

Program Verification: While Loops

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Exercises 15 : Reading and More

Record your time spent (in 0.1 hours) with brief tasks and durations in your learning log by hand writing!

- 1) Read [textB-ch04-ProgramVerification.pdf](#) (cont.)

Outline

Program Verification: While Loops

Key Points

Proving Partial Correctness - Example 1

Proving Partial Correctness - Example 2

Invariants on Sorting Algorithms

Proving Termination

Summary

Key Points

Key points to learn:

Partial correctness for while loops

- ▶ Determine whether a given formula is an invariant for a while loop.
- ▶ Find an invariant for a given while loop.
- ▶ Prove that a Hoare triple is satisfied under partial correctness for a program containing while loops.

Total correctness for while loops

- ▶ Determine whether a given formula is a variant for a while loop.
- ▶ Find a variant for a given while loop.
- ▶ Prove that a Hoare triple is satisfied under total correctness for a program containing while loops.

Proving Total Correctness of While Loops

- ▶ Partial correctness
- ▶ Termination

Proving Partial Correctness of While Loops

$\{P\}$
 $\{I\}$ implied (A)
while (B) {
 $\{ (I \wedge B) \}$ partial-**while**
 C
 $\{I\}$ <justify based on C – a subproof>
}
 $\{ (I \wedge (\neg B)) \}$ partial-**while**
 $\{Q\}$ implied (B)

Proof of implied (A): $(P \rightarrow I)$

Proof of implied (B): $((I \wedge (\neg B)) \rightarrow Q)$

I is called a **loop invariant**. We need to determine I !

What is a loop invariant?

A **loop invariant** is:

- ▶ A relationship among the variables. (A predicate formula involving the variables.)
- ▶ The word “invariant” means something that does not change.
- ▶ It is true before the loop begins.
- ▶ It is true at the start of every iteration of the loop and at the end of every iteration of the loop.
- ▶ It is true after the loop ends.

Proving partial correctness of while loops

Indicate the places in the program where the loop invariant is true.

$\llbracket (x \geq 0) \rrbracket$

$y = 1;$

$z = 0;$

while $(z \neq x)$ {

$z = z + 1;$

$y = y * z;$

}

$\llbracket (y = x!) \rrbracket$

Proving partial correctness of a while loop

Steps to follow:

- ▶ Find a loop invariant.
- ▶ Complete the annotations.
- ▶ Prove any implied's.

How do we find a loop invariant???

How do we find a loop invariant?

First, we need to understand the purpose of an invariant.

- ▶ The postcondition is the ultimate goal of our while loop.
- ▶ At every iteration, we are making progress towards the postcondition.
- ▶ The invariant is describing the progress we are making at every iteration.

Partial While - Example 1

Example 1:

Prove that the following triple is satisfied under partial correctness.

```
⌊(x ≥ 0)⌋  
y = 1;  
z = 0;  
while (z != x) {  
    z = z + 1;  
    y = y * z;  
}  
⌊(y = x!)⌋
```

Finding a loop invariant

Step 1: Write down the values of all the variables every time the while test is reached.

$\langle (x \geq 0) \rangle$

$y = 1;$

$z = 0;$

while $(z \neq x)$ {

$z = z + 1;$

$y = y * z;$

}

$\langle (y = x!) \rangle$

Finding a loop invariant

Step 2: Find relationships among the variables that are true for every while test. These are our candidate invariants.

Come up with some invariants in the next 2 minutes.

```
⌊(x ≥ 0)⌋  
y = 1;  
z = 0;  
while (z != x) {  
    z = z + 1;  
    y = y * z;  
}  
⌊(y = x!)⌋
```

x	z	y
5	0	1 = 0!
5	1	1 = 1!
5	2	2 = 2!
5	3	6 = 3!
5	4	24 = 4!
5	5	120 = 5!

CQ 1 Is this a loop invariant?

CQ 1: Is $(\neg(z = x))$ a loop invariant?

(A) Yes (B) No (C) I don't know...

```
⌈(x ≥ 0)⌋  
y = 1;  
z = 0;  
while (z != x) {  
    z = z + 1;  
    y = y * z;  
}  
⌋(y = x!)⌋
```

x	z	y
5	0	1 = 0!
5	1	1 = 1!
5	2	2 = 2!
5	3	6 = 3!
5	4	24 = 4!
5	5	120 = 5!

CQ 2 Is this a loop invariant?

