



如何将图模型整合到已有 关系型数据库中？



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第一部分

图数据库



图数据库

目前常见的使用图数据库场景有：

- 风控(欺诈检测/反洗钱)
- 图神经网络
- 知识图谱
- 社交网络

目前的问题

- 对于复杂的关系网络，传统关系型数据库无能为力
- 单独部署图数据库集群
- 部署运维两套数据库集群成本太高
- 在两个不同的数据库中数据一致性不能保障

探索方向

TiGraph 项目尝试验证在分布式关系型数据中无缝集成图模式：

- 同时包含关系型模型和图模型
- 同一个事务中操作图数据和关系型数据的能力
- 将图遍历作为 SQL 子查询(反之亦然)
- 在 SQL 中扩展出一个让 DBA 一眼就能学会的图遍历语法
- 对于 N 度人脉的场景性能对比

第二部分

关系型数据库

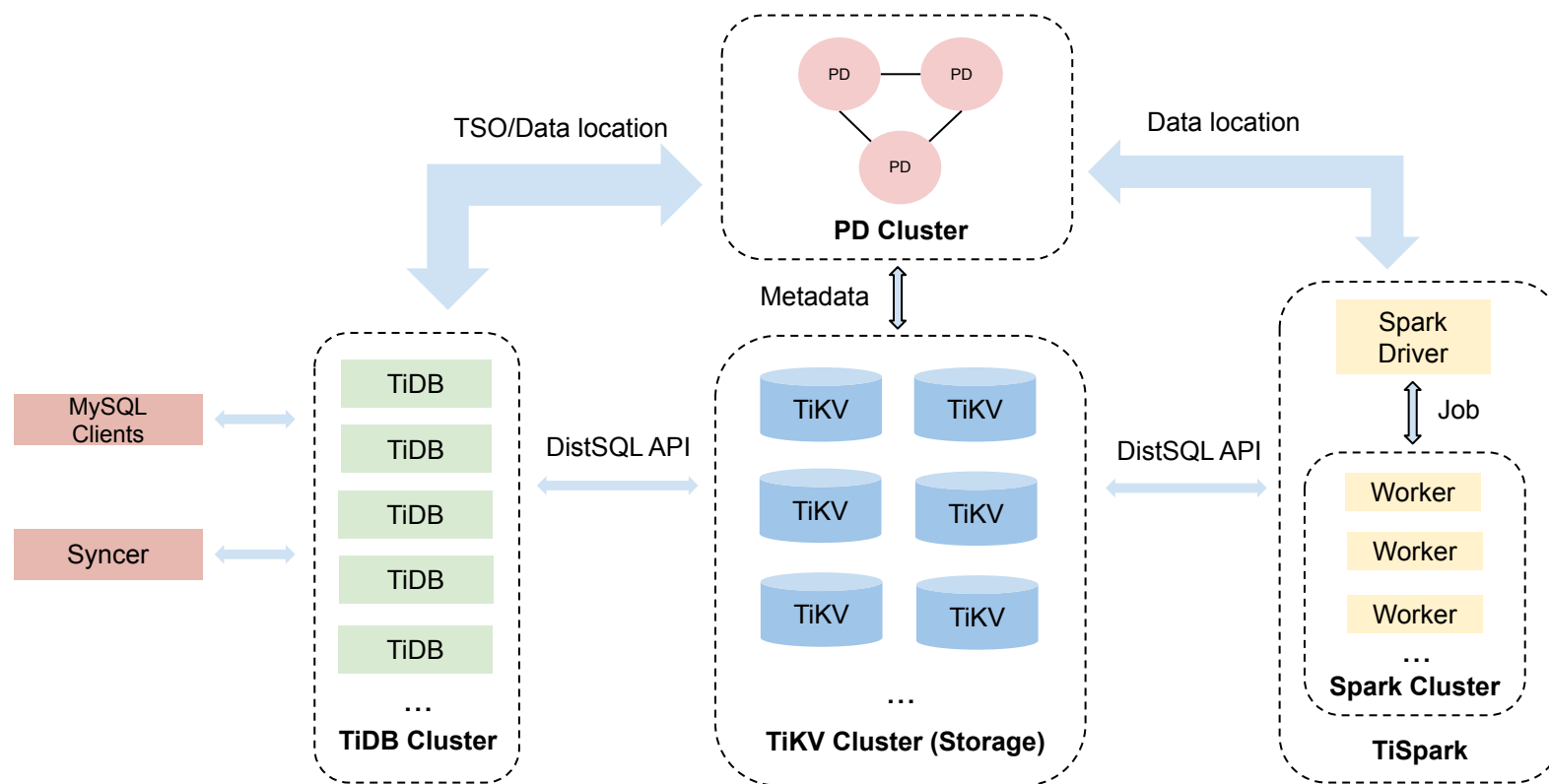


关系型数据库本质



- For a row in a Table, row data is encoded in key-value pairs with the format below:
 - t<<tableID>>_r<<rowID>> => [col1, col2, col3, col4]
- If there is secondary index with a column, the index data of a row is encoded in this way:
 - t<<tableID>>_i<<indexID>>_indexedColumnsValue => rowID
 - t<<tableID>>_i<<indexID>>_indexedColumnsValue_rowID => nil

TiDB 架构



“

《三国演义》



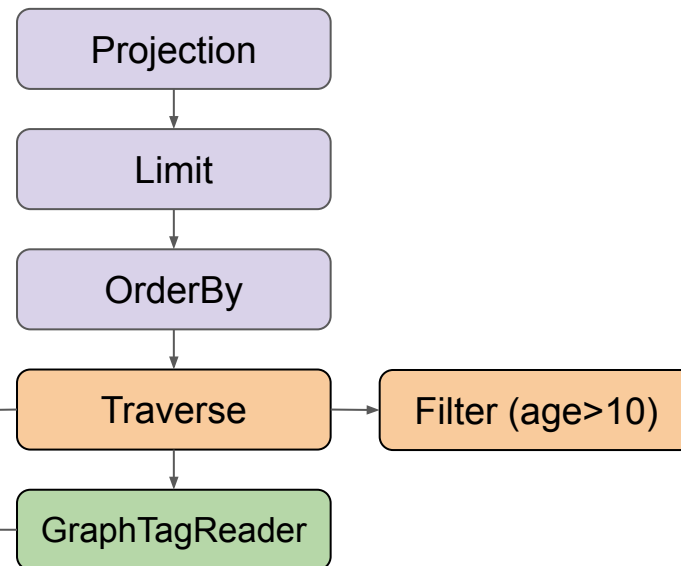
天下大势，合久必分，分久必合

图 + 关系型多模实现

```
SELECT * FROM people
  WHERE name="Henry"      # 从名字为 Henry 的顶点开始遍历
  TRAVERSE
    OUT(friend WHERE age>10). # 一度人脉 (年龄大于 10 岁的)
    OUT(friend).            # 二度人脉
    TAG(people)            # 输出这些朋友的 people 属性
  ORDER BY name           # 将这些二度人脉按照名字排序
  LIMIT 10;              # 取二度人脉的前 10 位
```

g_prefix	src_Vertex_ID	Edge_Type	Edge_ID	dst_Vertex_ID
----------	---------------	-----------	---------	---------------

