HW 3: DATA DEPENDENCES

DIHAN DAI

February 22, 2020

QUESTION 1

- (a) The validity condition for loop unrolling is that the dependence vector of permutation corresponding to moving the unrolled loop to the innermost loop is valid. Therefore,
 - (0,0,1,-1): unrolling t, i and k loops is valid.
 - (1, -2, 1, -1): unrolling i, j and k loops is valid.
 - (1, -1, 1, 0): unrolling i, j and k loops is valid.
- (b) The validity condition for loop permutation is that the dependence vector of the corresponding permutation is valid.
 - (0,0,1,-1): tijk, itjk, tjki, ijkt, jkti, jkti, tjik, ijtk, jtki, jitk, jtik
 - (1, -2, 1, -1): tijk, tikj, tjki, tjik, tkij, tk
ji, jtik, jtki, jkti, jkti, jitk, jikt
 - (1, -1, 1, 0): tijk, tikj, tjki, tjik, tkij, tkij, jtik, jtki, jkit, jkit, jitk, jikt, ktji, ktij, kjit, kjit
- (c) The vaidity condition for loop tiling is that the corresponding loops are fully permutable.
 - (0,0,1,-1): 3D tiling: tij, 2D tiling: ti, tj, ij,
 - (1, -2, 1, -1): 3D tiling: ijk, 2D tiling: tj, ij, ik, jk,
 - (1,-1,1,0): 3D tiling: ijk, 2D tiling: tj, ij, ik, jk,

2 DIHAN DAI

QUESTION 2

- (a) Yes. If there is an output dependences between (t_1, i_1, j_1) and (t_2, i_2, j_2) such that $(t_1, i_1, j_1) < (t_2, i_2, j_2)$, it must satisfy $(i_1 + 1, j_1 1) = (i_2 + 1, j_2 1)$. The dependence vectors are given by $(t_2 t_1, i_2 i_1, j_2 j_1)$ and the smallest of such vectors is (1, 0, 0). The leftmost nonzero element of the vector is positive, thus the dependence vector is valid.
- (b) Flow dependence: if there is a flow dependences between the write instance (t_w, i_w, j_w) and the read instance (t_r, i_r, j_r) . The dependences vector are given by $(t_r t_w, i_r i_w, j_r j_w)$ and they must satisfies
 - either $(i_r, j_r + 1) = (i_w + 1, j_w 1)$, where the smallest dependence vector is (0, 1, -2),
 - or $(i_r 1, j_r) = (i_w + 1, j_w 1)$, where the smallest dependence vector is (0, 2, -1). Both dependence vectors are valid because the leftmost nonzero elements of both vectors are positive.
 - Anti-dependence: if there is an anti-dependence between the read instance (t_r, i_r, j_r) and the write instance (t_w, i_w, j_w) . The dependences vector are given by $(t_w t_r, i_w i_r, j_w j_r)$ and they must satisfies
 - either $(i_r, j_r + 1) = (i_w + 1, j_w 1)$, where the smallest dependence vector is (1, -1, 2),
 - or $(i_r-1,j_r)=(i_w+1,j_w-1)$, where the smallest dependence vector is (1,-2,-1). Both dependence vectors are valid because the leftmost nonzero elements of both vectors are positive.

QUESTION 3

(i)

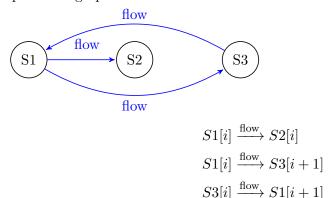
i loop:
$$S1[i, j, k] \to S3[i+1, j, k],$$

k loop: $S2[i, j, k] \to S1[i, j, k+1],$
j loop: $S3[i, j, k] \to S2[i, j+1, k].$

- (ii) No parallelization can be achieved since every loop carries dependence.
- (iii) No. It is because the dependences of every loop doesn't change no matter what permutation of the nested loops is.

QUESTION 4

(i) Dependence graph:



Statement 1 and 3 form a cycle, thus only statement 2 can be vectorized.

(ii)
$$\begin{aligned} & \textbf{for} \, (\, i = \! 1; \; i < \! 256; \; i \! + \! + \!) \\ & \{ \\ & a \, [\, i \,] \; = \; c \, [\, i \,] \! + \! 1; \\ & c \, [\, i + \! 1] \; = \; a \, [\, i \, - \! 1] \! + \! 1; \\ & \} \\ & b \, [\, 1 : \! 255 \,] \; = \; a \, [\, 1 : \! 255 \,] \, - \! 1; \end{aligned}$$

QUESTION 5

(i) Data dependences:

$$S[i, j, k] \to S[i + 1, j - 1, k - 1]$$
 flow