode;
}

// Build min heap
void buildMinHeap(struct MinH *minHeap) {
    int n = minHeap->size - 1;
    int i;

    for (i = (n - 1) / 2; i >= 0; --i)
        minHeapify(minHeap, i);
}

int isLeaf(struct MinHNode *root) {
    return !(root->left) && !(root->right);
}

```

```

struct MinH *createAndBuildMinHeap(char item[], int freq[], int size) {
    struct MinH *minHeap = createMinH(size);

    for (int i = 0; i < size; ++i)
        minHeap->array[i] = newNode(item[i], freq[i]);

    minHeap->size = size;
    buildMinHeap(minHeap);

    return minHeap;
}

```

```

struct MinHNode *buildHfTree(char item[], int freq[], int size) {
    struct MinHNode *left, *right, *top;
    struct MinH *minHeap = createAndBuildMinHeap(item, freq, size);

```

```

    while (!checkSizeOne(minHeap)) {
        left = extractMin(minHeap);
        right = extractMin(minHeap);

        top = newNode('$', left->freq + right->freq);

        top->left = left;
        top->right = right;

```

```

    insertMinHeap(minHeap, top);
    }
    return extractMin(minHeap);
}

```

```

void printHCodes(struct MinHNode *root, int arr[], int top) {
    if (root->left) {
        arr[top] = 0;
        printHCodes(root->left, arr, top + 1);
    }

```

```

    if (root->right) {
        arr[top] = 1;
        printHCodes(root->right, arr, top + 1);
    }
    if (isLeaf(root)) {
        cout << root->item << " | ";
        printArray(arr, top);
    }
}

```

// Wrapper function

```

void HuffmanCodes(char item[], int freq[], int size) {
    struct MinHNode *root = buildHfTree(item, freq, size);

```

```

    int arr[MAX_TREE_HT], top = 0;

```

```

    printHCodes(root, arr, top);
}

```

```

int main() {
    char arr[] = {'A', 'B', 'C', 'D'};
    int freq[] = {5, 1, 6, 3};

```

```

    int size = sizeof(arr) / sizeof(arr[0]);

```

```

    cout << "Char | Huffman code ";
    cout << "\n-----\n";
    HuffmanCodes(arr, freq, size);

```

```

}

```

OUTPUT

Char	Huffman code
C	0
B	100
D	101
A	11

RESULT:

Thus the C++ program for implementing huffman Coding Algorithm using greedy approach has been executed successfully

8)c)SHORTEST SUPERSTRING PROBLEM

AIM

To implement a C++ program for shortest superstring problem.

ALGORITHM

1. Create an auxiliary array of strings, temp[]. Copy contents of arr[] to temp[].
2. While the size of temp[] is greater than 1
3. In temp[], find the strings that result in maximum overlapping. Say we have two strings “s1” and “s2” that results in maximum overlapping among all the given strings in temp[].
4. Remove string “s1” and “s2” from the temp[] and add the resultant string after merging s1 and s2 in temp[].
5. Return temp[0], the only final string left in temp which is the required shortest superstring.

PROGRAM

```
#include <bits/stdc++.h>
using namespace std;
```

```
int Overlap(string s1,string s2, string &str)
{
    int max_len = INT_MIN;
    int len1 = s1.length();
    int len2 = s2.length();

    //suffix of s1 with prefix of s2
    for (int i = 1; i<=min(len1, len2); i++)
    {
        // Compare last i char of s1 with first i char of s2.
        if (s1.compare(len1-i, i, s2,0, i) == 0)
        {
            if (max_len<i)
            {
                max_len = i;
                str = s1 + s2.substr(i);
            }
        }
    }

    //suffix of s2 with prefix of s1
    for (int i = 1; i<=min(len1, len2); i++)
    {
        // Compare last i char of s1 with first i char of s2.
        if (s1.compare(0, i, s2,len2-i, i) == 0)
        {
            if (max_len<i)
            {
                max_len = i;
                str = s2 + s1.substr(i);
            }
        }
    }

    return max_len;
}

// Function to calculate
// smallest string that contains
// each string in the given
// set as substring.
string findShortestSuperstring(vector<string>arr)
{
    int len=arr.size();
    while(len>1)
    {
        int max_len = INT_MIN;
```

```

        int l, r;
        string max_string;

        // Maximum overlap
        for (int i = 0; i < len; i++)
        {
            for (int j = i + 1; j < len; j++)
            {
                string str;
                int res = Overlap(arr[i], arr[j], str);

                if (max_len < res)
                {
                    max_len = res;
                    max_string = str;
                    // indexes to replace
                    l = i, r = j;
                }
            }
        }
        len--;
        if (max_len == INT_MIN)
            arr[0] += arr[len];
        else
        {
            arr[l] = max_string;
            arr[r] = arr[l];
        }
    }
    return arr[0];
}

// Driver program
int main()
{
    vector<string> arr = {"catg", "ctaagt", "gcta", "ttca", "atgcatc"};
    cout << findShortestSuperstring(arr);

    return 0;
}

```

OUTPUT

“gctaagttcatgcatc”

RESULT:

Thus the C++ program for implementing shortest superstring problem using greedy approach has been executed successfully

8)d)MINIMUM SPANNING TREE

AIM:

To implement a C++ program for minimum spanning tree.

ALGORITHM

1. Initialize the minimum spanning tree with a vertex chosen at random.
2. Find all the edges that connect the tree to new vertices, find the minimum and add it to the tree
3. Keep repeating step 2 until we get a minimum spanning tree

PROGRAM

// Prim's Algorithm in C++

```
#include <cstring>
#include <iostream>
using namespace std;

#define INF 99999999

// number of vertices in graph
#define V 5

// create a 2d array of size 5x5
//for adjacency matrix to represent graph

int G[V][V] = {
    {0, 9, 75, 0, 0},
    {9, 0, 95, 19, 42},
    {75, 95, 0, 51, 66},
    {0, 19, 51, 0, 31},
    {0, 42, 66, 31, 0}};

int main() {
    int no_edge; // number of edge

    // create a array to track selected vertex
    // selected will become true otherwise false
    int selected[V];

    // set selected false initially
    memset(selected, false, sizeof(selected));

    // set number of edge to 0
    no_edge = 0;

    // the number of edge in minimum spanning tree will be
    // always less than (V - 1), where V is number of vertices in
    //graph

    // choose 0th vertex and make it true
    selected[0] = true;

    int x; // row number
    int y; // col number

    // print for edge and weight
    cout<< "Edge"
    << " : "
    << "Weight";
    cout<<endl;
    while (no_edge< V - 1) {
        //For every vertex in the set S, find the all adjacent vertices
        // , calculate the distance from the vertex selected at step 1.
        // if the vertex is already in the set S, discard it otherwise
        //choose another vertex nearest to selected vertex at step 1.

