## **Sorting** Chapter B8.

Caution! If you are expecting to sort efficiently on a parallel machine, whether its parallelism is small-scale or massive, you almost certainly want to use library routines that are specific to your hardware.

We include in this chapter translations into Fortran 90 of the general purpose serial sorting routines that are in Volume 1, augmented by several new routines that give pedagogical demonstrations of how parallel sorts can be achieved with Fortran 90 parallel constructions and intrinsics. However, we intend the above word "pedagogical" to be taken seriously: these new, supposedly parallel, routines are not likely to be competitive with machine-specific library routines. Neither do they compete successfully on serial machines with the all-serial routines provided (namely sort, sort2, sort3, indexx, and select).

```
SUBROUTINE sort_pick(arr)
USE nrtype
IMPLICIT NONE
REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr
   Sorts an array arr into ascending numerical order, by straight insertion. arr is replaced
   on output by its sorted rearrangement.
INTEGER(I4B) :: i,j,n
REAL(SP) :: a
n=size(arr)
do j=2,n
                              Pick out each element in turn.
    a=arr(j)
                              Look for the place to insert it.
    do i=j-1,1,-1
        if (arr(i) <= a) exit
        arr(i+1)=arr(i)
    end do
                              Insert it.
    arr(i+1)=a
end do
END SUBROUTINE sort_pick
```

Not only is sort\_pick (renamed from Volume 1's piksrt) not parallelizable, but also, even worse, it is an  $N^2$  routine. It is meant to be invoked only for the most trivial sorting jobs, say, N < 20.

Sorts an array arr into ascending numerical order by Shell's method (diminishing increment

Determine the starting increment.

Loop over the partial sorts.

Outer loop of straight insertion.

Inner loop of straight insertion.

USE nrtype IMPLICIT NONE

REAL(SP) :: v

inc=3\*inc+1
if (inc > n) exit

inc=inc/3
do i=inc+1,n

v=arr(i) j=i do

end do arr(j)=v

if (inc  $\leq$  1) exit

END SUBROUTINE sort\_shell

end do

end do

j=j-inc

n=size(arr)
inc=1

do

do

end do

SUBROUTINE sort\_shell(arr)

INTEGER(I4B) :: i,j,inc,n

REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr

if (arr(j-inc) <= v) exit
arr(j)=arr(j-inc)</pre>

if (j <= inc) exit

sort). arr is replaced on output by its sorted rearrangement.

```
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                                                                                                                                                                                                                                                                                                                                     Computing (ISBN 0-521-57439-0
```

```
The routine sort_shell is renamed from Volume 1's shell. Shell's Method, a diminishing increment sort, is not directly parallelizable. However, one can write a fully parallel routine (though not an especially fast one — see remarks at beginning of this chapter) in much the same spirit:
```

```
SUBROUTINE sort_byreshape(arr)
USE nrtype; USE nrutil, ONLY : swap
IMPLICIT NONE
REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr
   Sort an array arr by bubble sorting a succession of reshapings into array slices. The method
   is similar to Shell sort, but allows parallelization within the vectorized masked swap calls.
REAL(SP), DIMENSION(:,:), ALLOCATABLE :: tab
REAL(SP), PARAMETER :: big=huge(arr)
INTEGER(I4B) :: inc,n,m
n=size(arr)
inc=1
                                         Find the largest increment that fits.
dο
    inc=2*inc+1
    if (inc > n) exit
end do
do
                                          Loop over the different shapes for the reshaped
    inc=inc/2
    m=(n+inc-1)/inc
    allocate(tab(inc,m))
                                         Allocate space and reshape the array, big en-
    tab=reshape(arr, (/inc,m/) , (/big/) )
                                                     sures that fill elements stay at the
                                                     end.
          Bubble sort all the rows in parallel.
        call swap(tab(:,1:m-1:2),tab(:,2:m:2), &
            tab(:,1:m-1:2)>tab(:,2:m:2))
        call swap(tab(:,2:m-1:2),tab(:,3:m:2), &
            tab(:,2:m-1:2)>tab(:,3:m:2))
```

The basic idea is to reshape the given one-dimensional array into a succession of two-dimensional arrays, starting with "tall and narrow" (many rows, few columns), and ending up with "short and wide" (many columns, few rows). At each stage we sort all the rows in parallel by a bubble sort, giving something close to Shell's diminishing increments.

