

### **CS323** Lab 6

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### Outline

• Bison introduction

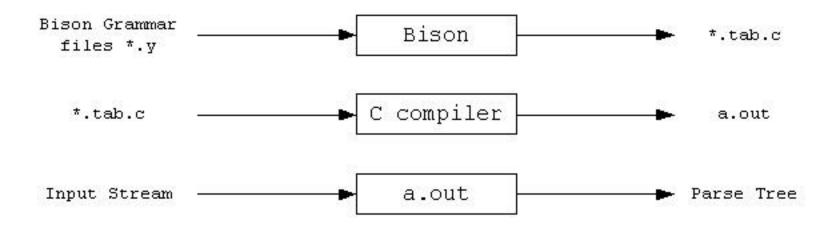
• Bison tutorial

### The Parser Generator Bison

• Bison generates a parser, which accepts the input token stream from Flex, to do syntax analysis, according to the specified CFG.

- Bison的前身为基于Unix的Yacc (Yet another compiler compiler)。Yacc的 发布时间比Lex还要早,其采用的语法分析技术的理论基础早在20世纪50年代就已经由Knuth逐步建立了起来,而Yacc本身则是贝尔实验室的S.C. Johnson基于这些理论在1975年到1978年写成的。
- 到了1985年,当时在UC Berkeley的一个研究生Bob Corbett在BSD下重写了Yacc,取名为Bison (美洲野牛,yak牦牛的近亲),后来GNU Project接管了这个项目,为其增加了许多新的特性,于是就有了我们今天所用的GNU Bison。

# Input/Output of Bison (Yacc)



### Structure of YACC Source Programs

- Declarations (声明)
  - Ordinary C declarations
  - Grammar tokens
- Translation rules (翻译规则)
  - Rule = a production + semantic action

declarations

%%

translation rules

%%

supporting C routines

- Supporting C routines (辅助性C语言例程)
  - Directly copied to y.tab.c
  - Can be invoked in the semantic actions
  - Other procedures such as error recovery routines may be provided

### **Translation Rules**

- The first head in the list of rules is taken as the start symbol of the grammar
- A semantic action is a sequence of C statements
  - \$\$ is a Bison's internal variable that holds the semantic value of the left-hand side of a production rule (i.e., the whole construct)
  - \$i holds the semantic value of ith grammar symbol of the body
- A semantic action is performed when we apply the associated production for reduction (归约,the reverse of rewrite)
  - Most of the time, the purpose of an action is to compute the semantic value of the whole construct from the semantic values of its parts, i.e., compute \$\$ using \$i's

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#### A Parser That Performs Calculation

- When processing arithmetic expressions, lexical anlayzer will recognize the following types of tokens: INT, ADD, SUB, MUL, DIV
  - When analyzing the input string "3 + 6 / 2", the lexer will output this string of tokens: INT ADD INT DIV INT
- After lexical analysis, the parser will check if the string of tokens
  produced by lexer has a valid structure or not according to the synatatic
  specification described by the following context-free grammar

```
Calc -> Exp
Exp -> Factor | Exp ADD Factor | Exp SUB Factor
Factor -> Term | Factor MUL Term | Factor DIV Term
Term -> INT
```

<sup>\*</sup>Red for non-terminals, Blue for terminals, Calc is the start symbol

# Valid and Invalid Inputs

#### Valid input expressions:

- 3
- 3 + 5 \* 4
- 3+6/2
- 3+2-1
- •

#### Invalid input expressions:

- -3
- 3+5\*
- ..

### Flex/Bison Code for Calculator

syntax.y

#### lex.l

```
%{
    #include "syntax.tab.h"
    #include "stdlib.h"

%}
%%
[0-9]+ { yylval = atoi(yytext); return INT; }
"+" { return ADD; }
"-" { return SUB; }
"*" { return MUL; }
"/" { return DIV; }
[ \n\r\t] {}
. { fprintf(stderr, "unknown symbol: %s\n", yytext); exit(1); }
```

#### yylval:

- Flex internal variable that is used to store the attribute of a recognized token
- Its data type is YYSTYPE (int by default)\*
- After storing values to yylval in Flex code, the values will be propagated to Bison (i.e., the syntax anlayzer part) and can be retrieved using \$n

```
응 {
    #include "lex.yy.c"
    void yyerror(const char*);
응 }
%token INT
%token ADD SUB MUL DIV
응응
Calc: /* to allow empty input */
    Exp { printf("= %d\n", $1); }
Exp: Factor
    | Exp ADD Factor \{ \$\$ = \$1 + \$3; \}
      Exp SUB Factor { $$ = $1 - $3; }
Factor: Term
    | Factor MUL Term { \$\$ = \$1 * \$3; }
    | Factor DIV Term { \$\$ = \$1 / \$3; }
Term: INT
응응
void yyerror(const char *s) {
    fprintf(stderr, "%s\n", s);
int main() {
    yyparse(); // will invoke yylex(
}
```

\*Can be customized by putting command like #define YYSTYPE char\* at the beginning of .l and .y files.

