

Lecture 17: Skiplists Conclusion.

→ Question: What is the expected max height of skiplist with n nodes?

"Back of Envelope"

• n nodes, on average how many have ...

• height ≥ 0 ?

• n

• height ≥ 1

• $n/2$

• height ≥ 2

• $n/4$

• height $\geq k$?

• $n/2^k$

→ when is this < 1 ?

$$\Rightarrow n < 2^k$$

$$\Rightarrow \log n < k$$

Hint that max. height may be $O(\log n)$.

★ Formalising things

→ H = rand var = max height in list of length n .

→ For each $k = 0, 1, 2, \dots$

L_k = length of list at height k ,
nodes with height $\geq k$.

Observe: $E(L_k) = \frac{n}{2^k}$

therefore, if $k = \log n + j$. ($j \geq 0$)

$$E(L_{\log n + j}) = \frac{n}{2^{\log n + j}} = \frac{n}{2^{\log n} \cdot 2^j} = \left[\frac{1}{2^j} \right]$$

To analyse $E(H)$, we can write ~~$E(H)$~~ H as sum of simpler variables.

two facts

- 1) $J_k \leq 1$ always
- 2) $J_k \leq L_k$. (if $L_k = 0$, so is J_k)

⇒ Def. $J_k = \begin{cases} 1 & \text{if } H \geq k \\ 0 & \text{otherwise} \end{cases}$

Therefore...

$$H = J_1 + J_2 + J_3 + \dots$$

$$\Rightarrow E(H) = E(J_1 + J_2 + \dots)$$

$$\Rightarrow E(H) = E(J_1) + E(J_2) + \dots$$

$$= E(J_1) + E(J_2) + \dots + E(J_{\log n}) + \dots$$

each ≤ 1

$$E(J_{\log n + k}) \leq E(L_{\log n + k})$$

$$\leq \underbrace{1 + 1 + \dots + 1}_{\log n \text{ times}} + \underbrace{\frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots}_1$$

$$\leq \log n + 1$$

Trick: Analyse search in reverse.

Rev. dirⁿ:

from a node at height h

if I can go up, go up.

if I can't go up, go left

Each step costs $O(1)$ in find procedure

$$\text{total \# of steps} = \underbrace{\# \text{ up}}_H + \underbrace{\# \text{ left}}_{??}$$

to understand this:

Observe: probability that I can go up is $1/2$.

⇒ prob of going left is $1/2$.



\Rightarrow probability of going left many times is small

Can show (See ODS §4.4)

$$E(\text{left steps}) = \log n + O(1)$$

So, expected time to find is

$$E(H) + E(\text{left steps}) = 2(\log n) + 1.$$

GOOGLE SEARCH \Rightarrow "concentration of measure"