\* \* \*

We now arrive at those routines, based on the Quicksort algorithm, that we actually intend for use with general N on serial machines:

```
SUBROUTINE sort(arr)
USE nrtype; USE nrutil, ONLY : swap, nrerror
TMPLTCTT NONE
REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr
INTEGER(I4B), PARAMETER :: NN=15, NSTACK=50
   Sorts an array arr into ascending numerical order using the Quicksort algorithm. arr is
   replaced on output by its sorted rearrangement.
   Parameters: NN is the size of subarrays sorted by straight insertion and NSTACK is the
   required auxiliary storage.
REAL(SP) :: a
INTEGER(I4B) :: n,k,i,j,jstack,l,r
INTEGER(I4B), DIMENSION(NSTACK) :: istack
n=size(arr)
jstack=0
1=1
r=n
do
    if (r-1 < NN) then
                                         Insertion sort when subarray small enough.
        do j=l+1,r
            a=arr(i)
            do i=j-1,1,-1
                if (arr(i) <= a) exit
                arr(i+1)=arr(i)
            end do
            arr(i+1)=a
        end do
        if (jstack == 0) RETURN
                                          Pop stack and begin a new round of partition-
        r=istack(jstack)
        l=istack(jstack-1)
        jstack=jstack-2
                                          Choose median of left, center, and right elements
    else
        k=(1+r)/2
                                             as partitioning element a. Also rearrange so
        call swap(arr(k),arr(l+1))
                                             that a(1) \leq a(1+1) \leq a(r).
        call swap(arr(1),arr(r),arr(1)>arr(r))
        call swap(arr(l+1),arr(r),arr(l+1)>arr(r))
        call swap(arr(1),arr(1+1),arr(1)>arr(1+1))
        i=1+1
                                          Initialize pointers for partitioning
        j=r
        a=arr(1+1)
                                          Partitioning element.
                                          Here is the meat.
        do
                                         Scan up to find element >= a.
                i=i+1
```

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```
if (arr(i) >= a) exit
            end do
                                          Scan down to find element <= a.
            do
                i=i-1
                if (arr(j) \le a) exit
            end do
                                          Pointers crossed. Exit with partitioning complete.
            if (j < i) exit
            call swap(arr(i),arr(j))
                                              Exchange elements.
        arr(l+1)=arr(j)
                                          Insert partitioning element.
        arr(j)=a
        jstack=jstack+2
          Push pointers to larger subarray on stack; process smaller subarray immediately.
        if (jstack > NSTACK) call nrerror('sort: NSTACK too small')
        if (r-i+1 >= j-1) then
            istack(jstack)=r
            istack(jstack-1)=i
            istack(jstack)=j-1
            istack(jstack-1)=1
            1=i
        end if
    end if
end do
END SUBROUTINE sort
```

call swap(...) ... call swap(...) One might think twice about putting all these external function calls (to nrutil routines) in the inner loop of something as streamlined as a sort routine, but here they are executed only once for each partitioning.

call swap(arr(i),arr(j)) This call is in a loop, but not the innermost loop. Most modern machines are very fast at the "context changes" implied by subroutine calls and returns; but in a time-critical context you might code this swap in-line and see if there is any timing difference.

