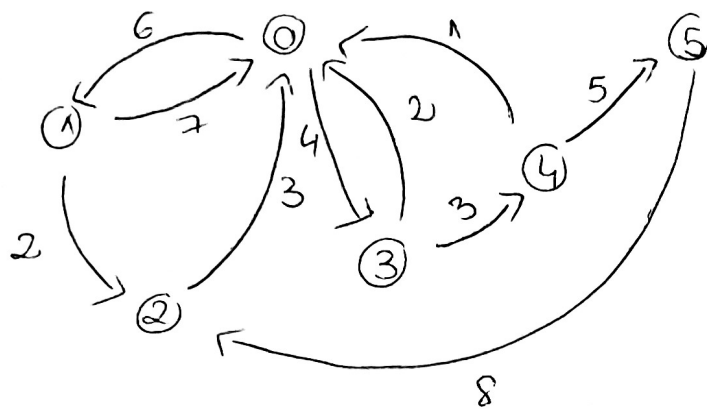


Problem 7. Floyd-Warshall algorithm.



D_i = cost (distance) matrix

P_i = path matrix

$$k=0: D_0 = \begin{pmatrix} 0 & 6 & \infty & 4 & \infty & \infty \\ 4 & 0 & 2 & \infty & \infty & \infty \\ 3 & \infty & 0 & \infty & \infty & \infty \\ 2 & \infty & \infty & 0 & 3 & \infty \\ 1 & \infty & \infty & \infty & 0 & 5 \\ \infty & \infty & 8 & \infty & \infty & 0 \end{pmatrix}; P_0 = \begin{pmatrix} 0 & 1 & 0 & 3 & 0 & 0 \\ 0 & 1 & 2 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 3 & 4 & 0 \\ 0 & 0 & 0 & 0 & 4 & 5 \\ 0 & 0 & 2 & 0 & 0 & 5 \end{pmatrix}$$

$k=1$: Using vertex 1 as an intermediate vertex

$$D_1 = \begin{pmatrix} 0 & 6 & \infty & 4 & \infty & \infty \\ 4 & 0 & 2 & 11 & \infty & \infty \\ 3 & 9 & 0 & 7 & \infty & \infty \\ 2 & 8 & \infty & 0 & 3 & \infty \\ 1 & 4 & \infty & 5 & 0 & 5 \\ \infty & \infty & 8 & \infty & \infty & 0 \end{pmatrix}; P_1 = \begin{pmatrix} 0 & 1 & 0 & 3 & 0 & 0 \\ 0 & 1 & 2 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 3 & 4 & 0 \\ 0 & 0 & 0 & 0 & 4 & 5 \\ 0 & 0 & 2 & 0 & 0 & 5 \end{pmatrix}$$

$k=2$: Using vertex 2 as an intermediate vertex

$$D_2 = \begin{pmatrix} 0 & 6 & 8 & 4 & \infty & \infty \\ 4 & 0 & 2 & 11 & \infty & \infty \\ 3 & 9 & 0 & 7 & \infty & \infty \\ 2 & 8 & 10 & 0 & 3 & \infty \\ 1 & 4 & 9 & 5 & 0 & 5 \\ \infty & \infty & 8 & \infty & \infty & 0 \end{pmatrix}; P_2 = \begin{pmatrix} 0 & 1 & 1 & 3 & 0 & 0 \\ 0 & 1 & 2 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 3 & 4 & 0 \\ 0 & 0 & 0 & 0 & 4 & 5 \\ 0 & 0 & 2 & 0 & 0 & 5 \end{pmatrix}$$

$k=3$: Using vertex 3 as an intermediate vertex

$$D_3 = \begin{pmatrix} 0 & 6 & 8 & 4 & \infty & \infty \\ 5 & 0 & 2 & 9 & \infty & \infty \\ 3 & 9 & 0 & 7 & \infty & \infty \\ 2 & 8 & 10 & 0 & 3 & \infty \\ 1 & 7 & 9 & 5 & 0 & 5 \\ 11 & 14 & 8 & 15 & \infty & 0 \end{pmatrix}, \quad P_3 = \begin{pmatrix} 0 & 1 & 1 & 3 & 0 & 0 \\ 2 & 1 & 2 & 2 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 3 & 4 & 0 \\ 0 & 0 & 0 & 0 & 4 & 5 \\ 2 & 2 & 2 & 2 & 0 & 5 \end{pmatrix}$$

$k=4$: Using vertex 4 as an intermediate vertex

$$D_4 = \begin{pmatrix} 0 & 6 & 8 & 4 & 7 & \infty \\ 5 & 0 & 2 & 9 & 12 & \infty \\ 3 & 9 & 0 & 7 & 10 & \infty \\ 2 & 8 & 10 & 0 & 3 & \infty \\ 1 & 7 & 9 & 5 & 0 & 5 \\ 11 & 14 & 8 & 15 & 18 & 0 \end{pmatrix}, \quad P_4 = \begin{pmatrix} 0 & 1 & 1 & 3 & 3 & 0 \\ 2 & 1 & 2 & 2 & 2 & 0 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 3 & 4 & 0 \\ 0 & 0 & 0 & 0 & 4 & 5 \\ 2 & 2 & 2 & 2 & 2 & 5 \end{pmatrix}$$

$k=5$: Using vertex 5 as an intermediate vertex

$$D_5 = \begin{pmatrix} 0 & 6 & 8 & 4 & 7 & 12 \\ 5 & 0 & 2 & 9 & 12 & 14 \\ 3 & 9 & 0 & 7 & 10 & 15 \\ 2 & 8 & 10 & 0 & 3 & 8 \\ 1 & 7 & 9 & 5 & 0 & 5 \\ 11 & 14 & 8 & 15 & 18 & 0 \end{pmatrix}, \quad P_5 = \begin{pmatrix} 0 & 1 & 1 & 3 & 3 & 3 \\ 2 & 1 & 2 & 2 & 2 & 2 \\ 0 & 0 & 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 3 & 4 & 4 \\ 0 & 0 & 0 & 0 & 4 & 5 \\ 2 & 2 & 2 & 2 & 2 & 5 \end{pmatrix}$$

$D_5(3,5)=8 \Rightarrow$ The minimum cost walk from 3 to 5 has the cost $D_5(3,5)=8$ and it is obtained from P_5 as follows:

$$n=3: P_5(3,5)=4, P_5(4,5)=5=t$$

The minimum cost walk: $3 \xrightarrow{3} 4 \xrightarrow{5} 5$