

Macroeconomics III

Stochastic Environment and Finite Markov Chains

Diego Rodrigues

SciencesPo
diego.desousarodrigues@sciencespo.fr

Fall 2023

1 Stochastic Environment

Environment

- The agents make decision at date $t = 0$;
- There are infinite periods;
- At each date $t \geq 0$ there is a realization of a stochastic event $s_t \in \mathcal{S}$;
- The history of these events until period t is denoted by $s^t = [s_0, s_1, s_2, \dots, s_t]$;
- A particular sequence of events occur with probability $\pi_t(s^t)$;
- The chance that a sequence of events s^τ occurs given the sequence s^t occurred is given by $\pi_t(s^\tau | s^t)$;
- There is a set of I agents that receive a perfectly anticipated endowment $y_t^i(s^t)$ in each period. This endowment is perishable, but the agents can trade among themselves;
- The individual savings can be positive or negative, but the aggregate one will always be zero;

- Agents have utility $u(c)$ that is increasing and concave, i.e., $u'(c) > 0$ and $u''(c) < 0$;
- Agents will draw a consumption plan for all his life $\{c_t^i(s^t)\}_{t=0}^{\infty} \equiv C^i$, which will be subject to a discount factor $\beta \in (0, 1)$ such that:

$$U(C^i) = \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) u[c_t^i(s^t)]$$

- First of all we will solve the **social planner's problem** and then verify the **conditions such that the solution of the competitive equilibrium is efficient**.

The Social Planner's problem

$$\max_{\{C^i\}_{i=1}^I} \sum_{i=1}^I \lambda^i U(C^i) \quad \text{s.a.} \quad \sum_{i=1}^I c_t^i(s^t) \leq \sum_{i=1}^I y_t^i(s^t), \quad \forall t, \forall s^t$$

- The problem of the Social Planner is then:

$$\mathcal{L} = \sum_{i=1}^I \lambda^i \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) u[c_t^i(s^t)] + \sum_{t=0}^{\infty} \sum_{s^t} \theta_t(s^t) \left[\sum_{i=1}^I (y_t^i(s^t) - c_t^i(s^t)) \right]$$

Let's derive the F.O.C.

- $[c_t^i(s^t)] : \lambda^i \beta^t \pi_t(s^t) u' [c_t^i(s^t)] = \theta_t(s^t), \quad \forall i \in I, \forall t, \forall s^t$

Let's derive the F.O.C.

- $[c_t^i(s^t)] : \lambda^i \beta^t \pi_t(s^t) u' [c_t^i(s^t)] = \theta_t(s^t), \quad \forall i \in I, \forall t, \forall s^t$
- By defining the aggregate endowment as $Y_t(s^t) \equiv \sum_{i=1}^I y_t^i(s^t)$ and assuming the **aggregate endowment between any two periods t and τ is the same**, we can show:
$$Y_t(s^t) = Y_\tau(s^\tau) \implies c_t^i(s^t) = c_\tau^i(s^\tau), \forall i \in I$$

Let's derive the F.O.C.

Divide the F.O.C. of the agent i by the F.O.C. of a generic agent 1:

$$\frac{\lambda^i \beta^t \pi_t(s^t) u' [c_t^i(s^t)] = \theta_t(s^t)}{\lambda^1 \beta^t \pi_t(s^t) u' [c_t^1(s^t)] = \theta_t(s^t)} \\ \implies u' [c_t^i(s^t)] = \frac{\lambda^1}{\lambda^i} u' [c_t^1(s^t)]$$

By assuming the derivative of the utility function admits inverse we can have:

$$c_t^i(s^t) = u'^{-1} \left(\frac{\lambda^1}{\lambda^i} u' [c_t^1(s^t)] \right) \\ \implies c_t^i(s^t) = h^i [c_t^1(s^t)]$$

The feasibility conditions is going to imply that:

$$\sum_{i=1}^I c_t^i(s^t) = \sum_{i=1}^I y_t^i(s^t) \implies \sum_{i=1}^I h^i [c_t^1(s^t)] = \sum_{i=1}^I y_t^i(s^t)$$

Let's derive the F.O.C.

