Spanning trees

Given g = (V, E) with |V| = n, the number of spanning trees of $G = n^{n-2}$.

- a spanning tree has |v|-1 edges.
- adding an edge creates an unique cycle

Every subgraph with n-1 edges and no cycles is a spanning tree.

Aim: to sample uniformly from the set of spanning trees of G.

- 1. Let X_t be a spanning tree S_1 . Choose any edge e from E $E(S_1)$ uniformly at random.
- 2. Let C be a unique cycle in $X_t \cup e$
- 3. Choose e' from E(C) uniformly.
- 4. Set $X_{t+1} = (X_t \{e'\}) \cup e$

Exercise: Write the transition matrix of this chain and show that it is irreducible, aperiodic and symmetric.

Knapsack Problems

m items each with value v_i and weight w_i for 1 < i < n. Find S = i, st $\arg \max_{\{i,j\}} \sum_i v_i = S$ and $\sum_i w_i < C$.

Space of feasible solutions $S = z : \langle w, z \rangle \leq C$

- Start at some $z \in S$
- Choose $j \in \{1, ..n\}$ uniformly
- $y = (z_1, z_2, \dots (1 z_j) \dots, z_m).$
- If $\langle w, y \rangle \leq C$,
 - $-X_{t+1} = y$
 - otherwise $X_{t+1} = z$

However, the walk is more efficient if we instead sample j $\pi(j) \propto e^{\beta \langle v, z \rangle}$ for some $\beta > 0$.

Metropolis algorithm

S is a discrete space. Q a symmetric transition matrix on S and π , the target distn.

Given X_t, X_{t+1} is chosen to follow MC,

- 1. Choose Y randomly according to Q, called the proposal. $P(Y = j | X_t = i) = q_{ij}$
- 2. Define the acceptance probability $\alpha = \min\{1, \pi_Y/\pi_i\}$
- 3. Let U ~ uniform(0,1). If $U \leq alpha$ then $X_{t+1} = Y$. Otherwise $X_{t+1} = X_t$

 $P_{ij} = q_{ij} * \alpha$

Proposition: If Q is irreducible, symmetric and π a prob distn. st $\pi > 0 \forall i \in S$, then the Metropolis chain is also irr, and reversible wrt π .

Proof:

$$Q \text{ is irr} \implies P \text{ irr}$$

$$\pi_{i}p_{ij} = \pi_{i}q_{ij}\alpha = q_{ij}\min\{\pi_{i}, \pi_{j}\} = \pi_{j}q_{ij}\min\{1, pi_{i}/pi_{j}\} = \pi_{j}p_{ji}$$