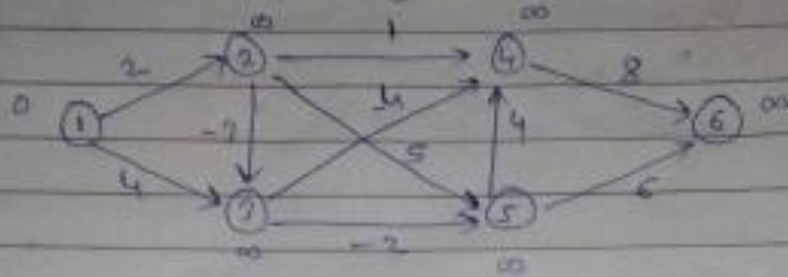


Tutorial - 7

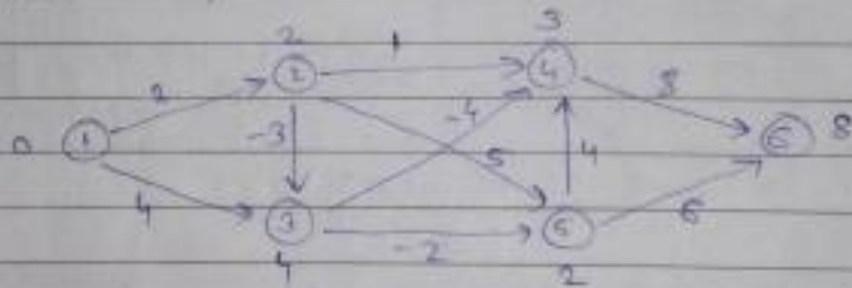
Q. find the shortest path from node 1 to every other node in the graph of fig. using the Bellman find Algorithm.



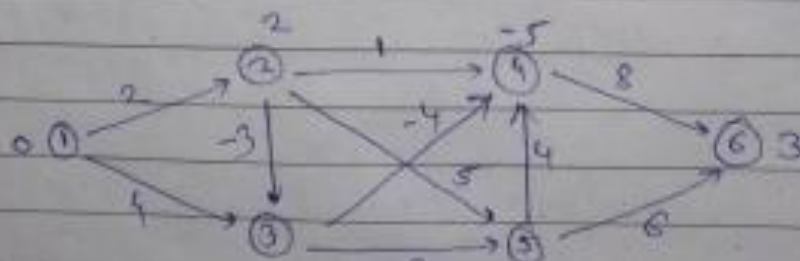
→ In Dp, we try all possible solution and pick up the best solution.
we repeat the relaxation process for $(|V|-1)$ times.

$$|V| = 6 \quad \therefore \text{repetition} = 6 - 1 = 5 \text{ times.}$$

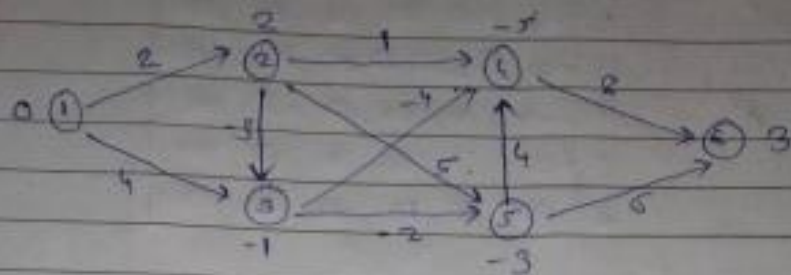
first time,



second time,



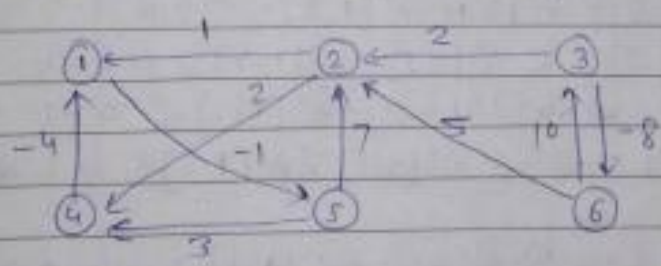
third time,



here we can see that after second step, values are not changing so this is our final solution.

distance	0	2	-1	-5	-3	3
node	1	2	3	4	5	6

q. find the shortest path using floyd-warshall algo.



→ here we have to find shortest distance from every vertex to other vertex.

