

Q1) Show the different inputs & outputs of the various phases of a compiler for the assignment statement  $C = (F-32) * (S/9)$

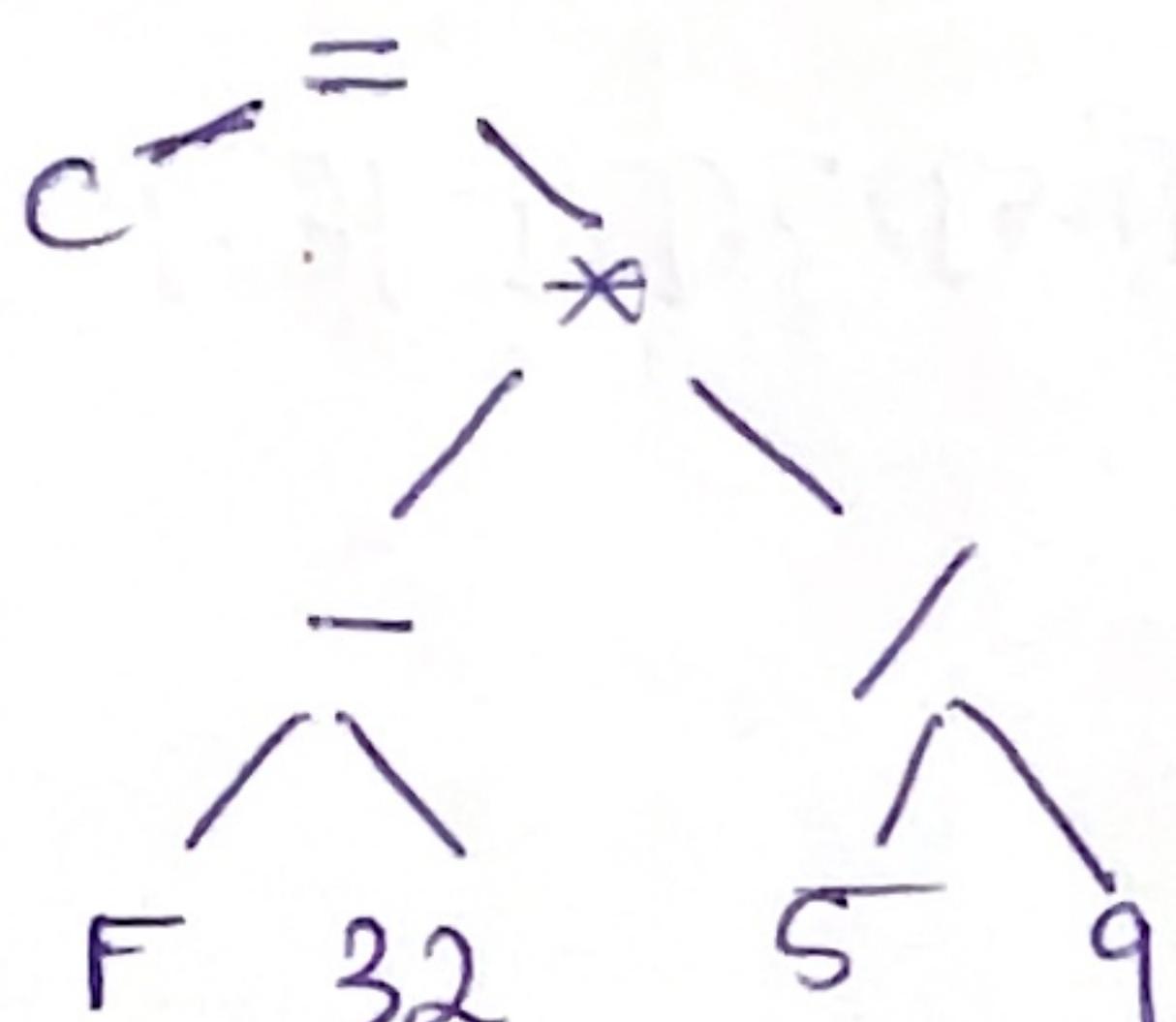
Ans → a) Lexical Analysis → Input → Source code as a sequence of characters.

Output → A stream of tokens.

Identifier (C), assign-op (=), left-paren, identifier (F), ~~assi~~ minus-op (-), number (32), ~~mult~~ right-paren, mult-op (\*), left-paren, number (S), div-op (/), number (9), right-paren

b) Syntax Analysis → Input → Token stream from lexical analysis.

Output → Parse Tree / Abstract Syntax Tree (AST) representing the expression <sup>structure</sup> syntax



c) Semantic Analysis → Input → AST from the syntax analysis phase.

Output → Annotated AST with type information and error checking:

- i) Check type compatibility (e.g. numeric types for arithmetic)
- ii) Validate variable declarations

d) Intermediate Code generation → Input → Annotated AST.

Output → Intermediate representation, often in three-address code (TAC) →

$$t_1 = F - 32$$

$$t_2 = S / 9$$

$$t_3 = t_1 * t_2$$

$$C = t_3$$

e) Code Optimisation → Input → Intermediate code

Output → Optimised intermediate code, removing redundancies:

$$t_1 = F - 32$$

$$C = t_1 * 0.5556$$

f) Code generation → Input → Optimised intermediate code

Output → Assembly or machine code:

Load F  
Sub B2  
Mul 0.5556  
Store C

g) Linking and loading → Input → Object code.

Output → Executable machine code ready for execution.

