

COSC-211

lecture-11

A stylized handwritten signature in black ink, featuring a large 'M' and a long, sweeping flourish.

- Dhruv
Marani

⇒ Accessing Elements in BST

◉ Helpful tool : print contents of BST in sorted order

1) One approach:

- find smallest element (left most)
- next smallest element is either the parent or smallest descendent of right child.

2) Alternative approach:

- start at root
- ① print everything in left subtree (I)
- ②. print root
- ③ print everything in right subtree (II)

Subroutine: starting @ v:

- print left subtree
- print v
- print right subtree

← JAVA CODE

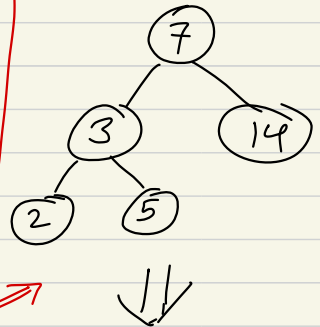
```

class Node<E> {
    // define it first.
    void printInOrder() {
        if (left != null) {
            left.printInOrder();
        }
        System.out.println(this.value);
        if (right != null) {
            right.printInOrder();
        }
    }
}
    
```

Annotations in image:

- Red arrows pointing to `left.printInOrder();` and `right.printInOrder();` are labeled `print "I"` and `print "J"` respectively.
- Red arrows pointing to `System.out.println(this.value);` are labeled `print "I"` and `print "J"` respectively.
- Red text below: `Keeping track of depth of tree.`

eg:



Example:

CALL STACK: ⇒ (2) 3 (5) 7 (14)

Call's	7.pio	3.pio 7.pio	2.pio 3.pio 7.pio	3.pio 7.pio	5.pio 3.pio 7.pio	3.pio 7.pio	7.pio	14.pio 7.pio	7.pio	∅
step:	1	2	3 print 2	4 print 3	5 print 5	6	7 print 7	8 print 14	9	=

⇒ Changing the processing order

→ (in-order): $\left. \begin{array}{l} \text{process left} \\ \text{process self} \\ \text{process right} \end{array} \right\}$

→ (pre-order): $\left. \begin{array}{l} \text{process self} \\ \text{process left} \\ \text{process right} \end{array} \right\}$

→ (post-order): $\left. \begin{array}{l} \text{process left} \\ \text{process right} \\ \text{process self} \end{array} \right\}$

Q1) What would be printed in pre & post order traversals.

Q2) Write a recursive method to find height
[done in assignment].

⇒ Last Time: AVL trees "height balanced"

→ for every $\left(\begin{array}{c} v \\ \swarrow \searrow \\ u \quad w \end{array} \right)$ $h(u)$ and $h(w)$ differ by ≤ 1

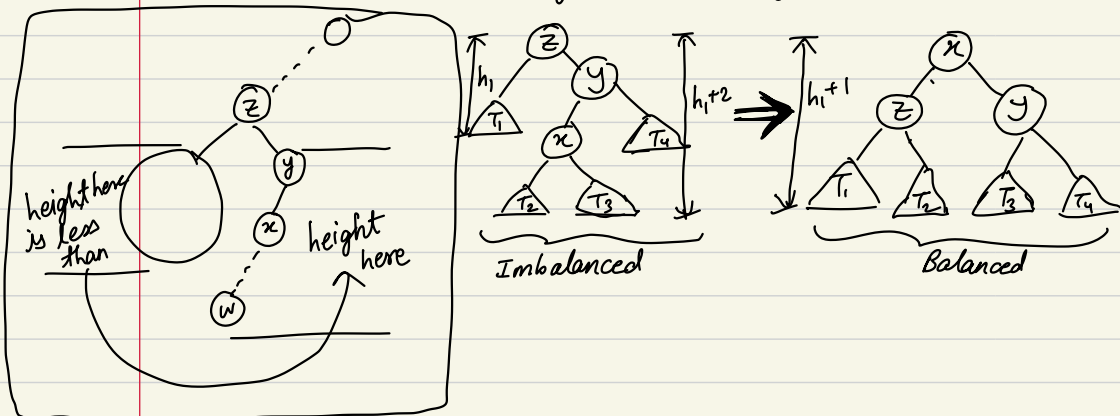
If T is AVL and with n nodes, then $h(T) = O(\log n)$

started: rules for maintaining balance on insertion and removal

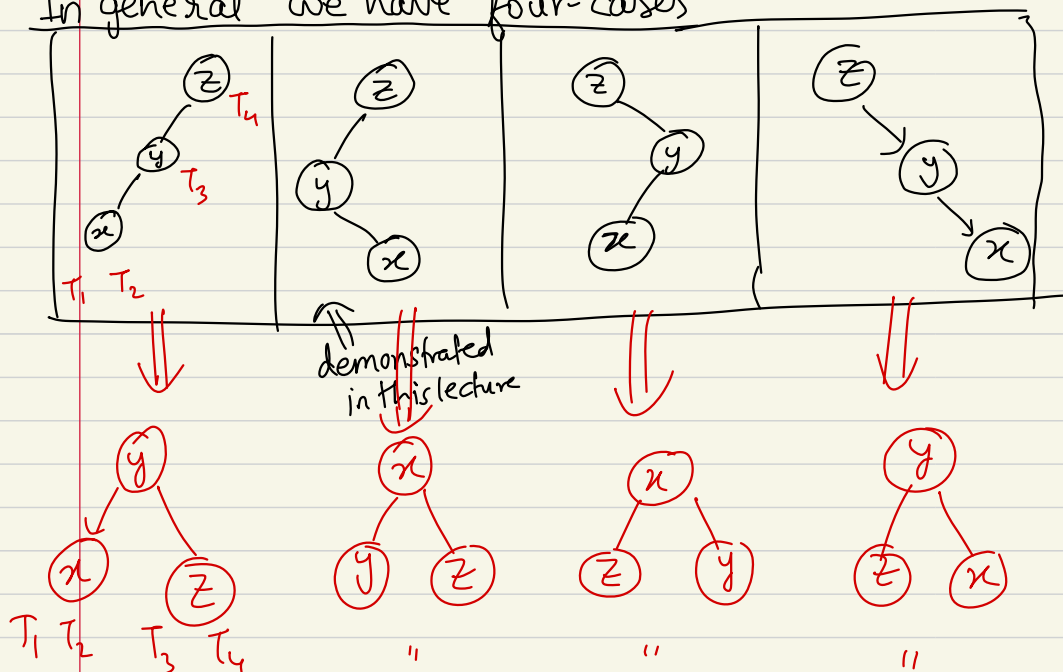
Idea for add: create new node "w"

(1) If no imbalance created, then no problem.

(2) Otherwise: z is the first ancestor of w that is imbalanced.



In general we have four-cases



⇒ How to maintain balance with node removal?