$A^0 =$

	1	2	3	4	5	6
1	0	∞	∞	∞	-1	∞
2	1	0	∞	2	∞	∞
3	∞	2	0	∞	∞	-8
4	-4	∞	∞	0	∞	∞
5	∞	7	∞	3	0	∞
6	∞	5	10	∞	∞	0

	1	2	3	4	5	6
1		0	∞	∞	-1	∞
2	1		0	∞	0	∞
3	∞	2		0	∞	-8
4	-4	∞	∞		0	-5
5	∞	7	∞	3		0
6	∞	5	10	∞	∞	

$$- A^0[2,3] = A^0[2,1] + A^0[1,3]$$

$$\infty < 1 + \infty = \infty$$

$$- A^0[2,4] = A^0[2,1] + A^0[1,4]$$

$$2 < 1 + \infty = \infty$$

$$- A^0[2,5] = A^0[2,1] + A^0[1,5]$$

$$\infty > 1 + (-1) = 0$$

$$- A^0[2,6] = A^0[2,1] + A^0[1,6]$$

$$\infty < 1 + (\infty) = \infty$$

$$A^0[3,2] = 2 < A^0[3,1] + A^0[1,2] = \infty + \infty$$

$$A^0[3,4] = \infty < A^0[3,1] + A^0[1,4] = \infty + \infty$$

$$A^0[3,5] = \infty > A^0[3,1] + A^0[1,5] = \infty - 1$$

$$A^0[3,6] = -8 < A^0[3,1] + A^0[1,6] = \infty + \infty$$

$$A^0[4,2] = \infty > A^0[4,1] + A^0[1,2] = -4 + \infty$$

$$A^0[4,3] = \infty > A^0[4,1] + A^0[1,3] = -4 + \infty$$

$$A^0[4,5] = \infty > A^0[4,1] + A^0[1,5] = -4 - 1$$

$$A^0[4,6] = \infty > A^0[4,1] + A^0[1,6] = -4 + \infty$$

$$A^0[5,2] = 7 < A^0[5,1] + A^0[1,2] = \infty + \infty$$

$$A^0[5,3] = \infty < A^0[5,1] + A^0[1,3] = \infty + \infty$$

$$A^0[5,4] = 3 < A^0[5,1] + A^0[1,4] = \infty + \infty$$

$$A^0[5,6] = \infty < A^0[5,1] + A^0[1,6] = \infty + \infty$$

$$\begin{aligned}
 A^0[G, 2] &= 5 < A^0[G, 1] + A^0[1, 2] = \infty + \infty \\
 A^0[G, 3] &= 10 < A^0[G, 1] + A^0[1, 3] = \infty + \infty \\
 A^0[G, 4] &= \infty < A^0[G, 1] + A^0[1, 4] = \infty + \infty \\
 A^0[G, 5] &= \infty > A^0[G, 1] + A^0[1, 5] = \infty - 1
 \end{aligned}$$

similarly,

$$A^2 = \begin{array}{c|cccccc} & 1 & 2 & 3 & 4 & 5 & 6 \\ \hline 1 & 0 & \infty & \infty & \infty & -1 & \infty \\ 2 & 1 & 0 & \infty & 2 & 0 & \infty \\ 3 & 3 & 2 & 0 & 4 & 2 & -8 \\ 4 & -4 & \infty & \infty & 0 & -5 & \infty \\ 5 & 8 & 7 & \infty & 3 & 0 & \infty \\ 6 & 6 & 5 & 10 & 7 & 5 & 0 \end{array}$$

$$A^3 = \begin{array}{c|cccccc} & 1 & 2 & 3 & 4 & 5 & 6 \\ \hline 1 & 0 & \infty & \infty & \infty & -1 & \infty \\ 2 & 1 & 0 & \infty & 2 & 0 & \infty \\ 3 & 3 & 2 & 0 & 4 & 2 & -8 \\ 4 & -4 & \infty & \infty & 0 & -5 & \infty \\ 5 & 8 & 7 & \infty & 3 & 0 & \infty \\ 6 & 6 & 5 & 10 & 7 & 5 & 0 \end{array}$$

