

# Languages and Automata

## Assignment 5, Tuesday 10<sup>th</sup> March, 2020

**Update 12 March:** please hand in via brightspace. Before submitting, make sure:

- the file is a PDF document
- your name and student number are included in the document (they might be printed).

**Deadline:** Friday 13<sup>th</sup> March, 2020, 17:00 (in Nijmegen!). This deadline is strict: submission in brightspace will close at that time.

**Goals:** After completing these exercises successfully you should be able to read context-free grammars, write down grammars for context-free languages and regular languages, and work with basic closure properties of context-free languages. The total number of points is 10.

There are 4 mandatory exercises, worth **10 points** in total. There are 3 more, extra hard, exercises. Be aware that these exercises are just for fun, you cannot earn any points with them.

### 1 Ambiguous grammars

Let  $\Sigma = \{a, b\}$ . Consider the following context-free grammar

$$G_1 = \boxed{S \rightarrow aS \mid Sb \mid ab \mid SS}$$

- Show that the grammar is ambiguous by giving two left-most derivations of  $abbb$ . (1pt)
- Give a regular expression for  $L_1 := \mathcal{L}(G_1)$ . (1pt)
- Give a non-ambiguous regular grammar for  $L_1$ . (1pt)

### 2 Constructing context-free grammars

For each of the following languages construct a context-free grammar that generates the language, and **explain why your answer is correct**. (4pt)

$$L_1 = \{a^n b^{n+m} a^m \mid n, m \geq 0\}$$

$$L_2 = \{w \in \{a, b, r\}^* \mid \text{every } a \text{ (in } w) \text{ is followed directly by } b, \text{ and } w \text{ ends with } brrr\}$$

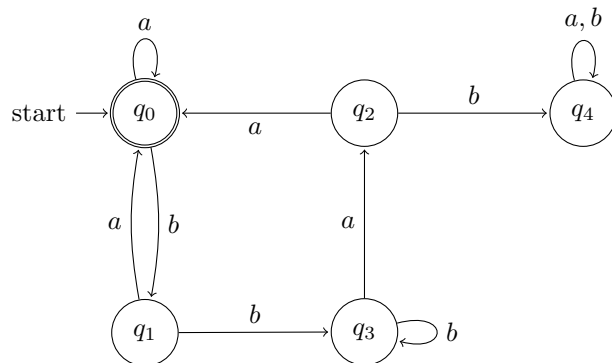
### 3 Closure properties

For each of the following statements, decide whether or not they are correct. Justify your answer with a proof (if the statement is correct) or a counterexample (if it is not correct). (2pt)

1. If  $L$  is context-free, then the complement  $\bar{L}$  is *not* context-free.
2. If  $L$  is context-free and  $K$  is regular, then  $L \cap K$  is regular.

## 4 From DFA to regular grammar

Consider the following DFA  $M$ :



Construct a context-free grammar that generates  $\mathcal{L}(M)$ .

**(1pt)**

## Fun Exercises

1. Construct a context-free grammar that generates the following language:

$$L_5 = \{w \in \{a, b\}^* \mid |w| = 2k + 1 \text{ and } w_1 = w_{k+1}\},$$

where  $w_i$  denotes the  $i$ -th symbol in a word  $w$ . That is,  $L_5$  consists of all words of odd length that have the same symbol in the first and middle positions.

2.  $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, \times, +, (, )\}$ . Construct a context-free grammar that generates the language

$$L_6 = \{w \in \Sigma^* \mid w \text{ is a well-formed arithmetical expressions}\}$$

NB.  $2 + 3 + 4 \times 5$  and  $((2 + 3) + 4) \times 5$  and  $((2 + 3)) + 4 \times 5$  are well-formed.  $2 + (3 + 4 \times 5$  and  $(2 + 3) + 4) \times 5$  and  $) ($  are not.

3. Suppose that  $L$  and  $L'$  are context-free languages. Show that both  $L^*$  and  $L L'$  are context-free languages.