CQ 2: Is $(z \leq x)$ a loop invariant?

(A) Yes (B) No (C) I don't know...

```
⌈(x ≥ 0)⌋  
y = 1;  
z = 0;  
while (z != x) {  
    z = z + 1;  
    y = y * z;  
}  
⌋(y = x!)⌋
```

x	z	y
5	0	1 = 0!
5	1	1 = 1!
5	2	2 = 2!
5	3	6 = 3!
5	4	24 = 4!
5	5	120 = 5!

CQ 3 Is this a loop invariant?

CQ 3: Is $(y = z!)$ a loop invariant?

(A) Yes (B) No (C) I don't know...

```
⌈(x ≥ 0)⌋  
y = 1;  
z = 0;  
while (z != x) {  
    z = z + 1;  
    y = y * z;  
}  
⌋(y = x!)⌋
```

x	z	y
5	0	1 = 0!
5	1	1 = 1!
5	2	2 = 2!
5	3	6 = 3!
5	4	24 = 4!
5	5	120 = 5!

CQ 4 Is this a loop invariant?

CQ 4: Is $(y = x!)$ a loop invariant?

(A) Yes (B) No (C) I don't know...

```
⌊(x ≥ 0)⌋  
y = 1;  
z = 0;  
while (z != x) {  
    z = z + 1;  
    y = y * z;  
}  
⌊(y = x!)⌋
```

x	z	y
5	0	1 = 0!
5	1	1 = 1!
5	2	2 = 2!
5	3	6 = 3!
5	4	24 = 4!
5	5	120 = 5!

CQ 5 Is this a loop invariant?

CQ 5: Is $((z \leq x) \wedge (y = z!))$ a loop invariant?

(A) Yes (B) No (C) I don't know...

```
⌈(x ≥ 0)⌋  
y = 1;  
z = 0;  
while (z != x) {  
    z = z + 1;  
    y = y * z;  
}  
⌋(y = x!)⌋
```

x	z	y
5	0	1 = 0!
5	1	1 = 1!
5	2	2 = 2!
5	3	6 = 3!
5	4	24 = 4!
5	5	120 = 5!

Finding a loop invariant

Step 3: Try each candidate invariant until we find one that works for our proof.

```
 $\langle\langle x \geq 0 \rangle\rangle$   
y = 1;  
z = 0;  
while (z != x) {  
    z = z + 1;  
    y = y * z;  
}  
 $\langle\langle y = x! \rangle\rangle$ 
```

x	z	y
5	0	1 = 0!
5	1	1 = 1!
5	2	2 = 2!
5	3	6 = 3!
5	4	24 = 4!
5	5	120 = 5!

How do we find an invariant?

A recap of the steps to find an invariant:

- ▶ Write down the values of all the variables every time the while test is reached.
- ▶ Find relationships among the variables that are true for every while test. These are our candidate invariants.
- ▶ Try each candidate invariant until we find one that works for our proof.

Partial While - Example 1 ($(z \leq x)$ as the invariant)

$\llbracket (x \geq 0) \rrbracket$	
$\llbracket (0 \leq x) \rrbracket$	implied (A)
$y = 1;$	
$\llbracket (0 \leq x) \rrbracket$	assignment
$z = 0;$	
$\llbracket (z \leq x) \rrbracket$	assignment
while ($z \neq x$) {	
$\llbracket ((z \leq x) \wedge (\neg(z = x))) \rrbracket$	partial- while
$\llbracket (z + 1 \leq x) \rrbracket$	implied (B)
$z = z + 1;$	
$\llbracket (z \leq x) \rrbracket$	assignment
$y = y * z;$	
$\llbracket (z \leq x) \rrbracket$	assignment
}	
$\llbracket ((z \leq x) \wedge (\neg(\neg(z = x)))) \rrbracket$	partial- while
$\llbracket (y = x!) \rrbracket$	implied (C)

CQ 7 Is there a proof for implied (A)?

We used $(z \leq x)$ as the invariant.

CQ 7: Is there a proof for implied (A)?

$$((x \geq 0) \rightarrow (0 \leq x))$$

(A) Yes

(B) No

(C) I don't know.

CQ 8 Is there a proof for implied (B)?

We used $(z \leq x)$ as the invariant.

CQ 8: Is there a proof for implied (B)?

$$(((z \leq x) \wedge (\neg(z = x))) \rightarrow (z + 1 \leq x))$$

(A) Yes

(B) No

(C) I don't know.