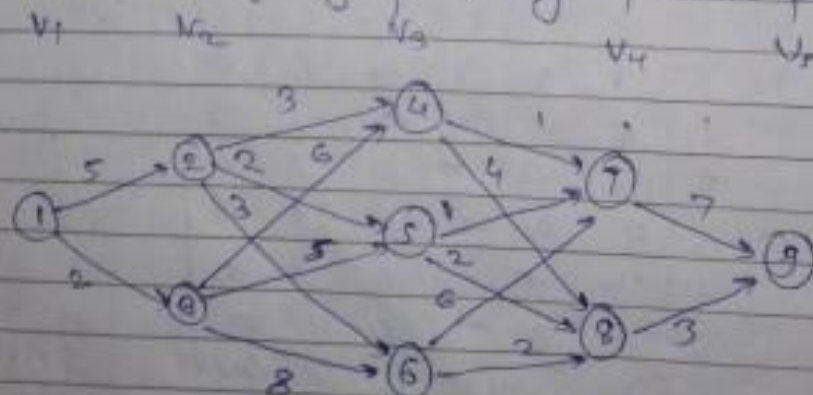
$$A^4 = \begin{array}{c|cccccc} & 1 & 2 & 3 & 4 & 5 & 6 \\ \hline 1 & 0 & \infty & \infty & \infty & -1 & \infty \\ 2 & -2 & 0 & \infty & 2 & -3 & \infty \\ 3 & 0 & 2 & 0 & 4 & -1 & -8 \\ 4 & -4 & \infty & \infty & 0 & -5 & \infty \\ 5 & -1 & 7 & \infty & 3 & 0 & \infty \\ 6 & 3 & 5 & 10 & 7 & 2 & 0 \end{array}$$

	1	2	3	4	5	6
1	0	6	∞	2	-1	∞
2	-4	0	∞	0	-3	∞
3	0	2	0	2	-1	-2
4	-6	2	∞	0	-5	∞
5	-1	7	∞	3	0	∞
6	1	5	10	5	2	0

	1	2	3	4	5	6
1	0	6	∞	2	-1	∞
2	-4	0	∞	0	-3	∞
3	-7	-3	0	-3	-6	-2
4	-6	2	∞	0	-5	∞
5	-1	7	∞	3	0	∞
6	1	5	10	5	2	0

This is shortest path from each vertex.

Q. Find the minimum cost path from s to t in multistage graph using dynamic programming.



step 1

v	1	2	3	4	5	6	7	8	9
cost									0
distance									9

step 2

for cost (4,7) = 7

step 3

for cost (4,8) = 3

v	1	2	3	4	5	6	7	8	9
cost							7	3	0
distance							9	9	9

step 4

for cost (3,4)

$$\begin{aligned} \text{cost}(3,4) &= \min \left\{ \begin{array}{l} c(4,7) + \text{cost}(4,7) \\ c(4,8) + \text{cost}(4,8) \end{array} \right\} \\ &= \min \{ 1+7, 6+3 \} \\ &= \min \{ 8, 7 \} \end{aligned}$$

$$\text{cost}(3,4) = 7$$

step 5

for cost (3,5)

$$\begin{aligned} \text{cost}(3,5) &= \min \left\{ \begin{array}{l} c(5,7) + \text{cost}(5,7) \\ c(5,8) + \text{cost}(5,8) \end{array} \right\} \\ &= \min \{ 1+7, 2+3 \} = 5 \end{aligned}$$

$$\text{cost}(3,5) = 5$$

v	1	2	3	4	5	6	7	8	9
cost				7	5		7	3	0
distance				8	8		9	9	9

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Step 6
for $\text{cost}(3,6) = \min \{ c(6,7) + \text{cost}(4,7), c(6,8) + \text{cost}(4,8) \}$
 $= \min \{ 6+7, 9+3 \} = 5$
 $\text{cost}(3,6) = 5$

Step 7
for $\text{cost}(2,2) = \min \{ c(2,4) + \text{cost}(3,4), c(2,5) + \text{cost}(3,5), c(2,6) + \text{cost}(3,6) \}$
 $= \min \{ 3+7, 2+5, 3+5 \}$
 $= \min \{ 10, 7, 8 \}$
 $\text{cost}(2,2) = 7$

Step 8
for $\text{cost}(2,3) = \min \{ c(3,4) + \text{cost}(3,4), c(3,5) + \text{cost}(3,5), c(3,6) + \text{cost}(3,6) \}$
 $= \min \{ 6+7, 5+5, 8+5 \}$
 $= \min \{ 13, 10, 13 \}$
 $\text{cost}(2,3) = 10$

Step 9
for $\text{cost}(1,1) = \min \{ c(1,2) + \text{cost}(2,2), c(1,3) + \text{cost}(2,3) \}$
 $= \min \{ 5+7, 2+10 \}$
 $= \min \{ 12, 12 \}$
 $\text{cost}(1,1) = 12$

