Topics List for Exam 2

• Queues (FIFO) #include <queue>

can't use vectors or arrays with queues!!!

• Know how to create an empty queue

```
deque<int> mydeck (3,100); // deque with 3 elements
                  list<int> mylist (2,200);
                                                 // list with 2 elements
                  queue<int> first;
                                       // empty queue
                  queue<int> second (mydeck); // queue initialized to copy of deque
                  queue<int,list<int> > third; // empty queue with list as underlying container
                  queue<int,list<int> > fourth (mylist);
                  cout << "size of first: " << first.size() << '\n';</pre>
                  cout << "size of second: " << second.size() << '\n';</pre>
                  cout << "size of third: " << third.size() << '\n';
cout << "size of fourth: " << fourth.size() << '\n';</pre>
• <u>know what push(), pop(), empty(), and size() do.</u> (front(), end())
                Void push(T value) adds to the back of the queue
                         myqueue.push (myint);
                Void pop() removes from the front of the queue
                         myqueue.pop();
                Bool empty() returns true or false
                         while (!myqueue.empty())
                          {
                                 sum += myqueue.front();
                                 myqueue.pop();
                Int size() returns the size of the queue
                         for (int i=0; i<5; i++) myints.push(i);</pre>
                         cout << "1. size: " << mvints.size();</pre>
```

- **Priority Queues** (ts first element is always the greatest of the elements it contains)
 - *Know how to create an empty priority queue.*

```
priority_queue<int> first; //empty
priority_queue<int> second (myints,myints+4);
priority_queue<int, std::vector<int>, std::greater<int> >third (myints,myints+4);
```

• Know what push(), pop(), top(), size() and empty() do, and be able to use them to write a function

Void push(const int value & val) This member function effectively calls the member function push_back of the *underlying container* object, and then reorders it to its location in the *heap* by calling the push_heap algorithm on the range that includes all the elements of the container.

Void **pop()** Removes the element on top of the priority_queue, effectively reducing its size by one. The element removed is the one with the highest value.

T **size()**; Returns the number of elements in the priority_queue.

T top(); Returns a constant reference to the *top element* in the priority gueue.

Example:

```
priority_queue<int> mypq;
  mypq.push(30);
  mypq.push(100);
  mypq.push(25);
  mypq.push(40);

cout<<mypq.size();

while (!mypq.empty()) {
      cout << ' ' << mypq.top();//prints top element
      mypq.pop(); //removes an element
}</pre>
```

- Be able to solve a problem using a priority queue.
- Our own implementation:
 - <u>Understand the implementation we used (binary max heap.)</u> (the max value is the root of the tree) (http://www.sci.brooklyn.cuny.edu/~mermelstein/Courses/CISC3130/PriorityQueue/03-OurOwnPri orityQueueImplementation/)

```
top()
    return arr[0]
        insert (x)
            arr.push_back(x);
        while (x is not at the root && x > parent(x))        swap;
pop()
        Swap arr[0] with arr[arr.size()-1]
        arr.pop_back();
        while(x is not a leaf && x<at least 1 of its children)
            Pick x's largest child and swap</pre>
```

Applications of Heaps:

- 1) Heap Sort: Heap Sort uses Binary Heap to sort an array in O(nLogn) time.
- **2)** Priority Queue: Priority queues can be efficiently implemented using Binary Heap because it supports insert(), delete() and extractmax(), decreaseKey() operations in O(logn) time.
- 3) Graph Algorithms: The priority queues are especially used in Graph Algorithms.
- **4)** Many problems can be efficiently solved using Heaps. See following for example.
 - a) K'th Largest Element in an array.
 - b) Sort an almost sorted array/
 - c) Merge K Sorted Arrays.
- Be able to run the insertElement() and removeMax() operations on a heap.

```
insertElement() O(h-->log base2 (n+1))
```

Complexity of the insertion operation is O(h), where **h** is heap's height. Taking into account completeness of the tree, $O(h) = O(\log n)$, where **n** is number of elements in a heap.

