

# Data Structures in C

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## Trees

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# Acknowledgement

- ❑ These lecture slides are based on slides by Professor Simon Hood
- ❑ Additional sources are cited separately

# Reading Assignment (required)

- Read Data Structures (recommended textbook)
  - Chapter 9 sections 9.1 – 9.5

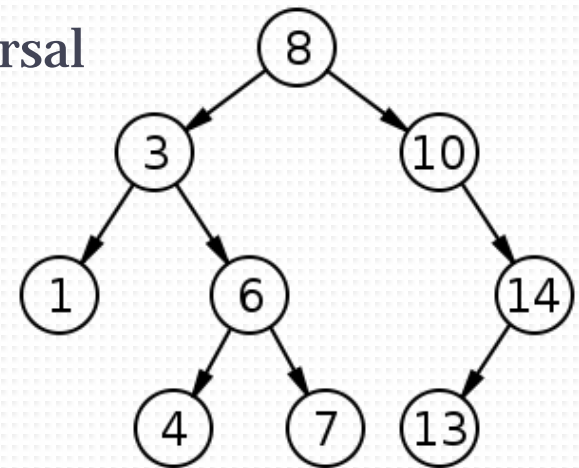
Note the textbook does a few things we might consider poor style, for example one-letter variable names and using int for Boolean values.



# Trees

- Next we'll learn about a data structure called a tree

- General trees
- Pre-order, post-order, and in-order traversal
- Binary Trees
  - Building
  - Traversing
  - Counting Nodes



- In principle a tree is not that different from a linked list
  - Imagine a linked list with **more than one next pointer**

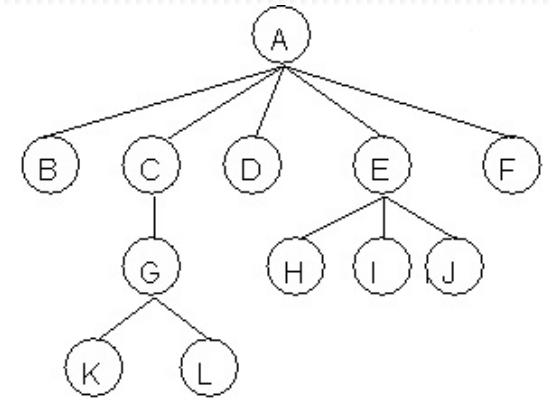
# Tree definition

- Formally, a tree is a set of nodes such that:
  1. There is one specially designated node called the root of the tree
  2. The remaining nodes are partitioned into  $m \geq 0$  disjoint sets  $T_1, T_2, \dots, T_m$
  3. Each of these sets ( $T_1, T_2$ , etc.) is also a tree

# Tree definition

## In other words...

- ❑ There is a root node
- ❑ All other nodes branch from the root node
- ❑ The branches can, in turn, branch again
- ❑ By convention, trees are drawn with the root at the top
- ❑ Normally each node stores some information
  - The tree shown above has letters of the alphabet at each node



# Tree terminology

- We use the terms **parent**, **child**, and **sibling** to refer to the nodes of a tree, but you may also see it written as branches, leaves, levels, and others
- For our structure to be a “tree” each node can only have one line leading in – though they can have any number leading out
  - Except for the root, which has no lines leading into it
- We’ll discuss **graphs** soon
  - A graph is more general than a tree, so many of the rules go away
  - A node of a graph can have two or more parent nodes, for example

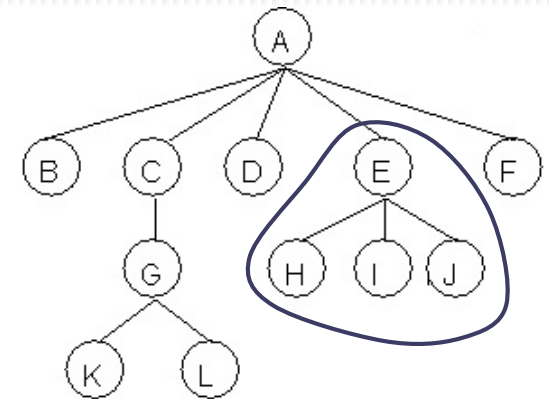
# Tree terminology

- The **degree** of a node is the number of subtrees that node has
  - We can think of it as the number of lines leaving the node
- The **depth** of a node is the number of branches which must be traversed on the path from the root to the node
  - How many lines to follow to reach that point
- The **height** of a tree is the number of levels in the tree
  - The height is one more than the tree's largest depth
- the depth (also called level) is concerned with the lines, while the height is concerned with the nodes



