

TP4-Exercicio1

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Trabalho Realizado Por:

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Exercício 1

1. Considere o seguinte programa, em Python anotado, para multiplicação de dois inteiros de precisão limitada a 16 bits. Assume-se que os inteiros são representáveis na teoria BitVecSort(16) do Z3.

```
assume m >= 0 and n >= 0 and r == 0 and x == m and y == n
0: while y > 0:
1:   if y & 1 == 1:
        y , r = y-1 , r+x
2:   x , y = x<<1 , y>>1
3: assert r == m * n
```

1. Usando indução verifique a terminação deste programa.

Estado Inicial

$$m \geq 0 \wedge n \geq 0 \wedge r == 0 \wedge x == m \wedge y == n \wedge pc$$

Transições do FOTS

$$\begin{aligned} & (pc = 0 \wedge pc' = 1 \wedge y \leq 0 \wedge m' = m \wedge n' = n \wedge r' = r \wedge x' = x \wedge y' = y) \\ & \quad \vee \\ & (pc = 0 \wedge pc' = 0 \wedge y > 0 \wedge m' = m \wedge n' = n \wedge r' = r + x \wedge x' = x << 1 \wedge y' = (y - 1) >> 1) \\ & \quad \vee \\ & (pc = 1 \wedge pc' = 1 \wedge m' = m \wedge n' = n \wedge r' = r \wedge x' = x \wedge y' = y) \end{aligned}$$

```
[25]: from z3 import *
from random import randint

#declarar
def declare(i):
    state = {}
    state['pc'] = Int('pc'+str(i))
    state['m'] = BitVec('m'+str(i),16)
    state['n'] = BitVec('n'+str(i),16)
    state['r'] = BitVec('r'+str(i),16)
```

```

state['x'] = BitVec('x'+str(i),16)
state['y'] = BitVec('y'+str(i),16)
return state

#iniciar
def init(state):
    return And(state['m'] == randint(0,10),
               state['n'] == randint(0,10),
               state['r'] == 0,
               state['x'] == state['m'],
               state['y'] == state['n'],
               state['pc'] == 0 )

#transicoes
def trans(curr,prox):

    aux = And(prox['m'] == curr['m'],
              prox['n'] == curr['n'],
              prox['r'] == curr['r'],
              prox['x'] == curr['x'],
              prox['y'] == curr['y'])

    #y <= 0 acabou ciclo
    t1 = And(curr['pc'] == 0, prox['pc'] == 1, curr['y'] <= 0, aux)

    #continua ciclo
    t2 = And(curr['pc'] == 0, prox['pc'] == 0,
             curr['y'] > 0,
             prox['m'] == curr['m'],
             prox['n'] == curr['n'],
             prox['r'] == curr['r'] + curr['x'],
             prox['x'] == (curr['x'] << 1),
             prox['y'] == ((curr['y']-1) >> 1))

    #fim
    t3 = And(curr['pc'] == 1, prox['pc'] == 1, aux)

    ###
    return Or(t1,t2,t3)

#gerar traco
def gera_traco(declare,init,trans,k):

```

```

s = Solver()
state = [declare(i) for i in range(k)]
s.add(init(state[0]))

for i in range(k-1):
    s.add(trans(state[i],state[i+1]))
if s.check()==sat:
    m=s.model()
    for i in range(k):
        print(i)
        for x in state[i]:
            print(x,"=",m[state[i][x]])

gera_traco(declare,init,trans,10)

```

```

0
pc = 0
m = 10
n = 3
r = 0
x = 10
y = 3
1
pc = 0
m = 10
n = 3
r = 10
x = 20
y = 1
2
pc = 0
m = 10
n = 3
r = 30
x = 40
y = 0
3
pc = 1
m = 10
n = 3
r = 30
x = 40
y = 0
4
pc = 1
m = 10
n = 3

```

```

r = 30
x = 40
y = 0
5
pc = 1
m = 10
n = 3
r = 30
x = 40
y = 0
6
pc = 1
m = 10
n = 3
r = 30
x = 40
y = 0
7
pc = 1
m = 10
n = 3
r = 30
x = 40
y = 0
8
pc = 1
m = 10
n = 3
r = 30
x = 40
y = 0
9
pc = 1
m = 10
n = 3
r = 30
x = 40
y = 0