for (int i=m iCurrNumOfItems, j=(i+1)/2-1; j>=0; i=j, j=(i+1)/2-1) {

```
if (BestOfTwo(i,j) == i) SwapItems(i,j);
else
        break;
m iCurrNumOfItems++;
return true;
                Pseudocodes:
                       insert(x)
                          arr.push back(x);
                          while(x is not at the root && x>parent(x)) swap
                       removeMax() O(h-->log base2 (n+1))
                          pop()
                             Swap arr[0] with arr[arr.size()-1]
                             arr.pop back();
                             while(x is not a leaf && x < at least 1 of its children)
                                      Pick x's largest child and swap
         • Know the running times for all of the priority queue functions.
                Using a vector:
                       push(x) --> call push back(x) --> O(1) amortized
                       pop() remove max el by shifting and find new max O(n)
                       top() every time you call push(), you check if the newest el is bigger O(1)
                Using a sorted vector:
                       push(x) needs to maintain a sorted order by shifting O(n)
                       pop() just call pop back O(1)
                       top() return back() O(1)
                Using binary heap: (we want)
                       push(x) O(n log n)
                       pop() O(log n)
                       top() O(1)
```

Sorting

• Know how to run at least one of the $O(n^2)$ time sorting algorithms

• bubble sort

}

Insertion sort

```
void insertionsort(int arr[], int N) {
    for(int i = 1; i < N; i++) {
        int j = i - 1;
        int temp = arr[i];
        while(j >= 0 && temp < arr[j]) {
            arr[j + 1] = arr[j];
            j--;;
        }
        arr[j + 1] = temp;
    }
}</pre>
```

• Know how to run the subroutines of the O(n log n) time sorting algorithms

```
• Merge sort (you have to know how to run merge())
               void merge(int *arr1, int n1,int *arr2, int n2, int *result) {
                       int *temp = new int [n1 + n2];
                       int index 1 = 0;
                       int index2 = 0;
                       int i = 0;
                       while(index 1 < n1 \&\& index 2 < n2) {
                               if(arr1[index1] <= arr2[index2]) {</pre>
                                       temp[i] = arr1[index1];
                                       index1++;
                               }
                               else {
                                 temp[i] = arr2[index2];
                                 index2++;
                               i++;
                       //only runs if arr1 is still full
                       while(index1 <n1){
                             temp[i] = arr1[index1];
                             index1++;
                            i++;
                       //only runs if arr2 is still full
                       while(index2 < n2){
                             temp[i] = arr2[index2];
                             index2++;
                             i++;
                       //copy all el from temp into result
                       for(int i = 0; i < n1 + n2j; i + +)
                             result[i] = temp[i];
                       Delete[temp];
```

```
• Quicksort (you have to know how to run partition())
                                    int partition (int arr[], int start, int end) {
                                           int pivotValue = arr [start];
                                           int pivotPosition = start;
                                            for (int pos =start+1; pos <= end; pos++){
                                                   if (arr[pos] < pivotValue) {
                                                          swap(arr[p i votPosition + 1], arr[pos]);
                                                          swap(arr[p ivotPosition ], arr[pivotPosition + 1]);
                                                          pivotPosition ++;
                                                   }
                                           return pivotPosition;
                                    void quickSort( int arr[ ] , int start, int end) {
                                           if (start< end) {
                                                   int p = partition(arr, start, end);
                                                   quickSort (arr, start, p - 1);
                                                   quickSort (arr, p + 1, end);
Heap sort
void heapSort(int arr[], int n)
    // Build max heap
    for (int i = n / 2 - 1; i >= 0; i--)
      heapify(arr, n, i);
    // Heap sort
    for (int i=n-1; i>=0; i--)
      swap(arr[0], arr[i]);
      // Heapify root element to get highest element at root again
      heapify(arr, i, 0);
    }
}
void heapSort(vector<int>& aVector) {
       priority_queue<int> pQueue;
       for(int i=0; i<aVector.size(); i++)
       pQueue.push(aVector[i]);
       for(int i=aVector.size()-1; i>=0; i--) {
       int maximumElement = pQueue.top();
       aVector[i] = maximumElement;
       pQueue.pop();
       }
```

}

```
}
```

```
Linear search:
int searchList (const int list[], int size, int value) {
    int index= O;
    int position= - 1;
    bool found= false;
    while (index< size && !found) {
        if (list[ index] == value) {
            found= true;
            position= index;
        }
        index++;
    }
    return position;
}</pre>
```

• Know the running times for all of the sorting algorithms discussed in class

sort	Best Case	Worst Case	Average
Insertion Sort	Already sorted, O(n)	Reverse sorted, O(n^2)	O(n^2)
Selection Sort	Already sorted, O(n^2)	Reverse sorted, O(n^2)	O(n^2)
Bubble Sort	Already sorted O(n)	Reverse sorted, O(n^2)	O(n^2)
Merge Sort	O(nlog(n))	Reverse sorted, O(nlog(n))	O(nlog(n))
Quicksort	O(nlog(n))	All elements equal, O(n^2)	O(nlog(n))
Heap Sort	O(n log n)	O(n log n)	O(n log n)

• Sets - contains unique values only!!!