# Tree terminology

- A **subtree** is a part of the tree starting at a node that's not the root
  - That node is the root of the subtree



- If the relative order of the subtrees is important, the tree is called an **ordered** tree
- On the other hand if the order is unimportant, the tree is called **oriented**
- The most important type of tree for us is the **Binary Tree**

# Binary Trees

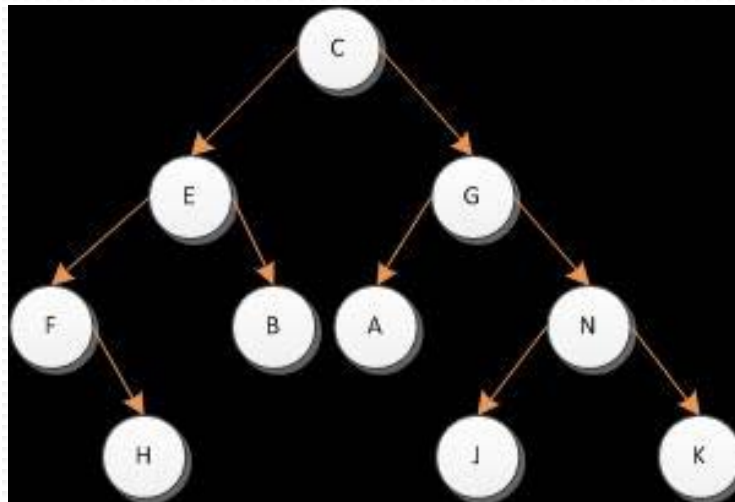
- ❑ A binary tree is a tree that has at most **two** children per node (a left and a right)
- ❑ Note that the order of the left child and right child will be important
- ❑ If there are two children for every node except the terminals (“leaves”), and the terminals have no children, the tree is **complete**

# Traversing a binary tree

- ❑ **Traversing** a tree means to “visit” all the nodes in some order
- ❑ There are three common ways to traverse a binary tree
  - Pre-order
  - In-order
  - Post-order
- ❑ We can use different traversals to interpret the tree's data in a different order, or to search for a specific data item in a different way

# Pre-order

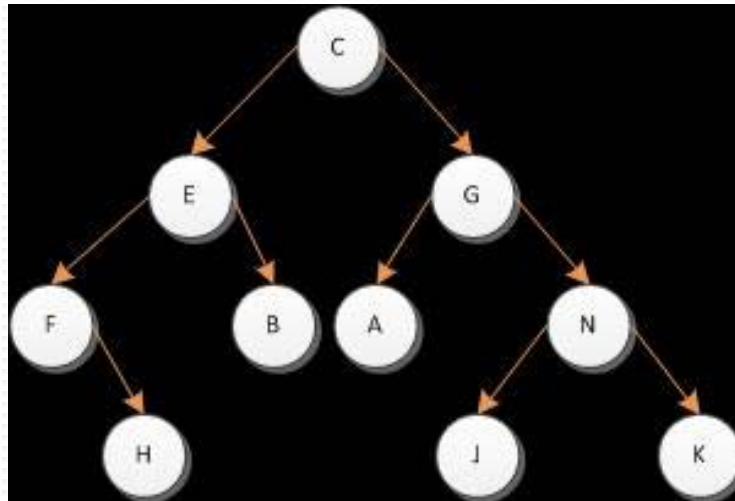
- Suppose we have the following tree



- Pre-order traversal says visit the root, then the left subtree (in pre-order), then the right subtree (in pre-order)
- The pre-order sequence is **C E F H B G A N J K**