```
SUBROUTINE sort2(arr,slave)
USE nrtype; USE nrutil, ONLY : assert_eq
USE nr, ONLY : indexx
IMPLICIT NONE
REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr,slave
   Sorts an array arr into ascending order using Quicksort, while making the corresponding
   rearrangement of the same-size array slave. The sorting and rearrangement are performed
   by means of an index array.
INTEGER(I4B) :: ndum
INTEGER(I4B), DIMENSION(size(arr)) :: index
ndum=assert_eq(size(arr), size(slave), 'sort2')
call indexx(arr,index)
                                 Make the index array.
arr=arr(index)
                                 Sort arr.
slave=slave(index)
                                 Rearrange slave.
END SUBROUTINE sort2
```

A close surrogate for the Quicksort partition-exchange algorithm can be coded, parallelizable, by using Fortran 90's pack intrinsic. On real compilers, unfortunately, the resulting code is not very efficient as compared with (on serial machines) the tightness of sort's inner loop, above, or (on parallel machines) supplied library sort routines. We illustrate the principle nevertheless in the following routine.

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```
RECURSIVE SUBROUTINE sort_bypack(arr)
USE nrtype; USE nrutil, ONLY: array_copy,swap
IMPLICIT NONE
REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr
   Sort an array arr by recursively applying the Fortran 90 pack intrinsic. The method is
   similar to Quicksort, but this variant allows parallelization by the Fortran 90 compiler.
REAL(SP) :: a
INTEGER(I4B) :: n,k,nl,nerr
INTEGER(I4B), SAVE :: level=0
LOGICAL, DIMENSION(:), ALLOCATABLE, SAVE :: mask
REAL(SP), DIMENSION(:), ALLOCATABLE, SAVE :: temp
n=size(arr)
if (n <= 1) RETURN
k=(1+n)/2
                                                 Pivot element is median of first, middle,
call swap(arr(1),arr(k),arr(1)>arr(k))
call swap(arr(k),arr(n),arr(k)>arr(n))
                                                    and last.
call swap(arr(1),arr(k),arr(1)>arr(k))
if (n <= 3) RETURN
level=level+1
                                                 Keep track of recursion level to avoid al-
if (level == 1) allocate(mask(n),temp(n))
                                                    location overhead.
a=arr(k)
mask(1:n) = (arr \le a)
                                                Which elements move to left?
mask(k) = .false.
call array_copy(pack(arr,mask(1:n)),temp,nl,nerr)
                                                           Move them.
mask(k) = .true.
                                                Move others to right.
temp(nl+2:n)=pack(arr,.not. mask(1:n))
temp(nl+1)=a
arr=temp(1:n)
call sort_bypack(arr(1:nl))
                                                And recurse.
call sort_bypack(arr(n1+2:n))
if (level == 1) deallocate(mask,temp)
level=level-1
END SUBROUTINE sort_bypack
     The following routine, sort_heap, is renamed from Volume 1's hpsort.
SUBROUTINE sort_heap(arr)
USE nrtype
USE nrutil, ONLY : swap
IMPLICIT NONE
REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr
   Sorts an array arr into ascending numerical order using the Heapsort algorithm. arr is
   replaced on output by its sorted rearrangement.
INTEGER(I4B) :: i,n
n=size(arr)
do i=n/2,1,-1
      The index i, which here determines the "left" range of the sift-down, i.e., the element to
      be sifted down, is decremented from n/2 down to 1 during the "hiring" (heap creation)
    call sift_down(i,n)
end do
do i=n,2,-1
      Here the "right" range of the sift-down is decremented from n-1 down to 1 during the
      "retirement-and-promotion" (heap selection) phase.
                                             Clear a space at the end of the array, and
    call swap(arr(1),arr(i))
    call sift_down(1,i-1)
                                                retire the top of the heap into it.
end do
CONTAINS
```

SUBROUTINE sift\_down(1,r)
INTEGER(I4B), INTENT(IN) :: 1,r

```
Carry out the sift-down on element arr(1) to maintain the heap structure.
INTEGER(I4B) :: j,jold
REAL(SP) :: a
a=arr(1)
jold=1
j=1+1
do
                                              "Do while j <= r:"
    if (j > r) exit
    if (j < r) then
        if (arr(j) < arr(j+1)) j=j+1
                                             Compare to the better underling.
    end if
    if (a \ge arr(j)) exit
                                             Found a's level. Terminate the sift-down. Oth-
                                                 erwise, demote a and continue.
    arr(jold)=arr(j)
    jold=j
    j=j+j
end do
                                              Put a into its slot.
arr(jold)=a
END SUBROUTINE sift_down
END SUBROUTINE sort_heap
```

