

# Network attacks



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# Hacking phases

1. Reconnaissance
2. Scanning
3. (enumeration)
- 4. Gaining Access**
  - a. Network Attacks**
  - b. System attacks
5. Maintaining access/escalating privileges
6. Clearing traces



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# MAC address

- “A media access control address (MAC address) **is a unique identifier assigned to a network interface controller (NIC)** for use as a network address in communications within a network segment.”
- “MAC addresses are primarily assigned by device manufacturers, and are therefore often referred to as the **burned-in address**, or as an **Ethernet hardware address**, hardware address”
- Is MAC unique?



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# MAC spoofing

```
pavle@pavle-ideapad:~$ macchanger -s wlp2s0
Current MAC: 70:c9:4e: (unknown)
Permanent MAC: 70:c9:4e: (unknown)
pavle@pavle-ideapad:~$ sudo ip link set dev wlp2s0 down
pavle@pavle-ideapad:~$ sudo macchanger -r wlp2s0
Current MAC: 70:c9:4e: (unknown)
Permanent MAC: 70:c9:4e: (unknown)
New MAC: ee:a5:b6:62:d4:c6 (unknown)
pavle@pavle-ideapad:~$ sudo ip link set dev wlp2s0 up
pavle@pavle-ideapad:~$ ifconfig wlp2s0
wlp2s0: flags=4163<UP,BROADCAST,RUNNING,MULTICAST> mtu 1500
    inet 192.168.1.109 netmask 255.255.255.0 broadcast 192.168.1.255
    ether ee:a5:b6:62:d4:c6 txqueuelen 1000 (Ethernet)
    RX packets 315963 bytes 425155682 (425.1 MB)
    RX errors 0 dropped 0 overruns 0 frame 0
    TX packets 189688 bytes 22451111 (22.4 MB)
    TX errors 0 dropped 0 overruns 0 carrier 0 collisions 0

pavle@pavle-ideapad:~$ sudo ip link set dev wlp2s0 down
pavle@pavle-ideapad:~$ sudo macchanger -p wlp2s0
Current MAC: ee:a5:b6:62:d4:c6 (unknown)
Permanent MAC: 70:c9:4e: (unknown)
New MAC: 70:c9:4e: (unknown)
pavle@pavle-ideapad:~$ sudo ip link set dev wlp2s0 up
pavle@pavle-ideapad:~$ ifconfig wlp2s0
wlp2s0: flags=4163<UP,BROADCAST,RUNNING,MULTICAST> mtu 1500
    inet 192.168.1.109 netmask 255.255.255.0 broadcast 192.168.1.255
    ether 70:c9:4e: txqueuelen 1000 (Ethernet)
    RX packets 315975 bytes 425157258 (425.1 MB)
    RX errors 0 dropped 0 overruns 0 frame 0
    TX packets 189691 bytes 22451596 (22.4 MB)
    TX errors 0 dropped 0 overruns 0 carrier 0 collisions 0
```



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# Switched network sniffing

- Switches fill the MAC address table
- MAC flooding:
  - Fill the switch MAC table with fake MAC addresses (macof tool: <https://kalilinuxtutorials.com/macof/>)
  - When the table is full, switch starts to flood packets and acts like a hub
  - solution: **switch port security** – allow only single MAC address per port



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# ARP protocol

```
▶ Frame 118: 42 bytes on wire (336 bits), 42 bytes captured (336 bits) on interface 0
▼ Ethernet II, Src: LiteonTe_ (70:c9: ), Dst: Broadcast (ff:ff:ff:ff:ff:ff)
  ▶ Destination: Broadcast (ff:ff:ff:ff:ff:ff)
  ▶ Source: LiteonTe_ (70:c9: )
  Type: ARP (0x0806)
▼ Address Resolution Protocol (request)
  Hardware type: Ethernet (1)
  Protocol type: IPv4 (0x0800)
  Hardware size: 6
  Protocol size: 4
  Opcode: request (1)
  Sender MAC address: LiteonTe_ (70:c9: )
  Sender IP address: 192.168.1.109
  Target MAC address: 00:00:00_00:00:00 (00:00:00:00:00:00)
  Target IP address: 192.168.1.103
```

```
▼ Ethernet II, Src: SamsungE_ (f4:0e: ), Dst: LiteonTe_ (70:c9: )
  ▶ Destination: LiteonTe_ (70:c9: )
  ▶ Source: SamsungE_ (f4:0e: )
  Type: ARP (0x0806)
▼ Address Resolution Protocol (reply)
  Hardware type: Ethernet (1)
  Protocol type: IPv4 (0x0800)
  Hardware size: 6
  Protocol size: 4
  Opcode: reply (2)
  Sender MAC address: SamsungE_ (f4:0e: )
  Sender IP address: 192.168.1.103
  Target MAC address: LiteonTe_ (70:c9: )
  Target IP address: 192.168.1.109
```