# In-order

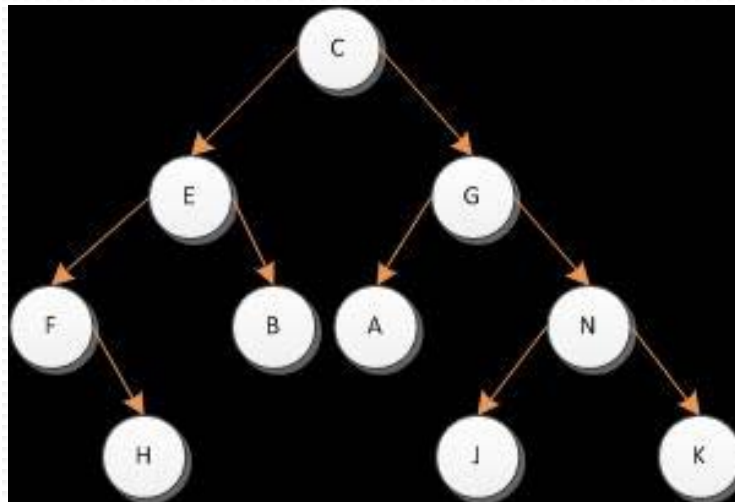
- Given the same tree, we can traverse it using in-order



- In-order traversal says visit the left subtree (in-order), root, then the right subtree (in-order)
- The in-order sequence is **F H E B C A G J N K**

# Post-order

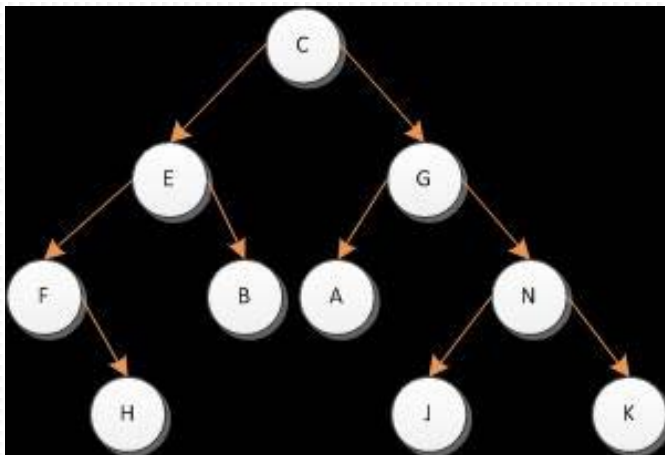
- We can also traverse the tree using post-order



- Post-order traversal says visit the left subtree (post-order), then the right subtree (post-order), then the root
- The post-order sequence is **H F B E A J K N G C**

# Other types of traversal

- Pre-order, in-order, and post-order are all considered types of **depth-first** traversal
  - Because the traversal tends to move “down” the tree before it moves across
- There is also **breadth-first** traversal, illustrated below

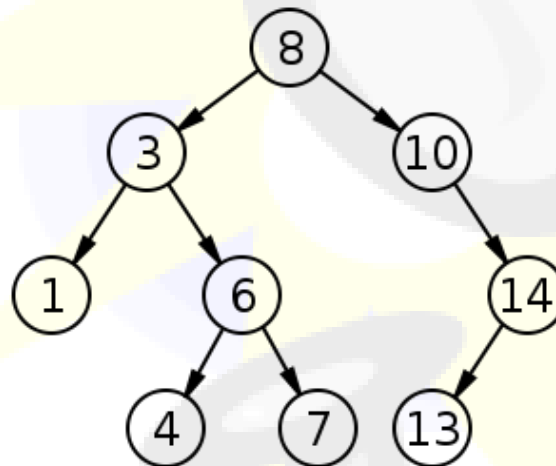


Breadth-first sequence:

**C E G F B A N H J K**

# Exercise 1

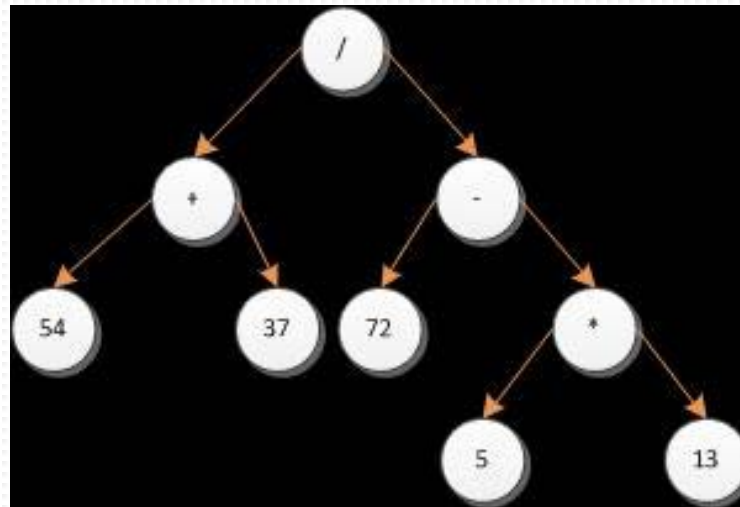
- For the binary tree shown below, write out the tree's data using pre-order, in-order, and post-order





# Math using trees

- If we want to calculate  $(54 + 37) / (72 - (5 * 13))$  we can represent this as a tree



- The pre-order traversal is  $/ + 54 37 - 72 * 5 13$
- The in-order traversal is  $54 + 37 / 72 - 5 * 13$
- The post-order traversal is  $54 37 + 72 5 13 * - /$

# Math using trees

- ❑ Interestingly, a computer can perform the mathematical operations in post-order more easily than our regular in-order!
- Create an empty stack
- Using post-order traversal, push numbers (operands) onto the stack
- If an operator is reached in post-order traversal, pop off the top two numbers, perform the operation on them, and push the result back on the stack
- ❑ This is also called a **postfix expression**, and is used by some calculators

## Exercise 2

- Draw a stack
- Use the post-order traversal described on the previous slide to calculate  $(54 + 37) / (72 - (5 * 13))$
- The post-order traversal is  $54\ 37\ +\ 72\ 5\ 13\ *\ -\ /$

# Binary Tree Implementation

The title 'Binary Tree Implementation' is centered in a large, white, sans-serif font. Below the title, there are several horizontal lines of varying lengths and colors (teal, light blue, and white) that create a layered effect. The bottom half of the slide features a light gray dotted pattern on a white background.

# Binary tree implementation

- We can represent a binary tree using a structure that's very similar to a linked list node structure

```
typedef struct treenode {  
    NodeData data; // struct inside a struct  
    struct treenode* left;  
    struct treenode* right;  
} TreeNode;
```

- Notice our data is stored in another struct—this will allow us to recycle our code by changing **NodeData** as needed
- Also notice there are two “next” pointers (that's why it's a binary tree)

# Binary tree implementation

- We might decide we just want our tree to store an **int**

```
typedef struct {  
    int num;  
} NodeData;
```

- ... or, we could store a word and associated frequency of that word

```
typedef struct{  
    char word[MAX_WORD_SIZE];  
    int freq;  
} NodeData;
```

# Binary tree implementation

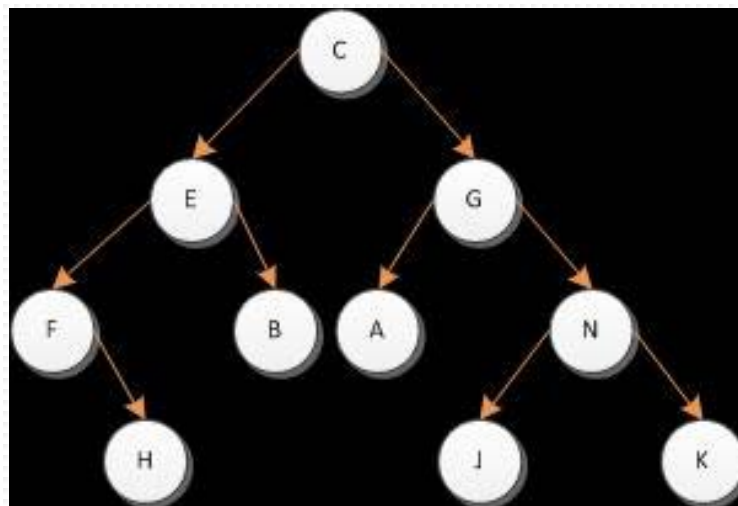
- And, of course, we need to know the root of the tree!
- Once we have access to the root, we have access to all the nodes via the left and right pointers

```
typedef struct {  
    TreeNode* root;  
} BinaryTree;
```

- Even though this structure is really just a node pointer we like to think of it as representing the whole tree

# Building a binary tree

- To create the binary tree shown below, we can supply data as a list of characters or strings
- C E F @ H @ @ B @ @ G A @ @ N J @ @ K @ @
- The @ represents an empty child node, we'll replace it with NULL in our structure





# Building a binary tree

- Note that we'll use the following definition of **NodeData**

```
typedef struct{  
    char word[MAX_WORD_SIZE];  
} NodeData;
```

- ... and that each node of our tree stores a **NodeData** structure as a variable (structure field) called **data**

# Building a binary tree

- ▣ This **recursive** function builds a tree from a text file containing a list of data specified by a FILE pointer (handle)

```
TreeNode* buildTree(FILE* in) {  
    char str[MAX_WORD_SIZE];  
    int cnt = fscanf(in, "%s", str);  
    if (cnt != 1 || strcmp(str, "@") == 0)  
        return NULL;  
    TreeNode* p = (TreeNode*) malloc(sizeof(TreeNode));  
    strcpy(p->data.word, str); // data not a ptr, use .  
    p->left = buildTree(in);  
    p->right = buildTree(in);  
    return p;  
}
```

# Building a binary tree

- We can now create a binary tree, `bt`, with the following code

```
FILE* in = fopen("btree.txt", "r");  
BinaryTree bt;  
bt.root= buildTree(in);
```

- With our tree built, let's see how we can **traverse** it
- Let's define a function called **visit()** that does the action to perform at each node (e.g. print the value in the node)

```
void visit(TreeNode* node) {  
    printf("%s ", node->data.word);  
}
```

# Traversing a binary tree

- We can visit each node in our tree using code like shown below
  - This is **pre-order** traversal

```
void preOrder(TreeNode* node) {  
    if (node != NULL) {  
        visit(node); // Process node, e.g print  
        preOrder(node->left);  
        preOrder(node->right);  
    }  
}
```

# Exercise 3 - Building a Binary Tree

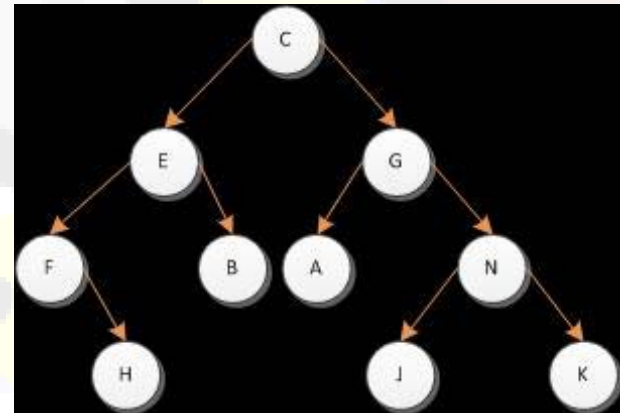
- ❑ Implement the binary tree described in the preceding slides
  - Use a separate C module (.c file) for your tree implementation
- ❑ Your main function should build a tree using the input data on the next slide
- ❑ After building, traverse the tree to print it out using the **preOrder()** function

## Exercise 3 - Building a Binary Tree (cont'd)

- Use the following data as your input file. Notice each word starts with the same letter as our example tree.

Cat  
Elephant  
Fox  
@  
Hyena  
@  
@  
Bullfrog  
@  
@  
Giraffe

Aardvark  
@  
@  
Narwhal  
Jaguar  
@  
@  
Kangaroo  
@  
@



## Exercise 4

- ❑ Starting with your program from Exercise 3, add functions **inOrder()** and **postOrder()** that traverse the tree in the other two common orders
  - Hint: Just change the order of the lines from `preOrder()`
- ❑ Call all three traversal functions and examine the output to see the differences

# Counting tree nodes

- We can count all the nodes in a binary tree with the following recursive function

```
int numNodes(TreeNode* root) {  
    if (root == NULL)  
        return 0;  
    return 1 + numNodes(root->left)  
        + numNodes(root->right);  
}
```



## Exercise 5

- ❑ Implement the function to count nodes in your program from Exercise 4
  - Call it from the main function and print out the result
- ❑ Note that you can also use the existing preOrder/inOrder/postOrder functions to count nodes, just by changing the **visit()** function
  - Will the three functions then give you different numbers, or the same number?
- ❑ Write a function called **visit2()** that counts nodes and use it to count the tree nodes in your program