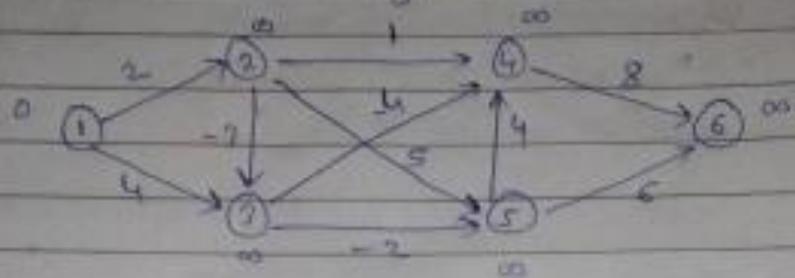


Tutorial - 7

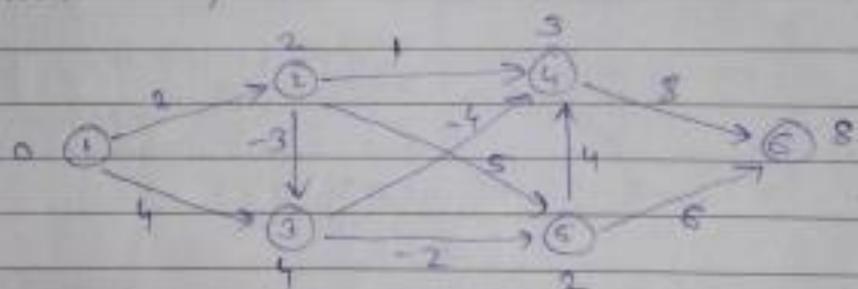
- Q. find the shortest path from node 1 to every other node in the graph of fig. using the Bellman Ford algorithm.



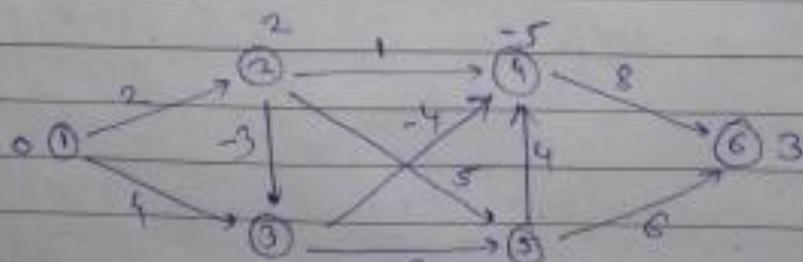
→ In Dp, we try all possible solution and pick up the best solution.
we repeat the relaxation process for $(|V|-1)$ times.

$$|V| = 6 \quad \therefore \text{relaxation} = 5 \cdot 1 = 5 \text{ times.}$$

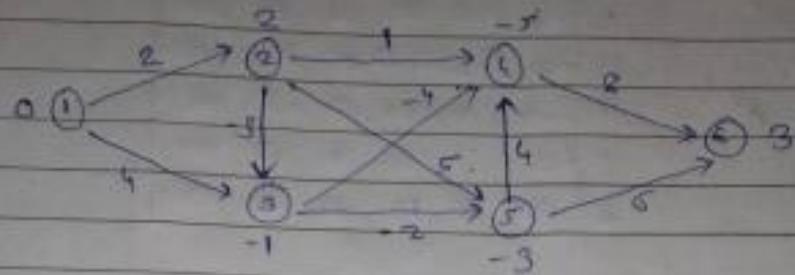
first time,



second time,



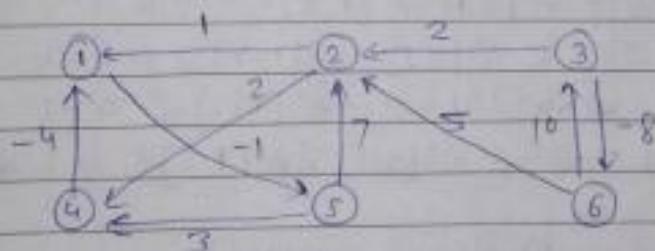
third time,



here we can see that after second step, values are not changing so this is our final solution.

| | | | | | | |
|----------|---|---|----|----|----|---|
| distance | 0 | 2 | -1 | -5 | -3 | 3 |
| node | 1 | 2 | 3 | 4 | 5 | 6 |

q. find the shortest path using floyd-warshall alg.