g_prefix	Vertex_ID	Tag_ID
----------	-----------	--------



<https://github.com/tigraph/tidb>

第三部分

DEMO

DEMO

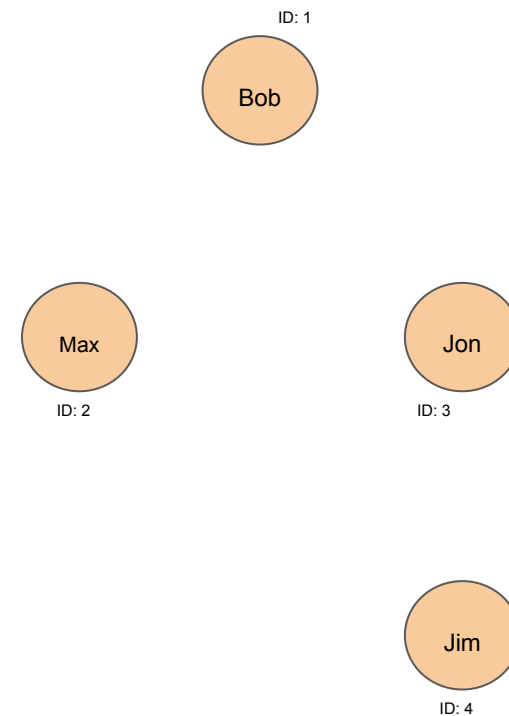
-- 创建 people 点

```
CREATE TAG people (id BIGINT, name VARCHAR(32));
```

-- 写入 4 个点

```
INSERT INTO people VALUES
```

```
(1, 'Bob'), (2, 'Max'), (3, 'Jon'), (4, 'Jim');
```



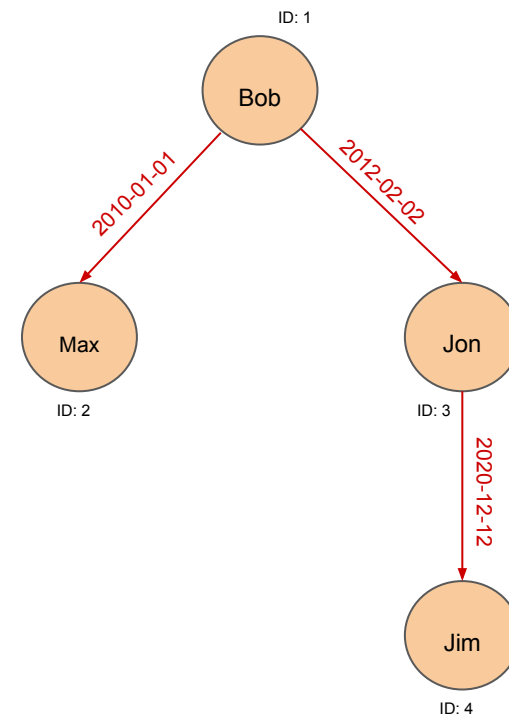
DEMO

-- 创建 follow 边

```
CREATE EDGE follow (src BIGINT, dst BIGINT, time DATE);
```

-- 写入 3 条边

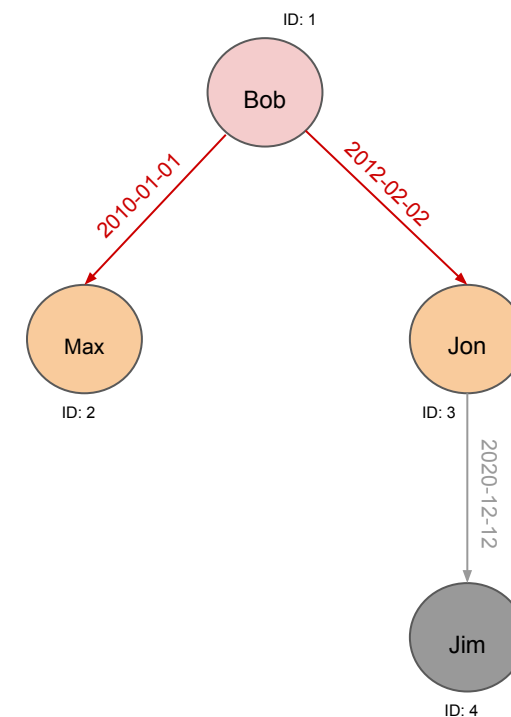
```
INSERT INTO follow VALUES (1, 2, '2010-01-01'),  
                           (1, 3, '2012-02-02'),  
                           (3, 4, '2020-12-12');
```



