

Assignment 02 of I2ML-s23

Q1 Use truth tables to prove the validity of the 3 hardest sequents you thought from the following list. Pay attention that $A, B, C \vdash D$ is equivalent to $A \wedge B \wedge C \rightarrow D$. (3 * 10 = 30 pts)

1. $p \vee q \vdash q \vee p$

2. $(p \wedge q) \wedge r \vdash p \wedge (q \wedge r)$

3. $p \rightarrow q \vdash \neg q \rightarrow \neg p$

4. $(p \wedge q) \vee (p \wedge r) \vdash p \wedge (q \vee r)$

5. $\vdash \neg p \vee q \rightarrow (p \rightarrow q)$

6. $\vdash (p \rightarrow q) \rightarrow \neg p \vee q$

7. $\neg(p \rightarrow q) \vdash q \rightarrow p$

8. $W \rightarrow X, Y \rightarrow Z \vdash W \vee Y \rightarrow X \vee Z$

9. $\vdash \neg(A \wedge B) \rightarrow (A \rightarrow \neg B)$

10. $p \wedge q \rightarrow r \vdash (p \rightarrow r) \vee (q \rightarrow r)$

11. $\vdash \neg(A \leftrightarrow \neg A)$

12. $A \leftrightarrow B \vdash \neg A \leftrightarrow \neg B$

Q2 Use truth tables to prove the validity of the 3 hardest formulas you thought from the following list. Pay attention that do not repeat with Q1. (3 * 10 = 30 pts)

- | | |
|--|---|
| 1. Commutativity of \wedge : $A \wedge B \leftrightarrow B \wedge A$ | 12. $\neg(A \wedge \neg A)$ |
| 2. Commutativity of \vee : $A \vee B \leftrightarrow B \vee A$ | 13. $\neg(A \rightarrow B) \leftrightarrow A \wedge \neg B$ |
| 3. Associativity of \wedge : $(A \wedge B) \wedge C \leftrightarrow A \wedge (B \wedge C)$ | 14. $\neg A \rightarrow (A \rightarrow B)$ |
| 4. Associativity of \vee : $(A \vee B) \vee C \leftrightarrow A \vee (B \vee C)$ | 15. $(\neg A \vee B) \leftrightarrow (A \rightarrow B)$ |
| 5. Distributivity of \wedge over \vee : $A \wedge (B \vee C) \leftrightarrow (A \wedge B) \vee (A \wedge C)$ | 16. $A \vee \perp \leftrightarrow A$ |
| 6. Distributivity of \vee over \wedge : $A \vee (B \wedge C) \leftrightarrow (A \vee B) \wedge (A \vee C)$ | 17. $A \wedge \perp \leftrightarrow \perp$ |
| 7. $(A \rightarrow (B \rightarrow C)) \leftrightarrow (A \wedge B \rightarrow C)$. | 18. $A \vee \neg A$ |
| 8. $(A \rightarrow B) \rightarrow ((B \rightarrow C) \rightarrow (A \rightarrow C))$ | 19. $\neg(A \leftrightarrow \neg A)$ |
| 9. $((A \vee B) \rightarrow C) \leftrightarrow (A \rightarrow C) \wedge (B \rightarrow C)$ | 20. $(A \rightarrow B) \leftrightarrow (\neg B \rightarrow \neg A)$ |
| 10. $\neg(A \vee B) \leftrightarrow \neg A \wedge \neg B$ | 21. $(A \rightarrow C \vee D) \rightarrow ((A \rightarrow C) \vee (A \rightarrow D))$ |
| 11. $\neg(A \wedge B) \leftrightarrow \neg A \vee \neg B$ | 22. $((A \rightarrow B) \rightarrow A) \rightarrow A$ |

Q3 For the given formula below, draw its parse tree and list the set of its all subformulas. (10 + 10 = 20 pts)

$$((p \rightarrow \sim q) \vee (p \wedge r) \rightarrow s) \vee \sim r$$

Q4 Finish 3 tasks for each parse tree in Figure A and Figure B: **1)** List the corresponding formula with fully dressed parentheses; **2)** List the shortest formula with proper precedence order and right associative discussed in lectures; **3)** Given a valuation/model and a parse tree of a formula, compute the truth value of the formula for that valuation/model (as done in a bottom-up fashion in Figure 1.7 on page 40 of textB). ($2 * (5+5+5) = 30$ pts)

4.1 **Figure A** in which *q evaluates to T and p and r evaluate to F*;

4.2 **Figure B** where we *let p be T, q be F and r be T*.