2) Provide regular expressions to characterise the following lexical items:

a) Identifiers → Regular expression →  $[a-zA-Z][a-zA-Z0-9]^*$   $([a-zA-Z0-9])^*$ ?

b) Numbers → Regular expression →  $(0|[1-9][0-9]^*)\cdot([0-9]^*[1-9])^*([eE][\pm])^*[1-9][0-9]^*$ ?

c) White space →  $[\t\n\r]+$

d) Arithmetic Operators → (+ | - | \* | / | div | mod)

e) Logical Constants → (true | false)

f) Logical Operations → ( $\neg$  |  $\wedge$  |  $\vee$  |  $\rightarrow$ ) (not | and | or)

g) Comparison Operators → ( $<$  |  $>$  |  $=$  |  $<$  |  $>$  |  $=$ )

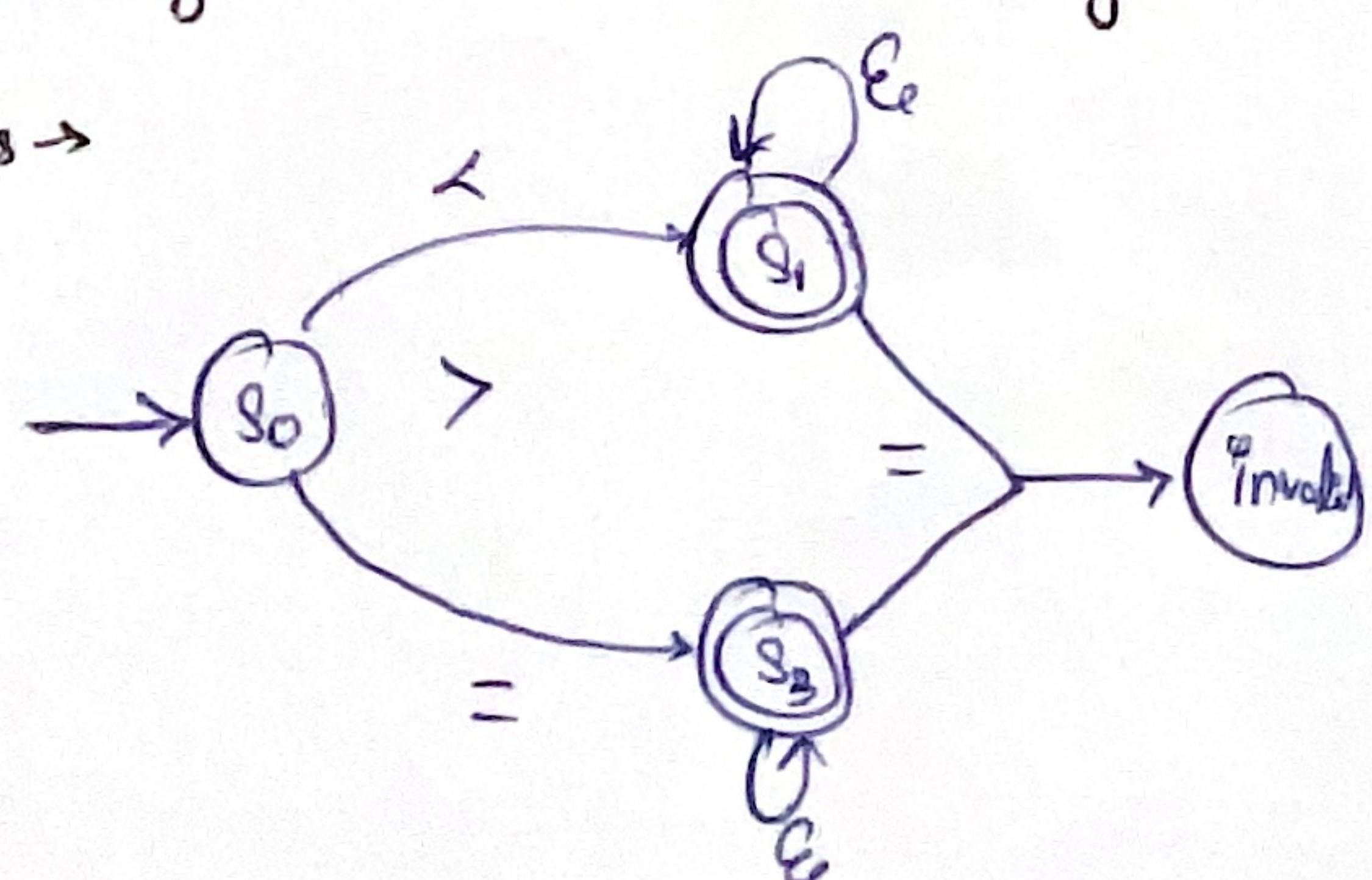
h) Parentheses → ( | )

i) Keywords → (if | then | else)

j) Literal Strings → "([^\"]\")\* (\\".["\"])\*")"