#### A Further Look at Semantic Actions

```
응 {
                                 syntax.y
    #include "lex.yy.c"
    void yyerror(const char*);
응 }
%token INT
%token ADD SUB MUL DIV
응응
Calc: /* to allow empty input */
      Exp { printf("= %d\n", $1);
Exp: Factor
    | Exp ADD Factor \{ \$\$ = \$1 + \$3; \}
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Factor: Term
    | Factor MUL Term { $$ = $1 * $3; }
    | Factor DIV Term { $$ = $1 / $3; }
Term: INT
응 응
void yyerror(const char *s) {
    fprintf(stderr, "%s\n", s);
int main() {
    yyparse();
```

When appyling this rule for reduction, the semantic analysis is successful (i.e., the start symbol calc can generate the input expression), we output the calculation result.

#### A Further Look at Semantic Actions

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    #include "lex.yy.c"
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Calc: /* to allow empty input */
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void yyerror(const char *s) {
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```

Intermediate calculation of semantic values of a language construct represented by the head symbol

#### A Further Look at Semantic Actions

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응응
void yyerror(const char *s) {
    fprintf(stderr, "%s\n", s);
int main() {
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```

Although no action is specified, Bison will still propagate the attribute of the only symbol INT to the head Term

Same as writing:  $\{\$\$ = \$1;\}$ 

```
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응응
void yyerror(const char *s) {
    fprintf(stderr, "%s\n", s);
int main() {
    yyparse();
```

Input: 3 + 5
Token string: INT ADD INT

3 INT

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Calc: /* to allow empty input */
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void yyerror(const char *s) {
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```

Input: 3 + 5
Token string: INT ADD INT

3 Term
3 INT

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void yyerror(const char *s) {
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int main() {
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```

Input: 3 + 5
Token string: INT ADD INT

3 Factor
3 Term
3 INT

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Input: 3 + 5
Token string: INT ADD INT

3 Exp
3 Factor
3 Term
3 INT

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```
Input: 3 + 5
Token string: INT ADD INT
```

```
3 Exp ADD

3 Factor

3 Term

3 INT
```

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```

```
Input: 3 + 5
Token string: INT ADD INT
```

```
3 Exp ADD

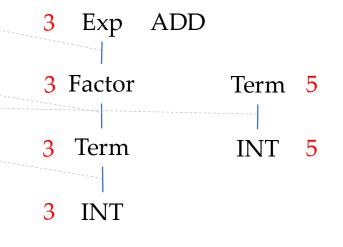
3 Factor

3 Term INT 5

3 INT
```

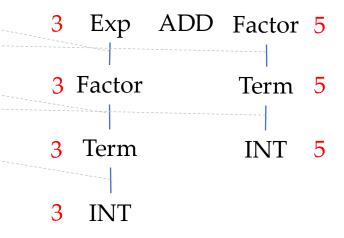
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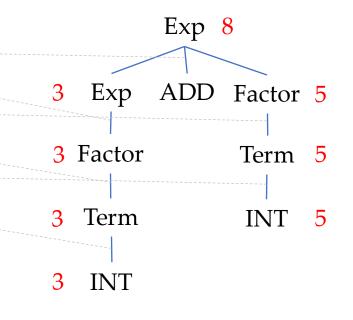
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응응
void yyerror(const char *s) {
    fprintf(stderr, "%s\n", s);
int main() {
   yyparse();
```

```
Input: 3 + 5
Token string: INT ADD INT
         Calc Print "=8"
         Exp 8
3 Exp
        ADD Factor 5
3 Factor
               Term 5
  Term
                INT 5
  INT
```

## Compile the Example

- Compile the Bison code into C code: a header file and an implementation file
  - bison -d syntax.y (why it is compilable without lex.yy.c?)
  - The command produces <u>syntax.tab.h</u> and <u>syntax.tab.c</u>
- Compile the flex source code
  - flex lex.1
  - The command produce <u>lex.yy.c</u>
- Putting things together
  - gcc syntax.tab.c -lfl -ly -o calc.out
  - The options -1f1 and -1y tell gcc to include Flex and Bison libraries

The code is available at lab5/calc. We provide a build target calc to facilitate the compilation.

### Run the Calculator

• In the runs below, we use pipes\* to pass the output of the first command to be the input of the second command

```
liu@liu-VirtualBox:~/Desktop/CS323-2022F/lab4/calc$ echo "3" | ./calc.out
= 3
liu@liu-VirtualBox:~/Desktop/CS323-2022F/lab4/calc$ echo "3+5*4" | ./calc.out
= 23
liu@liu-VirtualBox:~/Desktop/CS323-2022F/lab4/calc$ echo "3+6/2" | ./calc.out
= 6
liu@liu-VirtualBox:~/Desktop/CS323-2022F/lab4/calc$ echo "3+ 2-1" | ./calc.out
= 4
liu@liu-VirtualBox:~/Desktop/CS323-2022F/lab4/calc$ echo "-3" | ./calc.out
syntax error
liu@liu-VirtualBox:~/Desktop/CS323-2022F/lab4/calc$ echo "3+5*" | ./calc.out
syntax error
```

<sup>\*</sup> https://linuxhint.com/what-is-pipe-in-linux/

# Project Reminder

- Please start to design and implement your language if you haven't done so.
  - It is also fine if you choose to build a compiler for SPL or its simple variants.

- Milestone check: Nov. 18, during the lab session.
  - Please prepare test cases by yourself for the demo.
  - You should also prepare a report, which should at least contain: 1) the specification and core features of your language, 2) the design of your compiler, 3) the implementation progress.