V	1	2	3	4	5	6	7	8	9
cost	12	7	10	7	5	5	7	3	0
distance	2/9	5	5	2	2	2	3	3	9

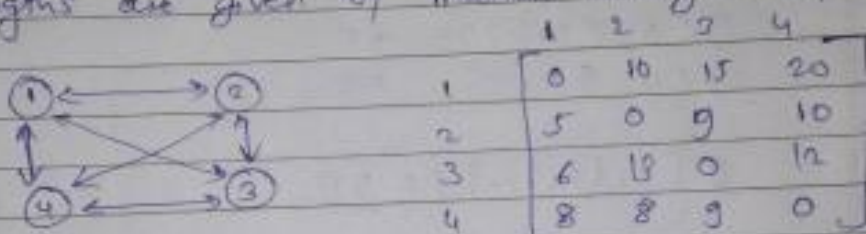
Here we get optimum minimum cost 12 but here we get this value for 2 different path.

$$\text{cost} = 12$$

path 1 :- $1 \rightarrow 2 \rightarrow 5 \rightarrow 8 \rightarrow 9$

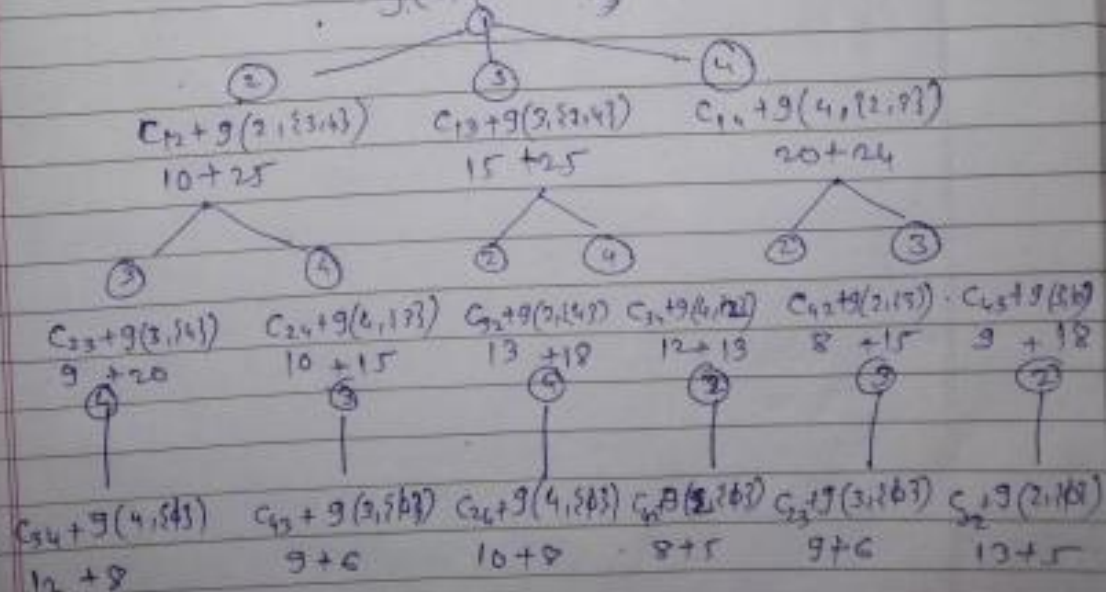
path 2 :- $1 \rightarrow 3 \rightarrow 5 \rightarrow 8 \rightarrow 9$

Consider the directed graph of fig. with edge lengths are given by matrix using T.V.S.P



$$\text{formula} \rightarrow g(i, S) = \min_{k \in S} \{ C_{ik} + g(k, S - \{k\}) \}$$

$$g(1, \{2, 3, 4\}) = 35$$



from the diagram,

$$g(2, \emptyset) = 5$$

$$g(3, \emptyset) = 6$$

$$g(4, \emptyset) = 8$$

$$g(2, \{3\}) = 15$$

$$g(2, \{4\}) = 18$$

$$g(3, \{2\}) = 18$$

$$g(3, \{4\}) = 20$$

$$g(4, \{2\}) = 19$$

$$g(4, \{3\}) = 15$$

$$g(2, \{3, 4\}) = 25$$

$$g(3, \{2, 4\}) = 25$$

$$g(4, \{2, 3\}) = 23$$

$$g(1, \{2, 3, 4\}) = \underline{35}$$

So minimum cost of TSP = 35