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# ARP table

```
pavle@pavle-ideapad:~$ arp
```

Address	HWtype	HWaddress	Flags	Mask	Iface
192.168.1.100	ether	d0:63:b4:00:59:40	C		wlp2s0
_gateway	ether	c0:c1:c0:c3:45:6f	C		wlp2s0
192.168.1.103	ether	f4:0e:22:f4:15:83	C		wlp2s0
192.168.1.101	ether	a4:50:46:11:5f:8f	C		wlp2s0



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# ARP poisoning

- ARP maps IP addresses to specific MAC addresses:
  - ARP query – (sent to MAC:FF:FF:FF:FF:FF:FF): What is the MAC address of the machine with IP address X?
  - ARP response from machine with IP address X – sends from the MAC address of X
- What if Y responds with MAC Y to the ARP query for IP X?
  - ARP spoofing – initiator gets the wrong MAC-IP mapping and sends packets destined to IP X to MAC Y.
- ARP poisoning tools: Ettercap, Cain & Abel
- IP DHCP snooping and ARP snooping for avoiding these attacks



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# DHCP protocol

- Dynamic Host Configuration Protocol

50	1275.0370591...	0.0.0.0	255.255.255.255	DHCP	342	DHCP Discover	- Transaction ID 0x5ef36f47
64	1277.8162619...	0.0.0.0	255.255.255.255	DHCP	342	DHCP Discover	- Transaction ID 0x5ef36f47
67	1278.8902579...	192.168.1.1	255.255.255.255	DHCP	590	DHCP Offer	- Transaction ID 0x5ef36f47
68	1278.8905763...	0.0.0.0	255.255.255.255	DHCP	342	DHCP Request	- Transaction ID 0x5ef36f47
69	1278.9032657...	192.168.1.1	255.255.255.255	DHCP	590	DHCP ACK	- Transaction ID 0x5ef36f47

▼ Option: (55) Parameter Request List

Length: 16

Parameter Request List Item: (1) Subnet Mask  
Parameter Request List Item: (28) Broadcast Address  
Parameter Request List Item: (2) Time Offset  
Parameter Request List Item: (3) Router  
Parameter Request List Item: (15) Domain Name  
Parameter Request List Item: (6) Domain Name Server  
Parameter Request List Item: (119) Domain Search  
Parameter Request List Item: (12) Host Name  
Parameter Request List Item: (44) NetBIOS over TCP/IP Name Server  
Parameter Request List Item: (47) NetBIOS over TCP/IP Scope  
Parameter Request List Item: (26) Interface MTU  
Parameter Request List Item: (121) Classless Static Route  
Parameter Request List Item: (42) Network Time Protocol Servers  
Parameter Request List Item: (249) Private/Classless Static Route (Microsoft)  
Parameter Request List Item: (33) Static Route  
Parameter Request List Item: (252) Private/Proxy autodiscovery

▼ Option: (53) DHCP Message Type (ACK)  
Length: 1  
DHCP: ACK (5)  
▼ Option: (54) DHCP Server Identifier  
Length: 4  
DHCP Server Identifier: 192.168.1.1  
▼ Option: (51) IP Address Lease Time  
Length: 4  
IP Address Lease Time: (86400s) 1 day  
▼ Option: (1) Subnet Mask  
Length: 4  
Subnet Mask: 255.255.255.0  
▼ Option: (3) Router  
Length: 4  
Router: 192.168.1.1  
▼ Option: (6) Domain Name Server  
Length: 4  
Domain Name Server: 109.122.98.6  
▶ Option: (255) End



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# DHCP starvation attack

- DoS type of attack
- Attacker sends fake requests with spoofed MAC addresses forcing DHCP server to lease all the IP addresses in the pool
- New users are not able to log in
- Tools: dhcpstarv, yersinia



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# Rogue DHCP server attack

- Attacker starves the real DHCP server
- Attacker acts as a DHCP server, sending default gateway IP address and/or DNS IP address
- DHCP protection – DHCP snooping feature on network devices – allows DHCP traffic only from the configured/authorized DHCP server