If the total endowment is the same we have $Y_t(s^t) = Y_\tau(s^\tau)$, where $Y_t(s^t) = \sum_{i=1}^I y_t^i(s^t)$ we can show by using the previous equation that:

$$\sum_{i=1}^I h^i [c_t^1(s^t)] = \sum_{i=1}^I h^i [c_\tau^1(s^\tau)]$$

The only way for the above to be true is such that $c_t^1(s^t) = c_\tau^1(s^\tau)$. Since agent 1 was chose arbitrarily, the above is valid for all agents $i \in I$

Agent's problem in an Arrow-Debreu structure

- Assume there are only two possible states of nature: 0 and 1, i.e., $\mathcal{S} = \{0, 1\}$
- Let the economy always begin in state 0 (i.e., the initial state is deterministic $\pi_0(0) = 1$)

Agent's problem in an Arrow-Debreu structure

- Assume there are only two possible states of nature: 0 and 1, i.e., $\mathcal{S} = \{0, 1\}$
- Let the economy always begin in state 0 (i.e., the initial state is deterministic $\pi_0(0) = 1$)
- The structure for an economy with 3 periods can be represented by the following tree:

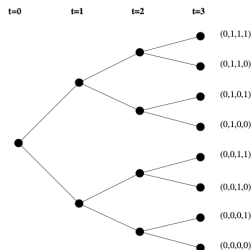


Figure 1: The Arrow-Debreu commodity space for a two-state.

Agent's problem in an Arrow-Debreu structure

- The **market structure is complete** in the sense that the agent can acquire **rights of consumption for any period and for any possible history**.

Agent's problem in an Arrow-Debreu structure

The problem of the agents is

$$\max_{\{c_t^i(s^t)\}_{t=0}^{\infty}} \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) u[c_t^i(s^t)]$$

subject to

$$\sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) y_t^i(s^t) \geq \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) c_t^i(s^t)$$

Agent's problem in an Arrow-Debreu structure

- Since the markets open only once, in $t = 0$, we will have just one Lagrange multiplier in the problem below:

$$\mathcal{L} = \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) u[c_t^i(s^t)] + \gamma^i \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) [y_t^i(s^t) - c_t^i(s^t)]$$

- The F.O.C.
 $[c_t^i(s^t)] : \beta^t \pi_t(s^t) u'[c_t^i(s^t)] = \gamma^i q_t^0(s^t), \quad \forall i \in I, \forall t, \forall s^t$

Agent's problem in an Arrow-Debreu structure

- Using the same idea we used for the Social Planner's problem we can divide the F.O.C. of the agent i by the F.O.C. of an agent 1 generic and obtain:

$$\frac{\beta^t \pi_t(s^t) u' [c_t^i(s^t)] = \gamma^i q_t^0(s^t)}{\beta^t \pi_t(s^t) u' [c_t^1(s^t)] = \gamma^1 q_t^0(s^t)} \implies u' [c_t^i(s^t)] = \frac{\gamma^i}{\gamma^1} u' [c_t^1(s^t)]$$

Agent's problem in an Arrow-Debreu structure

- Using the same idea we used for the Social Planner's problem we can divide the F.O.C. of the agent i by the F.O.C. of an agent 1 generic and obtain:

$$\frac{\beta^t \pi_t(s^t) u' [c_t^i(s^t)] = \gamma^i q_t^0(s^t)}{\beta^t \pi_t(s^t) u' [c_t^1(s^t)] = \gamma^1 q_t^0(s^t)} \implies u' [c_t^i(s^t)] = \frac{\gamma^i}{\gamma^1} u' [c_t^1(s^t)]$$

- In other words, by setting $\lambda^i = \frac{\gamma^i}{\gamma^1} \lambda^1$ and using those weights to solve the SP, we will reach the same allocation as the competitive equilibrium. In other words, **there are weights such that the allocation of the social planner is a competitive equilibrium.**

Agent's problem in an Arrow-Debreu structure

- So far, we obtained 4 important equations, which are:

$$\lambda^i \beta^t \pi_t(s^t) u' [c_t^i(s^t)] = \theta_t(s^t), \quad \forall i \in I, \forall t, \forall s^t \quad (\text{F.O.C. SP})$$

$$\beta^t \pi_t(s^t) u' [c_t^i(s^t)] = \gamma^i q_t^0(s^t), \quad \forall i \in I, \forall t, \forall s^t \quad (\text{F.O.C. AD})$$

$$\sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) [y_t^i(s^t) - c_t^i(s^t)] = 0 \quad (\text{BC})$$

$$Y_t(s^t) = \sum_{i=1}^I y_t^i(s^t) = \sum_{i=1}^I c_t^i(s^t) \quad (\text{F})$$

Agent's problem in an Arrow-Debreu structure

- By assuming $q_0^0(s^0) = 1$ we can then reach the Euler Equation:

Agent's problem in an Arrow-Debreu structure

- By assuming $q_0^0(s^0) = 1$ we can then reach the Euler Equation:

$$u' [c_0^i(s^0)] = \frac{\beta^t \pi_t(s^t) u' [c_t^i(s^t)]}{q_t^0(s^t)}$$

Agent's problem in an Arrow-Debreu structure

- By assuming $q_0^0(s^0) = 1$ we can then reach the Euler Equation:

$$u' [c_0^i(s^0)] = \frac{\beta^t \pi_t(s^t) u' [c_t^i(s^t)]}{q_t^0(s^t)}$$

- From the Euler Equation we can have idea about the prices in this economy:

$$q_t^0(s^t) = \beta^t \pi_t(s^t) \frac{u' [c_t^i(s^t)]}{u' [c_0^i(s^0)]}$$

Agent's problem in an Arrow-Debreu structure

- Suppose **there is no aggregate uncertainty**, $Y_t(s^t) = \bar{Y}$, $\forall t, \forall s^t$
- By the analysis we already made we can conclude $u'[c_t^i(s^t)] = u'[c_0^i(s^0)]$, which will lead us to a completely exogenous price:

$$q_t^0(s^t) = \beta^t \pi_t(s^t)$$

Agent's problem in an Arrow-Debreu structure

- Suppose $u(c) = \log(c) \implies u'(c) = \frac{1}{c}$
- **Show that the consumption of the agent is always a function of the total endowment.**

Agent's problem in an Arrow-Debreu structure

- By still considering the case where there is no aggregate uncertainty

$$q_t^0(s^t) = \beta^t \pi_t(s^t) \frac{u'[c_t^i(s^t)]}{u'[c_0^i(s^0)]}$$

$$\sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) y_t^i(s^t) = \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) c_t^i(s^t) \Rightarrow \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) y_t^i(s^t) = \sum_{t=0}^{\infty} \sum_{s^t \in S} \beta^t \pi_t(s^t) c_t^i(s^t)$$

Agent's problem in an Arrow-Debreu structure

- By still considering the case where there is no aggregate uncertainty

$$q_t^0(s^t) = \beta^t \pi_t(s^t) \frac{u'[c_t^i(s^t)]}{u'[c_0^i(s^0)]}$$

$$\sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) y_t^i(s^t) = \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) c_t^i(s^t) \Rightarrow \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) y_t^i(s^t) = \sum_{t=0}^{\infty} \sum_{s^t \in S} \beta^t \pi_t(s^t) c_t^i(s^t)$$

Define $\bar{C}^i \equiv \sum_{t=0}^{\infty} \sum_{s^t} c_t^i(s^t)$

Agent's problem in an Arrow-Debreu structure

- By still considering the case where there is no aggregate uncertainty

$$q_t^0(s^t) = \beta^t \pi_t(s^t) \frac{u'[c_t^i(s^t)]}{u'[c_0^i(s^0)]}$$

$$\sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) y_t^i(s^t) = \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) c_t^i(s^t) \Rightarrow \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) y_t^i(s^t) = \sum_{t=0}^{\infty} \sum_{s^t \in S} \beta^t \pi_t(s^t) c_t^i(s^t)$$

Define $\bar{C}^i \equiv \sum_{t=0}^{\infty} \sum_{s^t} c_t^i(s^t)$

$$= \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) c_t^i(s^t) = \bar{C}^i \sum_{t=0}^{\infty} \beta^t \sum_{s^t} \pi_t(s^t) = \bar{C}^i \cdot \frac{1}{1-\beta} \cdot 1$$

Agent's problem in an Arrow-Debreu structure

- By still considering the case where there is no aggregate uncertainty

$$q_t^0(s^t) = \beta^t \pi_t(s^t) \frac{u'[c_t^i(s^t)]}{u'[c_0^i(s^0)]}$$

$$\sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) y_t^i(s^t) = \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) c_t^i(s^t) \Rightarrow \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) y_t^i(s^t) = \sum_{t=0}^{\infty} \sum_{s^t \in S} \beta^t \pi_t(s^t) c_t^i(s^t)$$

Define $\bar{C}^i \equiv \sum_{t=0}^{\infty} \sum_{s^t} c_t^i(s^t)$

$$= \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) c_t^i(s^t) = \bar{C}^i \sum_{t=0}^{\infty} \beta^t \sum_{s^t} \pi_t(s^t) = \bar{C}^i \cdot \frac{1}{1-\beta} \cdot 1$$

$$\frac{\bar{C}^i}{1-\beta} = \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) y_t^i(s^t) \Rightarrow \bar{C}^i = (1-\beta) \sum_{t=0}^{\infty} \beta^t \sum_{s^t} \pi_t(s^t) y_t^i(s^t)$$

- Therefore, **the present value of consumption is equal to the expected value of the endowment of the agent.**

Agent's problem in a Sequential structure

- First of all define $Q(s_{t+1}|s^t)$ as being the price in t , given the history s^t of an unit of consumption good in period $t + 1$ contingent to the realization s_{t+1} .