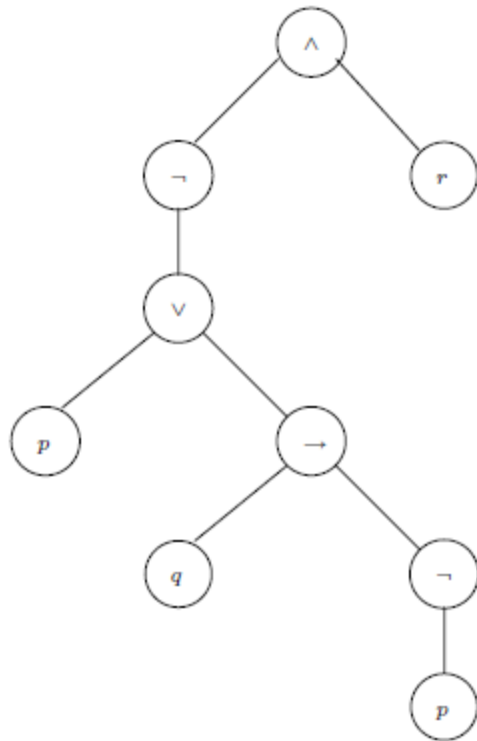


Figure A. In-order linear representation of the tree is a logical formula

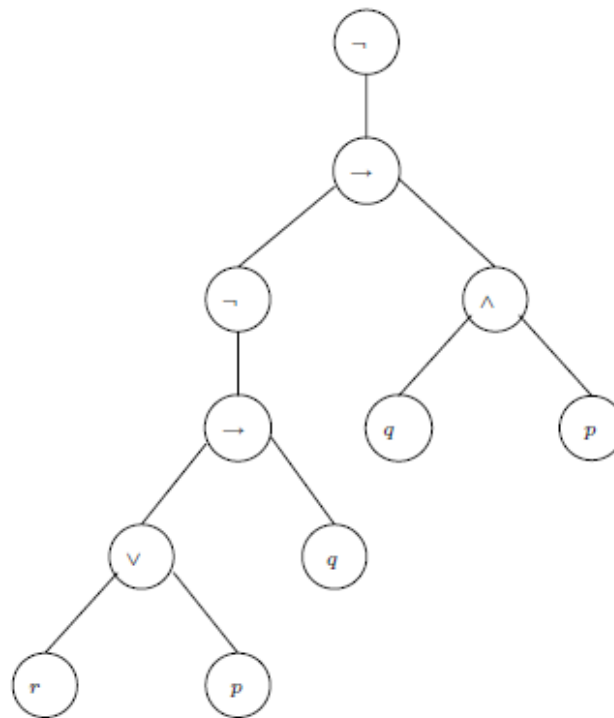


Figure B. A parse tree of a negated implication

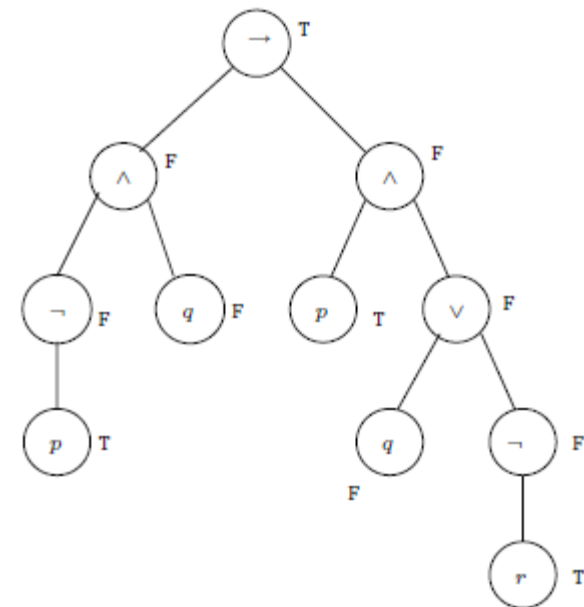


Figure 1.7. The evaluation of a logical formula under a given valuation
(On page 40 of textB)