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# IP and TCP based attacks

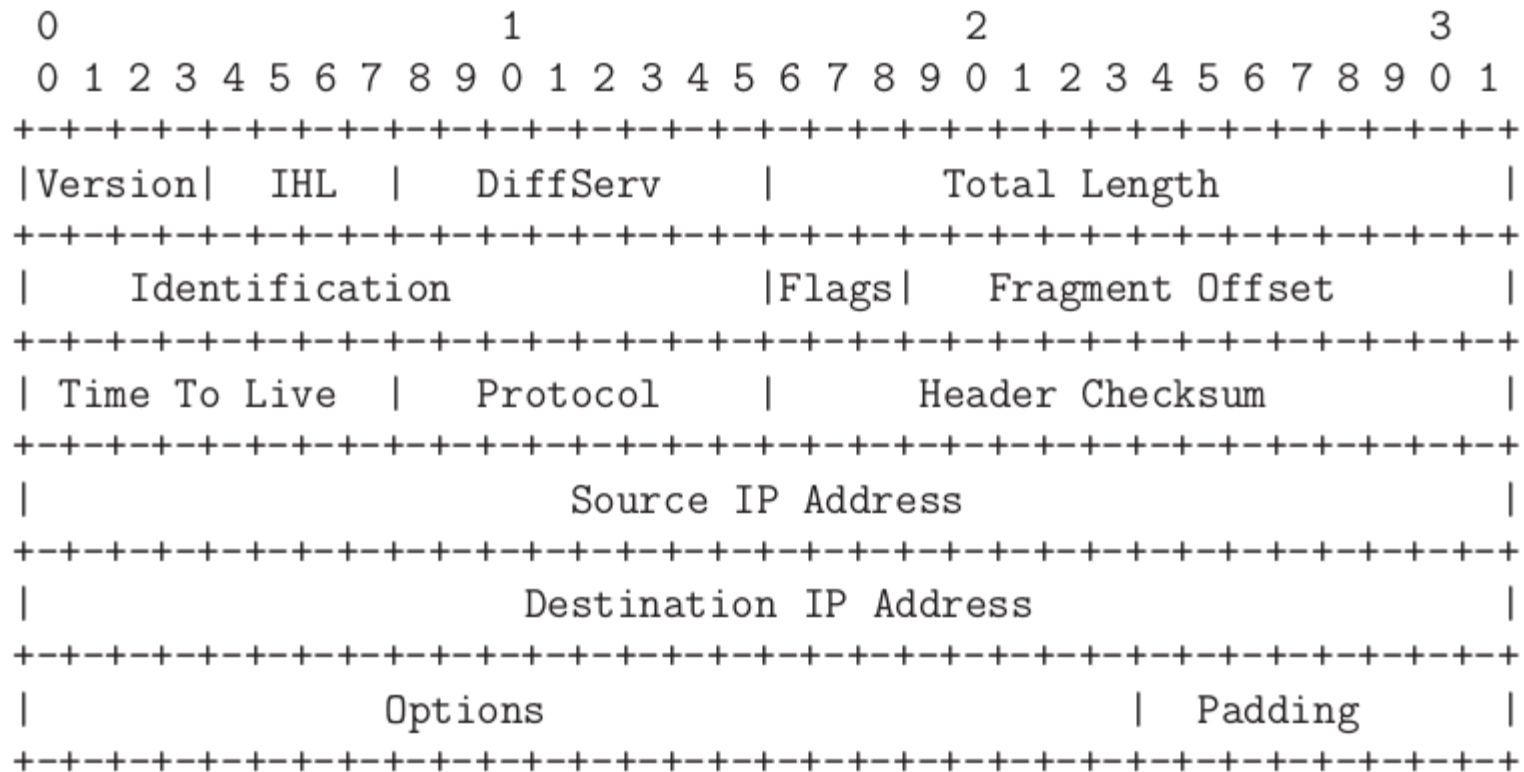


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# IPv4

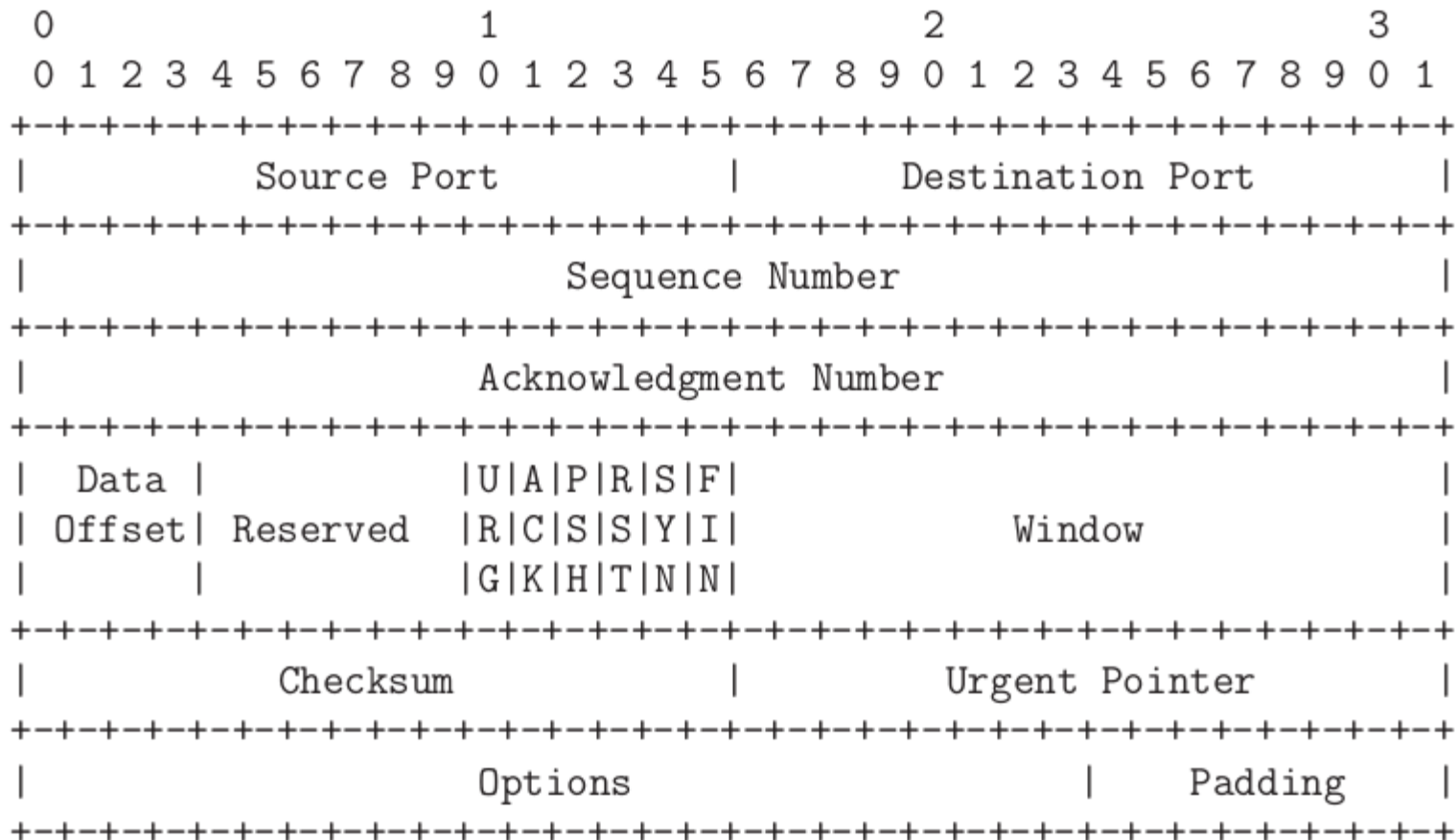


