

Asymptotic Notations

Asymptotic notations give us an idea about how good a given algorithm is compared to some other algorithm.

Let us see the mathematical definition of "order of" now.

Primarily there are three types of widely used asymptotic notations.

1. Big Oh notation (O)
2. Big Omega notation (Ω)
3. Big theta notation (Θ) \rightarrow Widely used one!

Big Oh notation

Big Oh notation is used to describe asymptotic upper bound.

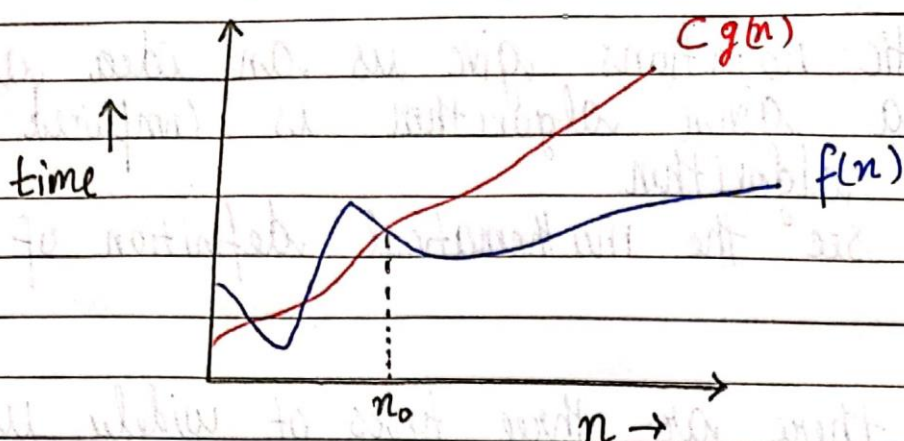
Mathematically, if $f(n)$ describes running time of an algorithm; $f(n)$ is $O(g(n))$ iff there exist positive constants C and n_0 such that $f(n) \leq Cg(n)$ for all $n \geq n_0$. (n-not)

$$0 \leq f(n) \leq Cg(n) \text{ for all } n \geq n_0$$

if a function is $O(n)$, it is automatically $O(n^2)$ as well!

\Downarrow
used to give upper bound on a function.

Graphic example for Big Oh (O)



Big Omega notation

Just like O notation provides an asymptotic upper bound, Ω notation provides asymptotic lower bound. Let $f(n)$ define running time of an algorithm;

$f(n)$ is said to be $\Omega(g(n))$ if there exists positive constants C and n_0 such that

$$0 \leq Cg(n) \leq f(n) \text{ for all } n \geq n_0$$

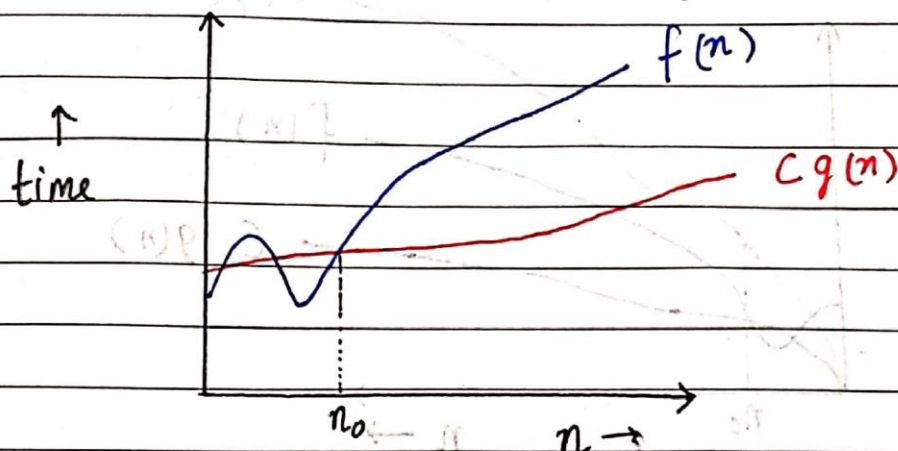


used to give lower bound on a function

if a function is $O(n^2)$ it is automatically $O(n)$ as well.



Graphic example for Big omega (Ω)



Big theta notation

Let $f(n)$ define running time of an algorithm

$f(n)$ is said to be $\Theta(g(n))$ iff $f(n)$ is $O(g(n))$ and $f(n)$ is $\Omega(g(n))$

Mathematically,

$$0 \leq f(n) \leq C_1 g(n) \quad \forall n \geq n_0 \rightarrow \text{Sufficiently large value of } n$$

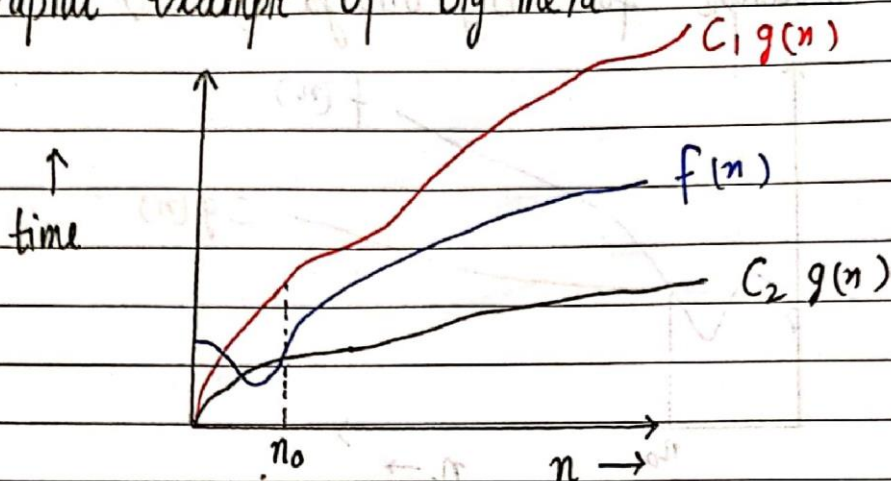
$$0 \leq C_2 g(n) \leq f(n) \quad \forall n \geq n_0 \rightarrow \text{Sufficiently large value of } n$$

Merging both the equations, we get:

$$0 \leq C_2 g(n) \leq f(n) \leq C_1 g(n) \quad \forall n \geq n_0$$

The equation simply means there exist positive constants C_1 and C_2 such that $f(n)$ is sandwiched between $C_2 g(n)$ and $C_1 g(n)$

Graphic example of Big theta



Which one of these to use?

Since Big theta gives a better picture of runtime for a given algorithm, most of the interviewers expect you to provide an answer in terms of Big theta when they say "Order of".

Quick Quiz ÷ Prove that $n^2 + n + 1$ is $O(n^3)$, $\Omega(n^2)$ and $\Theta(n^2)$ using respective definitions.

Increasing order of Common runtimes

$$1 < \log n < n < n \log n < n^2 < n^3 < 2^n < n^n$$

Better

Worse



Common runtimes from
better to worse