Q3) A grammar for branching statements is defined as

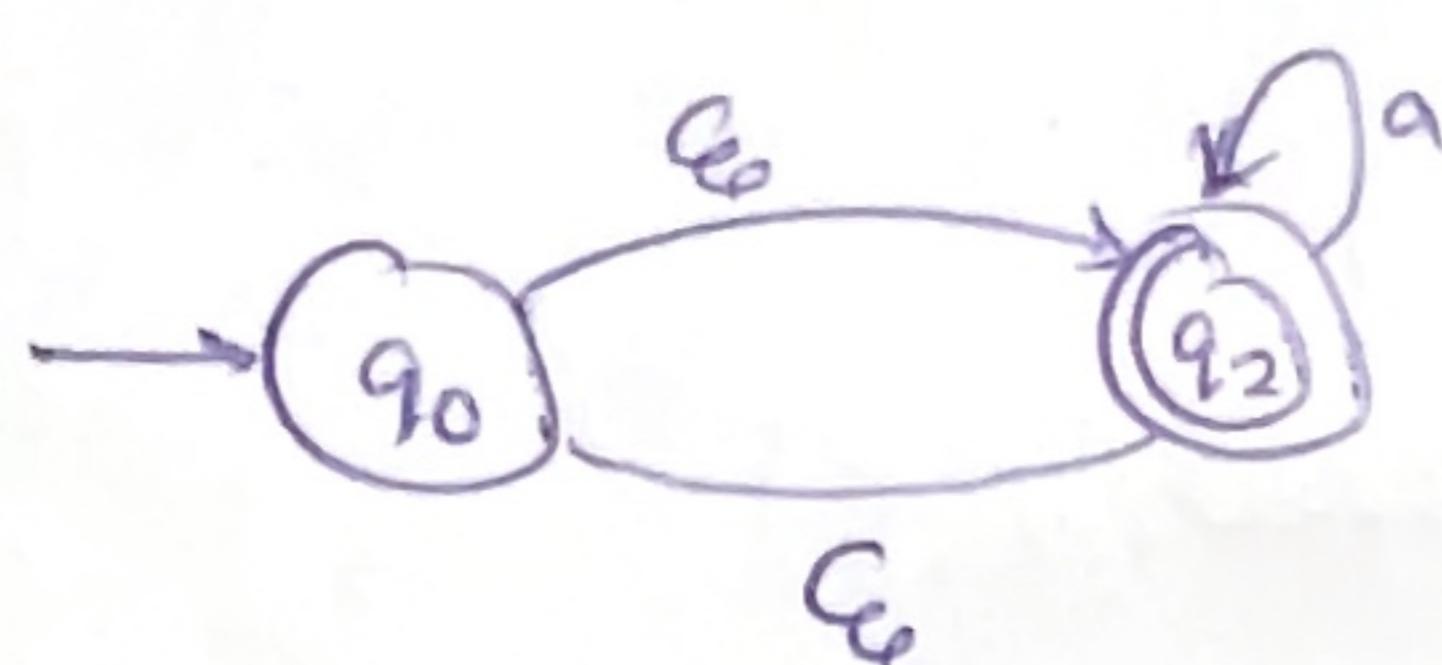
Ans →



4) Design the DFA for the given Regular Expression  $a(a|b)^*a$

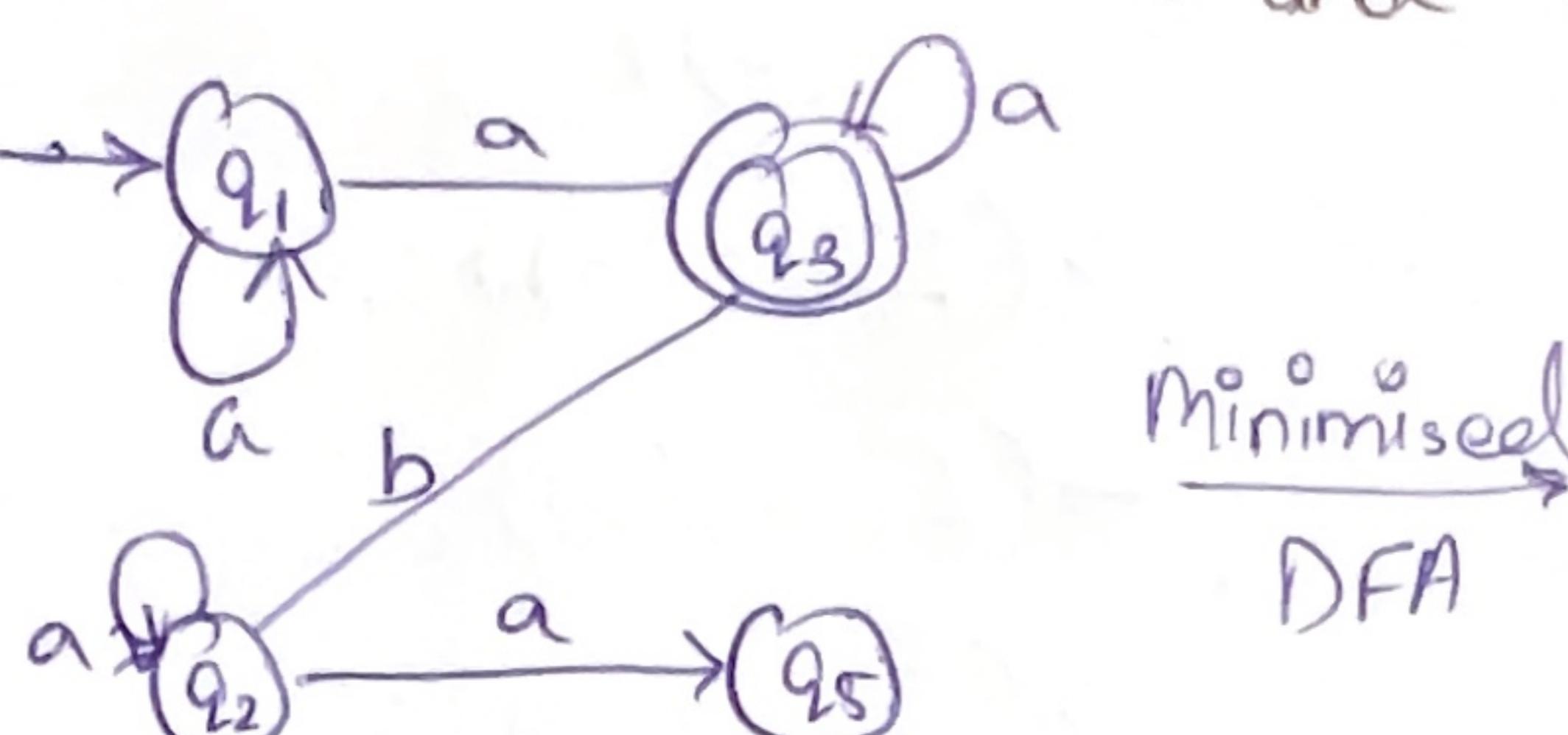
a) Convert the RG to  $\epsilon$ -NFA.

Ans →

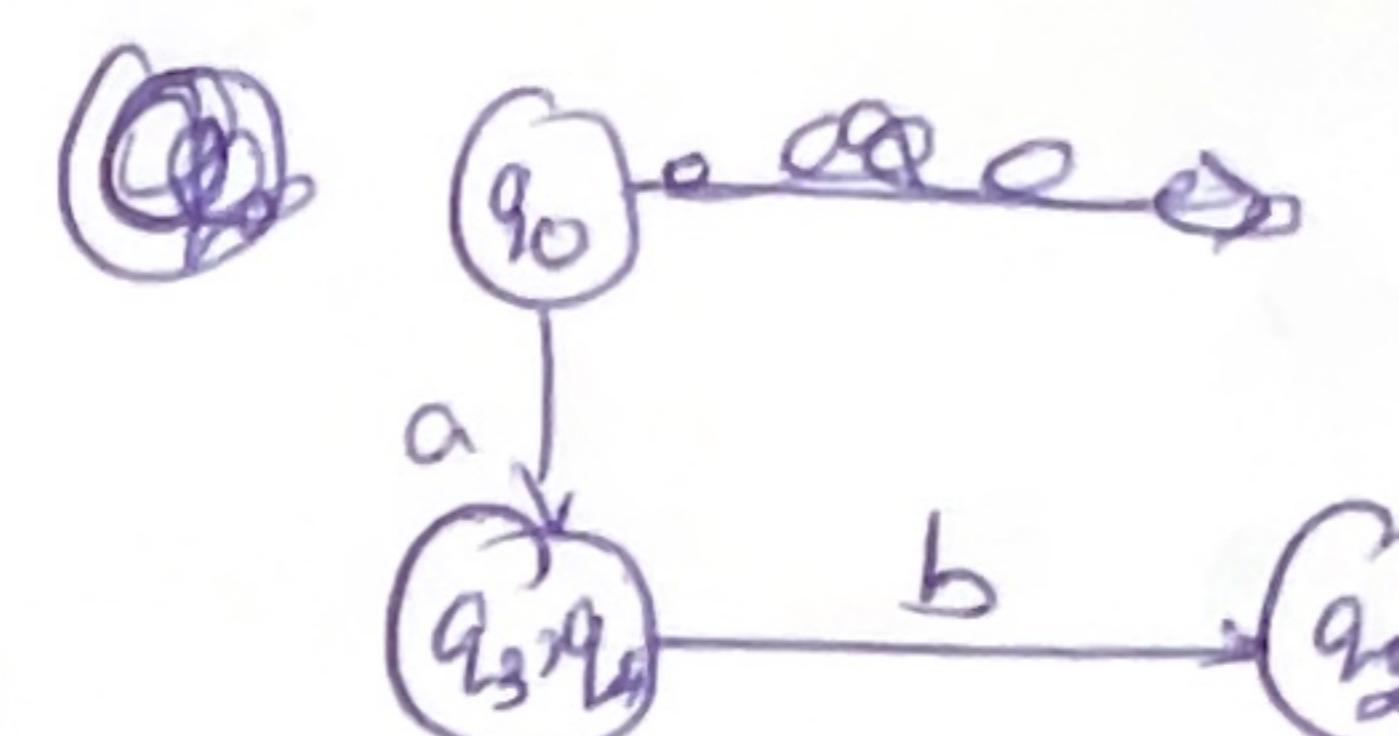


b) Convert the  $\epsilon$ -NFA to NFA and NFA to Minimized DFA.

Ans →



$\xrightarrow{\text{Minimized}}$



5) Divide the following C program into appropriate lexemes and count the number of tokens.

float limitedSquare (x)

{

float x;

/\* returns x-squared, but never more than 100 \*/

return (x < -10.0 || x >= 10.0) ? 100 : x\*x;

}

Ans → Lexeme Analysis.

- i) float → keyword ii) limitedSquare → Identifier iii) ( → Left parenthesis iv) x → Identifier
- v) ) → Right parenthesis vi) { → Left brace vii) float → keyword viii) x → Identifier ix) ; → Punctuation (Semicolon) x) /\* returns x-squared, but never more than 100 \*/ → Comment (Ignored by lexical analysis) xi) return → keyword xii) ( → Left parenthesis xiii) x → Identifier
- xiv) <= → Relational Operator xv) - → Arithmetic xvii) 10.0 → Number xviii) || → Logical
- xix) x = → Identifier xx) >= → Relational xxii) 10.0 → Number xxiii) ? → Conditional
- xxiv) 100 → Number xxv) : → Conditional xxvi) x → Identifier xxvii) \* → Arithmetic
- xxviii) x → Identifier xxix) ; → Punctuation xxviii) ? → Right Brace xxix) ) → Right Paren

Number of tokens → 29

6) Consider the context free grammar :  $S \rightarrow (L) \mid a$   $L \rightarrow L, S \mid S$  and the string  $((a,a),a,(a))$

a) Give the leftmost derivation for the string.

$$\text{Ans} \rightarrow S \rightarrow (L)$$

$$\rightarrow (L, S)$$

$$\rightarrow (L, S, S)$$

$$\rightarrow (S, S, S)$$

$$\rightarrow ((L), S, S)$$

$$\rightarrow ((L, S), S, S)$$

$$\rightarrow ((S, S), S, S)$$

$$\rightarrow ((a, S), S, S)$$

$$\rightarrow ((a, a), S, S)$$

$$\rightarrow ((a, a), a, S)$$

$$\rightarrow ((a, a), a, (L))$$

$$\rightarrow ((a, a), a, (S))$$

$$\rightarrow ((a, a), a, (a))$$

b) Give the rightmost derivation for the string.

$$\text{Ans} \rightarrow S \rightarrow (L)$$

~~$\rightarrow (L, S)$~~

$$\rightarrow (L, (L))$$

$$\rightarrow (L, (a))$$

$$\rightarrow (L, S, (a))$$

$$\rightarrow (L, a, (a))$$

$$\rightarrow (S, a, (a))$$

$$\rightarrow ((L), a, (a))$$

$$\rightarrow ((L, S), a, (a))$$

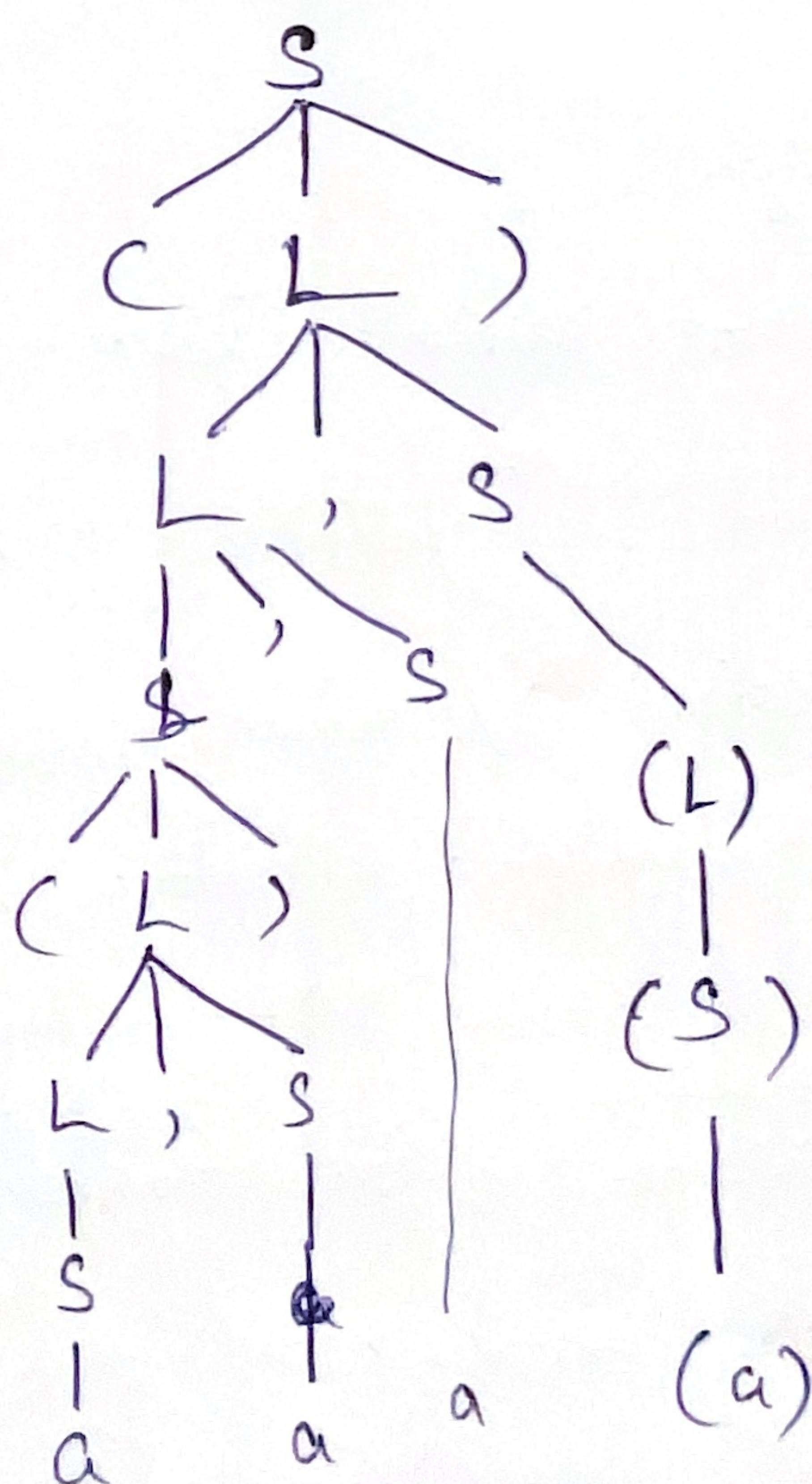
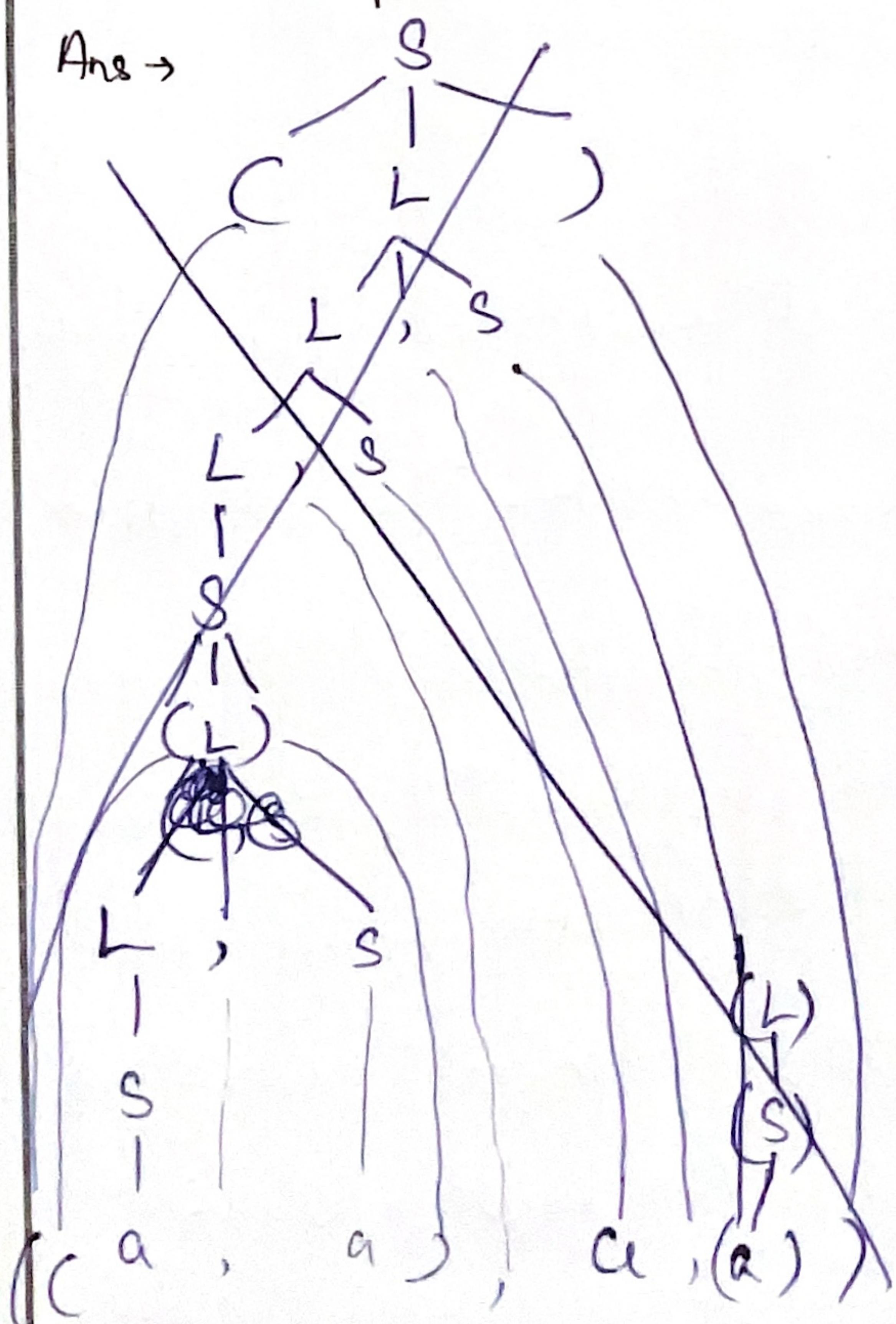
$$\rightarrow ((L, a), a, (a))$$