→ here we have to find shortest distance from every vertex to other vertex

| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|----|---|----|---|----|----|
| 1 | 0 | ∞ | ∞ | ∞ | -1 | ∞ |
| 2 | 1 | 0 | ∞ | 2 | ∞ | ∞ |
| 3 | ∞ | 2 | 0 | ∞ | ∞ | -8 |
| 4 | -4 | ∞ | ∞ | 0 | ∞ | ∞ |
| 5 | ∞ | 7 | ∞ | 3 | 0 | ∞ |
| 6 | ∞ | 5 | 10 | ∞ | ∞ | 0 |

| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|----|---|----|----|----|----|
| 1 | 0 | 0 | 0 | -1 | 0 | |
| 2 | 1 | 0 | 0 | 2 | 0 | 0 |
| 3 | 0 | 2 | 0 | 0 | 0 | -8 |
| 4 | -4 | 0 | 0 | 0 | -5 | 0 |
| 5 | 0 | 7 | 0 | 3 | 0 | 0 |
| 6 | 0 | 5 | 10 | 0 | 0 | 0 |

$$- A^{\circ}[2,3] \quad A^{\circ}[2,1] + A^{\circ}[1,3]$$

$$\infty < 1 + \infty = \infty$$

$$- A^{\circ}[2,4] \quad A^{\circ}[2,1] + A^{\circ}[1,4]$$

$$2 < 1 + \infty$$

$$- A^{\circ}[2,5] \quad A^{\circ}[2,1] + A^{\circ}[1,5]$$

$$\infty \rightarrow 1 + (-1) = 0$$

$$- A^{\circ}[2,6] \quad A^{\circ}[2,1] + A^{\circ}[1,6]$$

$$\infty < 1 + (\infty)$$

$$A^{\circ}[3,2] = 2 < A^{\circ}[3,1] + A^{\circ}[1,2] = \infty + \infty$$

$$A^{\circ}[3,4] = \infty < A^{\circ}[3,1] + A^{\circ}[1,4] = \infty + \infty$$

$$A^{\circ}[3,5] = \infty > A^{\circ}[3,1] + A^{\circ}[1,5] = \infty - 1$$

$$A^{\circ}[3,6] = -3 < A^{\circ}[3,1] + A^{\circ}[1,6] = \infty + \infty$$

$$A^{\circ}[4,2] = \infty \Rightarrow A^{\circ}[4,1] + A^{\circ}[1,2] = -4 + \infty$$

$$A^{\circ}[4,3] = \infty \Rightarrow A^{\circ}[4,1] + A^{\circ}[1,3] = -4 + \infty$$

$$A^{\circ}[4,5] = \infty \Rightarrow A^{\circ}[4,1] + A^{\circ}[1,5] = -4 - 1$$

$$A^{\circ}[4,6] = \infty \Rightarrow A^{\circ}[4,1] + A^{\circ}[1,6] = -4 + \infty$$

$$A^{\circ}[5,2] = 7 < A^{\circ}[5,1] + A^{\circ}[1,2] = \infty + \infty$$

$$A^{\circ}[5,3] = \infty < A^{\circ}[5,1] + A^{\circ}[1,3] = \infty + \infty$$

$$A^{\circ}[5,4] = 3 < A^{\circ}[5,1] + A^{\circ}[1,4] = \infty + \infty$$

$$A^{\circ}[5,6] = \infty < A^{\circ}[5,1] + A^{\circ}[1,6] = \infty + \infty$$

$$\begin{aligned}
 A^0[6,2] &= 5 < A^0[6,1] + A^0[1,2] = \infty + \infty \\
 A^0[6,3] &= 10 < A^0[6,1] + A^0[1,3] = \infty + \infty \\
 A^0[6,4] &= \infty < A^0[6,1] + A^0[1,4] = \infty + \infty \\
 A^0[6,5] &= \infty > A^0[6,1] + A^0[1,5] = \infty - 1
 \end{aligned}$$