Agent's problem in a Sequential structure

- First of all define $Q(s_{t+1}|s^t)$ as being the price in t , given the history s^t of an unit of consumption good in period $t + 1$ contingent to the realization s_{t+1} .
- The problem of the agents will be given by:

$$\max_{\{c_t^i(s^t), a_{t+1}^i(s^{t+1})\}_{t=0}^{\infty}} \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) u[c_t^i(s^t)]$$

subject to:

$$c_t^i(s^t) + \sum_{s_{t+1}} Q_t(s_{t+1} | s^t) a_{t+1}^i(s^{t+1}) \leq y_t^i(s^t) + a_t^i(s^t) \quad \forall t, \forall s^t$$

Agent's problem in a Sequential structure

- Let us define now a **natural limit for the debt** of the agent:

$$-a_t^i(s^t) \leq \sum_{\tau=t+1}^{\infty} \sum_{s^\tau | s^{t+1}} q_\tau^0(s^\tau) y_\tau^i(s^\tau)$$

- The **debt of the agent has to be lower than the present value of the income of the agent for the remaining of his/her life.**

Agent's problem in a Sequential structure

$$\mathcal{L} = \sum_{t=0}^{\infty} \sum_{s^t} \beta^t \pi_t(s^t) u[c_t^i(s^t)] + \sum_{t=0}^{\infty} \sum_{s^t} \eta_t^i(s^t) [y_t^i(s^t) + a_t^i(s^t) - c_t^i(s^t) - \sum_{s^{t+1}} Q_t(s_{t+1} | s^t) a_{t+1}^i(s^{t+1})]$$

In this case we will have the following F.O.C.

- The F.O.C.

$$[c_t^i(s^t)] : \beta^t \pi_t(s^t) u'[c_t^i(s^t)] = \eta_t^i(s^t)$$

$$[a_{t+1}^i(s^{t+1})] : \eta_t^i(s^t) Q_t(s_{t+1} | s^t) = \eta_{t+1}^i(s^{t+1})$$

Agent's problem in a Sequential structure

- So far, we obtained 4 important equations, which are:

$$Q_t(s_{t+1} | s^t) = \beta \pi_t(s^{t+1} | s^t) \frac{u' [c_{t+1}^i(s^{t+1})]}{u' [c_t^i(s^t)]} \quad (\text{EE})$$

$$c_t^i(s^t) + \sum_{s_{t+1}} Q_t(s_{t+1} | s^t) a_{t+1}^i(s^{t+1}) = y_t^i(s^t) + a_t^i(s^t) \quad \forall t, \forall s^t \quad (\text{BC})$$

$$Y_t(s^t) = \sum_{i=1}^I y_t^i(s^t) = \sum_{i=1}^I c_t^i(s^t) \quad (\text{F})$$

$$\sum_{i=1}^I a_t^i(s^t) = 0 \quad (\text{A})$$

Equivalence of equilibrium in both structures

Remember we have the following important equations in the two structures:

$$q_t^0(s^t) = \beta^t \pi_t(s^t) \frac{u'[c_t^i(s^t)]}{u'[c_0^i(s^0)]}$$

$$Q_t(s_{t+1} | s^t) = \beta \pi_t(s^{t+1} | s^t) \frac{u'[c_{t+1}^i(s^{t+1})]}{u'[c_t^i(s^t)]}$$

How do we guarantee that the allocations are the same ?

Equivalence of equilibrium in both structures

Remember we have the following important equations in the two structures:

$$q_t^0(s^t) = \beta^t \pi_t(s^t) \frac{u'[c_t^i(s^t)]}{u'[c_0^i(s^0)]}$$

$$Q_t(s_{t+1} | s^t) = \beta \pi_t(s^{t+1} | s^t) \frac{u'[c_{t+1}^i(s^{t+1})]}{u'[c_t^i(s^t)]}$$

How do we guarantee that the allocations are the same ?

Just set

Equivalence of equilibrium

$$\frac{q_{t+1}^0(s^{t+1})}{q_t^0(s^t)} \equiv Q_t(s_{t+1} | s^t)$$

Equivalence of equilibrium in both structures

Now we need to check whether the allocation in AD is comparable to that in a Sequential Market Equilibrium:

Set the following value and we will be done:

$$a_t^i(s^t) = \sum_{\tau=t}^{\infty} \sum_{s^\tau | s^t} q_\tau^t(s^\tau) [c_\tau^i(s^\tau) - y_\tau^i(s^\tau)]$$