$$\rightarrow ((S, a), a, (a))$$

$$\rightarrow ((a, a), a, (a))$$

c) Give a parse tree for the string.

$$\text{Ans} \rightarrow$$



d) Is the grammar ambiguous or unambiguous? Justify your answer.

Ans → A grammar is ambiguous if a single string can have multiple distinct parse trees on different derivations. This grammar is ambiguous because the string can be derived in multiple ways, leading to different parse trees.

e) Eliminate the left recursion from the grammar (if any).

Ans →  $S \rightarrow (L) | a$   
 $L \rightarrow L, S | S$

Rule  $L \rightarrow L, S$  is left-recursive because L appears on the left side immediately after L.

Rewrite rule for L

$$L \rightarrow SL'$$

$$L' \rightarrow, SL' | \epsilon$$

New Grammar

$$S \rightarrow (L) | a$$

$$L \rightarrow SL'$$

$$L' \rightarrow, SL' | \epsilon$$

f) Let G be the following context-free grammar, where E is the start symbol of G.

Ans →  $E \rightarrow E > E | E > E | E < E | E \leq E | E \ll E | E \gg E | E ! = E | id$

Construction of unambiguous grammar.

$$E \rightarrow EI$$

$$EI \rightarrow EI! = E2 | E2$$

$$E2 \rightarrow E3 << E2 | E3 >> E2 | E3$$

$$E3 \rightarrow E3 > E4 | E3 > E4 | E3 < E4 | E3 \leq E4 | E4$$

$$E4 \rightarrow id$$

Example derivation → Lets derive:  $id \ll id > id ! = id$

i)  $E \rightarrow EI \rightarrow E2 \rightarrow E3 << E2 \rightarrow E4 << E2 \rightarrow id \ll E2$

ii)  $E2 \rightarrow E3 > E4 \rightarrow E4 > E4 \rightarrow id > id$

iii)  $id \ll id > id$

iv)  $E2 \rightarrow EI ! = E2 \rightarrow E4 ! = E2 \rightarrow id ! = E2$

v)  $E2 \rightarrow E3 \rightarrow E4 \rightarrow id$

vi)  $id \ll id > id ! = id$

8) Let  $G$  be the following Context-Free Grammar.

$$P \rightarrow xQR$$

$$Q \rightarrow yz/z$$

$$R \rightarrow w/\epsilon$$

$$S \rightarrow y$$

Compute the FIRST and FOLLOW function for all the non-terminals present in the above grammar.