```

Usando o `bmc_eventually` verificamos a terminação do programa

```

[27]: def termina(state):
        return (state['pc'] == 1)

def bmc_eventually(declare,init,trans,prop,bound):
    for k in range (1,bound+1):
        s = Solver()
        # completar

```

```

state =[declare(i) for i in range(k)]
s.add(init(state[0]))
for i in range(k-1):
    s.add(trans(state[i],state[i+1]))
s.add(prop(state[k-1]))
if s.check()==sat:
    m=s.model()
    for i in range(k):
        print(i)
        for x in state[i]:
            print(x,"=",m[state[i][x]])
    return
print ("Property is valid up to traces of length "+str(k))

```

```
bmc_eventually(declare,init,trans,termina,20)
```

Property is valid up to traces of length 1

```

0
pc = 0
m = 8
n = 0
r = 0
x = 8
y = 0
1
pc = 1
m = 8
n = 0
r = 0
x = 8
y = 0

```

- b. Pretende-se verificar a correção parcial deste programa usando duas formas alternativas para lidar com programas iterativos: havoc e unfold. I. Usando o comando havoc e a metodologia WPC (weakest pre-condition) gere a condição de verificação que garanta a correção parcial.

```

assume m >= 0 and n >= 0 and r == 0 and x == m and y == n
0: while y > 0:
1:   if y & 1 == 1:
       y , r = y-1 , r+x
2:   x , y = x<<1 , y>>1
3: assert r == m * n

```

Usando o método de correção havoc obtemos a seguinte estrutura:

```

assume m >= 0 and n >= 0 and r == 0 and x == m and y == n
0: while y > 0:
    invariante
1:   if y & 1 == 1:

```

```

          y , r = y-1 , r+x
2:      x , y = x<<1 , y>>1
3: assert r == m * n

```

Logo com:

```

pre= m >= 0 and n >= 0 and r == 0 and x == m and y == n
pos= r == m * n
inv=

```

Temos:

```

assume pre;
assert inv;
havoc r,havoc x,havoc y;
((assume y>0 and inv;((assume y and 1==1;y==y-1;r==r+x)||
  assume not(y and 1==1););x==x<<1;y==y>>1;assert inv; assume False;assert pos;)||
(assume not(y>0) and inv;assert pos;))

```

#== havoc

```

pre->(inv and ForAll([r,x,y],
  ((assume y>0 and inv;((assume y and 1==1;y==y-1;r==r+x)||
    assume not(y and 1==1););x==x<<1;y==y>>1;assert inv; assume False;assert pos;)||
    assume not(y>0) and inv;assert pos;)))

```

#== false->..=TRUE

```

pre->(inv and ForAll([r,x,y],
  ((assume y>0 and inv;((assume y and 1==1;y==y-1;r==r+x)||
    assume not(y and 1==1););x==x<<1;y==y>>1;assert inv;))))and assume not(y>0) and inv;assert pos;

```

#== transformacao

```

pre->(inv and ForAll([r,x,y],
  (y>0 and inv->(((y and 1==1)-> inv; [y>>1/y] [x<<1/x] [r+x/r] [y-1/y]) and
  (not(y and 1==1)->inv; [y>>1/y] [x<<1/x])))) and (not(y>0) and inv) -> pos;)

```

- II. Usando a metodologia SPC (strongest pos-condition), para um parâmetro inteiro N , gere o fluxo que resulta do unfold do ciclo N vezes e construa a respetiva condição de verificação.

```

assume m >= 0 and n >= 0 and r == 0 and x == m and y == n
0: while y > 0:
1:   if y & 1 == 1:
      y , r = y-1 , r+x
2:   x , y = x<<1 , y>>1
3: assert r == m * n

