

3. Solution:

Edges of graph

$$A - B \rightarrow 7$$

$$B - C \rightarrow 8$$

$$D - E \rightarrow 15$$

$$F \Rightarrow 9 \rightarrow 11$$

$$A - D \rightarrow 5$$

$$B - E \rightarrow 7$$

$$D - F \rightarrow 6$$

$$E - G \rightarrow 9$$

$$B - D \rightarrow 9$$

$$C - E \rightarrow 5$$

$$E - F \rightarrow 8$$

Now arranging the edges in ascending order

$$A - D \rightarrow 5$$

$$C - E \rightarrow 5$$

$$D - F \rightarrow 6$$

$$A - B \rightarrow 7$$

$$B - E \rightarrow 7$$

$$B - C \rightarrow 8$$

$$E - F \rightarrow 8$$

$$E - G \rightarrow 9$$

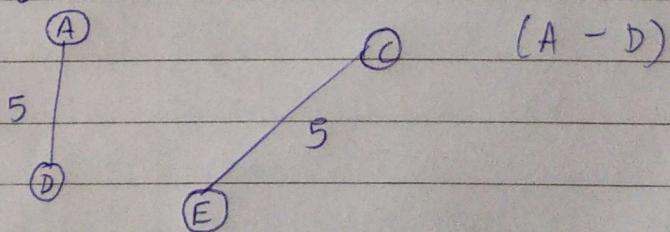
$$B - D \rightarrow 9$$

$$F - G \rightarrow 11$$

$$D - E \rightarrow 15$$

Now, creating a graph with edges having minimum length of no circuit forms and all the edges are included.

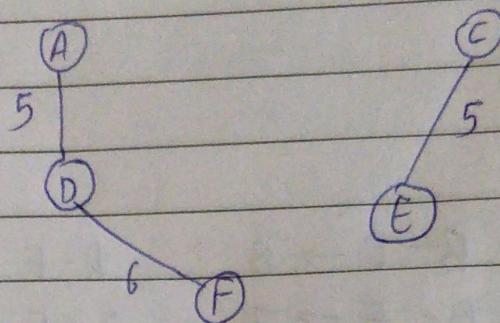
Step 1:



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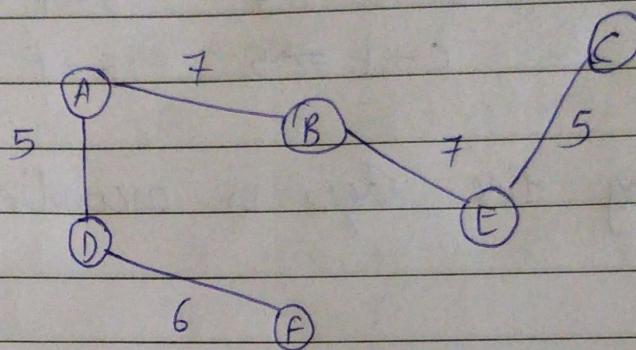
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Step 2:



(C - E)

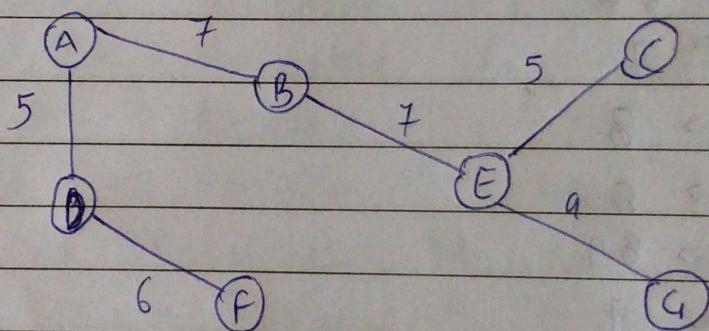
Step 3:



(D - F)

(B - C)

Step 4:

(E - F)  
(E - G)

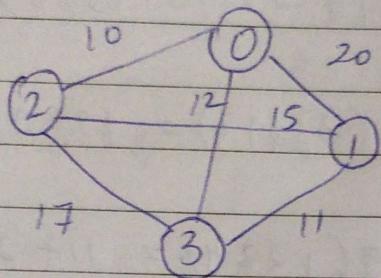
$$\text{Min. cost of the graph} = 5 + 6 + 7 + 7 + 5 + 9 \\ = 39.$$

4. Using formula

$$g(1, V - \{1\}) = \min_{2 \leq k \leq n} \{c_{ik} + g(k, V - \{1, k\})\}$$

for  $i \in S$

$$g(i, S) = \min \{c_{ij} + g(j, S - \{j\})\} \quad (j \in S)$$



$$\begin{matrix} & 0 & 1 & 2 & 3 \\ 0 & 0 & 20 & 10 & 12 \\ 1 & 20 & 0 & 15 & 11 \\ 2 & 10 & 15 & 0 & 17 \\ 3 & 12 & 11 & 17 & 0 \end{matrix}$$

Iteration 1:

$$g(1, \emptyset) = c_{10} = 20$$

$$g(2, \emptyset) = c_{20} = 10$$

$$g(3, \emptyset) = c_{30} = 12$$

Iteration 2: (intermediate)

$$g(1, \{2\}) = c_{12} + g(2, \emptyset) = 15 + 10 = 25$$

$$g(1, \{3\}) = c_{13} + g(3, \emptyset) = 11 + 12 = 23$$

$$g(2, \{1\}) = c_{21} + g(1, \emptyset) = 15 + 20 = 35$$

$$g(2, \{3\}) = c_{32} + g(3, \emptyset) = 17 + 12 = 29$$

$$g(3, \{1\}) = c_{31} + g(1, \emptyset) = 11 + 20 = 31$$

$$g(3, \{2\}) = c_{23} + g(2, \emptyset) = 17 + 10 = 27$$

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Iteration 3: (2<sup>nd</sup> intermediate)

$$g(1, \{2, 3\}) = \begin{cases} C_{12} + g(2, \{3\}) = 15 + 29 = 44 \\ g(3, \{2\}) \\ C_{13} + g(\cancel{3}, \{2\}) = 11 + 27 = 38 \quad \checkmark \min \end{cases}$$

$$g(2, \{1, 3\}) = \begin{cases} C_{21} + g(1, \{3\}) = 15 + 23 = 38 \quad \checkmark \min \\ C_{23} + g(3, \{1\}) = 17 + 31 = 48 \end{cases}$$

$$g(3, \{1, 2\}) = \begin{cases} C_{31} + g(1, \{2\}) = 11 + 25 = 36 \quad \checkmark \min \\ C_{32} + g(2, \{1\}) = 17 + 35 = 52 \end{cases}$$

Iteration 4:

$$g(0, \{1, 2, 3\}) = C_{01} + g(1, \{2, 3\}) = 20 + 38 + \cancel{58} = 58$$

$$C_{02} + g(2, \{1, 3\}) = 10 + 38 = 48$$

$$C_{03} + g(3, \{1, 2\}) = 12 + 36 = 48$$

∴ Min cost is 48.

5. Using finite automata,

$$P = RSTP$$

$$T = PST P S Q S T R R S T P Q P$$

O :

Define state : R :

RS :

RST :

RSTP :

Language : {P, Q, R, S, T}

	P	Q	R	S	T
O	0	0	1	0	0
R	0	0	1	2	0
RS	0	0	1	0	3
RST	4	0	1	0	0
RSTP	0	0	1	0	0

T : P S T P S Q S T R R S T P Q P

0 0 0 0 0 0 0 0 0 1 2 3 4

~~Present~~ Present

Hence we reach the final state 4 once.

∴ pattern RSTP is present in the text only one time.