# Towards Ptolemaic metric properties of the z-normalized Euclidean distance for MulTiS indexing

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May 13, 2024

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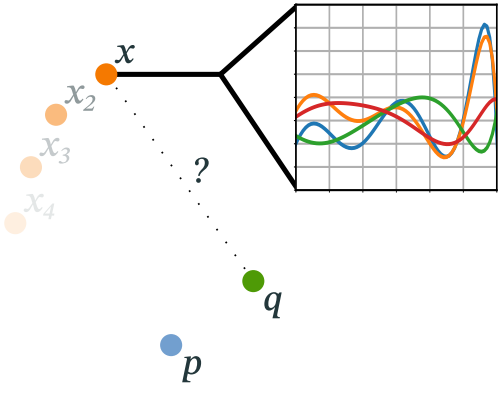
# Motivation: Comparing Time Series



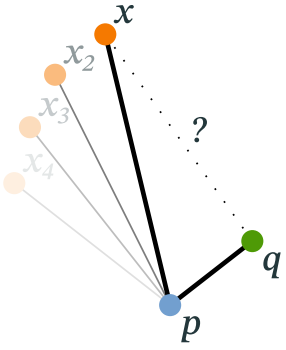

Compare time series using a (dis-)similarity measure $d$

Find time series similar to example $q$ in database $\mathbb{D}$

$$d : \mathbb{D} \times \mathbb{D} \to \mathbb{R}^+$$

$$\{x \in \mathbb{D} \mid d(q, x) \leq r\}$$

# Motivation: Comparing Time Series Efficiently

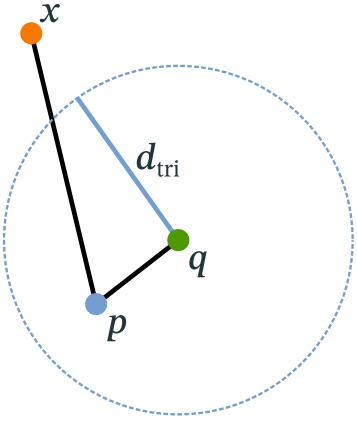



Cache distances to special pivot points $d(p, x_i)$ beforehand

Calculate $d(p, q)$ online

Determine $d(q, x) \leq r$

# Motivation: Metric Index

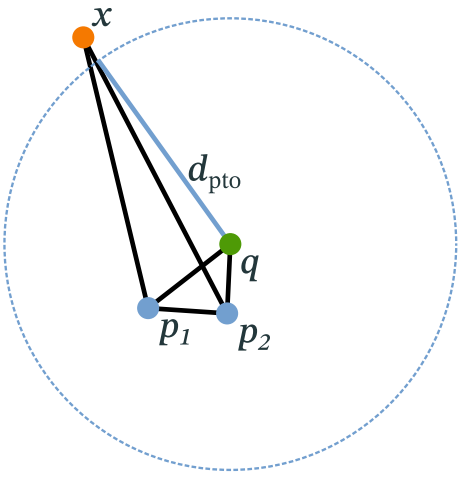



If $(\mathbb{D}, d)$ is a metric space:

Use the triangle inequality to access index

$$d(q, x) \geq d_{\text{tri}} = |d(p, x) - d(p, q)|$$

# Ptolemaic Index




If $(\mathbb{D}, d)$ is also Ptolemaic:

Even better approximation!

$$d(q,x) \geq d_{\text{pto}} = \frac{|\overline{xp_1} \cdot \overline{qp_2} - \overline{qp_1} \cdot \overline{xp_2}|}{\overline{p_1p_2}}$$

# Objective

Find a distance function for MulTiS that is

**metric and Ptolemaic** to speed up range queries!

# Structure

- Univariate z-normalized Euclidean distance
- Multivariate z-normalized Euclidean distance
- Indexability

# Univariate z-normalized Euclidean Distance

aka Matrix Profile distance

$$d_z(x, y) : \mathbb{R}^n \times \mathbb{R}^n \to \mathbb{R}^+$$

$$d_z(x, y) = \sqrt{\sum_i^n (\nu(x_i) - \nu(y_i))^2}$$

with $\quad \nu : x_i \mapsto \dfrac{x_i - \mathrm{E}(x)}{\sqrt{\mathrm{Var}(x)}}$

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$$\iff d_z(x, y) = \sqrt{2n}\sqrt{1 - \mathrm{Corr}(x, y)}$$

# Homomorphic to Euclidean Metric

$$d_z(x,y) = \quad d_{\text{eucl}}(\nu(x),\nu(y)) = \quad (d_{\text{eucl}} \circ \nu)(x,y)$$

$$\Rightarrow (\mathbb{R}^n, d_z) = \quad (\mathbb{R}^n, d_{\text{eucl}} \circ \nu) \sim \quad (\nu(\mathbb{R}^n), d_{\text{eucl}})$$

but $\nu(\mathbb{R}^n) \subset \mathbb{R}^n$, so $(\mathbb{R}^n, d_z) \subset (\mathbb{R}^n, d_{\text{eucl}})$