similarly,

$$A^2 = \begin{array}{|c|cccccc|} \hline & 1 & 2 & 3 & 4 & 5 & 6 \\ \hline 1 & 0 & \infty & \infty & \infty & -1 & \infty \\ 2 & 1 & 0 & \infty & 2 & 0 & \infty \\ 3 & 3 & 2 & 0 & 4 & 2 & -8 \\ 4 & -4 & \infty & \infty & 0 & -5 & \infty \\ 5 & 8 & 7 & \infty & 3 & 0 & \infty \\ 6 & 6 & 5 & 10 & 7 & 5 & 0 \\ \hline \end{array}$$

$$A^3 = \begin{array}{|c|cccccc|} \hline & 1 & 2 & 3 & 4 & 5 & 6 \\ \hline 1 & 0 & \infty & \infty & \infty & -1 & \infty \\ 2 & 1 & 0 & \infty & 2 & 0 & \infty \\ 3 & 0 & 2 & 0 & 4 & 2 & -\infty \\ 4 & -4 & \infty & \infty & 0 & -5 & \infty \\ 5 & 8 & 7 & \infty & 3 & 0 & \infty \\ 6 & 6 & 5 & 10 & 7 & 5 & 0 \\ \hline \end{array}$$

$$A^4 = \begin{array}{|c|cccccc|} \hline & 1 & 2 & 3 & 4 & 5 & 6 \\ \hline 1 & 0 & \infty & \infty & \infty & -1 & \infty \\ 2 & -2 & 0 & \infty & 2 & -3 & \infty \\ 3 & 0 & 2 & 0 & 4 & -1 & -8 \\ 4 & -4 & \infty & \infty & 0 & -5 & \infty \\ 5 & -1 & 7 & \infty & 3 & 0 & \infty \\ 6 & 3 & 5 & 10 & 7 & 2 & 0 \\ \hline \end{array}$$

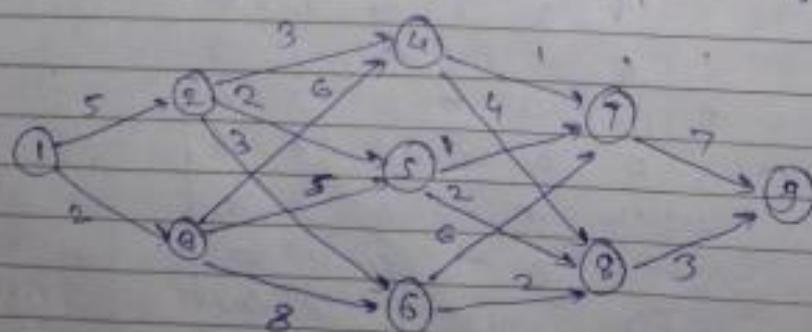
| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|----|---|----------|---|----|----------|
| 1 | 0 | 6 | ∞ | 2 | -1 | ∞ |
| 2 | -4 | 0 | ∞ | 0 | -3 | ∞ |
| 3 | 0 | 2 | 0 | 2 | -1 | -2 |
| 4 | -6 | 2 | ∞ | 0 | -5 | ∞ |
| 5 | -1 | 7 | ∞ | 3 | 0 | ∞ |
| 6 | 1 | 5 | 10 | 5 | 2 | 0 |

| | 1 | 2 | 3 | 4 | 5 | 6 |
|---|----|----|----------|----|----|----------|
| 1 | 0 | 6 | ∞ | 2 | -1 | ∞ |
| 2 | -4 | 0 | ∞ | 0 | -3 | ∞ |
| 3 | -7 | -3 | 0 | -3 | -6 | -8 |
| 4 | -6 | 2 | ∞ | 0 | -5 | ∞ |
| 5 | -1 | 7 | ∞ | 3 | 0 | ∞ |
| 6 | 1 | 5 | 10 | 5 | 2 | 0 |

This is shortest path from each vertex.