DEMO

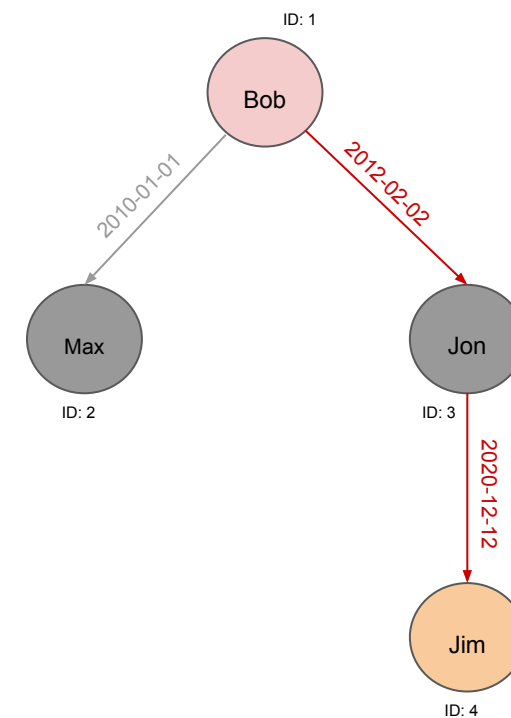
-- 查询 Bob 关注的人(Bob 的 1 度人脉)

```
SELECT * FROM people
  WHERE name = 'Bob'
  TRAVERSE
    OUT(follow).
  TAG(people);
```



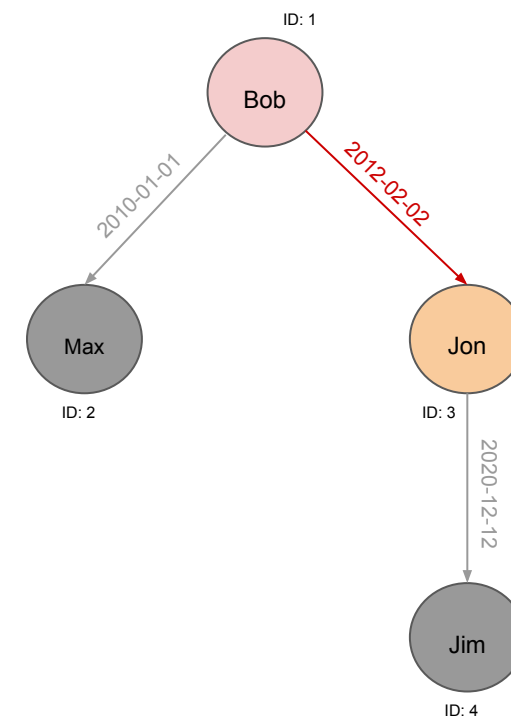
DEMO

```
-- 查询 Bob 的 2 度人脉
SELECT * FROM people
  WHERE name = 'Bob'
  TRAVERSE
    OUT(follow).
    OUT(follow).
  TAG(people);
```



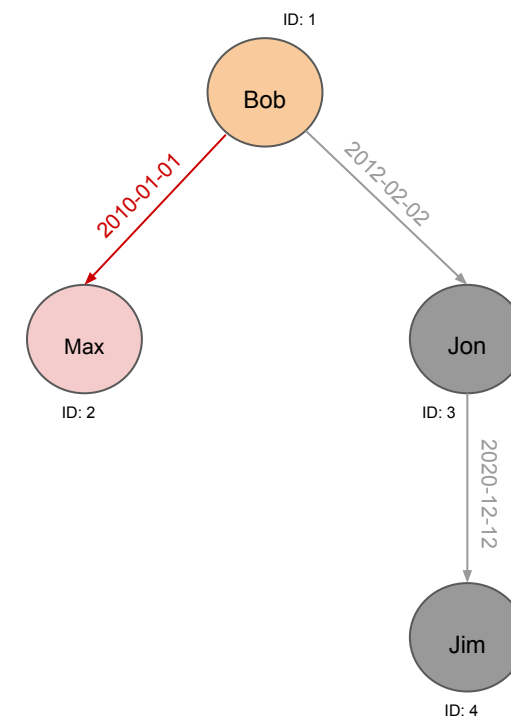
DEMO

```
-- 查询 Bob 在 2012 年关注的人
SELECT * FROM people
  WHERE name = 'Bob'
  TRAVERSE
    OUT(follow WHERE YEAR(time) = 2012).
  TAG(people);
```