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# TCP



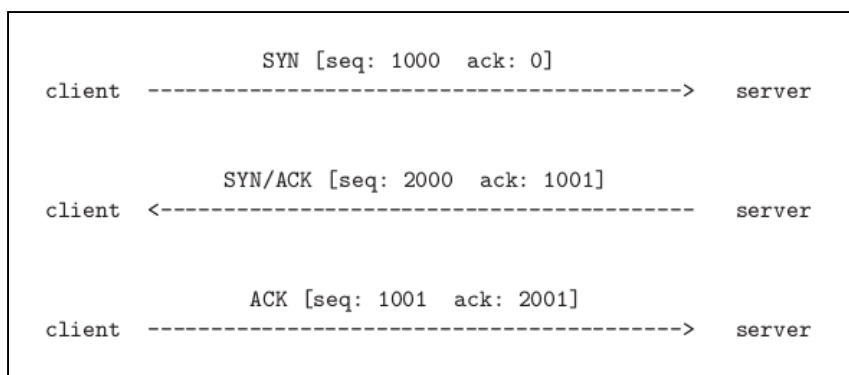
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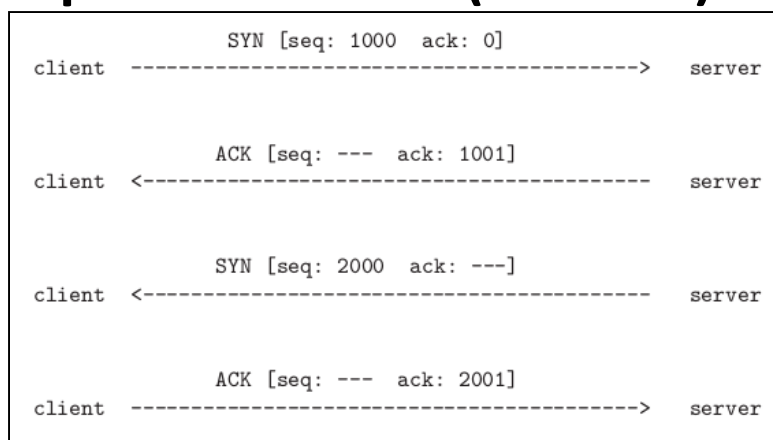
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# TCP handshake and split handshake