Another opportunity provided by Fortran 90 for a fully parallelizable sort, at least pedagogically, is to use the language's allowed access to the actual floating-point representation and to code a radix sort [1] on its bits. This is *not* efficient, but it illustrates some Fortran 90 language features perhaps worthy of study for other applications.

```
SUBROUTINE sort_radix(arr)
USE nrtype; USE nrutil, ONLY : array_copy,nrerror
IMPLICIT NONE
REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr
   Sort an array arr by radix sort on its bits.
INTEGER(I4B), DIMENSION(size(arr)) :: narr,temp
LOGICAL, DIMENSION(size(arr)) :: msk
INTEGER(I4B) :: k,negm,ib,ia,n,nl,nerr
 Because we are going to transfer reals to integers, we must check that the number of bits
  is the same in each:
ib=bit_size(narr)
ia=ceiling(log(real(maxexponent(arr)-minexponent(arr),sp))/log(2.0_sp)) &
    digits(arr)
if (ib /= ia) call nrerror('sort_radix: bit sizes not compatible')
                                                       Mask for all bits except sign bit.
negm=not(ishftc(1,-1))
n=size(arr)
narr=transfer(arr.narr.n)
                                                       Flip all bits on neg. numbers.
where (btest(narr,ib-1)) narr=ieor(narr,negm)
do k=0, ib-2
      Work from low- to high-order bits, and partition the array according to the value of the
   msk=btest(narr.k)
    call array_copy(pack(narr,.not. msk),temp,nl,nerr)
   temp(nl+1:n)=pack(narr,msk)
   narr=temp
end do
msk=btest(narr,ib-1)
                                                       The sign bit gets separate treat-
call array_copy(pack(narr,msk),temp,nl,nerr)
                                                           ment, since here 1 comes be-
temp(nl+1:n)=pack(narr,.not. msk)
                                                           fore 0.
narr=temp
where (btest(narr,ib-1)) narr=ieor(narr,negm)
                                                       Unflip all bits on neg. numbers.
arr=transfer(narr,arr,n)
END SUBROUTINE sort_radix
```

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\* \* 7



We overload the generic name indexx with two specific implementations, one for SP floating values, the other for I4B integers. (You can of course add more overloadings if you need them.)

```
SUBROUTINE indexx_sp(arr,index)
USE nrtype; USE nrutil, ONLY : arth,assert_eq,nrerror,swap
IMPLICIT NONE
REAL(SP), DIMENSION(:), INTENT(IN) :: arr
INTEGER(I4B), DIMENSION(:), INTENT(OUT) :: index
INTEGER(I4B), PARAMETER :: NN=15, NSTACK=50
   Indexes an array \mathtt{arr}, i.e., outputs the array \mathtt{index} of length N such that \mathtt{arr}(\mathtt{index}(j))
   is in ascending order for j=1,2,\ldots,N. The input quantity arr is not changed.
REAL(SP) :: a
INTEGER(I4B) :: n,k,i,j,indext,jstack,l,r
INTEGER(I4B), DIMENSION(NSTACK) :: istack
n=assert_eq(size(index),size(arr),'indexx_sp')
index=arth(1,1,n)
jstack=0
1=1
r=n
do
    if (r-1 < NN) then
        do j=l+1,r
            indext=index(j)
            a=arr(indext)
            do i=j-1,1,-1
                if (arr(index(i)) <= a) exit</pre>
                index(i+1)=index(i)
            end do
            index(i+1)=indext
        end do
        if (jstack == 0) RETURN
        r=istack(jstack)
        l=istack(jstack-1)
        jstack=jstack-2
    else
        k=(1+r)/2
        call swap(index(k),index(1+1))
        call icomp_xchg(index(1),index(r))
        call icomp_xchg(index(1+1),index(r))
        call icomp_xchg(index(l),index(l+1))
        i=1+1
        j=r
        indext=index(1+1)
        a=arr(indext)
        do
            dο
                if (arr(index(i)) >= a) exit
            end do
            do
                j=j-1
                if (arr(index(j)) <= a) exit</pre>
            end do
            if (j < i) exit
            call swap(index(i),index(j))
        end do
        index(l+1)=index(j)
        index(j)=indext
        jstack=jstack+2
        if (jstack > NSTACK) call nrerror('indexx: NSTACK too small')
        if (r-i+1 \ge j-1) then
```

