Big-O Cheat Sheet Download PDF

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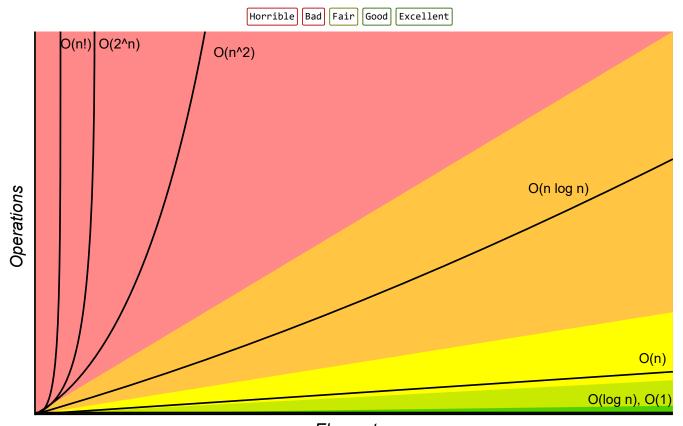
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Know Thy Complexities!

Hi there! This webpage covers the space and time Big-O complexities of common algorithms used in Computer Science. When preparing for technical interviews in the past, I found myself spending hours crawling the internet putting together the best, average, and worst case complexities for search and sorting algorithms so that I wouldn't be stumped when asked about them. Over the last few years, I've interviewed at several Silicon Valley startups, and also some bigger companies, like Google, Facebook, Yahoo, LinkedIn, and Uber, and each time that I prepared for an interview, I thought to myself "Why hasn't someone created a nice Big-O cheat sheet?". So, to save all of you fine folks a ton of time, I went ahead and created one. Enjoy! - Eric

Check out El Grapho, a graph data visualization library that supports millions of nodes and edges

Big-O Complexity Chart



Common Data Structure Operations

Data Structure	Time Con	Space Complexity							
	Average				Worst				Worst
	Access	Search	Insertion	Deletion	Access	Search	Insertion	Deletion	
<u>Array</u>	Θ(1)	Θ(n)	Θ(n)	Θ(n)	0(1)	0(n)	0(n)	0(n)	0(n)
<u>Stack</u>	Θ(n)	Θ(n)	Θ(1)	Θ(1)	0(n)	0(n)	0(1)	0(1)	0(n)
Queue	Θ(n)	Θ(n)	Θ(1)	Θ(1)	0(n)	0(n)	0(1)	0(1)	0(n)
Singly-Linked List	Θ(n)	Θ(n)	Θ(1)	Θ(1)	0(n)	0(n)	0(1)	0(1)	0(n)
Doubly-Linked List	Θ(n)	Θ(n)	Θ(1)	Θ(1)	0(n)	0(n)	0(1)	0(1)	0(n)
Skip List	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	0(n)	0(n)	0(n)	0(n)	0(n log(n))
<u>Hash Table</u>	N/A	Θ(1)	Θ(1)	Θ(1)	N/A	0(n)	0(n)	0(n)	0(n)
Binary Search Tree	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	0(n)	0(n)	0(n)	0(n)	0(n)
Cartesian Tree	N/A	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	N/A	0(n)	0(n)	0(n)	0(n)
B-Tree	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	O(log(n))	O(log(n))	0(log(n))	0(log(n))	0(n)
Red-Black Tree	$\Theta(\log(n))$	$\boxed{\Theta(\log(n))}$	$\Theta(\log(n))$	$\Theta(\log(n))$	O(log(n))	O(log(n))	0(log(n))	0(log(n))	0(n)
Splay Tree	N/A	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	N/A	0(log(n))	0(log(n))	0(log(n))	0(n)
AVL Tree	$\Theta(\log(n))$	$\boxed{\Theta(\log(n))}$	$\Theta(\log(n))$	$\boxed{\theta(\log(n))}$	O(log(n))	O(log(n))	O(log(n))	O(log(n))	0(n)
KD Tree	$\boxed{\Theta(\log(n))}$	$\boxed{\theta(\log(n))}$	$\Theta(\log(n))$	$\boxed{\theta(\log(n))}$	0(n)	0(n)	0(n)	0(n)	0(n)

Array Sorting Algorithms

Algorithm	Time Comp	olexity	Space Complexity		
	Best	Average	Worst	Worst	
Quicksort	$\Omega(n \log(n))$	Θ(n log(n))	O(n^2)	O(log(n))	
<u>Mergesort</u>	$\Omega(n \log(n))$	Θ(n log(n))	O(n log(n))	O(n)	
<u>Timsort</u>	$\Omega(n)$	Θ(n log(n))	O(n log(n))	O(n)	
<u>Heapsort</u>	$\Omega(n \log(n))$	Θ(n log(n))	0(n log(n))	0(1)	
Bubble Sort	$\Omega(n)$	Θ(n^2)	O(n^2)	0(1)	
Insertion Sort	$\Omega(n)$	Θ(n^2)	0(n^2)	0(1)	
Selection Sort	Ω(n^2)	Θ(n^2)	O(n^2)	0(1)	
Tree Sort	$\Omega(n \log(n))$	Θ(n log(n))	0(n^2)	0(n)	
Shell Sort	$\Omega(n \log(n))$	Θ(n(log(n))^2)	O(n(log(n))^2)	0(1)	
Bucket Sort	$\Omega(n+k)$	Θ(n+k)	0(n^2)	0(n)	
Radix Sort	$\Omega(nk)$	Θ(nk)	0(nk)	0(n+k)	
Counting Sort	$\Omega(n+k)$	Θ(n+k)	0(n+k)	0(k)	
<u>Cubesort</u>	$\Omega(n)$	Θ(n log(n))	O(n log(n))	0(n)	

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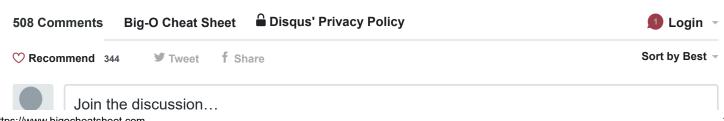


Contributors



Make this Page Better

Edit these tables!



https://www.bigocheatsheet.com



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Name



Michael Mitchell • 8 years ago

This is great. Maybe you could include some resources (links to khan academy, mooc etc) that would explain each of these concepts for people trying to learn them.

388 ^ | V 1 • Reply • Share >



Amanda Harlin → Michael Mitchell • 8 years ago

Yes! Please & thank you

85 ^ | V • Reply • Share >



Asim Ahmad → Amanda Harlin • 4 years ago

Can you Explain the Above Algorithm.??



Anonymous → Asim Ahmad • 3 years ago

Mr.

you can learn these algorithms easily in google by searching

Don't always ask or wait for someone to post things for you go out and search on internet You will find everything you want to learn

If you are a beginner in Data structures and algorithms then visit mycodeschool youtube channel and learn there

if you want more then email me at rise.d1105@gmail.com I will help you as much as I can 38 ^ | ~ 13 • Reply • Share >



Careerdrill → Anonymous • a year ago

https://www.careerdrill.com/



Trey Huffine → Asim Ahmad • a year ago

https://skilled.dev provides a detailed explanation



Cam Cecil → Michael Mitchell • 8 years ago

This explanation in 'plain English' helps: http://stackoverflow.com/qu...



Richard Wheatley - Cam Cecil • 6 years ago

this is plain english.



Tomás Gemes | xygrr → Richard Wheatley • 2 months ago

No, this is spooky english to some you know



Arjan Nieuwenhuizen → Michael Mitchell • 8 years ago • edited

Here are the links that I know of.

1 / 1111p.//auuiii.org/oourses/ai...

- #2) http://ocw.mit.edu/courses/...
- #3) https://www.udacity.com/cou...

probably as good or maybe better # 2, but I have not had a chance to look at it. http://ocw.mit.edu/courses/...

Sincerely,

Arjan

p.s.

https://www.coursera.org/co...

This course has just begun on coursera (dated 1 July 2013), and looks very good.

22 ^ | V • Reply • Share >



fireheron → Arjan Nieuwenhuizen • 8 years ago

Thank you Arjan. Espaecially the coursera.org one ;-)

5 ^ | V • Reply • Share >



@hangtwentyy → fireheron • 7 years ago

also this! http://opendatastructures.org

9 ^ | V • Reply • Share >



yth → @hangtwentyy • 6 years ago

thank you for sharing this.

1 ^ | V • Reply • Share >



Eduardo Sánchez - Michael Mitchell • 6 years ago

There is an amazing tutorial for Big O form Derek Banas in Youtube, that guy is amazing explaining!!!

see more

12 ^ | Y • Reply • Share >



Sudhanshu Mishra → Eduardo Sánchez • 5 years ago

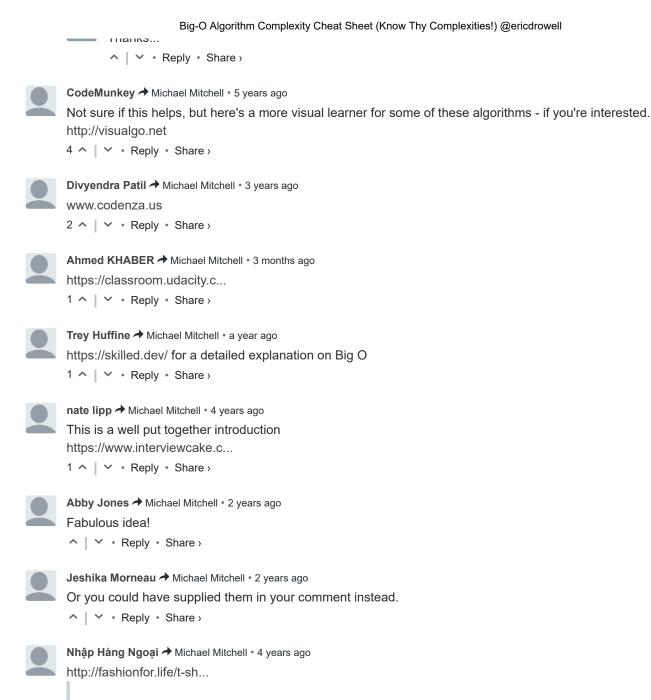
Cool! This is a more than adequate introduction! Thanks a ton for sharing!

1 ^ | V • Reply • Share >



Mohammed Hameed → Eduardo Sánchez • 2 years ago

Thanke





Blake Jennings • 8 years ago

i'm literally crying

101 ^ | Y • Reply • Share >



friend → Blake Jennings • 4 years ago

you give me a big o

5 ^ | V • Reply • Share >



Gokce Toykuyu • 9 years ago

Could we add some tree algorithms and complexities? Thanks. I really like the Red-Black trees;)

91 ^ | V • Reply • Share >



ericdrowell Mod → Gokce Toykuyu • 9 years ago

Excellent idea. I'll add a section that compares insertion, deletion, and search complexities for specific data structures

31 ^ | V • Reply • Share >



yash bedi → ericdrowell • 5 years ago

its been 4 years you haven't added that section:)

∧ | ∨ • Reply • Share >



Elliot Géhin → yash bedi • 5 years ago

It's up there Yash, bottom of the first table

1 ^ | Y • Reply • Share >



Ahsan Sukamuljo → yash bedi • a year ago

fuck off asshole

^ | ✓ 4 • Reply • Share >



Jonathan Neufeld → Gokce Toykuyu • 4 years ago

Where I come from we use trees on a regular rotation

3 ^ | V • Reply • Share >



Valentin Stanciu • 8 years ago

- 1. Deletion/insertion in a single linked list is implementation dependent. For the question of "Here's a pointer to an element, how much does it take to delete it?", single-linked lists take O(N) since you have to search for the element that points to the element being deleted. Double-linked lists solve this problem.
- 2. Hashes come in a million varieties. However with a good distribution function they are O(logN) worst case. Using a double hashing algorithm, you end up with a worst case of O(loglogN).
- 3. For trees, the table should probably also contain heaps and the complexities for the operation "Get Minimum".

63 ^ | V • Reply • Share >



Alexis Mas → Valentin Stanciu • 7 years ago

If you a list: A B C D, When you want to delete B, you can delete a node without iterating over the list.

- 1. B.data = C.data
- 2. B.next = C.next
- 3. delete C

If you can't copy data between nodes because its too expensive then yes, it's O(N)

6 ^ | V 2 • Reply • Share >



Miguel → Alexis Mas • 7 years ago

You still have to find the position in the list, which can only be done linearly.

8 ^ | V • Reply • Share >



Guest → Miguel • 7 years ago • edited

You still have to find the position in the list, which can only be done linearly.

3 ^ | Y • Reply • Share



Alexis Mas → Miguel • 7 years ago

Yes of course, If you need to search the node it's O(n), otherwise you can delete it as I stated before.

1 ^ | V • Reply • Share



Guest → Alexis Mas • 7 years ago

No need to find the position if you can delete it as Alexis mentioned

2 ^ | V 1 • Reply • Share >



Pingu App → Alexis Mas • 7 years ago

What if B is the last element in the list?

How would B's predecesor know that its next field should point to NULL and not to a futurely invalid memory address?

3 ^ | V • Reply • Share >



Alexis Mas → Pingu App • 7 years ago

In that case you can't deleted that way, you're forced to have a pointer to the previous item.

1 ^ | V 1 • Reply • Share >



Oscar Martinez Sanchez → Alexis Mas • a year ago

that's confusing, so if u want to delete the last element of the list, to maintain the delete method working u should receive two params, node to delete and previous node?

^ | ✓ • Reply • Share >



OmegaNemesis28 → Alexis Mas • 7 years ago • edited

To get to B - you HAVE to iterate over the list though. You can't just manipulate B without a pointer. So unless you do book-keeping and have pointers to specific nodes you intend to delete/manipulate, LinkLists are O(n) insert and delete.

3 ^ | V • Reply • Share >



Alexis Mas → OmegaNemesis28 • 7 years ago

Strictly speaking no, you don't. let's say you have this function.

public void delete(Node node)

That function doesn't care how did you got that node.

Did you got my point?

When you have a pointer to a node, and that node needs to be deleted you don't need to iterave over the list.

1 ^ | ~ 2 • Reply • Share >



Sam Lehman → Alexis Mas • 6 years ago

But in order to get to that pointer, you probably need to iterate through the list

2 ^ | V • Reply • Share >



OmegaNemesis28 → Alexis Mas • 7 years ago • edited

But that is MY point :p

You have to have the node FIRST. You have to iterate through the list before you can do that, unless you do book-keeping and happen to have said node. Reread what I said. "have pointers to specific nodes" Most of the time, you do not with LinkedLists. If you have a Linked List and want to delete index 5, you have to iterate to 5 and such. Your example was ABCD, our points are that you typically don't have the pointer to B just offhand. You have to obtain it first which will be O(n)

2 ^ | V • Reply • Share >



Chris B → OmegaNemesis28 • 6 years ago

Search and insert/delete are different operations. Insert/delete on an unsorted linked list is O(1). The fact that you might have to first search for the element that you want to delete is not considered relevant, as that functionality is covered by the O(n) search operation, not the O(1) insert/delete operations. A real world example of linked list insert/delete can be found in list del and list add of the Linux kernel source, those functions are only 2 and 4 lines of code, so should be easy to understand: http://lxr.free-electrons.c...

2 ^ | Y • Reply • Share >



pvlbzn → Alexis Mas • 5 years ago • edited

And you will introduce the side effect which will be hell to debug. Consider:

Singly linked list { A:1, B:2, C:3, D:4 } where is X:Y, y is a value, function `delete` which works as you described, function 'get' which returns pointer to the node by index.

// Take needed node C node t* node = get(list, 2) print(node->value) // prints 3 // Delete B delete(list, 1) // Try to access C again print(node->value) // well, enjoy your O(1) Don't. 2 ^ | V 1 • Reply • Share >



Reehz Ree • 6 years ago