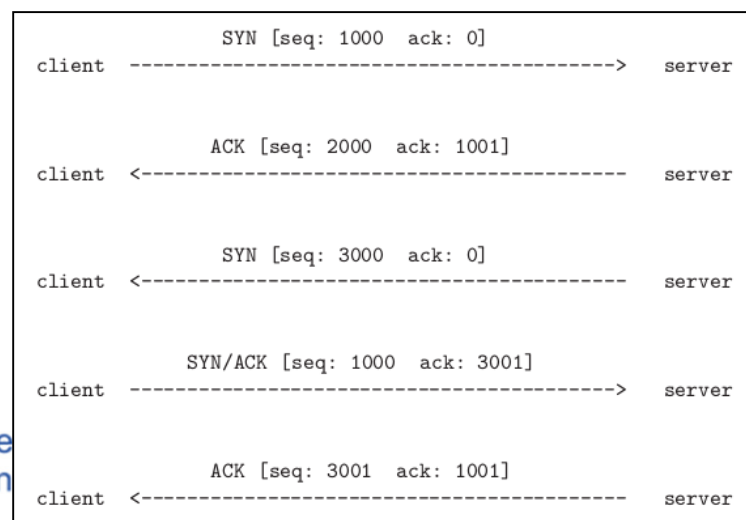
## 3-way handshake



## Split handshake (RFC 793)



- Who sends SYN/ACK?
- Split handshake can trick the IDS/FW software on the client side



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# Shrew DoS attack (1)

- The retransmission decision for a TCP segment is based on logic that operates at two different timescales:
  - When traffic congestion is low, the timescale used for determining the frequency of retransmission is RTT (Round Trip Time), which is typically of the order of a few tens of milliseconds.
  - However, when congestion is high, the frequency of retransmission is determined by the much longer RTO (Retransmission Timeout), which is generally of the order of a full second.



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# Shrew DoS attack (2)

- TCP AIMD congestion control:
  - Upon ACK:  $CWND = CWND + a$  (e.g.  $a = 1 \text{ SegmentSize}$ )
  - No ACK in RTT:  $CWND = CWND * b$  (e.g.  $b = 1/2$ )
  - No ACK in RTO – severe congestion  $CWND = 1$
- Initial value of RTO depends on RTT (RFC 2988).
- When RTT cannot be measured, the initial value for RTO value is set to 3s, minimum 1s.
- If no ACK is received within an RTO, the value of RTO doubles with each subsequent timeout.



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# Shrew DoS attack (3)

- Attacker creates a short burst which congests the link and causes regular TCP packets to be lost.
- Burst duration is approx. RTT. (so that it can't be measured)
- Next burst is after RTO (e.g. 1s)... (CWND=1)
- Next burst is after double RTO... (CWND=1)
- TCP sender stays in slow start state (CWND=1) while the average attacker traffic volume is low.



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# SYN Flooding

- Attacker repeatedly sends SYN TCP segments to every port on the server using a fake IP address.
- The server responds to each such attempt with a SYN/ACK segment from each open port and with an RST segment from each closed port.
- Attacker never sends back the expected ACK segment.
- As soon as a connection for a given port gets timed out, another SYN request arrives for the same port from the attacker.
- When a connection for a given port at the server gets into this state of receiving a never-ending stream of SYN segment the intruder has a sort of perpetual half-open connection with the victim host.
- Server can protect its resources by rate limiting all incoming SYN packets.