        int min = INF;
        x = 0;
```

```

y = 0;

for (int i = 0; i < V; i++) {
    if (selected[i]) {
        for (int j = 0; j < V; j++) {
            if (!selected[j] && G[i][j]) { // not in selected and there is an edge
                if (min > G[i][j]) {
                    min = G[i][j];
                    x = i;
                    y = j;
                }
            }
        }
    }
}

cout<< x << " - " << y << " : " << G[x][y];
cout<<endl;
    selected[y] = true;
    no_edge++;
}

return 0;
}

```

OUTPUT

```

Edge : Weight
0 - 1 : 9
1 - 3 : 19
3 - 4 : 31
3 - 2 : 51

```

RESULT:

Thus the C++ program for implementing minimum spanning tree using greedy approach has been executed successfully

8)e)SIEVE OF SUNDARAM

AIM

To implement a C++ program for Sieve of Sundaram

ALGORITHM

1. In general Sieve of Sundaram, produces primes smaller than $(2*x + 2)$ for given number x . Since we want primes smaller than n , we reduce $n-1$ to half. We call it $nNew$.
 $nNew = (n-1)/2$;
For example, if $n = 102$, then $nNew = 50$.
if $n = 103$, then $nNew = 51$
2. Create an array `marked[n]` that is going to be used to separate numbers of the form $i+j+2ij$ from others where $1 \leq i \leq j$
3. Initialize all entries of `marked[]` as false.
4. // Mark all numbers of the form $i + j + 2ij$ as true
5. // where $1 \leq i \leq j$
Loop for $i=1$ to $nNew$
 - a) $j = i$;
 - b) Loop While $(i + j + 2*i*j) \leq nNew$, then print $2*i + 1$ as first prime.
6. Remaining primes are of the form $2i + 1$ where i is index of NOT marked numbers. So print $2i + 1$ for all i such that `marked[i]` is false.

PROGRAM

```
// C++ program to print primes smaller than n using
// Sieve of Sundaram.
#include <bits/stdc++.h>
using namespace std;

// Prints all prime numbers smaller
int SieveOfSundaram(int n)
{
    // In general Sieve of Sundaram, produces primes smaller
    // than  $(2*x + 2)$  for a number given number  $x$ .
    // Since we want primes smaller than  $n$ , we reduce  $n$  to half
    int nNew = (n-1)/2;

    // This array is used to separate numbers of the form  $i+j+2ij$ 
    // from others where  $1 \leq i \leq j$ 
    bool marked[nNew + 1];

    // Initialize all elements as not marked
    memset(marked, false, sizeof(marked));

    // Main logic of Sundaram. Mark all numbers of the
    // form  $i + j + 2ij$  as true where  $1 \leq i \leq j$ 
    for (int i=1; i<=nNew; i++)
        for (int j=i; (i + j + 2*i*j) <= nNew; j++)
            marked[i + j + 2*i*j] = true;

    // Since 2 is a prime number
    if (n > 2)
        cout << 2 << " ";

    // Print other primes. Remaining primes are of the form
    //  $2*i + 1$  such that marked[i] is false.
    for (int i=1; i<=nNew; i++)
```

```
        if (marked[i] == false)
            cout << 2*i + 1 << " ";
    }

// Driver program to test above
int main(void)
{
    int n = 20;
    SieveOfSundaram(n);
    return 0;
}
```

OUTPUT

2 3 5 7 11 13 17 19

RESULT:

Thus the C++ program for implementing Sieve of Sundaram using greedy approach has been executed successfully

8)f)REMOVE INVALID PARENTHESIS

AIM

To implement a C++ program for removing invalid parenthesis.

ALGORITHM

1. To solve this, we will follow these steps –
2. Define a method called solve(), this will take pos, left, right, l, r, string array res and string temp.
3. if pos is same as size of s, then,
4. if left is same as 0 and right is same as 0, then,
5. if (m[temp] is non-zero) is false, then,
6. insert temp at the end of res
7. m[temp] := 1
8. return
9. if s[pos] is same as '(' and right > 0, then,
10. call solve(pos + 1, left, right - 1, l, r, res, temp + empty string)
11. Otherwise when s[pos] is same as ')' and left > 0, then –
12. call solve(pos + 1, left - 1, right, l, r, res, temp + empty string)
13. if s[pos] is same as '(', then,
14. call solve(pos + 1, left, right, l + 1, r, res, temp + "(")
15. Otherwise when s[pos] is same as ')' and l > r, then –
16. call solve(pos + 1, left, right, l, r + 1, res, temp + ')')
17. if s[pos] is not equal to '(' and s[pos] is not equal to ')', then,
18. call solve(pos + 1, left, right, l, r, res, temp + s[pos])

PROGRAM

```
#include <bits/stdc++.h>
using namespace std;
void print_vector(vector<auto> v){
    cout << "[";
    for(int i = 0; i<v.size(); i++){
        cout << v[i] << ", ";
    }
    cout << "]"<<endl;
}
class Solution {
public:
    string s;
    map <string,int> m;
    void solve(int pos, int left, int right,int l, int r, vector <string> &res, string temp=""){
        if(pos == s.size()){
            if(left==0 && right==0 && l==r){
                if(!m[temp])
                    res.push_back(temp);
                m[temp] = 1;
            }
        }
    }
};
```

```

        }
        return;
    }
    if(s[pos] == '(' && right>0 ){
        solve(pos+1,left,right-1,l,r,res,temp+"" );
    } else if(s[pos] == ')' && left>0) {
        solve(pos+1,left-1,right,l,r,res,temp+"" );
    }
    if(s[pos] == '(')solve(pos+1,left,right,l+1,r,res,temp+"(");
    else if(s[pos] == ')' && l>r)solve(pos+1,left,right,l,r+1,res,temp+")");
    if(s[pos]!='(' && s[pos]!=')')solve(pos+1,left,right,l,r,res,temp+s[pos]);
}
vector<string> removeInvalidParentheses(string s) {
    vector <string > res;
    int l = 0;
    int r=0;
    this->s = s;
    for(int i =0;i<s.size();i++){
        if(s[i] == '('){
            r++;
        } else if(s[i]==')') {
            if(r)r--;
            else l++;
        }
    }
    solve(0,l,r,0,0,res);
    return res;
}
};
main(){
    Solution ob;
    print_vector(ob.removeInvalidParentheses("()()()"));
}

```

OUTPUT

[()(), ()(),]

RESULT:

Thus the C++ program for implementing removing invalid parenthesis using greedy approach has been executed successfully

8)g) MAXIMUM RIBBON CUT PROBLEM

AIM

To implement a C++ program for Maximum ribbon cut problem.

ALGORITHM

1. A person X has a ribbon cut it by fulfilling the below conditions
2. After the cutting each ribbon piece should have length a, b or c.
3. After the cutting the number of ribbon pieces should be maximum.
4. $1 \leq n, a, b, c \leq 4000$ The numbers a, b and c can coincide.

PROGRAM

```
//cut the ribbon
#include<bits/stdc++.h>
using namespace std;

int main()
{
    int i,j,k,n;
    int a,b,c;
    int x,y,z;

    cout<<"Enter the length of ribbon"<<endl;
    cin>>n;

    cout<<"Enter the 3 values of lengths allowed"<<endl;
    cin>>a>>b>>c;

    //array to memoize values
    vector<int> dp(n+1);

    //initialize
    dp[0]=0;

    for(i=1;i<=n;i++)
    {
        x=y=z=-1;

        if(i>=a)
            x=dp[i-a];

        if(i>=b)
            y=dp[i-b];

        if(i>=c)
            z=dp[i-c];

        if(x!=-1 && y!=-1 && z!=-1)
            dp[i]=-1;

        else
            dp[i]=max(max(x,y),z)+1;
    }

    if(dp[n]==-1)
        cout<<"Not possible";

    else
        cout<<"Maximum number of pieces in which the ribbon can be cut is "<<endl<<dp[n];
        cout<<endl;
        return 0;
}
```

OUTPUT

Enter the length of ribbon

8

Enter the 3 values of lengths allowed

2

3

4

Maximum number of pieces in which the ribbon can be cut is

4

RESULT:

Thus the C++ program for implementing removing invalid parenthesis using greedy approach has been executed successfully

8)h)DIALS ALGORITHM

AIM

To implement a C++ program for Dials Algorithm.

ALGORITHM

1. Let the number of vertices be V , the number of edges be E , and let the source vertex be S . We assume the graph is represented as an adjacency list, where $\text{adjList}[v]$ gives the adjacent vertices of the vertex v . The adjacency list stores the vertices as well as the edge weights. Let the maximum edge weight be C .

2. Initialise the data structures. dist should contain all values as infinity, except for S . $\text{dist}[S]$ should be 0. $\text{buckets}[0]$ should contain S , and all other buckets should be empty. For this article, we are going to implement buckets as an array of hashsets (we will use dynamically sized arrays for this)

3. Iterate through the buckets array, and find the first bucket that is not empty. This will be our EXTRACT MIN step. As we have defined, $\text{buckets}[i]$ contains all the vertices at a distance i from the source. So, when we take the first non-empty bucket, we get the smallest value of i , which means those vertices have the smallest distance from the source.

4. IMPORTANT STEP: Note the index i over here. We will need it later. If $\text{buckets}[i]$ has many vertices, take the first one (we can take any). Let's call it v .

5. Go through the neighbours of v from the adjacency list. For each neighbour (let's call it nei): Check if $\text{dist}[v] + w < \text{dist}[\text{nei}]$. If true, move nei from its current bucket into its new bucket. Its distance from the source has changed, so its bucket will change too. Its old bucket will be $\text{buckets}[\text{dist}[\text{nei}]]$, as the buckets are indexed by distance. Its new bucket will be $\text{buckets}[\text{dist}[v] + w]$. buckets is an array of hashsets, so $\text{buckets}[x]$ will give a hashset for all values of x . Therefore, we will be able to delete and add elements to any bucket in constant time. Set $\text{dist}[\text{nei}] = \text{dist}[v] + w$.

6. Delete v from $\text{buckets}[i]$. Go to step 4, but instead of starting with index 0, continue from the previous index i . We can do this since we know that no vertex can have a distance smaller than i now. Any vertex that could have a distance smaller than i would have been present already. For the neighbours of v , whose distance we just reduced - their new distances can never be smaller than i , because edge weights are all positive. So, since $w > 0$, $\text{dist}[v] + w > \text{dist}[v]$. As $\text{dist}[v] = i$, $\text{dist}[\text{nei}] > i$. Hence, we can continue from i itself. (We do not go to $i+1$ because the bucket at i may have had many vertices, we have only considered one) We continue this until we reach the last bucket.

PROGRAM

```
#include <iostream>
#include <vector>
#include <unordered_set>
using namespace std;

// This class represents a graph using adjacency list representation
class Graph {
public:
    int V; // no of vertices
    // adjacency list of the graph. adj[u] contains the list of pairs (v, w)
    // where
    // v is a neighbour of u and w is the weight of the edge (u, v)
    vector<vector<pair<int, int>>> adj;
    Graph(int v) { // constructor
        V = v;
        // creates an adjacency list of size V+1, so that our vertices are
        // 1-indexed
        adj = vector<vector<pair<int, int>>>(V + 1);
    }
    // adds an edge from u to v, with weight w
    void addEdge(int u, int v, int w) { adj[u].push_back({v, w}); }
};

// Executes Dial's algorithm on the Graph g, and prints the shortest path
// from S to all other vertices. C is the maximum weight of any edge in the
// graph
void dialsAlgorithm(Graph g, int C, int S) {
```

```

// The maximum number of buckets possible
int maxBuckets = C * g.V;

// buckets[i] stores all vertices that have a distance of i from S
vector<unordered_set<int>> buckets(maxBuckets);

// dist[i] stores the shortest distance from S to i
//Initially, all distances are infinity or INT_MAX here.
vector<int> dist(g.V + 1, INT_MAX);

// initially, S is at distance 0 from itself
dist[S] = 0;
buckets[0].insert(S);

// The current bucket that we are at in the algorithm
int bucketPointer = 0;
while (true) {
    // iterate through the buckets until we find a non-empty bucket, or run
    // out of buckets
    while (bucketPointer < maxBuckets && buckets[bucketPointer].empty()) {
        bucketPointer++;
    }

    // if we ran out of buckets, then we are done
    if (bucketPointer >= maxBuckets)
        break;

    // otherwise, we have found a non-empty bucket, and we will process it
    // We can choose any vertex. We will choose the first one.
    int v = *buckets[bucketPointer].begin();

    // remove v from the bucket, as we won't need it again.
    buckets[bucketPointer].erase(v);

    // iterate through all the neighbours of v
    for (pair<int, int> neiPair : g.adj[v]) {

        int nei = neiPair.first; // the neighbour
        int w = neiPair.second; // the weight of the edge (v, nei)

        int altDist = dist[v] + w; // the distance from S to nei if we
        // travelled through v

        int currentDist = dist[nei]; // the current distance from S to nei

        // if we can improve the distance to nei by going through v, then we
        // will do so
        if (altDist < currentDist) {
            // if nei is not at infinity, it must be in some bucket. We will
            // remove it from that bucket
            if (currentDist != INT_MAX) {
                buckets[currentDist].erase(nei);
            }

            // insert nei into the bucket that corresponds to its new
            // distance
            buckets[altDist].insert(nei);

            // update the distance to nei
            dist[nei] = altDist;
        }
    }
}

// print the shortest distances from S to all other vertices
for (int i = 1; i <= g.V; i++) {
    cout << i << " ";
}
cout << endl;
for (int i = 1; i <= g.V; i++) {

```

```

        cout << dist[i] << " ";
    }
    // And, we are done!
}
int main() {
    Graph g(6);
    g.addEdge(1, 3, 2);
    g.addEdge(1, 2, 4);
    g.addEdge(2, 5, 2);
    g.addEdge(3, 2, 1);
    g.addEdge(3, 4, 4);
    g.addEdge(3, 5, 4);
    g.addEdge(4, 6, 1);
    g.addEdge(5, 4, 3);
    g.addEdge(5, 6, 3);
    dialsAlgorithm(g, 4, 1);
    return 0;
}

```

OUTPUT

```

1 2 3 4 5 6
0 3 2 6 5 7

```

RESULT:

Thus the C++ program for implementing Dials Algorithm using greedy approach has been executed successfully

8)i)FLIP THE WORLD

AIM

To implement a C++ program for Flip the World Algorithm

ALGORITHM

1. Reverse the given string str using STL function reverse().
2. Iterate the reversed string and whenever a space is found reverse the word before that space using the STL function reverse().

PROGRAM

```
// C++ program

#include <bits/stdc++.h>
using namespace std;

// Function to reverse the given string
string reverseString(string str)
{
    // Reverse str using inbuilt function
    reverse(str.begin(), str.end());

    // Add space at the end so that the
    // last word is also reversed
    str.insert(str.end(), ' ');

    int n = str.length();

    int j = 0;

    // Find spaces and reverse all words
    // before that
    for (int i = 0; i < n; i++) {
        // If a space is encountered
        if (str[i] == ' ') {
            reverse(str.begin() + j, str.begin() + i);

            // Update the starting index
            // for next word to reverse
            j = i + 1;
        }
    }

    // Remove spaces from the end of the
    // word that we appended
    str.pop_back();

    // Return the reversed string
    return str;
}

// Driver code
int main()
{
    string str = "I like this code";

    // Function call
    string rev = reverseString(str);

    // Print the reversed string
    cout << rev;
    return 0;
}
```

OUTPUT

code this like I

RESULT:

Thus the C++ program for implementing Flip the World Algorithm using greedy approach has been executed successfully

9)a)IMPLEMENTATION OF BACKTRACKING

AIM:

To Develop C++ program for implementing N Queen Problem process using backtracking

ALGORITHM:

1. Initialize: Start with an empty chessboard and a variable row set to 0 (representing the current row).
2. Base Case: If row is equal to the board size N, all queens are placed successfully. Print the solution (the positions of queens) and return.
3. Loop Through Columns: For each column in the current row, from 0 to N-1:
4. Check if placing a queen at the current row and column is safe (no other queens threaten it). If safe:
5. Place the queen in the current position.
6. Increment row to move to the next row.
7. Recursively call the algorithm with the updated row.
8. If the recursive call returns true, then a solution is found. Return true.
9. If the recursive call returns false, backtrack by removing the queen and trying the next column.
10. If no column in the current row leads to a solution, return false to backtrack further.

PROGRAM:

```
#include <iostream>
#include <vector>
using namespace std;
bool isSafe(int row, int col, const vector<int>& queenPositions) {
    for (int prevRow = 0; prevRow < row; ++prevRow) {
        int prevCol = queenPositions[prevRow];
        if (prevCol == col || abs(prevRow - row) == abs(prevCol - col)) {
            return false;
        }
    }
    return true;
}
void solveNQueens(int row, int n, vector<int>& queenPositions) {
    if (row == n) {
        for (int col : queenPositions) {
            cout << col << " ";
        }
        cout << endl;
        return;
    }
    for (int col = 0; col < n; ++col) {
        if (isSafe(row, col, queenPositions)) {
            queenPositions[row] = col;
            solveNQueens(row + 1, n, queenPositions);
        }
    }
}
int main() {
    int n;
    cout << "Enter the value of N: ";
    cin >> n;
    vector<int> queenPositions(n, -1);
    solveNQueens(0, n, queenPositions);
    return 0;
}
```

OUTPUT:

```
Enter the value of N: 8
0 4 7 5 2 6 1 3
0 5 7 2 6 3 1 4
0 6 3 5 7 1 4 2
0 6 4 7 1 3 5 2
```

```
1 3 5 7 2 0 6 4
1 4 6 0 2 7 5 3
1 4 6 3 0 7 5 2
1 5 0 6 3 7 2 4
1 5 7 2 0 3 6 4
1 6 2 5 7 4 0 3
1 6 4 7 0 3 5 2
1 7 5 0 2 4 6 3
2 0 6 4 7 1 3 5
2 4 1 7 0 6 3 5
2 4 1 7 5 3 6 0
2 4 6 0 3 1 7 5
2 4 7 3 0 6 1 5
```

RESULT:

Thus the C++ program for implementing N Queen Problem process using backtracking has been executed successfully

9)b)IMPLEMENTATION OF BRANCH & BOUND

AIM:

To develop C++ program to implement 0/1 Knapsack Problem using Branch and Bound.

ALGORITHM:

1. Create Priority Queue (Min Heap): Initialize a priority queue (min heap) to store nodes (candidates) for exploration. Each node contains information about the current item index, current weight, current value, and an upper bound (value of the node if all remaining items are taken).
2. Create Initial Node: Create an initial node with index 0 (first item), current weight 0, current value 0, and compute its upper bound using a greedy approach (considering fractional parts of items).
3. Enqueue Initial Node: Enqueue the initial node into the priority queue.
4. Loop while Priority Queue is Not Empty:
Dequeue a node from the priority queue.
If the node represents a feasible solution and has a higher value than the best solution found so far, update the best solution.
If the node's upper bound is greater than the best solution value:
Create two child nodes: one with the next item included and one without it.
For each child node, update its weight, value, and compute its upper bound.
Enqueue both child nodes into the priority queue if their bounds are greater than the best solution value.
5. Print the Best Solution Found.

PROGRAM:

```
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std;
struct Item {
    int value;
    int weight;
};
bool compare(Item a, Item b) {
    return (a.value * b.weight) > (b.value * a.weight);
}
double fractionalKnapsack(int capacity, vector<Item>& items) {
    sort(items.begin(), items.end(), compare);
    double totalValue = 0.0;
    for (const Item& item : items) {
        if (capacity >= item.weight) {
            totalValue += item.value;
            capacity -= item.weight;
        } else {
            totalValue += (static_cast<double>(capacity) / item.weight) * item.value;
            break;
        }
    }
    return totalValue;
}
int main() {
    int n, capacity;
    cout << "Enter the number of items: ";
    cin >> n;
    cout << "Enter the knapsack capacity: ";
    cin >> capacity;
    vector<Item> items(n);
    for (int i = 0; i < n; ++i) {
        cout << "Enter value and weight of item " << i + 1 << ": ";
        cin >> items[i].value >> items[i].weight;
    }
    double maxVal = fractionalKnapsack(capacity, items);
    cout << "Maximum value that can be obtained: " << maxVal << endl;
    return 0;
}
```


OUTPUT:

Enter the number of items: 2
Enter the knapsack capacity: 10
Enter value and weight of item 1:
3
4
Enter value and weight of item 2:
5
6
Maximum value that can be obtained: 8

RESULT:

Thus, the C++ program for 0/1 knapsack problem using branch and bound was executed and output was verified successfully.

10)a)HAMILTONIAN CIRCUIT

AIM:

To develop a C++ program to implement Hamiltonian circuit problem.

ALGORITHM:

1. Create a data structure (e.g., an adjacency matrix or adjacency list) to represent the graph and Initialize an empty path and an empty set to keep track of visited vertices.
2. Pick a starting vertex arbitrarily (you can start from any vertex since a Hamiltonian circuit visits every vertex).
3. At each step, choose a neighboring vertex that has not been visited yet and add it to the path.
Mark the chosen vertex as visited.
Recursively repeat this process for the chosen vertex.
If you find a vertex that has no unvisited neighbors, backtrack by removing it from the path and marking it as unvisited.
Continue this process until all vertices have been visited, and the path forms a Hamiltonian circuit.
4. Check if the path forms a valid Hamiltonian circuit by verifying that it starts and ends at the same vertex and visits each vertex exactly once.
5. If a Hamiltonian circuit is found, print or return the path. If not, report that no Hamiltonian circuit exists.

PROGRAM:

```
#include <iostream>
#include <vector>
using namespace std;
const int MAX_VERTICES = 10; // Maximum number of vertices in the graph
class HamiltonianCircuit {
private:
    int numVertices;
    vector<vector<int>>> graph;
    vector<int> path;
public:
    HamiltonianCircuit(int n) : numVertices(n) {
        graph.resize(numVertices, vector<int>(numVertices, 0));
        path.resize(numVertices, -1);
    }

    void addEdge(int u, int v) {
        graph[u][v] = 1;
        graph[v][u] = 1;
    }

    bool isSafe(int v, int pos) {
        if (!graph[path[pos - 1]][v])
            return false;

        for (int i = 0; i < pos; ++i)
            if (path[i] == v)
                return false;

        return true;
    }

    bool findHamiltonianCircuitUtil(int pos) {
        if (pos == numVertices) {
            // Hamiltonian circuit found, print the path
            for (int i : path)
                cout << i << " ";
            cout << path[0] << endl;
            return true;
        }

        for (int v = 1; v < numVertices; ++v) {
            if (isSafe(v, pos)) {
                path[pos] = v;
```

```

        if (findHamiltonianCircuitUtil(pos + 1))
            return true;
        path[pos] = -1; // Backtrack
    }
}

return false;
}

void findHamiltonianCircuit() {
    path[0] = 0; // Start from the first vertex
    if (!findHamiltonianCircuitUtil(1))
        cout << "No Hamiltonian circuit exists in the graph." << endl;
}
};

int main() {
    int numVertices, numEdges;
    cout << "Enter the number of vertices: ";
    cin >> numVertices;
    cout << "Enter the number of edges: ";
    cin >> numEdges;

    HamiltonianCircuit hc(numVertices);

    cout << "Enter the edges (u v):" << endl;
    for (int i = 0; i < numEdges; ++i) {
        int u, v;
        cin >> u >> v;
        hc.addEdge(u, v);
    }

    cout << "Hamiltonian Circuit:" << endl;
    hc.findHamiltonianCircuit();

    return 0;
}

```

OUTPUT:

```

Enter the number of vertices: 4
Enter the number of edges: 6
Enter the edges (u v):
0 1
0 2
0 3
1 2
1 3
2 3
Hamiltonian Circuit:
0 1 2 3 0

```

RESULT:

Thus the program to implement Hamiltonian Circuit problem was executed and output was verified successfully.

10)b)SUBSET SUM PROBLEM

AIM:

To develop a C++ program to implement subset sum problem.

ALGORITHM:

1. Let nums be an array or set of positive integers and sum be the target sum.
2. Create a 2D boolean array dp, where dp[i][j] represents whether there exists a subset of the first i elements in nums that can sum up to j.
3. Iterate through the elements of nums from 1 to n, where n is the number of elements in nums.
4. After completing the dynamic programming table dp, the value in dp[n][sum] will indicate whether there exists a subset of nums that sums up to the target sum.
5. If dp[n][sum] is true, a subset with the given sum exists; otherwise, it does not.
6. If dp[n][sum] is true, you can backtrack through the dp table to find the elements included in the subset that achieves the target sum

PROGRAM:

```
#include <iostream>
#include <vector>

bool isSubsetSum(std::vector<int>& nums, int sum) {
    int n = nums.size();
    std::vector<std::vector<bool>>> dp(n + 1, std::vector<bool>(sum + 1, false));

    // If the sum is 0, then it's always possible to have an empty subset.
    for (int i = 0; i <= n; i++) {
        dp[i][0] = true;
    }

    for (int i = 1; i <= n; i++) {
        for (int j = 1; j <= sum; j++) {
            if (nums[i - 1] <= j) {
                // If the current number can be included, check if there's a subset sum without it,
                // or with it (subtracting the current number from the sum).
                dp[i][j] = dp[i - 1][j] || dp[i - 1][j - nums[i - 1]];
            } else {
                // If the current number is greater than the sum being considered,
                // it can't be included in the subset.
                dp[i][j] = dp[i - 1][j];
            }
        }
    }

    // The final cell dp[n][sum] will contain the answer.
    return dp[n][sum];
}

int main() {
    std::vector<int> nums = {3, 34, 4, 12, 5, 2};
    int sum = 9;

    if (isSubsetSum(nums, sum)) {
        std::cout << "Subset with the given sum exists." << std::endl;
    } else {
        std::cout << "No subset with the given sum exists." << std::endl;
    }

    return 0;
}
```

OUTPUT:

Subset with the given sum exists.

RESULT

Thus, the C++ program to implement subset sum problem was executed and verified successfully.

10)c)COMBINATORIAL OPTIMIZATION PROBLEMS

AIM:

To develop a C++ program to implement combinational optimization problem.

ALGORITHM:

1. Define Your Problem by Identify the decision variables, constraints, and the objective function you want to optimize and determine the search space for your problem, which represents all possible solutions.
2. Initialize a solution or a population of solutions and set up any data structures or parameters needed for your algorithm.
3. Evaluate the fitness or objective value of each solution in your population based on the problem's criteria.
4. Define stopping criteria for your algorithm. Common criteria include a maximum number of iterations, a target fitness value, or a time limit.
5. Repeat until a termination criterion is met:
6. Apply a search operator or algorithm to generate new solutions or modify existing ones. This can include mutation, crossover, or other problem-specific operators.
7. Evaluate the fitness of the new solutions.
8. Select the best solutions to form the next generation or population.
9. After the optimization loop terminates, output the best solution found during the search.

PROGRAM:

```
#include <iostream>
#include <vector>
#include <algorithm>
#include <cmath>
#include <limits>

struct City {
    double x, y;
    // Calculate the Euclidean distance between two cities
    double distanceTo(const City& other) const {
        double dx = x - other.x;
        double dy = y - other.y;
        return sqrt(dx * dx + dy * dy);
    }
};

// Calculate the total tour length
double calculateTourLength(const std::vector<City>& cities, const std::vector<int>& tour) {
    double totalLength = 0.0;
    int numCities = cities.size();

    for (int i = 0; i < numCities - 1; ++i) {
        totalLength += cities[tour[i]].distanceTo(cities[tour[i + 1]]);
    }

    // Add the distance from the last city back to the starting city
    totalLength += cities[tour[numCities - 1]].distanceTo(cities[tour[0]]);

    return totalLength;
}

int main() {
    int numCities;
    std::cout << "Enter the number of cities: ";
    std::cin >> numCities;
    std::vector<City> cities(numCities);
    std::vector<int> tour(numCities);
    std::cout << "Enter the coordinates of each city:\n";
    for (int i = 0; i < numCities; ++i) {
        std::cout << "City " << i + 1 << " (x y): ";
        std::cin >> cities[i].x >> cities[i].y;
        tour[i] = i; // Initialize the tour with the order (0, 1, 2, ..., numCities-1)
    }
    double minTourLength = std::numeric_limits<double>::max();
    std::vector<int> minTour;
    do {
        double tourLength = calculateTourLength(cities, tour);
```

```

        if (tourLength < minTourLength) {
            minTourLength = tourLength;
            minTour = tour;
        }
    } while (std::next_permutation(tour.begin() + 1, tour.end()));

    std::cout << "Optimal Tour Order: ";
    for (int city : minTour) {
        std::cout << city << " ";
    }
    std::cout << "\nOptimal Tour Length: " << minTourLength << std::endl;
    return 0;
}

```

OUTPUT:

```

Enter the number of cities: 4
Enter the coordinates of each city:
City 1 (x y): 2 3
City 2 (x y): 4 5
City 3 (x y): 6 7
City 4 (x y): 5 6
Optimal Tour Order: 0 1 2 3
Optimal Tour Length: 11.3137

```

RESULT:

Thus the C++ program to implement the combinatorial optimization problem was executed and output was verified successfully.

10)d)PEOPLE HOLDING HANDS

AIM:

To write a C++ program to implement people holding hands.

ALGORITHM:

1. The areHoldingHands function takes a string arrangement as input, representing the arrangement of people.
2. It iterates through the characters in pairs and checks if they are holding hands according to the given rules (e.g., 'MF' or 'FM' indicates holding hands).
3. If any adjacent pair is not holding hands, it returns false.
4. If all adjacent pairs are holding hands, it returns true.
5. In the main function, you can change the arrangement string to represent your specific scenario.

PROGRAM:

```
#include <bits/stdc++.h>
using namespace std;
class Solution {
public:
    class UF{
    public:
        vector<int> parent;
        int count;
        UF(int N){
            count = N;
            parent = vector<int>(N);
            for (int i = 0; i < N; i++) {
                parent[i] = i;
            }
        }
        void unionn(int a, int b){
            int parA = getParent(a);
            int parB = getParent(b);
            if (parA == parB)
                return;
            count--;
            parent[parB] = parA;
        }
        int getParent(int i){
            if (parent[i] == i)
                return i;
            return parent[i] = getParent(parent[i]);
        }
    };
    int minSwapsCouples(vector<int>& row) {
        int n = row.size();
        int N = n / 2;
        UF uf(N);
        for (int gr = 0; gr < N; gr++) {
            int a = row[gr * 2];
            int b = row[gr * 2 + 1];
            uf.unionn(a / 2, b / 2);
        }
        return N - uf.count;
    }
};
main(){
    Solution ob;
    vector<int> v = {0,2,4,1,3,5};
    cout << (ob.minSwapsCouples(v));
}
```

OUTPUT:

2

RESULT:

Thus the C++ program to implement the people holding hands was executed and output was verified successfully.

10)e)KNIGHT'S TOUR PROBLEM

AIM:

To develop C++ program to implement Knight's tour problem.

ALGORITHM:

1. Create an N x N chessboard (usually N = 8 for a standard chessboard).
2. Initialize an empty path that will store the knight's moves.
3. Place the knight on an arbitrary square (the starting position).
4. Mark the starting square as visited.
5. Repeat the following steps until all squares on the chessboard have been visited: a. Calculate the accessibility of each unvisited neighbor square from the current position. Accessibility is the number of unvisited squares that can be reached from a particular square. b. Move the knight to the unvisited neighbor square with the lowest accessibility. c. Mark the new square as visited and add it to the path.
6. If all squares have been visited, the Knight's Tour is complete.
7. If the knight gets stuck (i.e., there are no unvisited neighbor squares with valid moves), backtrack to the previous square and try alternative moves until a solution is found or proven impossible.

PROGRAM:

```
#include <iostream>
#include <iomanip>
#define N 8
using namespace std;
int sol[N][N];

bool isValid(int x, int y, int sol[N][N]) { //check place is in range and not assigned yet
    return ( x >= 0 && x < N && y >= 0 && y < N && sol[x][y] == -1);
}

void displaySolution() {
    for (int x = 0; x < N; x++) {
        for (int y = 0; y < N; y++)
            cout << setw(3) << sol[x][y] << " ";
        cout << endl;
    }
}

int knightTour(int x, int y, int move, int sol[N][N], int xMove[N], int yMove[N]) {
    int xNext, yNext;
    if (move == N*N) //when the total board is covered
        return true;

    for (int k = 0; k < 8; k++) {
        xNext = x + xMove[k];
        yNext = y + yMove[k];
        if (isValid(xNext, yNext, sol)) { //check room is preoccupied or not
            sol[xNext][yNext] = move;
            if (knightTour(xNext, yNext, move+1, sol, xMove, yMove) == true)
                return true;
            else
                sol[xNext][yNext] = -1; //backtracking
        }
    }
    return false;
}

bool findKnightTourSol() {
    for (int x = 0; x < N; x++) //initially set all values to -1 of solution matrix
        for (int y = 0; y < N; y++)
            sol[x][y] = -1;
    //all possible moves for knight
    int xMove[8] = { 2, 1, -1, -2, -2, -1, 1, 2 };
    int yMove[8] = { 1, 2, 2, 1, -1, -2, -2, -1 };
    sol[0][0] = 0; //starting from room (0, 0)

    if (knightTour(0, 0, 1, sol, xMove, yMove) == false) {
        cout << "Solution does not exist";
    }
}
```



```
        return false;
    } else
        displaySolution();
    return true;
}

int main() {
    findKnightTourSol();
}
```

OUTPUT:

```
0 59 38 33 30 17 8 63
37 34 31 60 9 62 29 16
58 1 36 39 32 27 18 7
35 48 41 26 61 10 15 28
42 57 2 49 40 23 6 19
47 50 45 54 25 20 11 14
56 43 52 3 22 13 24 5
51 46 55 44 53 4 21 12
```

RESULT:

Thus the C++ program to implement the Knight's tour problem was executed and output was verified successfully.

10)f) SUDOKU SOLVER

AIM:

To Develop a C++ program to implement by solving Sudoku in a 3X3 matrix.

ALGORITHM:

1. Create a function that checks if the given matrix is valid Sudoku or not. Keep Hash map for the row, column and boxes.
2. Create a recursive function that takes a grid and the current row and column index.
3. If the index is at the end of the matrix, i.e. $i=N-1$ and $j=N$ then check if the grid is safe or not, if safe print the grid and return true else return false.
4. The other base case is when the value of column is N, i.e. $j = N$, then move to next row, i.e. $i++$ and $j = 0$.
5. If the current index is not assigned then fill the element from 1 to 9 and recur for all 9 cases with the index of next element, i.e. $i, j+1$. if the recursive call returns true then break the loop and return true.
6. If the current index is assigned then call the recursive function with the index of the next element, i.e. $i, j+1$

PROGRAM:

```
#include <iostream>
using namespace std;
// N is the size of the 2D matrix N*N
#define N 9
void print(int arr[N][N])
{
    for (int i = 0; i < N; i++)
    {
        for (int j = 0; j < N; j++)
            cout << arr[i][j] << " ";
        cout << endl;
    }
}

bool isSafe(int grid[N][N], int row,
            int col, int num)
{
    for (int x = 0; x <= 8; x++)
        if (grid[row][x] == num)
            return false;
    for (int x = 0; x <= 8; x++)
        if (grid[x][col] == num)
            return false;

    int startRow = row - row % 3,
        startCol = col - col % 3;

    for (int i = 0; i < 3; i++)
        for (int j = 0; j < 3; j++)
            if (grid[i + startRow][j +
                startCol] == num)
                return false;

    return true;
}

bool solveSudoku(int grid[N][N], int row, int col)
{
    if (row == N - 1 && col == N)
        return true;

    if (col == N) {
        row++;
        col = 0;
    }

    if (grid[row][col] > 0)
        return solveSudoku(grid, row, col + 1);
```

```

for (int num = 1; num <= N; num++)
{

    if (isSafe(grid, row, col, num))
    {
        grid[row][col] = num;
        if (solveSudoku(grid, row, col + 1))
            return true;
    }

    grid[row][col] = 0;
}
return false;
}

int main()
{
    int grid[N][N] = {
        {3, 0, 6, 5, 0, 8, 4, 0, 0},
        {5, 2, 0, 0, 0, 0, 0, 0, 0},
        {0, 8, 7, 0, 0, 0, 0, 3, 1},
        {0, 0, 3, 0, 1, 0, 0, 8, 0},
        {9, 0, 0, 8, 6, 3, 0, 0, 5},
        {0, 5, 0, 0, 9, 0, 6, 0, 0},
        {1, 3, 0, 0, 0, 0, 2, 5, 0},
        {0, 0, 0, 0, 0, 0, 0, 7, 4},
        {0, 0, 5, 2, 0, 6, 3, 0, 0} };
    if (solveSudoku(grid, 0, 0))
        print(grid);
    else
        cout << "no solution exists " << endl;

    return 0;
}

```

OUTPUT:

```

3 1 6 5 7 8 4 9 2
5 2 9 1 3 4 7 6 8
4 8 7 6 2 9 5 3 1
2 6 3 4 1 5 9 8 7
9 7 4 8 6 3 1 2 5
8 5 1 7 9 2 6 4 3
1 3 8 9 4 7 2 5 6
6 9 2 3 5 1 8 7 4
7 4 5 2 8 6 3 1 9

```

RESULT:

Thus the C++ program to implement the solving Sudoku in a 3X3 matrix was executed and output was verified successfully.

10)g) LETTER COMBINATIONS OF A PHONE NUMBER

AIM:

To Write a C++ program to implement letter combinations of a phone number.

ALGORITHM:

1. Create a list to store the generated combinations.
2. Create a recursive function to explore the combinations.
3. Initialize an empty string to represent the current combination.
4. Create a recursive function
5. If the current combination has the same length as the input digits, add it to the list of combinations.
6. Return from the function.
7. Get the letters associated with the current digit.
8. For each letter, append it to the current combination and recursively call the function with the next digit.
9. Backtrack by removing the last added letter to explore other possibilities.
10. Call the backtracking function with the initial parameters.
11. The list of combinations now contains all possible letter combinations.

PROGRAM:

```
#include<iostream>
#include<vector>
using namespace std;
vector<string> letterCombinations(string digits)
{
    vector<string> res;
    string charmap[10] = {"0", "1", "abc", "def", "ghi", "jkl", "mno", "pqrs", "tuv", "wxyz"};
    res.push_back("");
    for (int i = 0; i < digits.size(); i++)
    {
        vector<string> tempres;
        string chars = charmap[digits[i] - '0'];
        for (int c = 0; c < chars.size(); c++)
        {
            for (int j = 0; j < res.size(); j++)
            {
                tempres.push_back(res[j]+chars[c]);
            }
        }
        res = tempres;
    }
    return res;
}
int main()
{
    string digits;
    cout<<"Enter the number\n";
    cin>>digits;
    vector<string> res = letterCombinations(digits);
    for (auto &i : res)
    {
        for (auto &j : i)
            cout << j;
        cout << ' ';
    }
    cout << endl;
}
```

OUTPUT:

Input: digits = "23"

Output: "ad","ae","af","bd","be","bf","cd","ce","cf"

Example 2:

Input: digits = "2"

Output: "a","b","c"

RESULT:

Thus the C++ program to implement the letter combinations of a phone number was executed and output was verified successfully.

10)h)ZIGZAG CONVERSION

AIM:

Given a string and number of rows 'n'. Print the string formed by concatenating n rows when input string is written in row-wise Zig-Zag fashion.

ALGORITHM:

1. Create one module that can perform this kind of operation by taking the string and the number of rows.
2. To solve this, we will follow these steps:
3. when $n = 1$, then return s
4. create an array of strings arr of size n
5. $row := 0$, and $down := true$
6. for i in range 0 to size of string $- 1$
7. insert $s[i]$ at the end of string $arr[row]$
8. if $row = n - 1$, then $down := false$, otherwise when $row = 0$, then $down := true$
9. if $down$ is true, then increase row by 1, otherwise decrease row by 1
10. $ans :=$ blank string
11. for i in range 0 to $n - 1$:
12. $ans := ans + arr[i]$
13. return ans

PROGRAM:

```
#include <iostream>
#include <string>
using namespace std;
string convert(string s, int numRows) {
    if (numRows <= 1) {
        return s;
    }
    string result = "";
    int cycleLen = 2 * numRows - 2;
    for (int i = 0; i < s.length(); ++i) {
        for (int j = 0; j + i < s.length(); j += cycleLen) {
            result += s[j + i];
            if (i != 0 && i != numRows - 1 && j + cycleLen - i < s.length()) {
                result += s[j + cycleLen - i];
            }
        }
    }
    return result;
}

int main() {
    string input;
    cout<<"Enter the string\n";
    cin>>input;

    int numRows ;
    cout<<"Enter no of rows \n";
    cin>>numRows;
    string output = convert(input, numRows);
    cout << "Zigzag conversion: " << output << endl;
    return 0;
}
```

OUTPUT:

```
Enter the string
PAYPALISHIRING
Enter no of rows
3
Zigzag conversion: PAHNAPLSIIGYIR
Example 2:
Enter the string
PAYPALISHIRING
Enter no of rows
4
Zigzag conversion: PINALSIGYAHRPI
Example 3:
```

Enter the string
A
Enter no of rows
1
Zigzag conversion: A

RESULT:

Thus, the Print the string formed by concatenating n rows when input string is written in row-wise Zig-Zag fashion program was executed and output was verified successfully.

10)i)VALID SUDOKU

AIM:

To Develop a C++ program to implement by solving 9×9 matrix called a Sudoku.

ALGORITHM:

1. Check if the given Sudoku board has the column with unique numbers or not. Then we will check for the row.
2. Each 3*3 block contains all the numbers unique in it.
3. We will check each of the block row and block column if it contains any number duplicate we will return false otherwise return true.
4. Take input of a 2-D array for the Sudoku board.
5. A Boolean function to check the row whether the elements present in it are unique is not.
6. A Boolean function to check the column whether the elements present in it are unique is not.
7. A Boolean function to check the block whether the elements present in it are unique is not.

PROGRAM:

```
#include <iostream>
#include <vector>
using namespace std;
bool isValidSudoku(vector<vector<char>>& board){
    vector<vector<bool>> cols(9, vector<bool>(9, false));
    vector<vector<bool>> grid(9, vector<bool>(9, false));
    for(int i = 0; i < 9; ++i)
    {
        vector<bool> rows(9, false);

        for(int j = 0; j < 9; ++j)
        {
            if(!isdigit(board[i][j]))
                continue;
            int idx = board[i][j] - '1';
            if(rows[idx] == true)
                return false;
            rows[idx] = true;
            if(cols[j][idx] == true)
                return false;
            cols[j][idx] = true;
            int gridIdx = ((i/3) * 3) + (j/3);
            if(grid[gridIdx][idx] == true)
                return false;
            grid[gridIdx][idx] = true;
        }
    }
    return true;
}

int main(){
    vector<vector<char>> board1 = {
        {'5','3',.,.,.,'7',.,.,.,.},
        {'6',.,.,'1','9','5',.,.,.},
        {',','9','8',.,.,.,.,'6',.},
        {'8',.,.,.,'6',.,.,.,'3'},
        {'4',.,.,'8',.,'3',.,.,'1'},
        {'7',.,.,.,'2',.,.,.,'6'},
        {',','6',.,.,.,'2','8',.},
        {',.,.,'4','1','9',.,.,'5'},
        {',.,.,.,'8',.,.,'7','9'}
    };
    if(isValidSudoku(board1)){
        cout << "The Sudoku in board1 is Valid !" << endl;
    }
    else{
        cout << "The Sudoku in board1 is NOT Valid !" << endl;
    }
}
```

```

vector<vector<char>> board2 = {
    {'5','3','3',.,.,7',.,.,.},
    {'6',.,.,1,'9','5',.,.,.},
    {'9','8',.,.,.,6',.},
    {'8',.,.,6',.,.,3'},
    {'4',.,.,8',.,3',.,1},
    {'7',.,.,2',.,.,6'},
    {'6',.,.,.,2,'8',.},
    {'',.,.,4,'1','9',.,5'},
    {'',.,.,8',.,7,'9'}
};
if(isValidSudoku(board2)){
    cout << "The Sudoku in board2 is Valid !" << endl;
}
else{
    cout << "The Sudoku in board2 is NOT Valid !" << endl;
}

return 0;
} return true;
}
int main(){
    vector<vector<char>> sudoku= {
        {'5','3',.,.,7',.,.,.},
        {'6',.,.,1,'9','5',.,.,.},
        {'9','8',.,.,.,6',.},
        {'8',.,.,6',.,.,3'},
        {'4',.,.,8',.,3',.,1},
        {'7',.,.,2',.,.,6'},
        {'6',.,.,.,2,'8',.},
        {'',.,.,4,'1','9',.,5'},
        {'',.,.,8',.,7,'9'}
    };
    bool ans= isValidSudoku(sudoku);
    if(ans){
        cout<<"True"<<endl;
    } else {
        cout<<"false"<<endl;
    }
    return 0;
}

```

OUTPUT:

The Sudoku in board1 is Valid !
The Sudoku in board2 is NOT Valid !

RESULT:

Thus, the solving 9×9 matrix called a Sudoku program was executed and output was verified successfully.

10j) REVERSE PAIR

AIM:

To develop a C++ program to Implement Reverse Pair in an array.

ALGORITHM:

1. It takes as input a vector arr, along with indices left, mid, and right that represent a range within the array.
2. It initializes a variable count to 0 to count the reverse pairs within this range.
3. It takes as input a vector arr, along with indices left and right to represent the range of elements to be sorted and counted for reverse pairs.
4. It is the entry point for the reverse pairs counting algorithm.
5. It takes a vector nums as input and calls mergeSort on the entire array.
6. A sample vector arr is defined and The reversePairs function is called on this vector to count the reverse pairs.
7. The result is printed to the console.

PROGRAM:

```
#include <iostream>
#include <vector>
using namespace std;
int merge(vector<int>& arr, int left, int mid, int right) {
    int count = 0;
    int i = left;
    int j = mid + 1;
    while (i <= mid && j <= right) {
        if (arr[i] > 2LL * arr[j]) {
            count += mid - i + 1;
            j++;
        } else {
            i++;
        }
    }
    i = left;
    j = mid + 1;
    int k = 0;
    vector<int> temp(right - left + 1);
    while (i <= mid && j <= right) {
        if (arr[i] <= arr[j]) {
            temp[k++] = arr[i++];
        } else {
            temp[k++] = arr[j++];
        }
    }
    while (i <= mid) {
        temp[k++] = arr[i++];
    }
    while (j <= right) {
        temp[k++] = arr[j++];
    }
    for (i = left, k = 0; i <= right; i++, k++) {
        arr[i] = temp[k];
    }
    return count;
}
int mergeSort(vector<int>& arr, int left, int right) {
    int count = 0;
    if (left < right) {
        int mid = left + (right - left) / 2;
        count += mergeSort(arr, left, mid);
        count += mergeSort(arr, mid + 1, right);
        count += merge(arr, left, mid, right);
    }
    return count;
}
```

```

int reversePairs(vector<int>& nums) {
    return mergeSort(nums, 0, nums.size() - 1);
}
int main() {
    vector<int> arr = {1,3,2,3,1};
    int count = reversePairs(arr);
    cout << "Number of reverse pairs: " << count << endl;
    return 0;
}

```

OUTPUT:

Number of reverse pairs: 2

Explanation: The reverse pairs are:

(1, 4) --> nums[1] = 3, nums[4] = 1, 3 > 2 * 1

(3, 4) --> nums[3] = 3, nums[4] = 1, 3 > 2 * 1

Example 2:

Input: nums = [2,4,3,5,1]

Number of reverse pairs:3

Explanation: The reverse pairs are:

(1, 4) --> nums[1] = 4, nums[4] = 1, 4 > 2 * 1

(2, 4) --> nums[2] = 3, nums[4] = 1, 3 > 2 * 1

(3, 4) --> nums[3] = 5, nums[4] = 1, 5 > 2 * 1

RESULT:

Thus, the Reverse Pair in an array program was executed and output was verified successfully.