istack(jstack)=r

```
istack(jstack-1)=i
            r=j-1
        else
            istack(jstack)=j-1
            istack(jstack-1)=1
            1=i
        end if
    end if
end do
CONTAINS
SUBROUTINE icomp_xchg(i,j)
INTEGER(I4B), INTENT(INOUT) :: i,j
INTEGER(I4B) :: swp
if (arr(j) < arr(i)) then
    swp=i
    i=j
    j=swp
end if
END SUBROUTINE icomp_xchg
END SUBROUTINE indexx_sp
SUBROUTINE indexx_i4b(iarr,index)
USE nrtype; USE nrutil, ONLY : arth,assert_eq,nrerror,swap
IMPLICIT NONE
INTEGER(14B), DIMENSION(:), INTENT(IN) :: iarr
INTEGER(14B), DIMENSION(:), INTENT(OUT) :: index
INTEGER(I4B), PARAMETER :: NN=15, NSTACK=50
INTEGER(I4B) :: a
INTEGER(I4B) :: n,k,i,j,indext,jstack,l,r
INTEGER(I4B), DIMENSION(NSTACK) :: istack
n=assert_eq(size(index),size(iarr),'indexx_sp')
index=arth(1,1,n)
jstack=0
1=1
r=n
do
    if (r-1 < NN) then
        do j=l+1,r
            indext=index(j)
            a=iarr(indext)
            do i=j-1,1,-1
                if (iarr(index(i)) <= a) exit</pre>
                index(i+1)=index(i)
            end do
            index(i+1)=indext
        end do
        if (jstack == 0) RETURN
        r=istack(jstack)
        l=istack(jstack-1)
        jstack=jstack-2
    else
        k=(1+r)/2
        call swap(index(k),index(1+1))
        call icomp_xchg(index(1),index(r))
        call icomp_xchg(index(l+1),index(r))
        call icomp_xchg(index(1),index(1+1))
        i=1+1
        j=r
        indext=index(1+1)
        a=iarr(indext)
        do
            do
```

```
i=i+1
               if (iarr(index(i)) >= a) exit
           end do
           do
               j=j-1
               if (iarr(index(j)) <= a) exit</pre>
           end do
           if (j < i) exit
            call swap(index(i),index(j))
        end do
        index(l+1)=index(j)
        index(j)=indext
        jstack=jstack+2
        if (jstack > NSTACK) call nrerror('indexx: NSTACK too small')
        if (r-i+1 >= j-1) then
            istack(jstack)=r
           istack(jstack-1)=i
           r=j-1
        else
           istack(jstack)=j-1
           istack(jstack-1)=1
           1=i
        end if
   end if
end do
CONTAINS
SUBROUTINE icomp_xchg(i,j)
INTEGER(I4B), INTENT(INOUT) :: i,j
INTEGER(I4B) :: swp
if (iarr(j) < iarr(i)) then</pre>
   swp=i
   i=j
   j=swp
end if
END SUBROUTINE icomp_xchg
END SUBROUTINE indexx_i4b
SUBROUTINE sort3(arr,slave1,slave2)
USE nrtype; USE nrutil, ONLY : assert_eq
USE nr, ONLY : indexx
IMPLICIT NONE
REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr,slave1,slave2
   Sorts an array arr into ascending order using Quicksort, while making the corresponding
   rearrangement of the same-size arrays slave1 and slave2. The sorting and rearrangement
   are performed by means of an index array.
INTEGER(I4B) :: ndum
INTEGER(I4B), DIMENSION(size(arr)) :: index
ndum=assert_eq(size(arr), size(slave1), size(slave2), 'sort3')
call indexx(arr,index)
                                Make the index array.
arr=arr(index)
                                 Sort arr.
slave1=slave1(index)
                                 Rearrange slave1,
slave2=slave2(index)
                                 and slave2.
END SUBROUTINE sort3
```

```
FUNCTION rank(index)
USE nrtype; USE nrutil, ONLY : arth
IMPLICIT NONE
INTEGER(I4B), DIMENSION(:), INTENT(IN) :: index
INTEGER(I4B), DIMENSION(size(index)) :: rank
Given index as output from the routine indexx, this routine returns a same-size array
    rank, the corresponding table of ranks.
rank(index(:))=arth(1,1,size(index))
END FUNCTION rank
```

\* \* \*

Just as in the case of sort, where an approximation of the underlying Quicksort partition-exchange algorithm can be captured with the Fortran 90 pack intrinsic, the same can be done with indexx. As before, although it is in principle parallelizable by the compiler, it is likely not competitive with library routines.

```
RECURSIVE SUBROUTINE index_bypack(arr,index,partial)
USE nrtype; USE nrutil, ONLY: array_copy,arth,assert_eq
IMPLICIT NONE
REAL(SP), DIMENSION(:), INTENT(IN) :: arr
INTEGER(I4B), DIMENSION(:), INTENT(INOUT) :: index
INTEGER, OPTIONAL, INTENT(IN) :: partial
   Indexes an array arr, i.e., outputs the array index of length N such that \mathtt{arr}(\mathtt{index}(j))
   is in ascending order for j=1,2,\ldots,N. The method is to apply recursively the Fortran
   90 pack intrinsic. This is similar to Quicksort, but allows parallelization by the Fortran 90
   compiler. partial is an optional argument that is used only internally on the recursive calls.
REAL(SP) :: a
INTEGER(I4B) :: n,k,nl,indext,nerr
INTEGER(I4B), SAVE :: level=0
LOGICAL, DIMENSION(:), ALLOCATABLE, SAVE :: mask
INTEGER(I4B), DIMENSION(:), ALLOCATABLE, SAVE :: temp
if (present(partial)) then
   n=size(index)
else
    n=assert_eq(size(index),size(arr),'indexx_bypack')
    index=arth(1,1,n)
if (n <= 1) RETURN
k=(1+n)/2
call icomp_xchg(index(1),index(k))
                                                    Pivot element is median of first, mid-
call icomp_xchg(index(k),index(n))
                                                        dle, and last.
call icomp_xchg(index(1),index(k))
if (n <= 3) RETURN
level=level+1
                                                    Keep track of recursion level to avoid
if (level == 1) allocate(mask(n),temp(n))
                                                        allocation overhead.
indext=index(k)
a=arr(indext)
mask(1:n) = (arr(index) \le a)
                                                    Which elements move to left?
mask(k) = .false.
call array_copy(pack(index,mask(1:n)),temp,nl,nerr)
                                                            Move them.
mask(k) = .true.
temp(nl+2:n)=pack(index,.not. mask(1:n))
                                                    Move others to right.
temp(nl+1)=indext
index=temp(1:n)
call index_bypack(arr,index(1:nl),partial=1)
                                                    And recurse.
call index_bypack(arr,index(nl+2:n),partial=1)
if (level == 1) deallocate(mask,temp)
level=level-1
```

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```
CONTAINS
SUBROUTINE icomp_xchg(i,j)
IMPLICIT NONE
INTEGER(I4B), INTENT(INOUT) :: i,j
  Swap or don't swap integer arguments, depending on the ordering of their corresponding
  elements in an array arr.
INTEGER(I4B) :: swp
if (arr(j) < arr(i)) then
    swp=i
    i=j
    j=swp
end if
END SUBROUTINE icomp_xchg
END SUBROUTINE index_bypack
FUNCTION select(k,arr)
USE nrtype; USE nrutil, ONLY: assert, swap
IMPLICIT NONE
INTEGER(I4B), INTENT(IN) :: k
REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr
REAL(SP) :: select
   Returns the kth smallest value in the array arr. The input array will be rearranged to have
   this value in location arr(k), with all smaller elements moved to arr(1:k-1) (in arbitrary
   order) and all larger elements in arr(k+1:) (also in arbitrary order).
INTEGER(I4B) :: i,r,j,l,n
REAL(SP) :: a
n=size(arr)
call assert(k >= 1, k <= n, 'select args')
1=1
r=n
do
                                          Active partition contains 1 or 2 elements.
    if (r-1 \le 1) then
        if (r-l == 1) call swap(arr(l),arr(r),arr(l)>arr(r)) Active partition con-
        select=arr(k)
                                                                        tains 2 elements.
        RETURN
    else
                                          Choose median of left, center, and right elements
                                              as partitioning element a. Also rearrange so
        call swap(arr(i),arr(l+1))
                                              that arr(1) \le arr(1+1) \le arr(r).
        call swap(arr(1),arr(r),arr(1)>arr(r))
        call swap(arr(l+1),arr(r),arr(l+1)>arr(r))
        call swap(arr(1),arr(1+1),arr(1)>arr(1+1))
        i=1+1
                                          Initialize pointers for partitioning.
        j=r
        a=arr(1+1)
                                          Partitioning element.
                                          Here is the meat.
        do
                                          Scan up to find element > a.
                i=i+1
                if (arr(i) \ge a) exit
            end do
                                          Scan down to find element < a.
            do
                if (arr(j) \le a) exit
            end do
            if (j < i) exit
                                          Pointers crossed. Exit with partitioning complete.
            call swap(arr(i),arr(j))
                                              Exchange elements.
        end do
        arr(1+1)=arr(j)
                                          Insert partitioning element.
        arr(j)=a
        if (j \ge k) r = j-1
                                          Keep active the partition that contains the kth
                                              element.
```

```
if (j <= k) l=i
  end if
end do
END FUNCTION select</pre>
```

\* \* \*

The following routine, select\_inplace, is renamed from Volume 1's selip.

Volume 1's selip routine uses an entirely different algorithm, for the purpose of avoiding any additional memory allocation beyond that of the input array. Fortran 90 presumes a richer memory environment, so select\_inplace simply does the obvious (destructive) selection in scratch space. You can of course use the old selip if your in-core or in-cache memory is at a premium.

```
FUNCTION select_bypack(k,arr)
USE nrtype; USE nrutil, ONLY : array_copy,assert,swap
IMPLICIT NONE
INTEGER(I4B), INTENT(IN) :: k
REAL(SP), DIMENSION(:), INTENT(INOUT) :: arr
REAL(SP) :: select_bypack
   Returns the kth smallest value in the array arr. The input array will be rearranged to have
   this value in location arr(k), with all smaller elements moved to arr(1:k-1) (in arbitrary
   order) and all larger elements in arr(k+1:) (also in arbitrary order). This implementation
   allows parallelization in the Fortran 90 pack intrinsic.
LOGICAL, DIMENSION(size(arr)) :: mask
REAL(SP), DIMENSION(size(arr)) :: temp
INTEGER(I4B) :: i,r,j,l,n,nl,nerr
REAL(SP) :: a
n=size(arr)
call assert(k >= 1, k <= n, 'select_bypack args')</pre>
1=1
                                                         Initial left and right bounds.
r=n
do
                                                         Keep partitioning until desired el-
    if (r-1 \le 1) exit
                                                             ement is found.
    i=(1+r)/2
    call swap(arr(1),arr(i),arr(1)>arr(i))
                                                         Pivot element is median of first,
                                                             middle, and last.
    call swap(arr(i),arr(r),arr(i)>arr(r))
    call swap(arr(1),arr(i),arr(1)>arr(i))
    a=arr(i)
    mask(1:r) = (arr(1:r) \le a)
                                                         Which elements move to left?
    mask(i) = .false.
    call array_copy(pack(arr(l:r),mask(l:r)),temp(l:),nl,nerr)
                                                                        Move them.
```

```
mask(i) = .true.
   temp(j+1:r)=pack(arr(1:r),.not. mask(1:r))
                                                       Move others to right.
   temp(j)=a
   arr(1:r)=temp(1:r)
                                                       Reset bounds to whichever side
   if (k > j) then
       1=j+1
                                                           has the desired element.
    else if (k < j) then
        r=j-1
    else
       1=j
        r=j
    end if
end do
if (r-l == 1) call swap(arr(l),arr(r),arr(l)>arr(r))
                                                          Case of only two left.
select_bypack=arr(k)
END FUNCTION select_bypack
```



The above routine select\_bypack is parallelizable, but as discussed above (sort\_bypack, index\_bypack) it is generally not very efficient.

\* \* \*

The following routine, select\_heap, is renamed from Volume 1's hpsel.