• <u>Know how to create an empty set.</u> set<int> first;

Know what the operations insert(), erase(), find(), and count() do, and be able to use them to write a function to solve a problem. (You do **NOT** have to worry about the return type of insert of know how to use it. You must know how to use all other return values) pair<iterator,bool> insert(T key); (add key, if not already there) (#include <utility>) -iterator- el you just added or it's already there -bool- represents whether or not the insertion was successful (true, if el was just added; false-0, if el was already there) int erase(T key); (removes key from the set, if it's there; returns wheather el was removed) myset.erase (it); myset.erase (40); it = myset.find (60);myset.erase (it, myset.end()); set<T>::iterator find(T el); (returns iterator to el if found, and end() if not found) set<int> myset; set<int>::iterator it; // set some initial values: for (int i=1; i<=5; i++) myset.insert(i*10);// set: 10 20 30 40 50 it=mvset.find(20): myset.erase (it): myset.erase (myset.find(40)); cout << "myset contains:"; for (it=myset.begin(); it!=myset.end(); ++it)
 cout << ' ' << *it;</pre> cout << '\n'; int count(T elem); (like bool returns 1 =, if el is in a set; 0, if el is absent) if (myset.count(i)!=0)

• <u>Know how to iterate through a set using set<T>::iterator and const_iterator.</u>

else

cout << " is an element of myset.\n";</pre>

cout << " is not an element of myset.\n";</pre>

Maps

• Know how to create an empty map.

```
map<char,int> first; //empty map

first['a']=10;
first['b']=30;
first['c']=50;
first['d']=70;

map<char,int> second (first.begin(),first.end());

map<char,int> third (second);
```

• <u>Know what the operations insert(), erase(), find(), count(), and operator[] do, and be able to use them to write a function to solve a problem. (Just as with sets, you need not worry about the return type and value for insert()).</u>

Built-in functions:

```
insert(pair<key,value>p);
       -returns <iter,bool> (bool: true/false, depending on if key is there or not)
       -inserts the pair into the map
                map<char,int> mymap;
                // first insert function version (single parameter):
                mymap.insert (pair<char,int>('a',100));
                mymap.insert (pair<char,int>('z',200));
                pair<map<char,int>::iterator,bool> ret;
                ret = mymap.insert ( pair<char,int>('z',500) );
                if (ret.second==false) {
                       cout << "element 'z' already existed";
                       cout << " with a value of " << ret.first->second << '\n';
                }
                // second insert function version (with hint position):
                map<char,int>::iterator it = mymap.begin();
                mymap.insert (it, pair<char,int>('b',300)); // max efficiency inserting
                mymap.insert (it, pair<char,int>('c',400)); //no max efficiency inserting
                // third insert function version (range insertion):
                map<char,int> anothermap;
                anothermap.insert(mymap.begin(),mymap.find('c'));
                // showing contents:
                cout << "mvmap contains:\n";</pre>
                for (it=mymap.begin(); it!=mymap.end(); ++it)
                       cout << it->first << " => " << it->second << '\n';
                cout << "anothermap contains:\n";</pre>
                for (it=anothermap.begin(); it!=anothermap.end(); ++it)
                       std::cout << it->first << " => " << it->second << '\n';
int erase(T key); removes the pair from the map whose key is "key"
                                it=mymap.find('b');
                                mymap.erase (it);
                                                            // erasing by iterator
                                mymap.erase ('c');
                                                            // erasing by key
                                it=mymap.find ('e');
                                mymap.erase (it, mymap.end()); // erasing by range
int count(T key); returns 0 or 1, if the key is in a map or not
                        std::map<char,int> mymap;
                        char c:
                        mymap ['a']=101;
                        mymap ['c']=202;
                        mymap ['f']=303;
                        for (c='a'; c<'h'; c++)
                                      std::cout << c:
                                      if (mymap.count(c)>0)
                                      std::cout << " is an element of mymap.\n";
                                      std::cout << " is not an element of mymap.\n";
```

```
find()
```

```
iterator find (const key type& k);
const iterator find (const key type& k) const;
Searches the container for an element with a key equivalent to k and returns an
iterator to it if found, otherwise it returns an iterator to map::end.
                        map<char,int> mymap;
                        map<char,int>::iterator it;
                        mymap['a']=50;
                        mymap['b']=100;
                        mymap['c']=150;
                        mymap['d']=200;
                        it = mymap.find('b');
                        if (it != mymap.end());
                        mymap.erase (it);
                        // print content:
                        cout << "elements in mymap:" << '\n';
                        cout << "a => " << mymap.find('a')->second << '\n';
                        cout << "c => " << mymap.find('c')->second << '\n';
                        cout << "d => " << mymap.find('d')->second << '\n';
```