CQ 9 Is there a proof for implied (C)?

We used $(z \leq x)$ as the invariant.

CQ 9: Is there a proof for implied (C)?

$$(((z \leq x) \wedge (\neg(\neg(z = x)))) \rightarrow (y = x!))$$

(A) Yes

(B) No

(C) I don't know.

Partial While - Example 2

Example 2:

Prove that the following triple is satisfied under partial correctness.

```
⌊(x ≥ 0)⌋  
y = 1;  
z = 0;  
while (z < x) {  
    z = z + 1;  
    y = y * z;  
}  
⌊(y = x!)⌋
```

Let's try using $(y = z!)$ as the invariant in our proof.

Which invariant leads to a valid proof?

To check whether an invariant leads to a valid proof,
we need to check whether all of the implied's can be proved.

CQ 11 Is there a proof for implied (A)?

We used $(y = z!)$ as the invariant.

CQ 11: Is there a proof for implied (A)?

$$((x \geq 0) \rightarrow (1 = 0!))$$

(A) Yes

(B) No

(C) I don't know.

CQ 12 Is there a proof for implied (B)?

We used $(y = z!)$ as the invariant.

CQ 12: Is there a proof for implied (B)?

$$(((y = z!) \wedge (z < x)) \rightarrow (y * (z + 1) = (z + 1)!))$$

(A) Yes

(B) No

(C) I don't know.

CQ 13 Is there a proof for implied (C)?

We used $(y = z!)$ as the invariant.

CQ 13: Is there a proof for implied (C)?

$$(((y = z!) \wedge (\neg(z < x))) \rightarrow (y = x!))$$

(A) Yes

(B) No

(C) I don't know.

CQ 14 Is there a proof for implied (C)?

We used $((y = z!) \wedge (z \leq x))$ as the invariant.

CQ 14: Is there a proof for implied (C)?

$$(((y = z!) \wedge (z \leq x)) \wedge (\neg(z < x))) \rightarrow (y = x!))$$

(A) Yes

(B) No

(C) I don't know.

Selection sort

Algorithm. ↑ scans from left to right.

Invariants.

- Entries the left of ↑ (including ↑) fixed and in ascending order.
- No entry to right of ↑ is smaller than any entry to the left of ↑.



Insertion sort

Algorithm. ↑ scans from left to right.

Invariants.

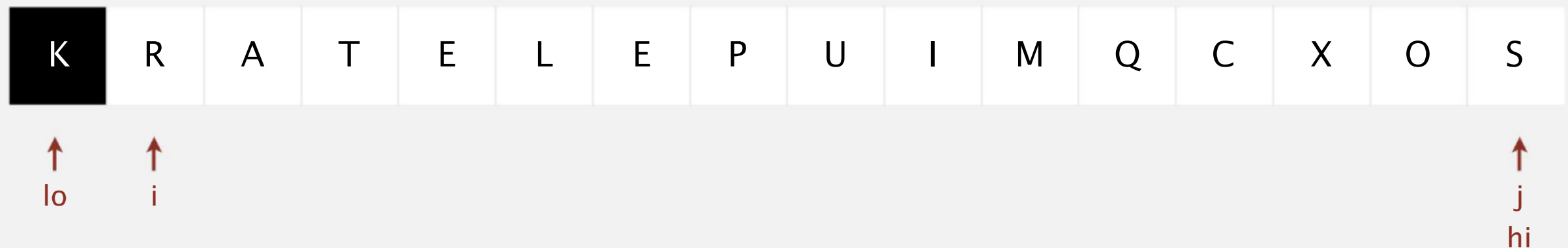
- Entries to the left of ↑ (including ↑) are in ascending order.
- Entries to the right of ↑ have not yet been seen.



Quicksort Partitioning

Repeat until i and j pointers cross.

- Scan i from left to right so long as $(a[i] < a[l_o])$.
- Scan j from right to left so long as $(a[j] > a[l_o])$.
- Exchange $a[i]$ with $a[j]$.