DEMO

```
-- 查询关注了 Max 了的人
SELECT * FROM people
  WHERE name = 'Max'
  TRAVERSE
    IN(follow).
  TAG(people);
```



第四部分

特性

一些主要特性

- ✓ DDL
 - ✓ Create tag
 - ✓ Create edge
 - ✓ Secondary index
- ✓ DML
 - ✓ Insert
 - ✓ Update
 - ✓ Delete
 - ✓ Insert on duplicate
 - ✓ Insert ignore
 - ✓ Secondary index
- ❑ Query
 - ✓ Secondary index
 - ✓ Selection
 - ✓ Order by
 - ❑ Group by

- ✓ Query
 - ✓ Limit
 - ✓ Expression Compatibility
 - ✓ Join
 - ✓ Subquery
 - ✓ Graph traverse
 - ✓ Out direction
 - ✓ In direction
 - ✓ Bidirection
 - ✓ Edge traverse
 - ✓ Distinct

Index

```
test3> desc SELECT * FROM people WHERE name = 'Bob' TRAVERSE OUT(follow).TAG(people);
```

id	estRows	task	access object	operator info
Traverse_7	1.00	root		
└─Point_Get_6	1.00	root	tag:people, index:name(name)	

用 Unique Index 走 point-get 优化

2 rows in set

Time: 0.011s

```
test3> select * from people;
```

id	name	register
1	Bob	2010-01-16 18:00:00
2	Max	2011-01-15 18:00:00
3	Jon	2012-01-14 18:00:00
4	Jim	2020-01-13 18:00:00

4 rows in set

Time: 0.118s

```
test3> desc SELECT * FROM people use index (register) WHERE register='2010-01-16 18:00:00' TRAVERSE OUT(follow).TAG(people);
```

id	estRows	task	access object	operator info
Traverse_9	0.00	root		
└─IndexLookUp_8	0.00	root		
└─IndexRangeScan_6(Build)	0.00	cop[tikv]	tag:people, index:register(register)	range:[2010-01-16 18:00:00,2010-01-16 18:00:00], keep order:false, stats:pseudo
└─GraphTagIDScan_7(Probe)	0.00	cop[tikv]	tag:people	keep order:false, stats:pseudo

用 IndexLookUp 避免“全表扫”

事务

```
(none)> use test4;
You are now connected to database "test4" as user "root"
Time: 0.001s
test4> CREATE TAG people (id BIGINT, name VARCHAR(32));
Query OK, 0 rows affected
Time: 0.005s
test4> CREATE EDGE follow (src BIGINT, dst BIGINT, time DATE);
Query OK, 0 rows affected
Time: 0.005s
test4> CREATE TABLE user (id BIGINT, name VARCHAR(32));
Query OK, 0 rows affected
Time: 0.005s
test4> BEGIN;
Query OK, 0 rows affected
Time: 0.000s
test4> INSERT INTO user VALUES (1, 'a'), (2, 'b');
Query OK, 2 rows affected
Time: 0.001s
test4> INSERT INTO people VALUES (1, 'a'), (2, 'b');
Query OK, 2 rows affected
Time: 0.001s
test4> INSERT INTO follow VALUES (1, 2, now());
Query OK, 1 row affected
Time: 0.001s
test4> select now();
+-----+
| now() |
+-----+
| 2021-01-16 19:34:35 |
+-----+
1 row in set
Time: 0.011s
test4> COMMIT;
Query OK, 0 rows affected
Time: 0.001s
test4>
```