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# IP address spoofing

- IP source address spoofing refers to an intruder using one or more forged source IP addresses to launch, say, a TCP SYN flood attack on a host in another network.
- Solution ingress filtering: RFC 2827 – do not allow routers to send out packets if their source IP address does not fall in the range assigned to that network.
- Also if there is NAT between the attacker and a victim, IP spoofing is not possible



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# DNS based attacks



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# DNS cache poisoning

- Putting wrong (IP address, symbolic name) mapping into the DNS servers cache.
- The attacker has to know (by sniffing) or guess destination port (16bit) and Transaction ID (16bit) which is randomly generated
- The attacker has to plant fake responses before the proper answer comes to the server
- Due to the birthday paradox the attacker does not have to do the brute force of all Transaction IDs – few hundred Transaction IDs are sufficient
- The attacker can increase the TTL of the fake DNS entry



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# DNS query/response messages

```
▶ Frame 24: 84 bytes on wire (672 bits), 84 bytes captured (672 bits) on interface 0
▶ Ethernet II, Src: LiteonTe_d0:8b:bf (70:c9:4e:d0:8b:bf), Dst: Cisco-Li_b6:6f:53 (00:1c:10:b6:6f:53)
▶ Internet Protocol Version 4, Src: 192.168.1.112, Dst: 109.122.98.6
▶ User Datagram Protocol, Src Port: 58744, Dst Port: 53
▼ Domain Name System (query)
  Transaction ID: 0xa911
  ▶ Flags: 0x0100 Standard query
  Questions: 1
  Answer RRs: 0
  Authority RRs: 0
  Additional RRs: 1
  ▼ Queries
    ▼ startpage.com: type A, class IN
      Name: startpage.com
      [Name Length: 13]
      [Label Count: 2]
      Type: A (Host Address) (1)
      Class: IN (0x0001)
  ▼ Additional records
    ▶ <Root>: type OPT
    [Response In: 25]
```

```
▶ Frame 25: 132 bytes on wire (1056 bits), 132 bytes captured (1056 bits) on interface
▶ Ethernet II, Src: Cisco-Li_b6:6f:53 (00:1c:10:b6:6f:53), Dst: LiteonTe_d0:8b:bf (70:
▶ Internet Protocol Version 4, Src: 109.122.98.6, Dst: 192.168.1.112
▶ User Datagram Protocol, Src Port: 53, Dst Port: 58744
▼ Domain Name System (response)
  Transaction ID: 0xa911
  ▶ Flags: 0x8180 Standard query response, No error
  Questions: 1
  Answer RRs: 3
  Authority RRs: 0
  Additional RRs: 1
  ▶ Queries
  ▼ Answers
    ▼ startpage.com: type A, class IN, addr 145.131.132.68
      Name: startpage.com
      Type: A (Host Address) (1)
      Class: IN (0x0001)
      Time to live: 30
      Data length: 4
      Address: 145.131.132.68
    ▶ startpage.com: type A, class IN, addr 145.131.132.84
    ▶ startpage.com: type A, class IN, addr 89.146.4.147
  ▼ Additional records
    ▶ <Root>: type OPT
    [Request In: 24]
    [Time: 0.034915497 seconds]
```