<https://algs4.cs.princeton.edu/lectures/demo/23DemoPartitioning.pdf>

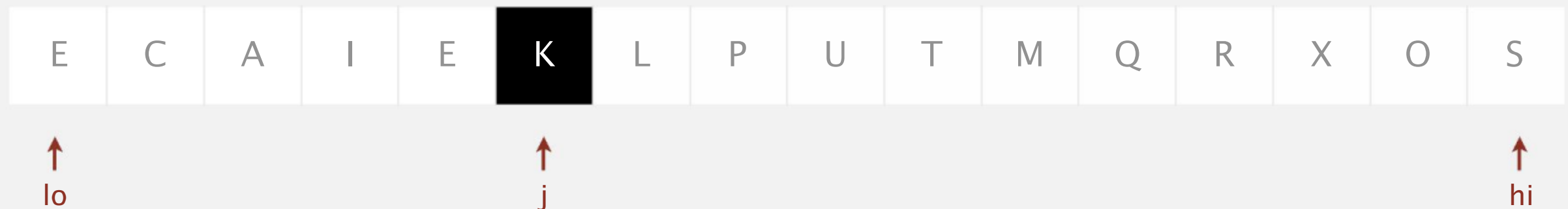
Quicksort Partitioning

Repeat until i and j pointers cross.

- Scan i from left to right so long as $(a[i] < a[lo])$.
- Scan j from right to left so long as $(a[j] > a[lo])$.
- Exchange $a[i]$ with $a[j]$.

When pointers cross.

- Exchange $a[lo]$ with $a[j]$.



partitioned!

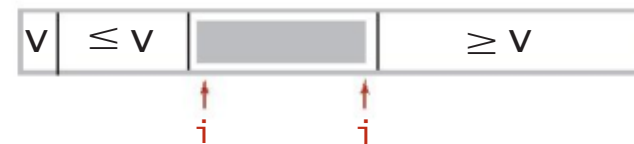
Quicksort: Java code for partitioning

```
private static int partition(Comparable[] a, int lo, int hi) {  
    int i = lo, j = hi+1;  
    while (true) {  
        while (less(a[++i], a[lo]))           find item on left to swap  
            if (i == hi) break;  
  
        while (less(a[lo], a[--j]))           find item on right to swap  
            if (j == lo) break;  
  
        if (i >= j) break;                     check if pointers cross  
        exch( a, i, j);                       swap  
    }  
  
    exch( a, lo, j);                          swap with partitioning item  
    return j;                                return index of item now known to be in place  
}
```

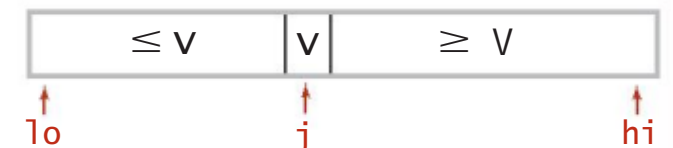
before



during




after



Quicksort: Java implementation

```
public class Quick {  
  
    private static int partition(Comparable[] a, int lo, int hi) {  
        /* see previous slide */  
    }  
  
    public static void sort (Comparable[] a) {  
        sort( a, 0, a.length - 1);  
    }  
  
    private static void sort (Comparable[] a, int lo, int hi) {  
        if (hi <= lo) return;  
        int j = partition( a, lo, hi);  
        sort( a, lo, j-1);  
        sort( a, j+1, hi);  
    }  
}
```

