

School of Engineering and Applied Science (SEAS), Ahmedabad University

CSE 400: Fundamentals of Probability in Computing

Lecture 10: Randomized Min-Cut Algorithm

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Topic: Randomized Min-Cut Algorithm

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Scribe: Lecture Summary

1 Outline

This lecture covers the fundamental concepts of the Min-Cut problem, comparing deterministic and randomized approaches. Key topics include:

- The Min-Cut Problem: Definition and Applications.
- Max-Flow Min-Cut Theorem.
- Deterministic Min-Cut (Stoer-Wagner Algorithm).
- Randomized Min-Cut (Karger's Algorithm).
- Performance Comparison and Success Probabilities.

2 The Min-Cut Problem

2.1 Why use Min-Cut?

Min-cut algorithms are utilized in various domains to solve problems related to network connectivity, reliability, and optimization:

- **Network Design:** Improves communication efficiency and optimizes network flow by finding the minimum capacity cut.
- **Communication Networks:** Helps understand network vulnerability to failures and aids in building robust, fault-tolerant systems.
- **VLSI Design:** Useful for partitioning circuits into smaller components to reduce interconnectivity complexity.

2.2 What is a Min-Cut?

- **Cut-set:** A set of edges whose removal breaks a graph into two or more connected components.
- **Min-cut Problem:** Given a graph $G = (V, E)$ with n vertices, find a minimum cardinality cut-set in G .

2.3 Edge Contraction

The primary operation in randomized min-cut algorithms is **edge contraction**:

- It removes an edge (u, v) and merges vertices u and v into a single vertex.
- All edges connecting u and v are eliminated.
- All other edges are retained; the resulting graph may have parallel edges but no self-loops.

3 Success and Failure in Min-Cut Runs

- **Successful Run:** An iteration of the algorithm that correctly identifies the minimum cut of the graph.
- **Unsuccessful Run:** An iteration where the algorithm fails to identify the minimum cut. This can happen if the algorithm happens to contract critical edges (edges belonging to the min-cut) early in the process.

4 Max-Flow Min-Cut Theorem

The theorem states: “In a flow network, the maximum amount of flow passing from the source (S) to the sink (T) is equal to the total weight of the edges in a minimum cut.”

- **Capacity of a Cut:** The sum of the capacity of edges oriented from a vertex in set X to a vertex in set Y .
- **Minimum Cut:** The cut in the network with the smallest possible capacity.
- **Maximum Flow:** The largest possible flow from source S to sink T .

5 Deterministic Min-Cut Algorithm

5.1 Stoer-Wagner Algorithm

This algorithm provides an exact solution. The core theorem states that for two vertices s and t , a minimum cut of G is the smaller of:

1. A minimum $s - t$ -cut of G .
2. A minimum cut of $G/\{s, t\}$ (the graph where s and t are merged).

Complexity: $O(V \cdot E + V^2 \log V)$.

6 Randomized Min-Cut Algorithm

6.1 Why Randomized Algorithms?

- Provide a probabilistic guarantee of success.
- Often provide accurate estimates with fewer iterations.
- Advantages: Efficiency, Parallelization, Approximation Guarantees, Robustness, and avoidance of worst-case instances.

6.2 Karger’s Randomized Algorithm

Karger’s algorithm uses random contraction steps to find a cut.

Theorem for Min-Cut Set: The algorithm outputs a min-cut set with probability at least:

$$P(\text{success}) \geq \frac{2}{n(n-1)}$$

Complexity: $O(V^2)$.

7 Comparison: Deterministic vs. Randomized

Feature	Deterministic (Stoer-Wagner)	Randomized (Karger's)
Guarantee	Always finds exact minimum cut	Approximate min-cut with high probability
Complexity	$O(V \cdot E + V^2 \log V)$	$O(V^2)$
Scalability	Higher complexity for large graphs	Generally more efficient for large graphs

8 Python Simulation

A simulation of the Randomized Min-Cut algorithm can be found in the lecture 10 Campuswire post (File: L8_RandomizedMin...ipynb).

End of Scribe