- Q. Find the minimum cost path from s to t in multistage graph using dynamic programming.

$v_1 \quad v_2 \quad v_3 \quad v_4 \quad v_5$



Step 1

| | | | | | | | | | |
|----------|---|---|---|---|---|---|---|---|---|
| v | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| cost | | | | | | | | 0 | |
| distance | | | | | | | | 9 | |

Step 2

$$\text{for cost}(4,7) = 7$$

Step 3

$$\text{for cost}(4,8) = 3$$

| | | | | | | | | | |
|----------|---|---|---|---|---|---|---|---|---|
| v | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| cost | | | | | | | 7 | 3 | 0 |
| distance | | | | | | | 9 | 9 | 9 |

Step 4

$$\text{for cost}(3,4)$$

$$\begin{aligned} \text{cost}(3,4) &= \min \{ c(4,7) + \text{cost}(7,4), \\ &\quad c(4,8) + \text{cost}(8,4) \} \\ &= \min \{ 1+7, 6+3 \} \\ &= \min \{ 8, 7 \} \end{aligned}$$

$$\text{cost}(3,4) = 7$$

Step 5

$$\text{for cost}(3,5)$$

$$\begin{aligned} \text{cost}(3,5) &= \min \{ c(5,7) + \text{cost}(7,4) + \text{cost}(4,3), \\ &\quad c(5,8) + \text{cost}(8,4) + \text{cost}(4,3) \} \\ &= \min \{ 1+7, 2+3 \} = 5 \end{aligned}$$

$$\text{cost}(3,5) = 5$$

| | | | | | | | | | |
|----------|---|---|---|---|---|---|---|---|---|
| v | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| cost | | | | 7 | 5 | | 7 | 3 | 0 |
| distance | | | | 8 | 8 | | 9 | 9 | 9 |

Step 6

$$\text{for } \text{cost}(3,6) = \min \{ c(6,7) + \text{cost}(7,8), \\ c(6,8) + \text{cost}(9,8) \} \\ = \min \{ 6+7, 9+3 \} = 5 \\ \text{cost}(3,6) = 5$$

Step 7

$$\text{for } \text{cost}(2,2) = \min \{ c(2,4) + \text{cost}(4,3,4), \\ c(2,5) + \text{cost}(7,5), \\ c(2,6) + \text{cost}(3,6) \} \\ = \min \{ 3+7, 2+5, 3+5 \} \\ = \min \{ 10, 7, 8 \} \\ \text{cost}(2,2) = 7$$

Step 8

$$\text{for } \text{cost}(2,3) = \min \{ c(3,4) + \text{cost}(3,5), \\ c(3,5) + \text{cost}(3,5), \\ c(3,6) + \text{cost}(3,6) \} \\ = \min \{ 6+7, 5+5, 2+5 \} \\ = \min \{ 13, 10, 13 \} \\ \text{cost}(2,3) = 10$$

Step 9

$$\text{for } \text{cost}(1,1) = \min \{ c(1,2) + \text{cost}(2,2), \\ c(1,2) + \text{cost}(2,3) \} \\ = \min \{ 5+7, 2+10 \} \\ = \min \{ 12, 12 \} \\ \text{cost}(1,1) = 12$$

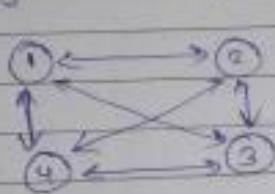
| V | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
|----------|-----|---|----|---|---|---|---|---|---|
| cost | 12 | 7 | 10 | 7 | 5 | 5 | 7 | 3 | 0 |
| distance | 2/9 | 5 | 5 | 2 | 2 | 2 | 3 | 9 | 9 |