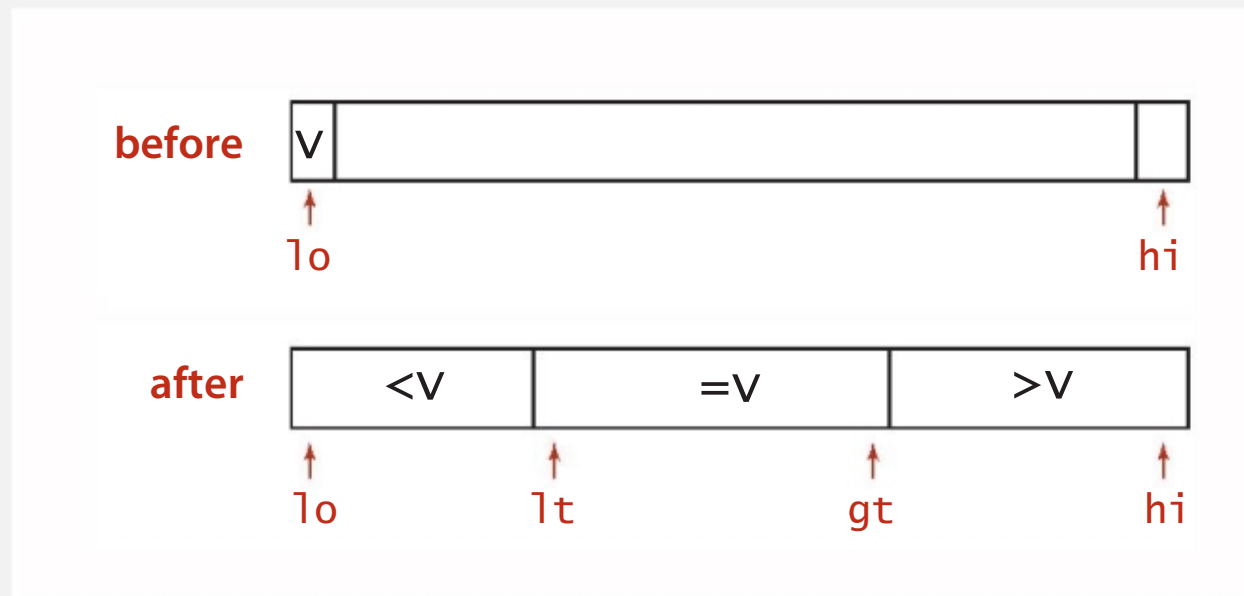
Call overloaded
sort method



3-way Partitioning

Goal. Partition array into **three** parts so that:

- Entries between lt and gt equal to the partition item.
- No larger entries to left of lt .
- No smaller entries to right of gt .



Dutch national flag problem. [Edsger Dijkstra]

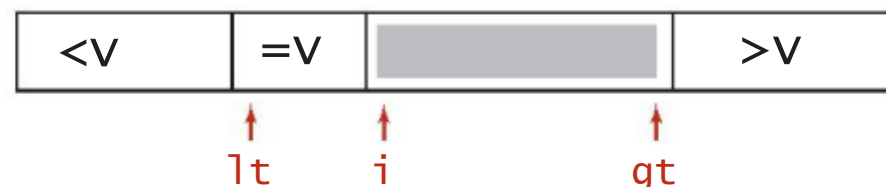
- Conventional wisdom until mid 1990s: not worth doing.
- Now incorporated into C library `qsort()` and Java 6 system sort.

Dijkstra 3-way partitioning

- Let v be partitioning item $a[lo]$.
- Scan i from left to right.
 - ($a[i] < v$): exchange $a[lt]$ with $a[i]$; increment both lt and i
 - ($a[i] > v$): exchange $a[gt]$ with $a[i]$; decrement gt
 - ($a[i] == v$): increment i

