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## Lecture-09: More on BSTs and Midterm #1 (Hand Out)

### ⊙ How to add (y)?

- use find like procedure to see if  $y$  is already in tree
- if not, let  $v$  be the last node searched. (node  $v$  has at most 1 child)
- if  $y < v.x$ , make new node for  $y$  and set as left child of  $x$ .
- if  $y > v.x$ , make new node for  $y$  and set  $y$  as right child of  $x$
- how long does this take?

$$- \quad \mathcal{O}(\text{depth of tree}) = \mathcal{O}(\text{height of root})$$

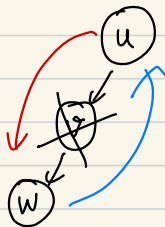
### ⊙ How to remove (y)?

- first find node  $v$  storing  $y$
- 3 cases:
  - (1)  $v$  is leaf
  - (2)  $v$  has one child
  - (3)  $v$  has 2 children

Case 1: If  $v == v.\text{parent}.\text{left}$ , set  $v.\text{parent}.\text{left} = \text{null}$   
If  $v == v.\text{parent}.\text{right}$ , set  $v.\text{parent}.\text{right} = \text{null}$

Case 2 (1 child): Assume  $v.\text{left} \neq \text{null}$ ,  $v.\text{right} == \text{null}$   
set  $v.\text{left}.\text{parent} = v.\text{parent}$  [sim. other case]

Assume  $v == v.\text{parent}.\text{left}$  [sim. other case]  
set  $v.\text{parent}.\text{left} = v.\text{left}$



Case 3 on next page

### Case 3 (2 children)

{ choice: replace the value of  $v$  with smallest ~~no~~ value larger than  $v$ .

→ What node  $w$  stores this value?

- go to right child of  $v$
- go to left descendants until no left left.  
→ this node is  $w$ .

- copy the value  $w$  to  $v$  ( $v.n = w.n$ )

- remove the node  $w$  (at most 1 child) → either case 1/2

- how long does this take?

$$- O(\text{depth of tree}) = O(\text{height of root})$$

### Questions:

1. how many element/nodes can a binary tree of height  $h$  store?  $(2^{h+1} - 1)$

2. Min Height of BT with nodes?  $(\log(n+1) - 1)$

3. What is max height of BT with  $n$ -nodes?  
 $(n-1)$

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