• <u>Know how to iterate through a map using map<T1, T2>::iterator and const_iterator, and know how to handle each element that comes from the iterator (the key is the first component of the pair, and the value is the second component of the pair).</u> #include<utility>

• Binary Search Trees

```
    How to create TreeNodes.
        struct BSTnode {
            int key;
            BSTnode *left, *right;
        }

    Know how to search for an item in a binary search tree.
        BSTnode* search (BSTnode* root, int key) {
            if (root == NULL) return NULL;
            if (root→ key == key) return root;
            if (key<root→ key) return search (root→ left, key);
            return search(root→ right, key);
        </pre>

        void insert(BSTnode*& root, int key) {
```

```
if (root == NULL){
    root == new BSTnode;
    root→ key = key;
    root→ right = root→ left = NULL;
}
else if(key<root→ key)
    insert (root→ left, key);
else insert (root→ right, key);
}</pre>
```

• <u>Know how to visit all nodes in a binary search tree and perform some operation (like printing out all nodes, adding up all nodes, etc.)</u>

```
void postorder (BSTnode* root){
        if (root == NULL) return;
         portorder (root\rightarrow left);
        postorder (root\rightarrow right);
         cout << root \rightarrow key << endl;
}
                                                     void inorder (BSTnode* root){
                                                              if (root == NULL) return;
                                                              inorder (root\rightarrow left);
                                                              cout << root \rightarrow key << endl;
                                                              inorder (root\rightarrow right);
                                                      }
void preorder (BSTnode* root){
        if (root == NULL) return;
         cout << root \rightarrow key << endl;
         preorder (root\rightarrow left);
        preorder (root\rightarrow right);
}
                                                     void deallocateTree (BSTnode* root){
                                                              if (root == NULL) return;
                                                              deallocateTree (root\rightarrow left);
                                                              deallocateTree (root\rightarrow right);
                                                              delete root;
                                                      }
int sumAll (BSTnode* root){
         if (root == NULL) return 0;
         Return sumAll(root-->left) + root\rightarrow key + sumAll (root\rightarrow right)
}
```

Test questions:

- 1. From given definition for binary set nodes, write a function that sums all nodes and prints them out
- 2. Binary max heap→ what happens if you insert an element. Show steps what it looks like before & after
- 3. Declare a map. Use insert, search.
- 4. Definitions: binary heap, max heap; binary search tree
- 5. Write functions using sets and maps

- 6. Sort array using bubble or insertion sort
- 7. Run merge on 2 sorted arrays
- 8. Use sorting on set, map, heap, priority queue, binary search tree
- 9. Answer ?: what a running time would be, if we implement sorts with vector/map/queue
- 10. Know how binary max heap work
- 11. From a given array, define if it's a binary max heap or not
- 12. Draw binary max heap, then run insert & delete
- 13. Run insert & delete on an array and define what an array looks like and what a tree looks like
- 14. Run a partition algo on an array
- 15. Use find or count to find an el in a map, don't just use [] in cout
- 16. Trace through push() and pop() from a given array down the tree (tree push O(log n))
- 17. Draw a tree from an array→ what happens after using push() and pop()
- 18. Define if given is a legal max heap
- 19. Know how to run push() and pop() on binary max heap

```
Side notes:
```

In a map insert() returns a pair

Answer to #4:

- 1. Binary tree is a structure with nodes, where each node has <=2 children
- 2. A complete binary tree where all levels, except for the last, have to be filled
- 3. A binary max heap is a complete binary tree, where each node obeys the max heap property for every node x, $x \le parent(x)$ (x = binary tree) its children