```
SUBROUTINE select_heap(arr,heap)
USE nrtype; USE nrutil, ONLY : nrerror, swap
USE nr, ONLY : sort
IMPLICIT NONE
REAL(SP), DIMENSION(:), INTENT(IN) :: arr
REAL(SP), DIMENSION(:), INTENT(OUT) :: heap
   Returns in heap, an array of length M, the largest M elements of the array arr of length
   N, with heap(1) guaranteed to be the the Mth largest element. The array arr is not
   altered. For efficiency, this routine should be used only when M \ll N.
INTEGER(I4B) :: i,j,k,m,n
m=size(heap)
n=size(arr)
if (m > n/2 .or. m < 1) call nrerror('probable misuse of select_heap')
heap=arr(1:m)
call sort(heap)
                                     Create initial heap by overkill! We assume m \ll n.
                                     For each remaining element...
do i=m+1,n
    if (arr(i) > heap(1)) then
                                         Put it on the heap?
        heap(1)=arr(i)
        j=1
        do
                                     Sift down.
            k=2*i
            if (k > m) exit
            if (k /= m) then
                if (heap(k) > heap(k+1)) k=k+1
            if (heap(j) <= heap(k)) exit</pre>
            call swap(heap(k),heap(j))
            j=k
        end do
    end if
end do
END SUBROUTINE select_heap
```

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\* \* \*

```
FUNCTION eclass(lista,listb,n)
USE nrtype; USE nrutil, ONLY : arth,assert_eq
TMPLTCTT NONE
INTEGER(I4B), DIMENSION(:), INTENT(IN) :: lista,listb
INTEGER(I4B), INTENT(IN) :: n
INTEGER(I4B), DIMENSION(n) :: eclass
   Given M equivalences between pairs of n individual elements in the form of the input arrays
   lista and listb of length \dot{M}, this routine returns in an array of length {\tt n} the number
   of the equivalence class of each of the n elements, integers between 1 and n (not all such
   integers used).
INTEGER :: j,k,l,m
m=assert_eq(size(lista), size(listb), 'eclass')
eclass(1:n)=arth(1,1,n)
                                      Initialize each element its own class.
                                       For each piece of input information...
do l=1,m
    j=lista(1)
    do
                                       Track first element up to its ancestor.
        if (eclass(j) == j) exit
        j=eclass(j)
    end do
   k=listb(1)
    do
                                       Track second element up to its ancestor.
        if (eclass(k) == k) exit
        k=eclass(k)
    if (j \neq k) eclass(j)=k
                                      If they are not already related, make them so.
end do
do j=1,n
                                       Final sweep up to highest ancestors.
    do
        if (eclass(j) == eclass(eclass(j))) exit
        eclass(j)=eclass(eclass(j))
    end do
end do
END FUNCTION eclass
FUNCTION eclazz(equiv,n)
USE nrtype; USE nrutil, ONLY: arth
IMPLICIT NONE
INTERFACE
    FUNCTION equiv(i,j)
    USE nrtype
    IMPLICIT NONE
   LOGICAL(LGT) :: equiv
INTEGER(I4B), INTENT(IN) :: i,j
    END FUNCTION equiv
END INTERFACE
INTEGER(I4B), INTENT(IN) :: n
INTEGER(I4B), DIMENSION(n) :: eclazz
   Given a user-supplied logical function equiv that tells whether a pair of elements, each
   in the range 1...n, are related, return in an array of length n equivalence class numbers
   for each element.
INTEGER :: i,j
eclazz(1:n)=arth(1,1,n)
do i=2,n
                                              Loop over first element of all pairs.
                                              Loop over second element of all pairs.
    do j=1,i-1
                                              Sweep it up this much.
        eclazz(j)=eclazz(eclazz(j))
        if (equiv(i,j)) eclazz(eclazz(eclazz(j)))=i
          Good exercise for the reader to figure out why this much ancestry is necessary!
    end do
end do
                                              Only this much sweeping is needed finally.
do i=1.n
    eclazz(i)=eclazz(eclazz(i))
end do
END FUNCTION eclazz
```

## CITED REFERENCES AND FURTHER READING:

Knuth, D.E. 1973, Sorting and Searching, vol. 3 of The Art of Computer Programming (Reading, MA: Addison-Wesley), §5.2.5. [1]