Ans →

	First	Follow
P	{x}	{\$}
Q	{y, z}	{w, y}
R	{w, $\epsilon$ }	{y}
S	{y}	{\$}

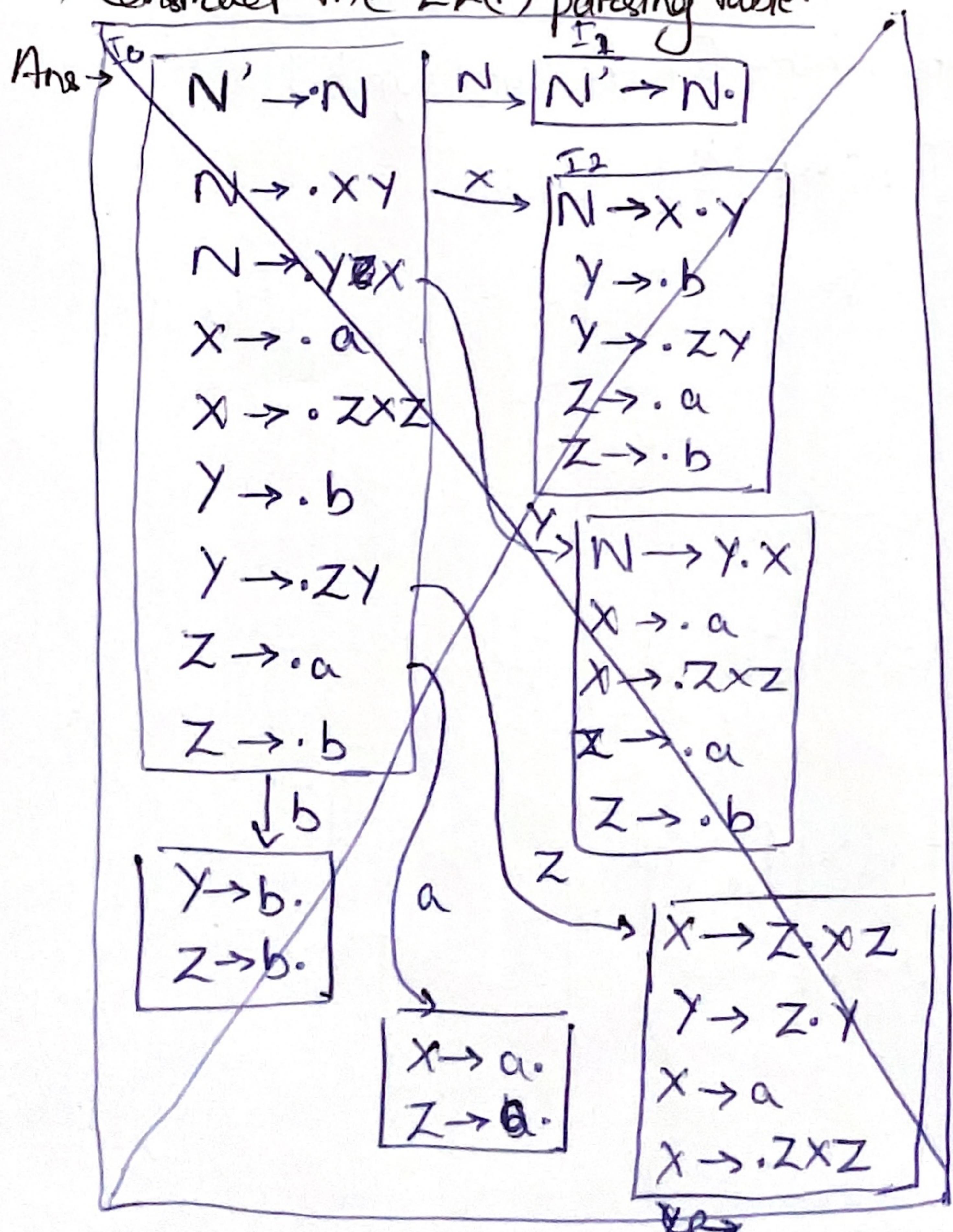
9) Consider the following grammar.  $N \rightarrow XY|YX$

$$X \rightarrow a|ZXZ$$

$$Y \rightarrow b|ZY$$

$$Z \rightarrow a|b$$

a) Construct the LL(1) parsing table.



	First	Follow
N	{a, b}	\$
X	{a, b}	{b, a}
Y	{b, a}	{\$, a, b}
Z	{a, b}	{a, b}

LL(1) Table →

	a	b
N	$N \rightarrow XY YX$	$N \rightarrow YX XY$
X	$X \rightarrow a ZXZ$	$ZXZ$
Y	$ZY$	$b ZY$
Z	$a$	$b$

b) Is this grammar LL(1)?

Ans → No. Conflicts exists in the parsing table.

Multiple entries for same non-terminal on terminal pair.

c) Verify the given grammar is ambiguous or not.