Here we get optimum minimum cost 12 but here we get this value for 2 different paths.

$$\boxed{\text{cost} = 12}$$

path 1 :- $1 \rightarrow 2 \rightarrow 5 \leftarrow 8 \rightarrow 9$
 path 2 :- $1 \rightarrow 3 \rightarrow 5 \rightarrow 8 \rightarrow 9$

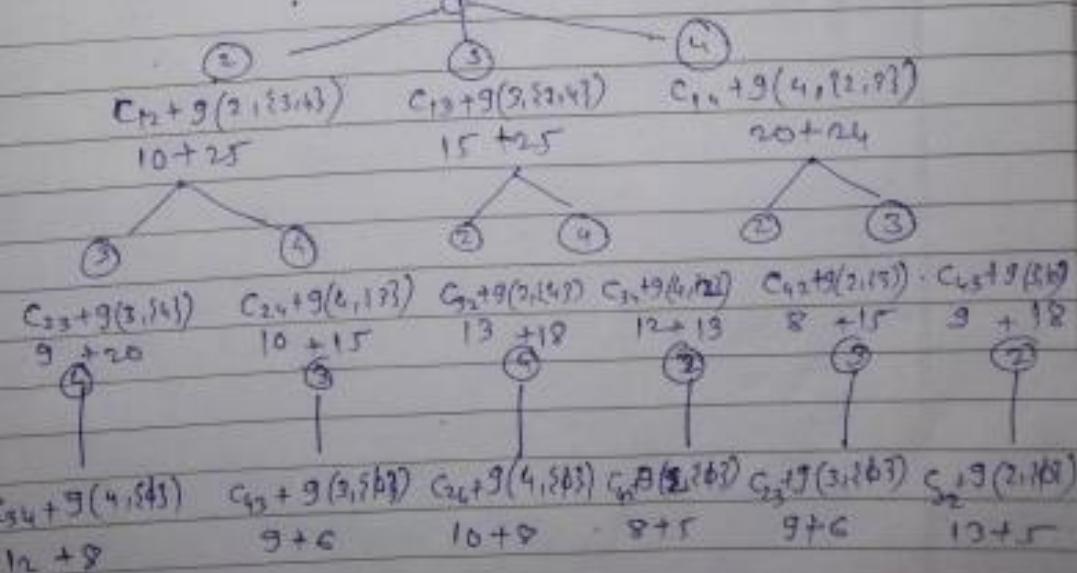
Consider the directed graph of fig. with edge lengths are given by matrix using TVSP



| | 1 | 2 | 3 | 4 |
|---|---|----|----|----|
| 1 | 0 | 10 | 15 | 20 |
| 2 | 5 | 0 | 9 | 10 |
| 3 | 6 | 10 | 0 | 12 |
| 4 | 8 | 8 | 9 | 0 |

formula $\rightarrow g(i,j) = \min_{k \in S} \{ C_{ik} + g(k, j - \{ k \}) \}$

$$g(1, \{2, 3, 4\}) = 35$$



from tree diagram.

$$g(2, \emptyset) = 5$$

$$g(3, \emptyset) = 6$$

$$g(4, \emptyset) = 8$$

$$g(2, \{3\}) = 15$$

$$g(2, \{4\}) = 18$$

$$g(3, \{2\}) = 18$$

$$g(3, \{4\}) = 20$$

$$g(4, \{2\}) = 13$$

$$g(4, \{3\}) = 15$$

$$g(2, \{3, 4\}) = 25$$

$$g(3, \{2, 4\}) = 25$$

$$g(4, \{2, 3\}) = 23$$

$$g(1, \{2, 3, 4\}) = \underline{35}$$

so minimum cost of TUSP = 35