

```

printf ("Top element : %d\n", *stack);
// Popping elements from the stack
printf ("Popped element : %d\n", pop(&stack));
printf ("Popped element : %d\n", pop(&stack));
// check if the stack is empty
printf ("Is the stack empty? %s\n",
isEmpty(&stack) ? "Yes" : "No");

return 0;
}

```

\* queues } abstract data structure (core data structure)  
stack

2023/06/21

### Tutorials (8)

① What is a circular queue?

A circular queue also known as "Ring Buffer" is a data structure that follows first-in-first-out (FIFO) principle.

It is implemented as an array/queue with a fixed size, where the last position is connected to the first position / first index forming a circular behaviour.

② What are the characteristics of circular queue?

Pointers in stack : Top/peer

- front and rear as pointers : In circular queues two pointers ~~front~~ maintain the order of the circular queue.
- The circular behaviour : When the data structure has reached its capacity the next

element will be inserted into a vacant position in the front of the data structure.

- Fixed size : The circular queues are initialized with a fixed size which prevents them from growing dynamically.

③ Give applications of circular queue?

- Cache management.
- Printer spooling
- Buffer management
- Event handling.
- \* Traffic management
- \* Memory management
- \* OS / CPU scheduling.

④ What is the algorithm of circular queue?





ALGO tute 8 question 4 answer;

The algorithm for a circular queue typically involves the following operations:

1. Initialize the circular queue:

Create an array of a fixed size to hold the elements of the queue.

Initialize two pointers, "front" and "rear," both set to -1 (indicating an empty queue).

Check if the queue is empty:

If  $\text{front} == -1$  and  $\text{rear} == -1$ , the queue is empty.

Check if the queue is full:

If  $(\text{rear} + 1) \% \text{size} == \text{front}$ , the queue is full, where "size" represents the maximum capacity of the queue.

Enqueue (insert) an element into the circular queue:

2. Check if the queue is full.

If it is, return an error or perform appropriate error handling.

If the queue is empty ( $\text{front} == -1$  and  $\text{rear} == -1$ ), set both front and rear pointers to 0.

Otherwise, increment the rear pointer by 1:  $\text{rear} = (\text{rear} + 1) \% \text{size}$ .

Insert the new element at the rear position in the array:  $\text{queue}[\text{rear}] = \text{element}$ .

Dequeue (remove) an element from the circular queue:

3. Check if the queue is empty.

If it is, return an error or perform appropriate error handling.

Store the value of the front element to return it later:  $\text{element} = \text{queue}[\text{front}]$ .

If  $\text{front} == \text{rear}$  (indicating the queue has only one element), set both front and rear pointers back to -1.

Otherwise, increment the front pointer by 1:  $\text{front} = (\text{front} + 1) \% \text{size}$ .

Get the front element of the circular queue without removing it:

4. Check if the queue is empty.

If it is, return an error or perform appropriate error handling.

Return the value of the front element:  $\text{element} = \text{queue}[\text{front}]$ .

Get the size of the circular queue:

If  $\text{front} == -1$  and  $\text{rear} == -1$ , the queue is empty, so return 0.

If  $\text{rear} \geq \text{front}$ , the size of the queue is  $(\text{rear} - \text{front} + 1)$ .

If  $\text{rear} < \text{front}$ , the size of the queue is  $(\text{size} - \text{front} + \text{rear} + 1)$ .