$\Rightarrow (\mathbb{R}^n, d_z)$ is a Ptolemaic pseudometric

# Multivariate Extension

$$d_m(X, Y) : (\mathbb{R}^{n \times m}) \times (\mathbb{R}^{n \times m}) \to \mathbb{R}^+$$

$$d_m(X, Y) = \sqrt{\sum_i^k (d_z(X_i, Y_i))^2} = ||\vec{d_z}(X, Y)||$$

# Properties of the Multivariate Extension

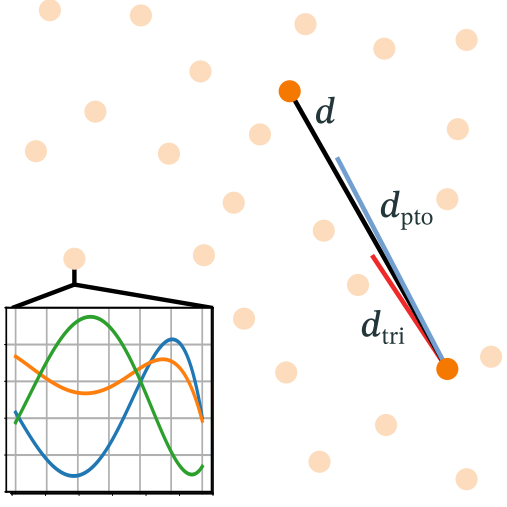
**is a pseudometric**

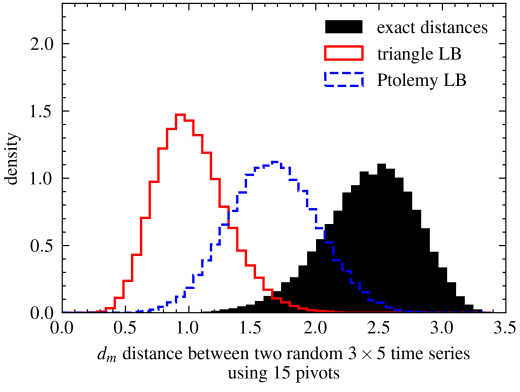
Use norm subadditivity ($||X + Y|| \leq ||X|| + ||Y||$)

**is probably Ptolemaic**

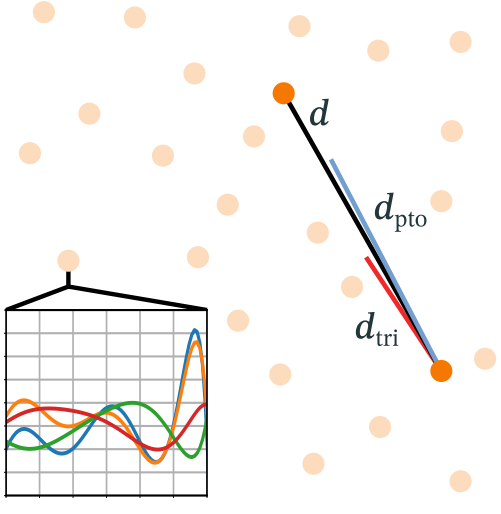
only numerical evidence yet

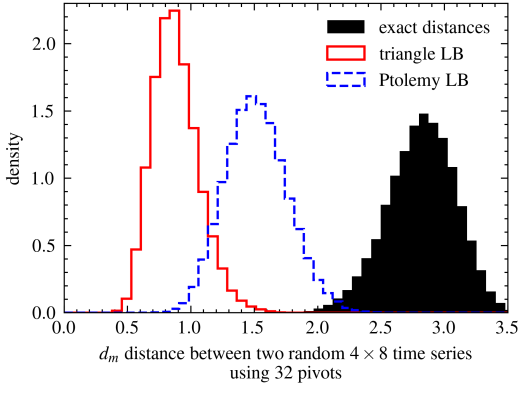
# Indexing Improvements: Low Dimensionality



# Indexing Improvements: High Dimensionality



# Conclusion

## Summary

- The multivariate z-normalized Euclidean distance is a Ptolemaic pseudometric
- Ptolemy's inequality improves indexability of such metric spaces

## Future Work

- Analytical proof
- Integrate into indexing structures, test on real data
- Pivot selection

# Backup Slides

# Experiment: Counterexample Search

- generate $5 \cdot 10^7$ MulTiS quadruplets
- with $X \in \mathbb{R}^{4\times5}$ (4 observables, 5 time slices)
- draw each element of the matrix from a uniform distribution
  1. $\{X_{ij} \in \mathbb{Z} \mid -20 \leq X_{ij} \leq 20\}$
  2. $\{X_{ij} \in \mathbb{R} \mid -20 \leq X_{ij} \leq 20\}$

# Experiment: Counterexample Search

- generate random points $\{X_{ij} \in \mathbb{R} \mid -20 \leq X_{ij} \leq 20\}$
- for each measured distance:
  - choose 15 (32) random pivots
  - only use the pivots that generate the highest lower bounds
  - calculate actual distance, best triangle lower bound, best Ptolemaic lower bounds

# Metric Space

metric space $(M, d)$ with time series $x, y, z \in M$

$d(x, y) = 0 \iff x = y$ (identity of indiscernibles)

$d(x, y) = d(y, x)$ (symmetry)

$d(x, z) \leq d(x, y) + d(y, z)$ (triangle inequality)

# The “fourth” metric axiom

Do we need $d(x, y) > 0$?

$$0 = d(x, x) \leq d(x, y) + d(y, x) \qquad \text{(triangle inequality)}$$

$$0 = d(x, x) \leq 2d(x, y) \qquad \text{(symmetry)}$$

# Ptolemaic Spaces

quadratic form distance $\sqrt{(x-y)^T A(x-y)}$

Jensen–Shannon distance

$$\sqrt{H(P/2+Q/2)-H(P)/2-H(Q)/2}$$

triangular distance $\sqrt{\sum_i \frac{(P_i-Q_i)^2}{P_i+Q_i}}$

cosine distance $\sqrt{1-\cos \measuredangle XY}$

$(M, \sqrt{d})$ for any metric space $(M, d)$