```

Desenrolando o ciclo em if's ficamos com: (Desenrolamos no máximo 16 vezes pois é o tamanho máximo do BitVec)

```

assume m >= 0 and n >= 0 and r == 0 and x == m and y == n
if (y > 0):
    if y & 1 == 1:
        y , r = y-1 , r+x
    x , y = x<<1 , y>>1
    if (y > 0):
        if y & 1 == 1:
            y , r = y-1 , r+x
        x , y = x<<1 , y>>1
        if (y > 0):
            if y & 1 == 1:
                y , r = y-1 , r+x
            x , y = x<<1 , y>>1
            if (y > 0):
                if y & 1 == 1:
                    y , r = y-1 , r+x
                x , y = x<<1 , y>>1

    (...)

    if (y > 0):
        if y & 1 == 1:
            y , r = y-1 , r+x
        x , y = x<<1 , y>>1
        assert not (y > 0)

assert r == m * n

```

Como tem de ser em single assignment (SA)

```

assume m >= 0 and n >= 0 and r0 == 0 and x0 == m and y0 == n
if (y0 > 0):
    if y0 & 1 == 1:
        y1 , r1 = y0-1 , r0+x0
    else:
        r16 = r0
    x1 , y2 = x0<<1 , y1>>1
    if (y2 > 0):
        if y2 & 1 == 1:
            y3 , r2 = y2-1 , r1+x1
        else:
            r16 = r1
        x2 , y4 = x1<<1 , y3>>1
        if (y4 > 0):
            if y4 & 1 == 1:
                y5 , r3 = y4-1 , r2+x2
            else:

```

```

        r16 = r2
    x3 , y6 = x2<<1 , y5>>1

    (...)

    if (y30 > 0):
        if y30 & 1 == 1:
            y31 , r16 = y30-1 , r15+x15
        else:
            r16 = r15
            x16 , y32 = x15<<1 , y31>>1
            assert not (y32 > 0)

    else:
        r16 = r15

    else:
        r16 = r2

    else:
        r16 = r1

else:
    r16 = r0

assert r16 == m * n

```

-Passo intermédio

```

assume m >= 0 and n >= 0 and r0 == 0 and x0 == m and y0 == n
assume (y0 > 0);
    assume y0 & 1 == 1;
        y1 , r1 = y0-1 , r0+x0
    ||
    assume not y0 & 1 == 1;
        r16 = r0
    x1 , y2 = x0<<1 , y1>>1
    assume (y2 > 0);
        assume y2 & 1 == 1;
            y3 , r2 = y2-1 , r1+x1
        ||
        assume not (y2 & 1 == 1);
            r16 = r1
    x2 , y4 = x1<<1 , y3>>1
    assume (y4 > 0);
        assume y4 & 1 == 1;
            y5 , r3 = y4-1 , r2+x2
        ||
        assume not (y4 & 1 == 1);
            r16 = r2
    x3 , y6 = x2<<1 , y5>>1

```



```

(...)

assume (y30 > 0);
    assume y30 & 1 == 1;
        y31 , r16 = y30-1 , r15+x15
    ||
    assume not (y30 & 1 == 1);
        r16 = r15
    x16 , y32 = x15<<1 , y31>>1
    assert not (y32 > 0);
    ||
    assume not (y30 > 0);
        r16 = r15
||
    assume not (y4 > 0);
        r16 = r2
||
    assume not (y2 > 0);
        r16 = r1
||
assume not (y0 > 0);
    r16 = r0

assert r16 == m * n

```

-normalizado com 17 fluxos usar a strongest pre condition

```

assume m >= 0 and n >= 0 and r0 == 0 and x0 == m and y0 == n
assume (y0 > 0);
assume y0 & 1 == 1;
y1 , r1 = y0-1 , r0+x0
||
assume not y0 & 1 == 1;
r16 = r0
x1 , y2 = x0<<1 , y1>>1
assume (y2 > 0);
assume y2 & 1 == 1;
y3 , r2 = y2-1 , r1+x1
||
assume not (y2 & 1 == 1);
r16 = r1
x2 , y4 = x1<<1 , y3>>1
assume (y4 > 0);
assume y4 & 1 == 1;
y5 , r3 = y4-1 , r2+x2
||
assume not (y4 & 1 == 1);
r16 = r2
x3 , y6 = x2<<1 , y5>>1