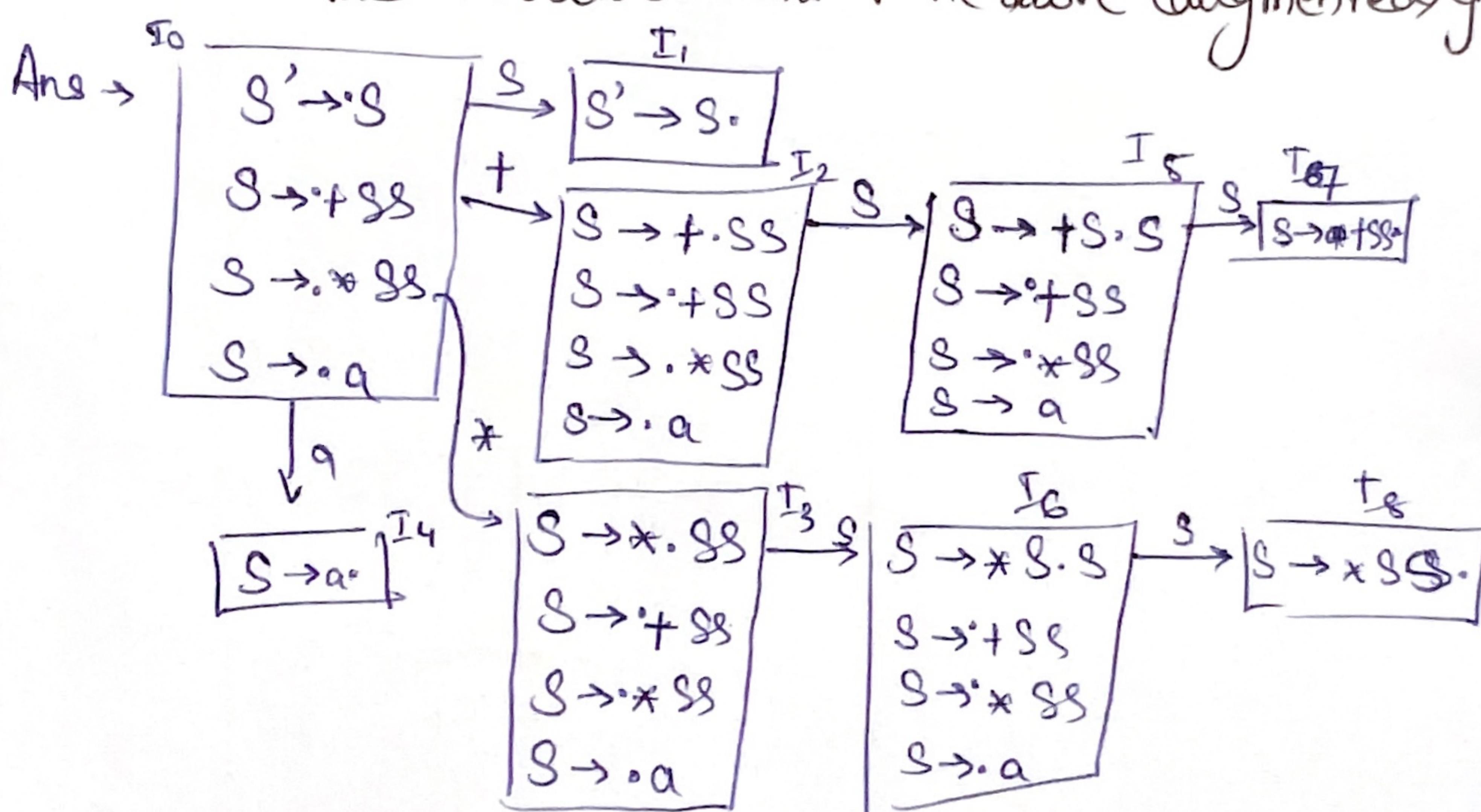
Ans → Let's derive a string ababa.

$N \rightarrow XY$	$N \rightarrow YX$	$N \rightarrow Y \text{ } \boxed{X}$
$\rightarrow ZXZY$	$\rightarrow ZYX$	$\rightarrow ZYX$
$\rightarrow ab$	$\rightarrow ZZYX$	$\rightarrow ZYZXZ$
	$\rightarrow ZZZYX$	$\rightarrow abab$
	$\rightarrow ababa$	

The grammar is not ambiguous as it didn't have same string generated from different parse trees.

10) Consider the following grammar:  $S \rightarrow +SS \mid *SS \mid a$  and the string  $+a^*aa$ .

a) Construct the SLR sets of items of the above (augmented) grammar.



b) Compute the GOTO function for these sets of items. Show the parsing table for this grammar.

Is this grammar SLR?

Ans →

I <sub>0</sub>	I <sub>2</sub>	GOTO	Action			
I <sub>1</sub>	I <sub>2</sub>	*	#	a	\$	accept
					s	3
					s	4
I <sub>2</sub>	I <sub>2</sub>	I <sub>3</sub>	I <sub>4</sub>		s	5
I <sub>3</sub>	I <sub>2</sub>	I <sub>3</sub>	I <sub>4</sub>		s	6
I <sub>4</sub>	I <sub>3</sub>	I <sub>3</sub>	I <sub>4</sub>	I <sub>4</sub>	s	7
I <sub>5</sub>	I <sub>2</sub>	I <sub>5</sub>	I <sub>4</sub>		s	8
I <sub>6</sub>	I <sub>2</sub>	I <sub>3</sub>	I <sub>4</sub>		s	
I <sub>7</sub>	s					
I <sub>8</sub>	s					

Yes the grammar is SLR(1)  
as it has no conflicts in parsing table.