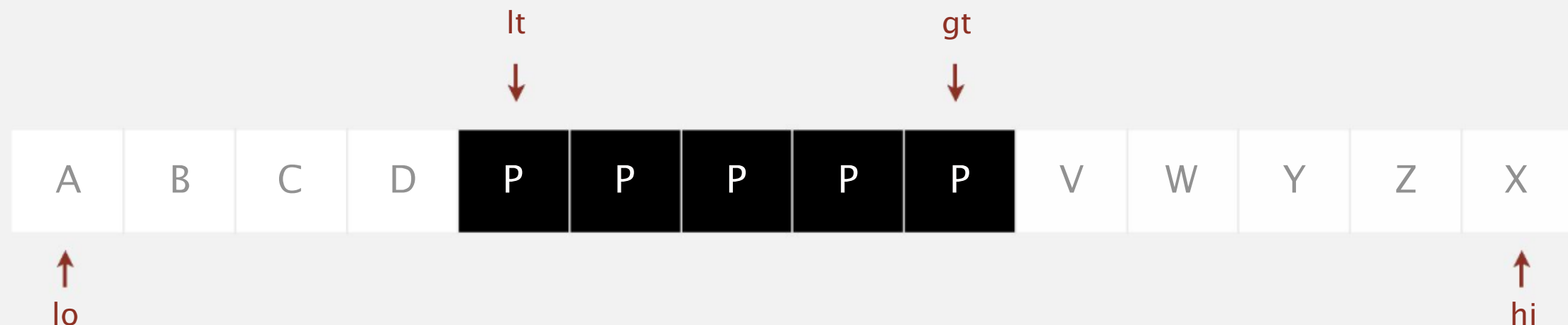


invariant

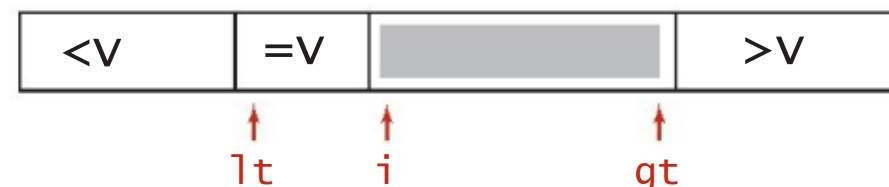


Dijkstra 3-way partitioning

- Let v be partitioning item $a[lo]$.
- Scan i from left to right.
 - ($a[i] < v$): exchange $a[lt]$ with $a[i]$; increment both lt and i
 - ($a[i] > v$): exchange $a[gt]$ with $a[i]$; decrement gt
 - ($a[i] == v$): increment i

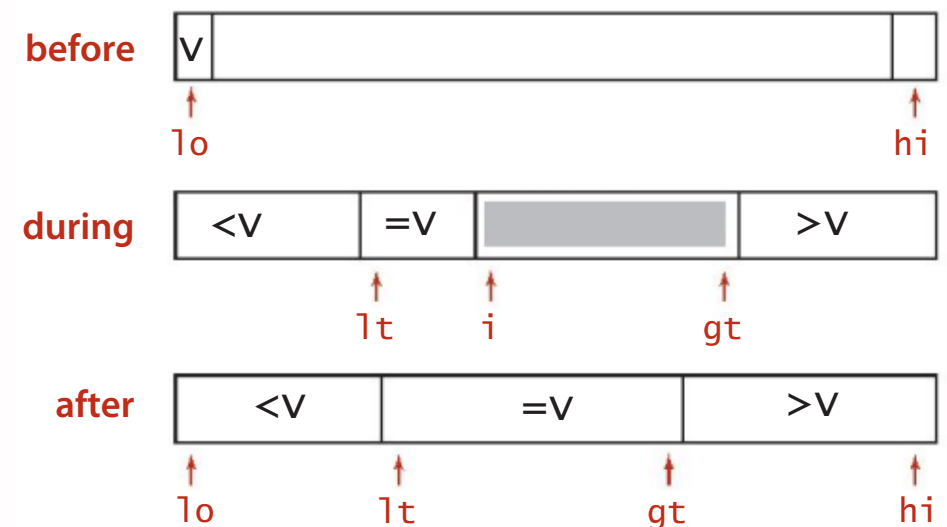


invariant



3-way quicksort: Java implementation

```
private static void sort (Comparable[] a, int lo, int hi) {  
    if (hi <= lo) return;  
  
    int lt = lo, gt = hi;  
    Comparable v = a[lo];  
    int i = lo;  
  
    while (i <= gt) {  
        int cmp = a[i].compareTo(v);  
        if (cmp < 0) exch( a, lt++, i++);  
        else if (cmp > 0) exch( a, i, gt--);  
        else i++;  
    }  
  
    sort( a, lo, lt - 1);  
    sort( a, gt + 1, hi);  
}
```



Proving Termination

Find an integer expression that

- ▶ is non-negative before the loop starts, at every iteration of the loop, and after the loop ends.
- ▶ decreases by at least 1 at every iteration of the loop.

This integer expression is called a **variant** (something that changes).

The loop must terminate because a non-negative integer can decrease by 1 a finite number of times.

Example 2: Finding a variant

Example 2:

Prove that the following program terminates.

```
y = 1;  
z = 0;  
while (z < x) {  
    z = z + 1;  
    y = y * z;  
}
```

How do we find a variant? The loop guard ($z < x$) helps.

Example 2: Proof of Termination

Consider the variant $(x - z)$.

Before the loop starts, $(x - z) \geq 0$ because the precondition is $(x \geq 0)$ and the second assignment mutates z to be 0.

During every iteration of the loop, $(x - z)$ decreases by 1 because x does not change and z increases by 1.

Thus, $x - z$ will eventually reach 0.

When $x - z = 0$, the loop guard $z < x$ will terminate the loop.

Summary

Key points learnt:

Partial correctness for while loops

- ▶ Determine whether a given formula is an invariant for a while loop.
- ▶ Find an invariant for a given while loop.
- ▶ Prove that a Hoare triple is satisfied under partial correctness for a program containing while loops.

Total correctness for while loops

- ▶ Determine whether a given formula is a variant for a while loop.
- ▶ Find a variant for a given while loop.
- ▶ Prove that a Hoare triple is satisfied under total correctness for a program containing while loops.