在一个事务中写 TABLE 和 TAG 数据

```
test4> select now();
+-----+
| now() |
+-----+
| 2021-01-16 19:35:12 |
+-----+
1 row in set
Time: 0.011s
test4> SELECT * FROM user;
+----+-----+
| id | name |
+----+-----+
0 rows in set
Time: 0.011s
test4> SELECT * FROM people;
+----+-----+
| id | name |
+----+-----+
0 rows in set
Time: 0.030s
test4> SELECT * FROM user;
+----+-----+
| id | name |
+----+-----+
| 1 | a |
| 2 | b |
+----+-----+
2 rows in set
Time: 0.011s
test4> SELECT * FROM user;
+----+-----+
| id | name |
+----+-----+
| 1 | a |
| 2 | b |
+----+-----+
2 rows in set
Time: 0.011s
test4> SELECT * FROM people WHERE name = 'a' TRAVERSE OUT(follow).TAG(people);
+----+-----+
| id | name |
+----+-----+
| 2 | b |
+----+-----+
1 row in set
```

其他客户端无法读取事务中未提交的数据

事务提交后，写入的数据才可见

- 支持分布式事务。
- 可以在同一个事务里面同时操作 TABLE 和 TAG。

子查询

```
mysql> CREATE TABLE orders(id bigint PRIMARY KEY AUTO_INCREMENT, user_id bigint, item_id bigint);
```

```
mysql> CREATE TAG user(id bigint, name varchar(64));
```

```
mysql> CREATE EDGE contact(src bigint, dst bigint);
```

```
mysql> INSERT INTO USER VALUES (1, "UserA"), (2, "UserA"), (3, "UserA"), (4, "UserA"), (5, "UserA"), (6, "UserA");
```

```
mysql> INSERT INTO contact VALUES (1, 2), (2, 3), (3, 4), (2, 5), (5, 6), (4, 7);
```

```
mysql> INSERT INTO orders(user_id, item_id) VALUES (4, 19230), (6, 4323);
```

```
mysql> -- 查找某一个用户的三度人脉的下单记录
```

```
mysql> SELECT * FROM orders WHERE user_id IN (SELECT id FROM USER WHERE id=1 traverse OUT(contact).OUT(contact).OUT(contact).TAG(USER) LIMIT 10);
```

id	user_id	item_id
1	4	19230
2	6	4323

2 rows in set (0.003s)

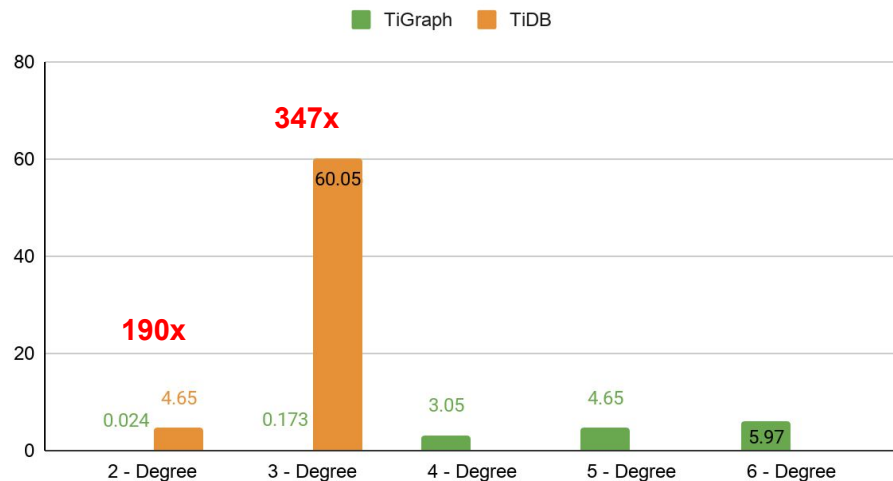
图遍历结果作为子查询

第五部分

Benchmark

Benchmark

Query time (s)



- Workload
 - 24vc + 64G
 - Unistore
 - N 度人脉
 - 数据规模
 - 100k vertices or 100k rows
 - 30 - 100 friends (6.5m edges)
- Notes
 - TiDB 4-degree > 7h (cancelled)

Id	User	Host	db	Command	Time	State	Info
3	root	127.0.0.1	test	Query	26637	autocommit	SELECT cou
11	root	127.0.0.1	test	Query	0	autocommit	show proc

关系型数据库跑 4 度人脉
至少提升 $26637/3.05 = 8733$ 倍

Thank You

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