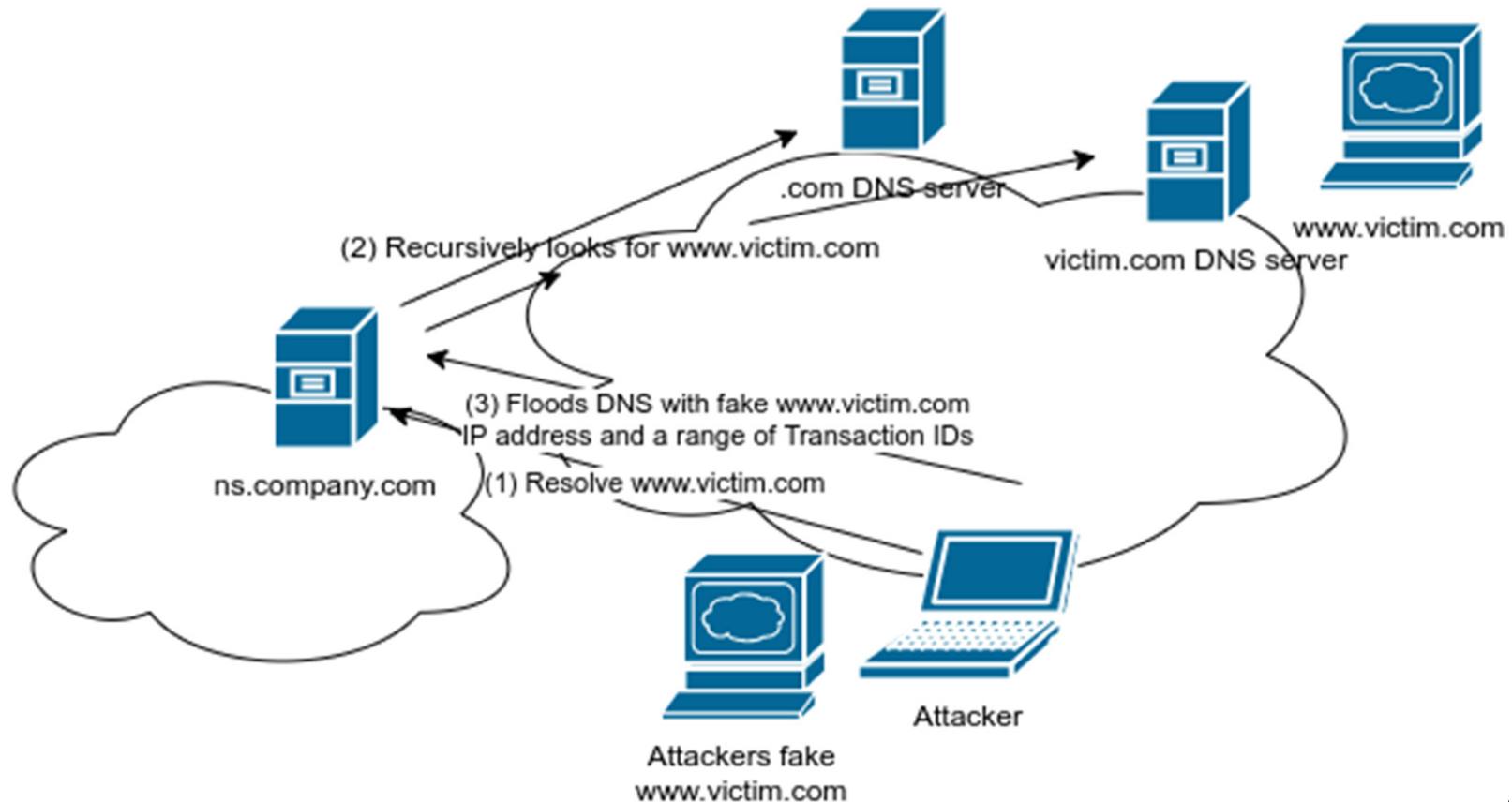


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# DNS cache poisoning



- If Additional section is used, the IP address of the victim.com DNS can be poisoned as well



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# Kaminsky DNS poisoning attack

- DNS server accepts records which it has not asked for
- Attacker asks for the resolution of: 1.victim.com, 2.victim.com,...,x.victim.com which do not exist
- Attacker tries to poison the cache of these non-existent entries before the answer from ns.victim.com comes as before
- In the additional records of all these attempts it puts the fake IP address of the ns.victim.com
- Measures against: randomize port and Transaction ID numbers & bailiwick check
- Look at:  
[https://www.youtube.com/watch?time\\_continue=940&v=qftKfFVHVuY&feature=emb\\_title](https://www.youtube.com/watch?time_continue=940&v=qftKfFVHVuY&feature=emb_title)

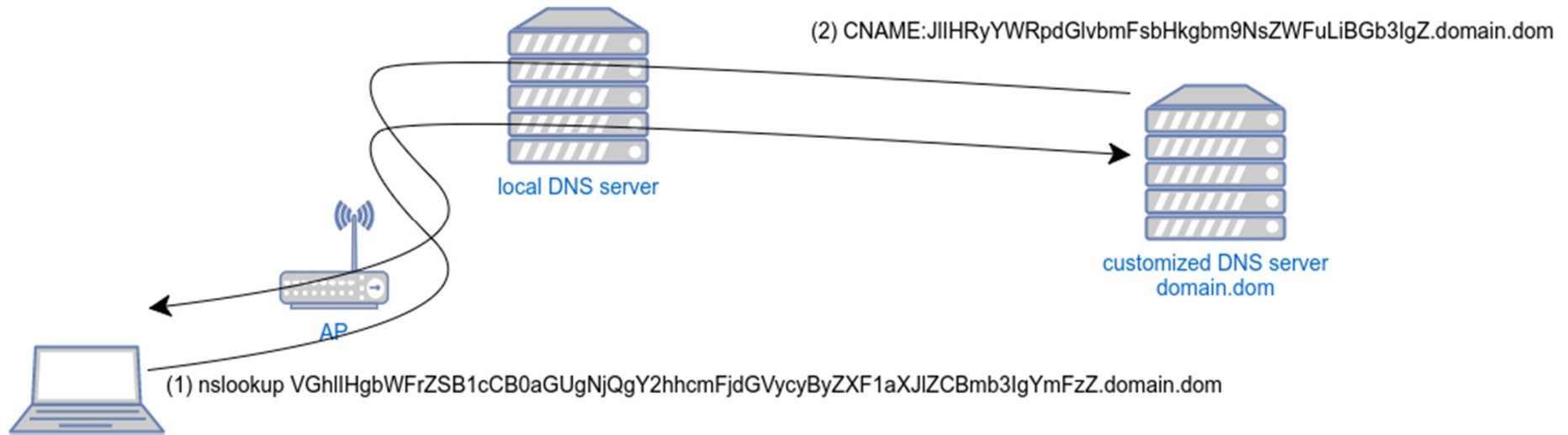


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# DNS tunneling/data exfiltration



- DNS traffic passes the firewall
- Use DNS to transport other traffic
- First use to bypass paid WiFi access
- Can be used for botnet C&C



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# DNS tunneling tools

- High throughput tools:
  - DNScat: <http://tadek.pietraszek.org/projects/DNScat/>
  - DNS2TCP: <https://www.aldeid.com/wiki/Dns2tcp>
  - Iodine: <https://code.kryo.se/iodine/>
- Low throughput tools:
  - DNS messenger
  - MULTIGRAIN
  - Wekby



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# DNS fast fluxing

- A method to hide botnet
- Quickly change the IP address for a single symbolic name to hide the malicious servers.
- Used for phishing botnets
- Single flux: Quickly change DNS records
- Double flux: Bots are proxies towards the attackers server



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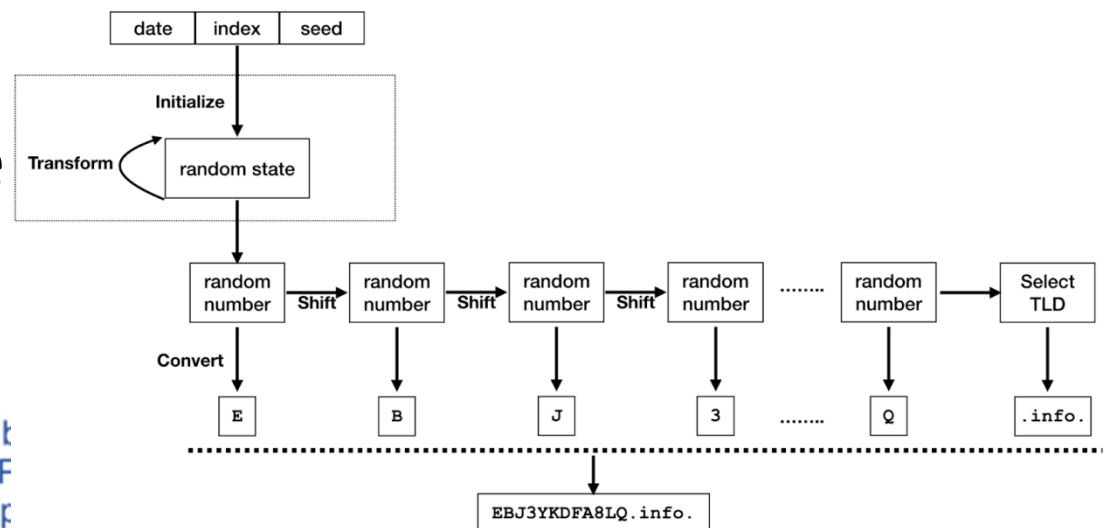
# Domain generation algorithm

- No domain hardcoded into the bot malware source
- Domains for C&C dynamically created and changed periodically
- Seed for domain creation can be timestamp, message on a popular social network,...

Domain examples:

t3622c4773260c097e2e9b26705212ab85.ws  
u83ccf36d9f02e9ea79a9d16c0336677e4.to  
v02bec0c090508bc76b3ea81dfc2198a71.in  
wa9e4628c334324e181e40f33f878c153f.hk  
xdcc5481252db5f38d5fc18c9ad3b2f7fd.cn

Words made of the two word from a dictionary combination, etc.



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