```

(...)

```
assume (y30 > 0);
assume y30 & 1 == 1;
y31 , r16 = y30-1 , r15+x15
||
assume not (y30 & 1 == 1);
r16 = r15
x16 , y32 = x15<<1 , y31>>1
assert not (y32 > 0) and r16 == m * n
```

||

(...)

||

```
assume m >= 0 and n >= 0 and r0 == 0 and x0 == m and y0 == n
assume (y0 > 0);
assume y0 & 1 == 1;
y1 , r1 = y0-1 , r0+x0
||
assume not y0 & 1 == 1;
r16 = r0
x1 , y2 = x0<<1 , y1>>1
assume (y2 > 0);
assume y2 & 1 == 1;
y3 , r2 = y2-1 , r1+x1
||
assume not (y2 & 1 == 1);
r16 = r1
x2 , y4 = x1<<1 , y3>>1
assume not (y4 > 0);
r16 = r2
assert r16 == m * n
```

||

```
assume m >= 0 and n >= 0 and r0 == 0 and x0 == m and y0 == n
assume (y0 > 0);
assume y0 & 1 == 1;
```

```

y1 , r1  = y0-1 , r0+x0
||
assume not y0 & 1 == 1;
r16 = r0
x1 , y2 = x0<<1 , y1>>1
assume not (y2 > 0);
r16 = r2
assert r16 == m * n

```

||

```

assume m >= 0 and n >= 0 and r0 == 0 and x0 == m and y0 == n
assume not (y0 > 0);
r16 = r0
assert r16 == m * n

```

C. Codifique, em SMT's e em ambos os casos, a verificação da correcção parcial.

i) Usando o Havoc com as deduções previamente feitas, temos:

```

[11]: def prove(f):
    s = Solver()
    s.add(Not(f))
    r = s.check()
    if r == unsat:
        print("Proved")
    else:
        print("Failed to prove")
        m = s.model()
        for v in m:
            print(v, '=', m[v])

m= BitVec('m',16)
n= BitVec('n',16)
r= BitVec('r',16)
x= BitVec('x',16)
y= BitVec('y',16)
pre=And( m >=0,n >=0,r == 0,x == m,y == n)
pos= r == m*n
inv=And(y>=0,y<=n,x == m + r)

d1=Implies(And(Not(y==0),1==1),substitute(substitute(substitute(substitute(inv,(y,y>>1))),(x,x<<1))),(x,x<<1)))
d2=Implies(Not(And(Not(y==0),1==1)),substitute(substitute(inv,(y,y>>1))),(x,x<<1)))
f1=inv
f2=ForAll([r,x,y],Implies(And(y>0,inv),And(d1,d2)))

```

```
f3=Implies(And(Not(y>0),inv),pos)
prove(Implies(pre,And(f1,f2,f3)))
```

Proved

ii) Usando o unfold com as deduções previamente feitas temos:

```
[12]: r = []
x = []
y1 = []
y2 = []

for i in range (17):
    r.append(BitVec('r'+str(i),16))
    x.append(BitVec('x'+str(i),16))
    y1.append(BitVec('y1'+str(i),16))
    y2.append(BitVec('y2'+str(i),16))

m = (BitVec('m'+str(i),16))
n = (BitVec('n'+str(i),16))

preCond = And(m >=0,n >=0,r[0] == 0,x[0] == m,y1[0] == n)
posCond = r[16] == m*n

def condicao(num):
    if num == 0:
        return Implies(And(preCond, Not(y1[num] > 0),r[16] == r[num]), posCond)

    aux=And([And(r[i+1] == r[i] + x[i],
                y2[i+1] == y1[i] - 1,
                x[i+1] == x[i] << 1,
                y1[i+1] == y2[i+1] >> 1)
            for i in range(num)])

    if num == 16:
        ret = Implies(And(preCond,aux),And(Not(y1[num] > 0), posCond))
    else:
        ret = Implies(And(preCond,aux,Not(y1[num] > 0),r[16] == r[num]),
            ↪posCond)
        return Or(ret, condicao(num - 1))

prove(condicao(16))
```

Proved

[]: