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7	$\begin{array}{c} 6.6.2 \text{ Stirling numbers of the second kind} \\ 6.7 \text{ Fast Walsh-Hadamard Transform} \\ 6.8 \text{ Simplex Algorithm} \\ 6.9 \text{ Subset Convolution} \\ 6.9.1 \text{ Construction} \\ 6.10 \text{ Schreier-Sims Algorithm} \\ 6.11 \text{ Berlekamp-Massey Algorithm} \\ 6.12 \text{ Fast Linear Recurrence} \\ 6.13 \text{ Miller Rabin} \\ 6.14 \text{ Pollard's Rho} \\ 6.15 \text{ Meissel-Lehmer Algorithm} \\ 6.16 \text{ Discrete Logarithm} \\ 6.17 \text{ Quadratic Residue} \\ 6.18 \text{ Gaussian Elimination} \\ 6.19 \text{ Characteristic Polynomial} \\ 6.20 \ \mu \text{ function} \\ 6.21 \text{ Partition Function} \\ 6.22 \ \frac{1\pi}{4} \ \ \text{Enumeration} \\ 6.23 \text{ De Bruijn Sequence} \\ 6.24 \text{ Extended GCD} \\ 6.25 \text{ Euclidean Algorithms} \\ 6.26 \text{ Chinese Remainder Theorem} \\ 6.27.1 \text{ Kirchhoff's Theorem} \\ 6.27.2 \text{ Tutte's Matrix} \\ 6.27.3 \text{ Cayley's Formula} \\ 6.27.4 \text{ Erdős-Gallai Theorem} \\ 6.27.4 \text{ Erdős-Gallai Theorem} \\ \end{array}$	13 13 14 14 14 15 15 15 16 16 17 17 17 17 17 17 17 17 17 17	<pre>return EOF;     p = buffer; } return *p++; } template <typename t=""> inline bool rit(T&amp; x) {     char c = 0; bool flag = false;     while (c = getchar(), (c &lt; '0' &amp;&amp; c != '-')    c &gt; '9') if (c = -1) return false;     c == '-' ? (flag = true, x = 0) : (x = c - '0');     while (c = getchar(), c &gt;= '0' &amp;&amp; c &lt;= '9') x = x * 10 + c - '0';     if (flag) x = -x;     return true; }  1.3 Increase stack size  const int size = 256 &lt;&lt; 20; register long rsp asm("rsp"); char *p = (char*)malloc(size) + size, *bak = (char*)rsp;    asm("movq %0, %%rsp\n"::"r"(p)); // main</typename></pre>	

# 2 Flows, Matching

# 2.1 Dinic's Algorithm

```
struct Edge {
 int to, cap, rev;
 Edge(int t, int c, int r) : to(t), cap(c), rev(r) {}
int Flow(vector<vector<Edge>> g, int s, int t) {
 int n = g.size(), res = 0;
 vector<int> lev(n, -1), iter(n);
 while (true) {
   vector<int> que(1, s);
    fill(lev.begin(), lev.end(), -1);
    fill(iter.begin(), iter.end(), 0);
    lev[s] = 0;
    for (int it = 0; it < que.size(); ++it) {</pre>
      int x = que[it];
      for (Edge &e : g[x]) {
        if (e.cap > 0 \&\& lev[e.to] == -1) {
          lev[e.to] = lev[x] + 1;
          que.push_back(e.to);
     }
   if (lev[t] == -1) break;
   auto Dfs = [\&] (auto dfs, int x, int f = 1000000000) {
      if (x == t) return f;
      int res = 0;
      for (int &it = iter[x]; it < g[x].size(); ++it) {</pre>
        Edge &e = g[x][it];
        if (e.cap > 0 \& lev[e.to] == lev[x] + 1) {
          int p = dfs(dfs, e.to, min(f - res, e.cap));
          res += p;
e.cap -= p;
          g[e.to][e.rev].cap += p;
      if (res == 0) lev[x] = -1;
      return res;
   res += Dfs(Dfs, s);
  return res;
```

#### 2.2 Minimum-cost flow

```
struct Edge {
  int to, cap, rev, w;
 Edge(int t, int c, int r, int w) : to(t), cap(c), rev(r), w(w
pair<int, int> Flow(vector<vector<Edge>> g, int s, int t) {
 int N = g.size();
 vector<int> dist(N), ed(N), pv(N);
 vector<bool> inque(N);
  int flow = 0, cost = 0;
 while (true) {
   dist.assign(N, kInf);
   inque.assign(N, false);
   pv.assign(N, -1);
   dist[s] = 0;
   queue<int> que;
   que.push(s);
   while (!que.empty()) {
      int x = que.front(); que.pop();
      inque[x] = false;
      for (int i = 0; i < g[x].size(); ++i) {</pre>
        Edge &e = g[x][i];
        if (e.cap > 0 \&\& dist[e.to] > dist[x] + e.w) {
          dist[e.to] = dist[x] + e.w;
          pv[e.to] = x;
          ed[e.to] = i;
          if (!inque[e.to]) {
            inque[e.to] = true;
            que.push(e.to);
       }
     }
    if (dist[t] == kInf) break;
    int f = kInf:
    for (int x = t; x != s; x = pv[x]) f = min(f, g[pv[x]][ed[x])
    ]].cap);
    for (int x = t; x != s; x = pv[x]) {
```

```
Edge &e = g[pv[x]][ed[x]];
    e.cap -= f;
    g[e.to][e.rev].cap += f;
}
flow += f;
    cost += f * dist[t];
}
return make_pair(flow, cost);
}
```

### 2.3 Gomory-Hu Tree

```
int g[maxn];
vector<edge> GomoryHu(int n){
  vector<edge> rt;
  for(int i=1;i<=n;++i)g[i]=1;
  for(int i=2;i<=n;++i){
    int t=g[i];
    flow.reset(); // clear flows on all edge
    rt.push_back({i,t,flow(i,t)});
    flow.walk(i); // bfs points that connected to i (use edges not fully flow)
    for(int j=i+1;j<=n;++j){
        if(g[j]==t && flow.connect(j))g[j]=i; // check if i can reach j
    }
}
return rt;
}</pre>
```

# 2.4 Stoer-Wagner Minimum Cut

```
int w[kN][kN], g[kN], del[kN], v[kN];
void AddEdge(int x, int y, int c) {
   w[x][y] += c;
   w[y][x] += c;
pair<int, int> Phase(int n) {
   fill(v, v + n, 0), fill(g, g + n, 0);
   int s = -1, t = -1;
   while (true) {
     int c = -1;
     for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;</pre>
        if (c == -1 || g[i] > g[c]) c = i;
     if (c == -1) break;
     v[c] = 1, s = t, t = c;
     for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
        g[i] += w[c][i];
   return make_pair(s, t);
int GlobalMinCut(int n) {
   int cut = kInf;
   fill(del, 0, sizeof(del));
   for (int i = 0; i < n - 1; ++i) {
     int s, t; tie(s, t) = Phase(n)
     del[t] = 1, cut = min(cut, g[t]);
     for (int j = 0; j < n; ++j) {
  w[s][j] += w[t][j];</pre>
        w[j][s] += w[j][t];
     }
   return cut;
}
```

#### 2.5 Kuhn-Munkres Algorithm

```
int64_t KuhnMunkres(vector<vector<int>> W) {
   int N = W.size();
   vector<int>> fl(N, -1), fr(N, -1), hr(N), hl(N);
   for (int i = 0; i < N; ++i) {
     hl[i] = *max_element(W[i].begin(), W[i].end());
   }
   auto Bfs = [&](int s) {
     vector<int>> slk(N, kInf), pre(N);
     vector<bool>> vl(N, false), vr(N, false);
     queue<int>> que;
     que.push(s);
   vr[s] = true;
   auto Check = [&](int x) -> bool {
     if (vl[x] = true, fl[x] != -1) {
        que.push(fl[x]);
        return vr[fl[x]] = true;
}
```

```
while (x != -1) swap(x, fr[fl[x] = pre[x]]);
    return false:
  while (true) {
    while (!que.empty()) {
      int y = que.front(); que.pop();
      for (int x = 0, d = 0; x < N; ++x) {
        if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] - W[x][y])
          if (pre[x] = y, d) slk[x] = d;
          else if (!Check(x)) return;
    int d = kInf;
for (int x = 0; x < N; ++x) {
      if (!vl[x] \&\& d > slk[x]) d = slk[x];
    for (int x = 0; x < N; ++x) {
      if (vl[x]) hl[x] += d;
      else slk[x] -= d;
      if (vr[x]) hr[x] -= d;
    for (int x = 0; x < N; ++x) {
      if (!vl[x] && !slk[x] && !Check(x)) return;
 }
for (int i = 0; i < N; ++i) Bfs(i);
int64_t res = 0;
for (int i = 0; i < N; ++i) res += W[i][fl[i]];</pre>
return res;
```

# 2.6 Maximum Matching on General Graph

```
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
 for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
 g[u].push_back(v);
  g[v].push_back(u);
int Find(int u) {
  return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
 static int tk = 0;
  tk++:
  x = Find(x), y = Find(y);
  for (; ; swap(x, y)) {
    if (x != n) {
      if (v[x] == tk) return x;
      v[x] = tk;
      x = Find(pre[match[x]]);
   }
 }
pre[x] = y, y = match[x];
    if (s[y] == 1) q.push(y), s[y] = 0;
    if (fa[x] == x) fa[x] = 1;
    if (fa[y] == y) fa[y] = 1;
    x = pre[y];
 }
bool Bfs(int r, int n) {
 for (int i = 0; i \le n; ++i) fa[i] = i, s[i] = -1;
  while (!q.empty()) q.pop();
  q.push(r);
  s[r] = 0;
  while (!q.empty()) {
    int x = q.front(); q.pop();
    for (int u : g[x]) {
      if (s[u] == -1) {
        pre[u] = x, s[u] = 1;
        if (match[u] == n) {
          for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
```

```
last = match[b], match[b] = a, match[a] = b;
           return true;
         }
         q.push(match[u]);
         s[match[u]] = 0;
       } else if (!s[u] && Find(u) != Find(x)) {
         int 1 = LCA(u, x, n);
         Blossom(x, u, 1);
         Blossom(u, x, 1);
    }
  }
  return false;
int Solve(int n) {
   int res = 0;
   for (int x = 0; x < n; ++x) {
     if (match[x] == n) res += Bfs(x, n);
   return res:
}}
```

# 2.7 Maximum Weighted Matching on General Graph

```
| struct WeightGraph {
   static const int inf = INT_MAX;
   static const int maxn = 514;
   struct edge {
     int u, v, w;
     edge(){}
     edge(int u, int v, int w): u(u), v(v), w(w) {}
   int n, n_x;
   edge g[maxn * 2][maxn * 2];
   int lab[maxn * 2];
   int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa[maxn *
   int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[maxn * 2];
   vector<int> flo[maxn * 2];
   int e_delta(const edge &e) { return lab[e.u] + lab[e.v] - g[e
   .u][e.v].w * 2; }
void update_slack(int u, int x) { if (!slack[x] || e_delta(g[
     u][x]) < e_delta(g[slack[x]][x])) slack[x] = u; }
   void set_slack(int x) {
     slack[x] = 0;
     for (int u = 1; u \le n; ++u)
       if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
         update_slack(u, x);
   void q_push(int x) {
     if (x \le n) q.push(x);
     else for (size_t i = 0; i < flo[x].size(); i++) q_push(flo[</pre>
     x][i]);
  void set_st(int x, int b) {
     st[x] = b;
     if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)
      set_st(flo[x][i], b);
   int get_pr(int b, int xr) {
     int pr = find(flo[b].begin(), flo[b].end(), xr) - flo[b].
     begin();
     if (pr % 2 == 1) {
       reverse(flo[b].begin() + 1, flo[b].end());
       return (int)flo[b].size() - pr;
     return pr;
  }
   void set_match(int u, int v) {
     match[u] = g[u][v].v;
     if (u <= n) return;</pre>
     edge e = g[u][v];
     int xr = flo_from[u][e.u], pr = get_pr(u, xr)
     for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo[u][i</pre>
     ^ 1]);
     set match(xr. v):
     rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].end());
   void augment(int u, int v) {
     for (; ; ) {
       int xnv = st[match[u]];
       set_match(u, v);
       if (!xnv) return:
       set_match(xnv, st[pa[xnv]]);
```

```
u = st[pa[xnv]], v = xnv;
  }
}
int get_lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0:
void add_blossom(int u, int lca, int v) {
  int b = n + 1;
  while (b <= n_x && st[b]) ++b;</pre>
  if (b > n_x) ++n_x;
lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[match[x]]),
   q_push(y);
  reverse(flo[b].begin() + 1, flo[b].end())
  for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].push\_back(x), flo[b].push\_back(y = st[match[x]]),
   q_push(y);
  set_st(b, b);
  for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].w = 0;
  for (int x = 1; x \le n; ++x) flo_from[b][x] = 0;
  for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)
       if (g[b][x].w == 0 \mid \mid e_delta(g[xs][x]) < e_delta(g[b][
   x]))
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
  if (flo_from[xs][x]) flo_from[b][x] = xs;</pre>
  set slack(b):
}
void expand_blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)</pre>
    set_st(flo[b][i], flo[b][i]);
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b, xr);
  for (int i = 0; i < pr; i += 2) {
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
  for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    S[xs] = -1, set_slack(xs);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
   S[nu] = 0, q_push(nu);
else if (S[v] == 0) {
    int lca = get_lca(u, v);
    if (!lca) return augment(u,v), augment(v,u), true;
    else add_blossom(u, lca, v);
  }
  return false;
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
  memset(slack + 1, 0, sizeof(int) * n_x);
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
    if (st[x] == x \&\& !match[x]) pa[x] = 0, S[x] = 0, q_push(
  if (q.empty()) return false;
  for (; ; ) {
    while (q.size()) {
      int u = q.front(); q.pop();
```

```
if (S[st[u]] == 1) continue;
         for (int v = 1; v <= n; ++v)
if (g[u][v].w > 0 && st[u] != st[v]) {
              if (e_delta(g[u][v]) == 0) {
                if (on_found_edge(g[u][v])) return true;
              } else update_slack(u, st[v]);
       int d = inf;
       for (int b = n + 1; b \le n_x; ++b)
         if (st[b] == b && S[b] == 1) d = min(d, lab[b] / 2);
       for (int x = 1; x <= n_x; ++x)
         if (st[x] == x \&\& slack[x]) {
            if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x]));
            else if (S[x] == 0) d = min(d, e_delta(g[slack[x]][x
      ]) / 2);
       for (int u = 1; u \le n; ++u) {
         if (S[st[u]] == 0) {
            if (lab[u] <= d) return 0;</pre>
         lab[u] -= d;
} else if (S[st[u]] == 1) lab[u] += d;
       for (int b = n + 1; b \le n_x; ++b)
          if (st[b] == b) {
            if (S[st[b]] == 0) lab[b] += d * 2;
            else if (S[st[b]] == 1) lab[b] -= d * 2;
         }
       q = queue<int>();
       for (int x = 1; x <= n_x; ++x)
  if (st[x] == x && slack[x] && st[slack[x]] != x &&</pre>
      e_delta(g[slack[x]][x]) == 0)
            if (on_found_edge(g[slack[x]][x])) return true;
       for (int b = n + 1; b \le n_x; ++b)
         if (st[b] == b \&\& S[b] == 1 \&\& lab[b] == 0)
      expand_blossom(b);
     return false;
   pair<long long, int> solve() {
     memset(match + 1, 0, sizeof(int) * n);
     n x = n:
     int n_matches = 0;
     long long tot_weight = 0;
     for (int u = 0; u <= n; ++u) st[u] = u, flo[u].clear();</pre>
     int w_max = 0;
     for (int u = 1; u <= n; ++u)</pre>
       for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
         w_max = max(w_max, g[u][v].w);
     for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
     while (matching()) ++n_matches;
     for (int u = 1; u \le n; ++u)
       if (match[u] && match[u] < u)</pre>
         tot_weight += g[u][match[u]].w;
     return make_pair(tot_weight, n_matches);
   void add_edge(int ui, int vi, int wi) { g[ui][vi].w = g[vi][
      ui].w = wi; }
   void init(int _n) {
     n = _n;
     for (int u = 1; u <= n; ++u)</pre>
       for (int v = 1; v <= n; ++v)
         g[u][v] = edge(u, v, 0);
  }
};
 2.8
        Minimum Cost Circulation
```

```
struct Edge { int to, cap, rev, cost; };
vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
bool mark[kN];
int NegativeCycle(int n) {
  memset(mark, false, sizeof(mark));
  memset(dist, 0, sizeof(dist));
  int upd = -1;
  for (int i = 0; i <= n; ++i) {
    for (int j = 0; j < n; ++j) {
      int idx = 0;
      for (auto &e : g[j]) {
        if (e.cap > 0 && dist[e.to] > dist[j] + e.cost) {
          dist[e.to] = dist[j] + e.cost;
          pv[e.to] = j, ed[e.to] = idx;
          if (i == n) {
```

```
while (!mark[upd]) mark[upd] = true, upd = pv[upd];
              return upd;
         idx++;
      }
    }
   return -1:
}
 int Solve(int n) {
   int rt = -1, ans = 0;
   while ((rt = NegativeCycle(n)) >= 0) {
     memset(mark, false, sizeof(mark));
     vector<pair<int, int>> cyc;
     while (!mark[rt]) {
       cyc.emplace_back(pv[rt], ed[rt]);
       mark[rt] = true;
       rt = pv[rt];
     reverse(cyc.begin(), cyc.end());
     int cap = kInf;
     for (auto &i : cyc) {
       auto &e = g[i.first][i.second];
       cap = min(cap, e.cap);
     for (auto &i : cyc) {
  auto &e = g[i.first][i.second];
       e.cap -= cap;
       g[e.to][e.rev].cap += cap;
ans += e.cost * cap;
     }
   return ans;
|}
```

#### 2.9Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x,y,l,u), connect  $x\to y$  with capacity u-l. 3. For each vertex v, denote by in(v) the difference between the sum
  - of incoming lower bounds and the sum of outgoing lower bounds. 4. If in(v) > 0, connect  $S \to v$  with capacity in(v), otherwise,
  - connect  $v \to T$  with capacity -in(v).
    - To maximize, connect t - $\rightarrow s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution.
    - Otherwise, the maximum flow from s to t is the answer. To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f + f' \neq \sum_{v \in V, in(v) > 0}^{\infty} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e + f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x, y) \in M$ ,  $x \to y$  otherwise. 2. DFS from unmatched vertices in X. 3.  $x \in X$  is chosen iff x is unvisited.

  - 4.  $y \in Y$  is chosen iff y is visited.
- Maximum density induced subgraph

  - 1. Binary search on answer, suppose we're checking answer T 2. Construct a max flow model, let K be the sum of all weights 3. Connect source  $s \to v, \, v \in G$  with capacity K

  - 4. For each edge (u, v, w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
  - For  $v \in G$ , connect it with sink  $v \to t$  with capacity  $K + 2T (\sum_{e \in E(v)} w(e)) 2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight
  - Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
  - 1. If  $p_v > 0$ , create edge (s, v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$
  - 2. Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.

    3. The mincut is equivalent to the maximum profit of a subset of
  - projects.

$$\sum_{x} c_{x} x + \sum_{y} c_{y} \bar{y} + \sum_{xy} c_{xy} x \bar{y} + \sum_{xyx'y'} c_{xyx'y'} (x \bar{y} + x' \bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$  and create edge (s,y) with capacity  $c_n$
- Create edge (x, y) with capacity  $c_{xy}$
- 3. Create edge (x, y) and edge (x', y') with capacity  $c_{xyx'y'}$ .

# Data Structure

#### <ext/pbds> 3.1

```
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef tree<int, null_type, std::less<int>, rb_tree_tag,
     tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
  tree_set s;
  s.insert(71); s.insert(22);
  assert(*s.find_by_order(0) == 22); assert(*s.find_by_order(1)
      == 71):
  assert(s.order_of_key(22) == 0); assert(s.order_of_key(71) ==
      1);
  s.erase(22);
  assert(*s.find_by_order(0) == 71); assert(s.order_of_key(71)
     == 0);
  // mergable heap
  heap a, b; a.join(b);
  // persistant
  rope<char> r[2];
  r[1] = r[0];
  std::string st = "abc";
  r[1].insert(0, st.c_str());
  r[1].erase(1, 1);
  std::cout << r[1].substr(0, 2) << std::endl;</pre>
  return 0;
```

#### Li Chao Tree 3.2

```
namespace lichao {
struct line {
  long long a, b;
   line(): a(0), b(0) {}
   line(long long a, long long b): a(a), b(b) {}
   long long operator()(int x) const { return a * x + b; }
line st[maxc * 4];
 int sz, lc[maxc * 4], rc[maxc * 4];
 int gnode() {
  st[sz] = line(1e9, 1e9);
lc[sz] = -1, rc[sz] = -1;
   return sz++;
}
void init() {
  sz = 0;
void add(int l, int r, line tl, int o) {
   bool lcp = st[o](l) > tl(l);
   bool mcp = st[o]((l + r) / 2) > tl((l + r) / 2);
   if (mcp) swap(st[o], tl);
   if (r - l == 1) return;
   if (lcp != mcp) {
     if (lc[o] == -1) lc[o] = gnode();
     add(1, (1 + r) / 2, t1, lc[o]);
   } else {
     if (rc[o] == -1) rc[o] = gnode();
     add((1 + r) / 2, r, tl, rc[o]);
long long query(int l, int r, int x, int o) {
  if (r - l == 1) return st[o](x);
   if (x < (l + r) / 2) {
     if (lc[o] == -1) return st[o](x);
     return min(st[o](x), query(l, (l + r) / 2, x, lc[o]));
   } else {
     if (rc[o] == -1) return st[o](x);
     return min(st[o](x), query((l + r) / 2, r, x, rc[o]));
| }}
```

### 3.3 Link-Cut Tree

```
struct node {
  node *ch[2], *fa, *pfa;
 int sum, v, rev, id;
node(int s, int id): id(id), v(s), sum(s), rev(0), fa(nullptr
     ), pfa(nullptr) {
    ch[0] = nullptr;
    ch[1] = nullptr;
  int relation() {
    return this == fa->ch[0] ? 0 : 1;
  void push() {
    if (!rev) return;
    swap(ch[0], ch[1]);
    if (ch[0]) ch[0]->rev ^= 1;
    if (ch[1]) ch[1]->rev ^= 1;
    rev = 0;
  void pull() {
    sum = v
    if (ch[0]) sum += ch[0]->sum;
    if (ch[1]) sum += ch[1]->sum;
  void rotate() {
    if (fa->fa) fa->fa->push();
    fa->push(), push(), swap(pfa, fa->pfa);
    int d = relation();
    node *t = fa;
    if (t->fa) t->fa->ch[t->relation()] = this;
    fa = t -> fa, t -> ch[d] = ch[d \land 1];
    if (ch[d ^ 1]) ch[d ^ 1]->fa = t;
ch[d ^ 1] = t, t->fa = this;
    t->pull(), pull();
 }
  void splay() {
    while (fa) {
      if (!fa->fa) {
        rotate();
         continue;
      fa->fa->push(), fa->push();
      if (relation() == fa->relation()) fa->rotate();
      else rotate(), rotate();
 }
  void evert() { access(), splay(), rev ^= 1; }
  void expose() {
    splay(), push();
    if (ch[1]) {
      ch[1]->fa = nullptr, ch[1]->pfa = this;
      ch[1] = nullptr, pull();
    }
  bool splice() {
    splay();
    if (!pfa) return false:
    pfa->expose(), pfa->ch[1] = this, fa = pfa;
    pfa = nullptr, fa->pull();
    return true;
  void access() {
    expose();
    while (splice());
  int query() { return sum; }
namespace lct {
node *sp[maxn];
void make(int u, int v) {
 // create node with id u and value v
  sp[u] = new node(v, u);
void link(int u, int v) {
  // u become v's parent
  sp[v]->evert();
  sp[v]->pfa = sp[u];
void cut(int u, int v) {
  // u was v's parent
  sp[u]->evert();
  sp[v]->access(), sp[v]->splay(), sp[v]->push();
  sp[v]->ch[0]->fa = nullptr;
  sp[v]->ch[0] = nullptr;
  sp[v]->pull();
```

```
void modify(int u, int v) {
  sp[u]->splay();
  sp[u]->v = v;
  sp[u]->pull();
}
int query(int u, int v) {
  sp[u]->evert(), sp[v]->access(), sp[v]->splay();
  return sp[v]->query();
int find(int u) -
  sp[u]->access();
  sp[u]->splay();
  node *p = sp[u];
  while (true) {
    p->push();
    if (p->ch[0]) p = p->ch[0];
    else break;
  3
  return p->id;
```

# 4 Graph

# 4.1 Heavy-Light Decomposition

```
void dfs(int x, int p) {
   dep[x] = \sim p ? dep[p] + 1 : dep[x];
   sz[x] = 1;
   to[x] = -1;
   fa[x] = p;
   for (const int &u : g[x]) {
     if (u == p) continue;
     dfs(u, x);
     sz[x] += sz[u];
     if (to[x] == -1 \mid | sz[to[x]] < sz[u]) to[x] = u;
  }
}
void hld(int x, int t) {
   static int tk = 0;
   fr[x] = t;
   dfn[x] = tk++;
   if (!~to[x]) return;
   hld(to[x], t);
   for (const int &u : g[x]) {
     if (u == fa[x] || u == to[x]) continue;
     hld(u, u);
   }
vector<pair<int, int>> get(int x, int y) {
   int fx = fr[x], fy = fr[y];
   vector<pair<int, int>> res;
   while (fx != fy) {
     if (dep[fx] < dep[fy]) {</pre>
       swap(fx, fy);
       swap(x, y);
     res.emplace_back(dfn[fx], dfn[x] + 1);
     x = fa[fx];
     fx = fr[x];
   res.emplace\_back(min(dfn[x], dfn[y]), max(dfn[x], dfn[y]) +
   int lca = (dep[x] < dep[y] ? x : y);
   return res;
}
```

# 4.2 Centroid Decomposition

```
void get_center(int now) {
  v[now] = true; vtx.push_back(now);
  sz[now] = 1; mx[now] = 0;
  for (int u : G[now]) if (!v[u]) {
    get_center(u);
    mx[now] = max(mx[now], sz[u]);
    sz[now] += sz[u];
  }
}
void get_dis(int now, int d, int len) {
  dis[d][now] = cnt;
  v[now] = true;
  for (auto u : G[now]) if (!v[u.first]) {
    get_dis(u, d, len + u.second);
  }
}
```

#### 4.3 Minimum Mean Cycle

```
// d[i][j] == 0 if {i,j} !in E
long long d[1003][1003],dp[1003][1003];
pair<long long,long long> MMWC(){
  memset(dp,0x3f,sizeof(dp))
  for(int i=1;i<=n;++i)dp[0][i]=0;</pre>
  for(int i=1;i<=n;++i){</pre>
    for(int j=1; j<=n;++j){</pre>
      for(int k=1;k<=n;++k){</pre>
         dp[i][k]=min(dp[i-1][j]+d[j][k],dp[i][k]);
    }
  long long au=1ll<<31,ad=1;
  for(int i=1;i<=n;++i){</pre>
    if(dp[n][i]==0x3f3f3f3f3f3f3f3f3f)continue;
    long long u=0,d=1;
for(int j=n-1;j>=0;--j){
      if((dp[n][i]-dp[j][i])*d>u*(n-j)){
         u=dp[n][i]-dp[j][i];
         d=n-j;
      }
    if(u*ad<au*d)au=u,ad=d;</pre>
  long long g=__gcd(au,ad);
  return make_pair(au/g,ad/g);
```

#### 4.4 Minimum Steiner Tree

```
namespace steiner {
// Minimum Steiner Tree - O(N * 3^K + N^2 * 2^K)
// z[i] = the weight of the i-th vertex
const int maxn = 64, maxk = 10;
const int inf = 1e9;
int w[maxn][maxn], z[maxn], dp[1 << maxk][maxn], off[maxn];</pre>
void init(int n) {
  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < n; ++j) w[i][j] = inf;
    z\Gamma i = 0;
    w[i][i] = 0;
 }
void add_edge(int x, int y, int d) {
  w[x][y] = min(w[x][y], d);
  w[y][x] = min(w[y][x], d);
void build(int n) {
  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < n; ++j) {
    w[i][j] += z[i];
      if (i != j) w[i][j] += z[j];
    }
  }
  for (int k = 0; k < n; ++k) {
    for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) w[i][j] = min(w[i][j], w[i][k
] + w[k][j] - z[k]);</pre>
  }
int solve(int n, vector<int> mark) {
  build(n);
  int k = (int)mark.size();
  assert(k < maxk);</pre>
  for (int s = 0; s < (1 << k); ++s) {
```

```
for (int i = 0; i < n; ++i) dp[s][i] = inf;</pre>
   for (int i = 0; i < n; ++i) dp[0][i] = 0;
   for (int s = 1; s < (1 << k); ++s) {
     if (__builtin_popcount(s) == 1) {
       int x = __builtin_ctz(s);
for (int i = 0; i < n; ++i) dp[s][i] = w[mark[x]][i];</pre>
       continue:
     for (int i = 0; i < n; ++i) {
       for (int sub = s & (s - 1); sub; sub = s & (sub - 1)) {
         dp[s][i] = min(dp[s][i], dp[sub][i] + dp[s \land sub][i] -
      z[i]);
     for (int i = 0; i < n; ++i) {
       off[i] = inf;
       for (int j = 0; j < n; ++j) off[i] = min(off[i], dp[s][j]
       + w[j][i] - z[j]);
     for (int i = 0; i < n; ++i) dp[s][i] = min(dp[s][i], off[i</pre>
     ]);
  }
   int res = inf;
   for (int i = 0; i < n; ++i) res = min(res, dp[(1 << k) - 1][i
      1);
   return res;
| }}
```

# 4.5 Directed Minimum Spanning Tree

```
template <typename T> struct DMST {
  T g[maxn][maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
    for(int i = 0; i < maxn; ++i) {</pre>
       for(int j = 0; j < maxn; ++j) g[i][j] = inf;
      vis[i] = inc[i] = false;
    }
  }
  void addedge(int u, int v, T w) {
    g[u][v] = min(g[u][v], w);
  T operator()(int root, int _n) {
    if (dfs(root) != n) return -1;
    T ans = 0;
    while (true) {
       for (int i = 1; i \le n; ++i) fw[i] = inf, fr[i] = i;
       for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
         for (int j = 1; j <= n; ++j) {
  if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
             fw[i] = g[j][i];
             fr[i] = j;
           }
        }
      }
       int x = -1;
       for (int i = 1; i <= n; ++i) if (i != root && !inc[i]) {</pre>
         int j = i, c = 0;
         while (j != root && fr[j] != i && c <= n) ++c, j = fr[j]
         if (j == root || c > n) continue;
         else { x = i; break; }
      if (!~x) {
         for (int i = 1; i \le n; ++i) if (i != root \&\& !inc[i])
     ans += fw[i];
         return ans;
       for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
      do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true; }
     while (y != x);
       inc[x] = false;
       for (int k = 1; k \le n; ++k) if (vis[k]) {
         for (int j = 1; j <= n; ++j) if (!vis[j]) {</pre>
           if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
if (g[j][k] < inf && g[j][k] - fw[k] < g[j][x]) g[j][</pre>
     x] = g[j][k] - fw[k];
        }
      }
    return ans;
  int dfs(int now) {
```

```
int r = 1;
vis[now] = true;
for (int i = 1; i <= n; ++i) if (g[now][i] < inf && !vis[i
]) r += dfs(i);
return r;
}
};</pre>
```

# 4.6 Maximum Clique

```
struct MaxClique {
  // change to bitset for n > 64.
   int n, deg[maxn];
  uint64_t adj[maxn], ans;
  vector<pair<int, int>> edge;
   void init(int n_) {
     fill(adj, adj + n, 0ull);
     fill(deg, deg + n, 0);
     edge.clear();
  }
  void add_edge(int u, int v) {
     edge.emplace_back(u, v);
     ++deg[u], ++deg[v];
  }
  vector<int> operator()() {
     vector<int> ord(n);
     iota(ord.begin(), ord.end(), 0);
sort(ord.begin(), ord.end(), [&](int u, int v) { return deg
     [u] < deg[v]; });
     vector<int> id(n);
     for (int i = 0; i < n; ++i) id[ord[i]] = i;</pre>
     for (auto e : edge) {
       int u = id[e.first], v = id[e.second];
       adj[u] |= (1ull << v);
       adj[v] |= (1ull << u);
     uint64_t r = 0, p = (1ull << n) - 1;
     dfs(r, p);
     vector<int> res;
     for (int i = 0; i < n; ++i) {
       if (ans >> i & 1) res.push_back(ord[i]);
     return res;
#define pcount __builtin_popcountll
  void dfs(uint64_t r, uint64_t p) {
     if (p == 0) {
       if (pcount(r) > pcount(ans)) ans = r;
       return;
     if (pcount(r | p) <= pcount(ans)) return;</pre>
     int x = __builtin_ctzll(p & -p);
     uint64_t c = p & \simadj[x];
     while (c > 0) {
       // bitset._Find_first(); bitset._Find_next();
       x = __builtin_ctzll(c & -c);
       r |= (1ull << x);
       dfs(r, p & adj[x]);
       r &= ~(1ull << x);
       p \&= \sim (1ull << x);
       c ^= (1ull << x);
  }
|};
```

# 4.7 Tarjan's Algorithm

```
void dfs(int x, int p) {
   dfn[x] = low[x] = tk++;
   int ch = 0;
   st.push(x); // bridge
   for (auto e : g[x]) if (e.first != p) {
      if (!ins[e.second]) { // articulation point
        st.push(e.second);
      ins[e.second] = true;
   }
   if (~dfn[e.first]) {
      low[x] = min(low[x], dfn[e.first]);
      continue;
   }
   dfs(u.first, x);
   if (low[u.first] >= low[x]) { // articulation point
      cut[x] = true;
   while (true) {
      int z = st.top(); st.pop();
   }
}
```

```
bcc[z] = sz;
    if (z == e.second) break;
}
sz++;
}
if (ch == 1 && p == -1) cut[x] = false;
if (dfn[x] == low[x]) { // bridge
    while (true) {
        int z = st.top(); st.pop();
        bcc[z] = sz;
        if (z == x) break;
}
}
```

#### 4.8 Dominator Tree

```
| vector<int> BuildDominatorTree(vector<vector<int>> q, int s) {
   int N = g.size();
   vector<vector<int>> rdom(N), r(N);
   vector < int > dfn(N, -1), rev(N, -1), fa(N, -1), sdom(N, -1),
      dom(N, -1), val(N, -1), rp(N, -1);
   int stamp = 0;
   auto Dfs = [\&](auto dfs, int x) -> void {
     rev[dfn[x] = stamp] = x;
     fa[stamp] = sdom[stamp] = val[stamp] = stamp;
     for (int u : g[x]) {
       if(dfn[u] == -1) {
         dfs(dfs, u);
         rp[dfn[u]] = dfn[x];
       r[dfn[u]].push_back(dfn[x]);
    }
  };
   function<int(int, int)> Find = [&](int x, int c) {
     if (fa[x] == x) return c ? -1 : x;
     int p = Find(fa[x], 1);
     if (p == -1) return c ? fa[x] : val[x];
     if (sdom[val[x]] > sdom[val[fa[x]]]) val[x] = val[fa[x]];
     fa[x] = p;
     return c ? p : val[x];
   auto Merge = [\&](int x, int y) \{ fa[x] = y; \};
  Dfs(Dfs, s);
   for (int i = stamp - 1; i >= 0; --i) {
     for (int u : r[i]) sdom[i] = min(sdom[i], sdom[Find(u, 0)])
     if (i) rdom[sdom[i]].push_back(i);
     for (int u : rdom[i]) {
       int p = Find(u, 0);
if (sdom[p] == i) dom[u] = i;
       else dom[u] = p;
     if (i) Merge(i, rp[i]);
  }
  vector<int> res(N, -2);
   res[s] = -1;
   for (int i = 1; i < stamp; ++i) {</pre>
     if (sdom[i] != dom[i]) dom[i] = dom[dom[i]];
   for (int i = 1; i < stamp; ++i) res[rev[i]] = rev[dom[i]];</pre>
   return res;
13
```

# 4.9 Virtual Tree

```
void VirtualTree(vector<int> v) {
  v.push_back(0);
  sort(v.begin(), v.end(), [&](int i, int j) { return dfn[i] <</pre>
     dfn[j]; });
  v.resize(unique(v.begin(), v.end()) - v.begin());
  vector<int> stk;
  for (int u : v) {
    if (stk.empty()) {
      stk.push_back(u);
       continue;
    int p = GetLCA(u, stk.back());
    if (p != stk.back()) {
      while (stk.size() >= 2 && dep[p] <= dep[stk[stk.size() -</pre>
     2]]) {
         int x = stk.back();
        stk.pop_back();
        AddEdge(x, stk.back());
```

```
if (stk.back() != p) {
   AddEdge(stk.back(), p);
   stk.pop_back();
   stk.push_back(p);
}

stk.push_back(u);
}

for (int i = 0; i + 1 < stk.size(); ++i) AddEdge(stk[i], stk[i + 1]);
}</pre>
```

# 4.10 Vizing's Theorem

```
namespace vizing { // returns edge coloring in adjacent matrix
      G. 1 - based
 int C[kN][kN], G[kN][kN];
void clear(int N) {
  for (int i = 0; i <= N; i++) {</pre>
     for (int j = 0; j \le N; j++) C[i][j] = G[i][j] = 0;
  }
}
void solve(vector<pair<int, int>> &E, int N, int M) {
  int X[kN] = {}, a;
auto update = [&](int u) {
     for (X[u] = 1; C[u][X[u]]; X[u]++);
   auto color = [&](int u, int v, int c) {
     int p = G[u][v];
     G[u][v] = G[v][u] = c;
     C[u][c] = v, C[v][c] = u;
     C[u][p] = C[v][p] = 0;
     if (p) X[u] = X[v] = p;
     else update(u), update(v);
     return p;
  };
   auto flip = [&](int u, int c1, int c2) {
     int p = C[u][c1];
     swap(C[u][c1], C[u][c2]);
     if (p) G[u][p] = G[p][u] = c2;
     if (!C[u][c1]) X[u] = c1;
     if (!C[u][c2]) X[u] = c2;
     return p;
  for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {</pre>
     int u = E[t].first, v0 = E[t].second, v = v0, c0 = X[u], c
     vector<pair<int, int>> L;
     int vst[kN] = {};
     while (\bar{[G[u][v0]}) {
       L.emplace_back(v, d = X[v]);
if (!C[v][c]) for (a = (int)L.size() - 1; a >= 0; a--) c
      = color(u, L[a].first, c);
       else if ((C[u][d]) for (a = (int)L.size() - 1; a >= 0; a >= 0
      --) color(u, L[a].first, L[a].second);
       else if (vst[d]) break
       else vst[d] = 1, v = C[u][d];
     if (!G[u][v0]) {
       for (; v; v = flip(v, c, d), swap(c, d));
       if (C[u][c0]) {
          for (a = (int)L.size() - 2; a >= 0 && L[a].second != c;
         for (; a \ge 0; a - -) color(u, L[a].first, L[a].second);
       } else t--;
j }}
```

#### 4.11 System of Difference Constraints

Given m constrains on n variables  $x_1, x_2, \ldots, x_n$  of form  $x_i - x_j \leq w$  (resp,  $x_i - x_j \geq w$ ), connect  $i \to j$  with weight w. Then connect  $0 \to i$  for all i with weight 0 and find the shortest path (resp, longest path) on the graph. dis(i) will be the maximum (resp, minimum) solution to  $x_i$ .

# 5 String

# 5.1 Knuth-Morris-Pratt Algorithm

```
vector<int> kmp(const string &s) {
  vector<int> f(s.size(), 0);
  // f[i] = length of the longest prefix (excluding s[0:i])
  such that it coincides with the suffix of s[0:i] of the
  same length
```

```
// i + 1 - f[i] is the length of the smallest recurring
     period of s[0:i]
   int k = 0;
   for (int i = 1; i < (int)s.size(); ++i) {</pre>
     while (k > 0 \& s[i] != s[k]) k = f[k - 1];
     if (s[i] == s[k]) ++k;
     f[i] = k;
  }
   return f;
vector<int> search(const string &s, const string &t) {
   // return 0-indexed occurrence of t in s
   vector<int> f = kmp(t), res;
for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
     while (k > 0 \& (k == (int)t.size() || s[i] != t[k])) k = f
     [k - 1];
     if (s[i] == t[k]) ++k;
     if (k == (int)t.size()) res.push_back(i - t.size() + 1);
   return res;
}
```

# 5.2 Z Algorithm

```
int z[maxn];
// z[i] = LCP of suffix i and suffix 0
void z_function(const string& s) {
    memset(z, 0, sizeof(z));
    z[0] = (int)s.length();
    int l = 0, r = 0;
    for (int i = 1; i < s.length(); ++i) {
        z[i] = max(0, min(z[i - l], r - i + 1));
        while (i + z[i] < s.length() && s[z[i]] == s[i + z[i]]) {
        l = i; r = i + z[i];
        ++z[i];
    }
}
</pre>
```

# 5.3 Manacher's Algorithm

```
int z[maxn];
int manacher(const string& s) {
   string t = ".";
   for (int i = 0; i < s.length(); ++i) t += s[i], t += '.';
   int l = 0, r = 0, ans = 0;
   for (int i = 1; i < t.length(); ++i) {
      z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
      while (i - z[i] >= 0 && i + z[i] < t.length() && t[i - z[i] ]] == t[i + z[i]]) ++z[i];
      if (i + z[i] > r) r = i + z[i], l = i;
   }
   for (int i = 1; i < t.length(); ++i) ans = max(ans, z[i] - 1);
   return ans;
}</pre>
```

#### 5.4 Aho-Corasick Automaton

ql = qr = 0; q[qr++] = root;

```
struct AC {
  static const int maxn = 1e5 + 5;
  int sz, ql, qr, root;
  int cnt[maxn], q[maxn], ed[maxn], el[maxn], ch[maxn][26], f[
  int gnode() {
    for (int i = 0; i < 26; ++i) ch[sz][i] = -1;
    f[sz] = -1;
    ed[sz] = 0;
    cnt[sz] = 0;
    return sz++;
  void init() {
    sz = 0;
    root = gnode();
  int add(const string &s) {
    for (int i = 0; i < s.length(); ++i) {
      if (ch[now][s[i] - 'a'] == -1) ch[now][s[i] - 'a'] =
     gnode();
      now = ch[now][s[i] - 'a'];
    ed[now] = 1;
    return now:
  void build_fail() {
```

```
while (al < ar) {
       int now = q[ql++];
       for (int i = 0; i < 26; ++i) if (ch[now][i] != -1) {
         int p = ch[now][i], fp = f[now];
         while (fp != -1 \&\& ch[fp][i] == -1) fp = f[fp];
         int pd = fp != -1 ? ch[fp][i] : root;
         f[p] = pd;
         el[p] = ed[pd] ? pd : el[pd];
         q[qr++] = p;
    }
  }
  void build(const string &s) {
    build_fail();
     int now = root;
    for (int i = 0; i < s.length(); ++i) {
      while (now != -1 \& ch[now][s[i] - 'a'] == -1) now = f[
      now = now != -1 ? ch[now][s[i] - 'a'] : root;
       ++cnt[now];
    for (int i = qr - 1; i >= 0; --i) cnt[f[q[i]]] += cnt[q[i]]
     ]];
  long long solve(int n) {
    build_fail();
    vector<vector<long long>> dp(sz, vector<long long>(n + 1,
     for (int i = 0; i < sz; ++i) dp[i][0] = 1;
     for (int i = 1; i <= n; ++i) {
       for (int j = 0; j < sz; ++j) {
         for (int k = 0; k < 2; ++k) {
           if (ch[j][k] != -1) {
             if (!ed[ch[j][k]])
               dp[j][i] += dp[ch[j][k]][i - 1];
           } else {
             int z = f[j];
             while (z != root \&\& ch[z][k] == -1) z = f[z];
             int p = ch[z][k] == -1 ? root : ch[z][k];
            if (ch[z][k] == -1 || !ed[ch[z][k]]) dp[j][i] += dp
     [p][i - 1];
        }
    return dp[0][n];
|};
```

#### 5.5 Suffix Automaton

```
struct SAM {
 static const int maxn = 5e5 + 5;
  int nxt[maxn][26], to[maxn], len[maxn];
  int root, last, sz;
 int gnode(int x) {
    for (int i = 0; i < 26; ++i) nxt[sz][i] = -1;
    to[sz] = -1;
   len[sz] = x;
   return sz++;
 void init() {
   sz = 0;
    root = gnode(0);
   last = root;
 void push(int c) {
   int cur = last;
    last = gnode(len[last] + 1);
    for (; ~cur && nxt[cur][c] == -1; cur = to[cur]) nxt[cur][c
    ] = last;
    if (cur == -1) return to[last] = root, void();
    int link = nxt[cur][c];
    if (len[link] == len[cur] + 1) return to[last] = link, void
    ();
    int tlink = gnode(len[cur] + 1);
    for (; ~cur && nxt[cur][c] == link; cur = to[cur]) nxt[cur
    ][c] = tlink;
    for (int i = 0; i < 26; ++i) nxt[tlink][i] = nxt[link][i];</pre>
   to[tlink] = to[link];
   to[link] = tlink;
   to[last] = tlink;
 void add(const string &s) {
   for (int i = 0; i < s.size(); ++i) push(s[i] - 'a');</pre>
```

```
bool find(const string &s) {
     int cur = root;
     for (int i = 0; i < s.size(); ++i) {</pre>
       cur = nxt[cur][s[i] - 'a'];
       if (cur == -1) return false;
     return true;
  }
   int solve(const string &t) {
     int res = 0, cnt = 0;
     int cur = root;
     for (int i = 0; i < t.size(); ++i) {</pre>
       if (~nxt[cur][t[i] - 'a']) {
         ++cnt;
         cur = nxt[cur][t[i] - 'a'];
       } else {
         for (; ~cur && nxt[cur][t[i] - 'a'] == -1; cur = to[cur
         if (\simcur) cnt = len[cur] + 1, cur = nxt[cur][t[i] - 'a'
     ];
         else cnt = 0, cur = root;
       }
       res = max(res, cnt);
     return res;
  }
};
```

### 5.6 Suffix Array

```
|// sa[i]: sa[i]-th suffix is the i-th lexigraphically smallest
      suffix.
   <code>lcp[i]: longest common prefix of suffix sa[i] and suffix sa[</code>
      i - 17.
namespace sfx {
vector<int> Build(const string &s) {
   int n = s.size();
   vector<int> str(n * 2), sa(n * 2), c(max(n, 256) * 2), x(max(n, 256)), p(n), q(n * 2), t(n * 2);
   for (int i = 0; i < n; ++i) str[i] = s[i];</pre>
   auto Pre = [&](int *sa, int *c, int n, int z) {
     memset(sa, 0, sizeof(int) * n);
     memcpy(x.data(), c, sizeof(int) * z);
  };
   auto Induce = [&](int *sa, int *c, int *s, int *t, int n, int
      z) {
     memcpy(x.data() + 1, c, sizeof(int) * (z - 1));
     for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] - 1]) sa[
     x[s[sa[i] - 1]]++] = sa[i] - 1;
memcpy(x.data(), c, sizeof(int) * z);
for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i] - 1])
       sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
  };
   auto SAIS = [&](auto self, int *s, int *sa, int *p, int *q,
     int *t, int *c, int n, int z) -> void {
     bool uniq = t[n - 1] = true;
     int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n, last =
     memset(c, 0, sizeof(int) * z);
     for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
     for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
     if (unia) {
       for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
       return:
     for (int i = n - 2; i >= 0; --i) t[i] = (s[i] == s[i + 1] ?
       t[i + 1] : s[i] < s[i + 1]);
     Pre(sa, c, n, z);
     for (int i = 1; i <= n - 1; ++i) if (t[i] && !t[i - 1]) sa
      [--x[s[i]]] = p[q[i] = nn++] = i;
     Induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) if (sa[i] && t[sa[i]] && !t[sa[
      i] - 1]) {
       bool neq = last < 0 | | memcmp(s + sa[i], s + last, (p[q[
      sa[i]] + 1] - sa[i]) * sizeof(int));
       ns[q[last = sa[i]]] = nmxz += neq;
     self(self, ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz +
      1):
     Pre(sa, c, n, z);
     for (int i = nn - 1; i >= 0; --i) sa[--x[s[p[nsa[i]]]]] = p
      [nsa[i]];
     Induce(sa, c, s, t, n, z);
   SAIS(SAIS, str.data(), sa.data(), p.data(), q.data(), t.data
      (), c.data(), n + 1, 256);
   return vector<int>(sa.begin() + 1, sa.begin() + n + 1);
```

#### 5.7 Palindromic Tree

```
struct PalindromicTree {
  int link[kN], len[kN], dp[kN], nxt[kN][26], sz, sf;
int gnode(int l, int fl = -1) {
    len[sz] = 1;
    link[sz] = fl;
    fill(nxt[sz], nxt[sz] + 26, -1);
    return sz++;
  void Init() {
    sz = 0;
    sf = 1;
    gnode(-1, 0);
    gnode(0, 0);
  void Push(const string &s, int pos) {
    int cur = sf, z = s[pos] - 'a'
    while (pos - 1 - len[cur] < 0 | | s[pos - 1 - len[cur]] != s
     [pos]) cur = link[cur];
    if (nxt[cur][z] != -1) {
      sf = nxt[cur][z];
    } else {
      int ch = gnode(len[cur] + 2);
      nxt[cur][z] = sf = ch;
      if (len[ch] == 1) {
        link[ch] = 1;
      } else {
        cur = link[cur];
        while (pos - 1 - len[cur] < 0 || s[pos - 1 - len[cur]]</pre>
     != s[pos]) cur = link[cur];
        link[ch] = nxt[cur][z];
      }
    dp[sf] += 1;
  long long Build(const string &s) {
    for (int i = 0; i < s.size(); ++i) Push(s, i);</pre>
    for (int i = sz - 1; i >= 0; --i) dp[link[i]] += dp[i];
    long long res = 0;
    for (int i = 0; i < sz; ++i) res = max(res, 1LL * dp[i] *
     len[i]);
    return res;
} plt;
```

#### 5.8 Circular LCS

```
string s1, s2;
int dp[kN * 2][kN];
int nxt[kN * 2][kN];
void reroot(int px) {
 int py = 1;
  while (py <= m && nxt[px][py] != 2) py++;</pre>
  nxt[px][py] = 1;
  while (px < 2 * n \&\& py < m) {
    if (nxt[px + 1][py] == 3) px++, nxt[px][py] = 1;
    else if (nxt[px + 1][py + 1] == 2) px++, py++, nxt[px][py]
     = 1:
    else py++;
  while (px < 2 * n && nxt[px + 1][py] == 3) px++, nxt[px][py]
    = 1;
int track(int x, int y, int e) { // use this routine to find
    LCS as string
  int ret = 0;
  while (y != 0 \&\& x != e) {
    if (nxt[x][y] == 1) y--;
```

```
else if (nxt[x][y] == 2) ret += (s1[x] == s2[y]), x--, y--;
     else if (nxt[x][y] == 3) x--;
   return ret;
 int solve(string a, string b) {
   n = a.size(), m = b.size();
s1 = "#" + a + a, s1 = '#' + b;
   for (int i = 0; i \le 2 * n; i++) {
     for (int j = 0; j <= m; j++) {
  if (j == 0) { nxt[i][j] = 3; continue; }</pre>
        if (i == 0) { nxt[i][j] = 1; continue; }
        dp[i][j] = -1;
        if (dp[i][j] < dp[i][j - 1]) dp[i][j] = dp[i][j - 1], nxt
      [i][j] = 1;
        if (dp[i][j] < dp[i - 1][j - 1] + (s1[i] == s2[j])) dp[i]
      [j] = dp[i - 1][j - 1] + (s1[i] == s2[j]), nxt[i][j] = 2;
        if (dp[i][j] < dp[i - 1][j]) dp[i][j] = dp[i - 1][j], nxt</pre>
      [i][j] = 3;
     }
   }
   int ret = dp[n][m];
   for (int i = 1; i < n; i++) reroot(i), ret = max(ret, track(n
       + i, m, i));
   return ret;
}
```

#### 5.9 Lexicographically Smallest Rotation

```
| string rotate(const string &s) {
   int n = s.length();
   string t = s + s;
   int i = 0, j = 1;
   while (i < n && j < n) {
      int k = 0;
      while (k < n && t[i + k] == t[j + k]) ++k;
      if (t[i + k] <= t[j + k]) j += k + 1;
      else i += k + 1;
      if (i == j) ++j;
   }
   int pos = (i < n ? i : j);
   return t.substr(pos, n);
}</pre>
```

#### 6 Math

### 6.1 Fast Fourier Transform

```
namespace fft {
struct cplx {
  double re, im;
  cplx(): re(0), im(0) {}
  cplx(double r, double i): re(r), im(i) {}
  cplx operator+(const cplx &rhs) const { return cplx(re + rhs.
     re, im + rhs.im); }
  cplx operator-(const cplx &rhs) const { return cplx(re - rhs.
     re, im - rhs.im); }
  cplx operator*(const cplx &rhs) const { return cplx(re * rhs.
    re - im * rhs.im, re * rhs.im + im * rhs.re); }
  cplx conj() const { return cplx(re, -im); }
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
  for (int i = 0; i \leftarrow maxn; ++i)
    omega[i] = cplx(cos(2 * pi * i / maxn), sin(2 * pi * i / maxn))
void bitrev(vector<cplx> &v, int n) {
  int z = __builtin_ctz(n) - 1;
  for (int i = 0; i < n; ++i) {
    int x = 0;
    for (int j = 0; (1 << j) < n; ++j) x ^= (i >> j & 1) << (z
     - i):
    if (x > i) swap(v[x], v[i]);
 }
void fft(vector<cplx> &v, int n) {
  bitrev(v, n);
  for (int s = 2; s <= n; s <<= 1) {
    int z = s \gg 1;
    for (int i = 0; i < n; i += s) {
      for (int k = 0; k < z; ++k) {
        cplx x = v[i + z + k] * omega[maxn / s * k];
```

v[i + z + k] = v[i + k] - x;

```
v[i + k] = v[i + k] + x;
     }
  }
 void ifft(vector<cplx> &v, int n) {
  fft(v, n);
   reverse(v.begin() + 1, v.end());
   for (int i = 0; i < n; ++i) v[i] = v[i] * cplx(1. / n, 0);
 vector<long long> convolution(const vector<int> &a, const
      vector<int> &b) {
   // Should be able to handle N <= 10^5, C <= 10^4
   int sz = 1;
   while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
   vector<cplx> v(sz);
   for (int i = 0; i < sz; ++i) {
     double re = i < a.size() ? a[i] : 0;
     double im = i < b.size() ? b[i] : 0;</pre>
     v[i] = cplx(re, im);
   fft(v, sz);
  for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);</pre>
     cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj()) * cplx
      (0, -0.25);
     if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i].conj
()) * cplx(0, -0.25);
     v[i] = x;
   ifft(v, sz);
   vector<long long> c(sz);
   for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);
   return c;
vector<int> convolution_mod(const vector<int> &a, const vector<</pre>
      int> &b, int p) {
   int sz = 1;
   while (sz < (int)a.size() + (int)b.size() - 1) sz <<= 1;</pre>
   vector<cplx> fa(sz), fb(sz);
for (int i = 0; i < (int)a.size(); ++i)</pre>
     fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
   for (int i = 0; i < (int)b.size(); ++i)</pre>
     fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
  fft(fa, sz), fft(fb, sz);
double r = 0.25 / sz;
   cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
   for (int i = 0; i \leftarrow (sz >> 1); ++i) {
     int j = (sz - i) & (sz - 1);
     cplx a1 = (fa[i] + fa[j].conj());
     cplx a2 = (fa[i] - fa[j].conj()) * r2;
     cplx b1 = (fb[i] + fb[j].conj()) * r3;
cplx b2 = (fb[i] - fb[j].conj()) * r4;
     if (i != j) {
       cplx c1 = (fa[j] + fa[i].conj());
       cplx c2 = (fa[j] - fa[i].conj()) * r2;
       cplx d1 = (fb[j] + fb[i].conj()) * r3;
cplx d2 = (fb[j] - fb[i].conj()) * r4;
       fa[i] = c1 * d1 + c2 * d2 * r5;
       fb[i] = c1 * d2 + c2 * d1;
     fa[j] = a1 * b1 + a2 * b2 * r5;
     fb[j] = a1 * b2 + a2 * b1;
   fft(fa, sz), fft(fb, sz);
   vector<int> res(sz);
   for (int i = 0; i < sz; ++i) {
     long long a = round(fa[i].re);
     long long b = round(fb[i].re);
     long long c = round(fa[i].im);
     res[i] = (a + ((b \% p) \ll 15) + ((c \% p) \ll 30)) \% p;
   return res:
| }}
```

### Number Theoretic Transform

```
vector<int> omega;
void Init() {
  omega.resize(kN + 1);
   long long x = fpow(kRoot, (Mod - 1) / kN);
   omega[0] = 1;
  for (int i = 1; i <= kN; ++i) {
  omega[i] = 1LL * omega[i - 1] * x % kMod;</pre>
void Transform(vector<int> &v, int n) {
```

```
BitReverse(v, n);
  for (int s = 2; s <= n; s <<= 1) {
    int z = s \gg 1;
    for (int i = 0; i < n; i += s) {
      for (int k = 0; k < z; ++k) {
        int x = 1LL * v[i + k + z] * omega[kN / s * k] % kMod;
        v[i + k + z] = (v[i + k] + kMod - x) % kMod;
        (v[i + k] += x) \% = kMod;
   }
 }
}
void InverseTransform(vector<int> &v, int n) {
  Transform(v, n);
  for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i]);
  const int kInv = fpow(n, kMod - 2);
  for (int i = 0; i < n; ++i) v[i] = 1LL * v[i] * inv % kMod;
```

# 6.3 NTT Prime List

```
Prime
             Root
                     Prime
                                   Root
7681
             17
                     167772161
12289
                     104857601
             11
                     985661441
40961
65537
                     998244353
786433
             10
                     1107296257
                                   10
                     2013265921
5767169
             3
                                   31
7340033
             3
                     2810183681
                                   11
23068673
                     2885681153
             3
                                   3
469762049
                     605028353
```

#### 6.4Formal Power Series

```
Poly Inverse(Poly f) {
  int n = f.size()
  Poly q(1, fpow(f[0], kMod - 2));
  for (int s = 2;; s <<= 1) {
     if (f.size() < s) f.resize(s);</pre>
    Poly fv(f.begin(), f.begin() + s);
    Poly fq(q.begin(), q.end());
    fv.resize(s + s);
    fq.resize(s + s);
    ntt::Transform(fv, s + s);
    ntt::Transform(fq, s + s);
    for (int i = 0; i < s + s; ++i) {
   fv[i] = 1LL * fv[i] * fq[i] % kMod * fq[i] % kMod;</pre>
    ntt::InverseTransform(fv, s + s);
    Poly res(s);
    for (int i = 0; i < s; ++i) {
      res[i] = kMod - fv[i];
       if (i < (s >> 1)) {
         int v = 2 * q[i] % kMod;
         (res[i] += v) >= kMod ? res[i] -= kMod : 0;
      }
    q = res;
    if (s >= n) break;
  q.resize(n);
  return q;
Poly Divide(const Poly &a, const Poly &b) {
  int n = a.size(), m = b.size(), k = 2;
  while (k < n - m + 1) k <<= 1;
  Poly ra(k), rb(k);
  for (int i = 0; i < min(n, k); ++i) ra[i] = a[n - 1 - i];
for (int i = 0; i < min(m, k); ++i) rb[i] = b[m - 1 - i];
  auto rbi = Inverse(rb);
  auto res = Multiply(rbi, ra);
  res.resize(n - m + 1);
  reverse(res.begin(), res.end());
  return res;
Poly Modulo(const Poly &a, const Poly &b) {
  if (a.size() < b.size()) return a;</pre>
  auto dv = Multiply(Divide(a, b), b);
  assert(dv.size() == a.size());
  for (int i = 0; i < dv.size(); ++i) {</pre>
    dv[i] = (a[i] + kMod - dv[i]) % kMod;
  while (!dv.empty() && dv.back() == 0) dv.pop_back();
  return dv;
Poly Derivative(const Poly &f) {
  int n = f.size();
```

vector<int> res(n - 1);

```
for (int i = 0; i < n - 1; ++i) { res[i] = 1LL * f[i + 1] * (i + 1) % kMod;
  return res;
Poly Integral(const Poly &f) {
  int n = f.size();
  vector<int> res(n + 1);
for (int i = 0; i < n; ++i) {
  res[i + 1] = 1LL * f[i] * fpow(i + 1, kMod - 2) % kMod;</pre>
  return res;
Poly Evaluate(const Poly &f, const vector<int> &x) {
  if (x.empty()) return Poly();
  int n = x.size();
  vector<Poly> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = {kMod - x[i], 1};
for (int i = n - 1; i > 0; --i) up[i] = Multiply(up[i * 2],
      up[i * 2 + 1]);
  vector<Poly> down(n * 2);
  down[1] = Modulo(f, up[1]);
  for (int i = 2; i < n * 2; ++i) down[i] = Modulo(down[i >>
      1], up[i]);
  vector<int> y(n);
  for (int i = 0; i < n; ++i) y[i] = down[i + n][0];
  return y;
Poly Interpolate(const vector<int> &x, const vector<int> &y) {
  int n = x.size();
  vector<Poly> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = \{kMod - x[i], 1\}; for (int i = n - 1; i > 0; --i) up[i] = Multiply(up[i * 2],
      up[i * 2 + 1]);
  vector<int> a = Evaluate(Derivative(up[1]), x);
for (int i = 0; i < n; ++i) {
    a[i] = 1LL * y[i] * fpow(a[i], kMod - 2) % kMod;</pre>
  vector<Poly> down(n * 2);
  for (int i = 0; i < n; ++i) down[i + n] = {a[i]};
for (int i = n - 1; i > 0; --i) {
     auto lhs = Multiply(down[i * 2], up[i * 2 + 1]);
auto rhs = Multiply(down[i * 2 + 1], up[i * 2]);
     assert(lhs.size() == rhs.size());
     down[i].resize(lhs.size());
     for (int j = 0; j < lhs.size(); ++j) {
  down[i][j] = (lhs[j] + rhs[j]) % kMod;</pre>
  return down[1];
Poly Log(Poly f) {
  int n = f.size();
  if (n == 1) return {0};
  auto d = Derivative(f);
  f.resize(n - 1);
  d = Multiply(d, Inverse(f));
  d.resize(n - 1);
  return Integral(d);
Poly Exp(Poly f) {
  int n = f.size();
  Poly q(1, 1);
  f[0] += 1;
  for (int s = 1; s < n; s <<= 1) {
   if (f.size() < s + s) f.resize(s + s);
   Poly g(f.begin(), f.begin() + s + s);</pre>
     Poly h(q.begin(), q.end());
     h.resize(s + s);
     h = Log(h);
     for (int i = 0; i < s + s; ++i) {
        g[i] = (g[i] + kMod - h[i]) % kMod;
     g = Multiply(g, q);
     g.resize(s + s);
     q = g;
  }
  assert(q.size() >= n);
  q.resize(n);
  return q;
Poly SquareRootImpl(Poly f) {
  if (f.empty()) return {0};
  int z = QuadraticResidue(f[0], kMod), n = f.size();
  constexpr int kInv2 = (kMod + 1) >> 1;
  if (z == -1) return {-1};
```

```
vector<int> q(1, z);
   for (int s = 1; s < n; s <<= 1) {
     if (f.size() < s + s) f.resize(s + s);</pre>
     vector<int> fq(q.begin(), q.end());
     fq.resize(s + s);
     vector<int> f2 = Multiply(fq, fq);
     f2.resize(s + s);
     for (int i = 0; i < s + s; ++i) {
       f2[i] = (f2[i] + kMod - f[i]) % kMod;
     f2 = Multiply(f2, Inverse(fq));
     f2.resize(s + s);

for (int i = 0; i < s + s; ++i) {

  fq[i] = (fq[i] + kMod - 1LL * f2[i] * kInv2 % kMod) %
     q = fq;
  }
   q.resize(n);
   return q;
Poly SquareRoot(Poly f) {
   int n = f.size(), m = 0;
   while (m < n \&\& f[m] == 0) m++;
   if (m == n) return vector<int>(n);
   if (m & 1) return {-1};
   auto s = SquareRootImpl(vector<int>(f.begin() + m, f.end()));
   if (s[0] == -1) return \{-1\};
   vector<int> res(n);
   for (int i = 0; i < s.size(); ++i) res[i + m / 2] = s[i];
   return res;
}
```

#### 6.5 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for  $\beta$  being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by  $Q_k$  the polynomial such that  $F(Q_k)=0\pmod{x^{2^k}}$ ,

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

#### 6.6 General Purpose Numbers

# 6.6.1 Bernoulli numbers

EGF of Bernoulli numbers is  $B(x) = \frac{x}{e^x - 1}$ .

#### 6.6.2 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$
 
$$S(n,1) = S(n,n) = 1$$
 
$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^{n}$$

#### 6.7 Fast Walsh-Hadamard Transform

- 1. XOR Convolution
  - $f(A) = (f(A_0) + f(A_1), f(A_0) f(A_1))$
  - $f^{-1}(A) = (f^{-1}(\frac{A_0 + A_1}{2}), f^{-1}(\frac{A_0 A_1}{2}))$
- 2. OR Convolution
  - $f(A) = (f(A_0), f(A_0) + f(A_1))$
  - $f^{-1}(A) = (f^{-1}(A_0), f^{-1}(A_1) f^{-1}(A_0))$
- 3. AND Convolution
  - $f(A) = (f(A_0) + f(A_1), f(A_1))$
  - $f^{-1}(A) = (f^{-1}(A_0) f^{-1}(A_1), f^{-1}(A_1))$

Description: maximize  $\mathbf{c}^T \mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Returns  $-\infty$  if

```
6.8 Simplex Algorithm
```

infeasible and  $\infty$  if unbounded.

```
const double eps = 1e-9;
const double inf = 1e+9;
int n. m:
vector<vector<double>> d;
vector<int> p, q;
void pivot(int r, int s) {
   double inv = 1.0 / d[r][s];
  for (int i = 0; i < m + 2; ++i) {
  for (int j = 0; j < n + 2; ++j) {</pre>
       if (i != r \&\& j != s) d[i][j] -= d[r][j] * d[i][s] * inv;
   for (int i = 0; i < m + 2; ++i) if (i != r) d[i][s] *= -inv;
   for (int j = 0; j < n + 2; ++j) if (j != s) d[r][j] *= +inv;
  d[r][s] = inv;
   swap(p[r], q[s]);
bool phase(int z) {
   int x = m + z;
  while (true) {
     int s = -1;
     for (int i = 0; i <= n; ++i) {
       if (!z \&\& q[i] == -1) continue
       if (s == -1 || d[x][i] < d[x][s]) s = i;
     if (d[x][s] > -eps) return true;
    int r = -1;
for (int i = 0; i < m; ++i) {
       if (d[i][s] < eps) continue;
if (r == -1 || d[i][n + 1] / d[i][s] < d[r][n + 1] / d[r</pre>
     ][s]) r = i;
     if (r == -1) return false;
     pivot(r, s);
  }
}
vector<double>> solve(const vector<vector<double>> &a, const
     vector<double> &b, const vector<double> &c) {
  m = b.size(), n = c.size();
  d = vector<vector<double>>(m + 2, vector<double>(n + 2));
  for (int i = 0; i < m; ++i) {
    for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
  p.resize(m), q.resize(n + 1);
  for (int i = 0; i < m; ++i) p[i] = n + i, d[i][n] = -1, d[i][
     n + 1] = b[i];
   for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
  q[n] = -1, d[m + 1][n] = 1;
   int r = 0;
  for (int i = 1; i < m; ++i) if (d[i][n + 1] < d[r][n + 1]) r
  if (d[r][n + 1] < -eps) {
     pivot(r, n);
     if (!phase(1) || d[m + 1][n + 1] < -eps) return vector<</pre>
      double>(n, -inf);
     for (int i = 0; i < m; ++i) if (p[i] == -1) {
       int s = min_element(d[i].begin(), d[i].end() - 1) - d[i].
      begin();
       pivot(i, s);
    }
  if (!phase(0)) return vector<double>(n, inf);
   vector<double> x(n);
   for (int i = 0; i < m; ++i) if (p[i] < n) \times [p[i]] = d[i][n +
     1];
  return x;
| }
6.9 Subset Convolution
Description: h(s) = \sum_{s' \subset s} f(s')g(s \setminus s')
```

```
if (s >> j & 1) {
    a[i][s] += a[i][s ^ (1 << j)];
    b[i][s] += b[i][s ^ (1 << j)];</pre>
    }
  }
}
vector<vector<int>>> c(n + 1, vector<int>(m));
for (int s = 0; s < m; ++s) {
  for (int i = 0; i <= n; ++i) {
    for (int j = 0; j \le i; ++j) c[i][s] += a[j][s] * b[i - j]
   ][s];
  }
for (int i = 0; i <= n; ++i) {
  for (int j = 0; j < n; ++j) {
    for (int s = 0; s < m; ++s) {
      if (s >> j & 1) c[i][s] -= c[i][s ^ (1 << j)];
}
vector<int> res(m);
for (int i = 0; i < m; ++i) res[i] = c[__builtin_popcount(i)</pre>
   ][i];
return res;
```

#### 6.9.1 Construction

Standard form: maximize  $\mathbf{c}^T\mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Dual LP: minimize  $\mathbf{b}^T\mathbf{y}$  subject to  $A^T\mathbf{y} \geq \mathbf{c}$  and  $\mathbf{y} \geq 0$ .  $\bar{\mathbf{x}}$  and  $\bar{\mathbf{y}}$  are optimal if and only if for all  $i \in [1, n]$ , either  $\bar{x}_i = 0$  or  $\sum_{j=1}^m A_{ji}\bar{x}_{jj} = c_i$  holds and for all  $i \in [1, m]$  either  $\bar{y}_i = 0$  or  $\sum_{j=1}^n A_{ij}\bar{x}_{jj} = b_j$  holds.

```
1. In case of minimization, let c_i' = -c_i
2. \sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j
```

- 3.  $\sum_{1 \le i \le n} A_{ji} x_i = b_j$ 
  - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$   $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

#### 6.10 Schreier-Sims Algorithm

```
namespace schreier {
vector<vector<int>>> bkts, binv;
vector<vector<int>> lk;
vector<int> operator*(const vector<int> &a, const vector<int> &
    b) {
  vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i) res[i] = b[a[i]];</pre>
  return res;
}
vector<int> inv(const vector<int> &a) {
  vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i) res[a[i]] = i;
  return res;
int filter(const vector<int> &g, bool add = true) {
  n = (int)bkts.size();
  vector < int > p = g;
  for (int i = 0; i < n; ++i) {
    assert(p[i] >= 0 \&\& p[i] < (int)lk[i].size());
    int res = lk[i][p[i]];
    if (res == -1) {
      if (add) {
        bkts[i].push_back(p);
        binv[i].push_back(inv(p));
        lk[i][p[i]] = (int)bkts[i].size() - 1;
      return i;
    p = p * binv[i][res];
  return -1;
bool inside(const vector<int> &g) { return filter(g, false) ==
void solve(const vector<vector<int>>> &gen, int _n) {
  bkts.clear(), bkts.resize(n);
  binv.clear(), binv.resize(n);
  lk.clear(), lk.resize(n);
  vector<int> iden(n);
  iota(iden.begin(), iden.end(), 0);
```

```
for (int i = 0; i < n; ++i) {
    k[i].resize(n, -1);</pre>
    bkts[i].push_back(iden);
    binv[i].push_back(iden);
    lk[i][i] = 0;
  for (int i = 0; i < (int)gen.size(); ++i) filter(gen[i]);</pre>
  queue<pair<pair<int, int>, pair<int, int>>> upd;
  for (int i = 0; i < n; ++i) {
    for (int j = i; j < n; ++j) {
      for (int k = 0; k < (int)bkts[i].size(); ++k) {</pre>
        for (int l = 0; l < (int)bkts[j].size(); ++l)</pre>
           upd.emplace(make_pair(i, k), make_pair(j, l));
    }
 }
  while (!upd.empty()) {
    auto a = upd.front().first;
    auto b = upd.front().second;
    upd.pop();
    int res = filter(bkts[a.first][a.second] * bkts[b.first][b.
     second]);
    if (res == -1) continue;
    pair<int, int> pr = make_pair(res, (int)bkts[res].size() -
     1);
    for (int i = 0; i < n; ++i) {
  for (int j = 0; j < (int)bkts[i].size(); ++j) {</pre>
         if (i <= res) upd.emplace(make_pair(i, j), pr);</pre>
         if (res <= i) upd.emplace(pr, make_pair(i, j));</pre>
   }
 }
long long size() {
  long long res = 1;
  for (int i = 0; i < n; ++i) res = res * bkts[i].size();</pre>
  return res;
```

# 6.11 Berlekamp-Massey Algorithm

```
template <int P>
vector<int> BerlekampMassey(vector<int> x) {
  vector<int> cur, ls;
 int lf = 0, ld = 0;
for (int i = 0; i < (int)x.size(); ++i) {
    int t = 0;
    for (int j = 0; j < (int)cur.size(); ++j)</pre>
      (t += 1LL * cur[j] * x[i - j - 1] % P) %= P;
    if (t == x[i]) continue;
    if (cur.empty()) {
      cur.resize(i + 1);
      lf = i, ld = (t + P - x[i]) % P;
      continue:
    int k = 1LL * fpow(ld, P - 2, P) * (t + P - x[i]) % P;
    vector<int> c(i - lf - 1);
    c.push_back(k);
    for (int j = 0; j < (int)ls.size(); ++j)
    c.push_back(1LL * k * (P - ls[j]) % P);</pre>
    if (c.size() < cur.size()) c.resize(cur.size());</pre>
    for (int j = 0; j < (int)cur.size(); ++j)
      c[j] = (c[j] + cur[j]) % P;
    if (i - lf + (int)ls.size() >= (int)cur.size()) {
      ls = cur, lf = i;
      ld = (t + P - x[i]) \% P;
    cur = c:
  return cur;
```

#### 6.12 Fast Linear Recurrence

```
res.resize(n + 1);
return res;
};
vector<int> p(n + 1), e(n + 1);
p[0] = e[1] = 1;
for (; k > 0; k >>= 1) {
   if (k & 1) p = Combine(p, e);
   e = Combine(e, e);
}
int res = 0;
for (int i = 0; i < n; ++i) (res += 1LL * p[i + 1] * s[i] % P
   ) %= P;
return res;
}</pre>
```

#### 6.13 Miller Rabin

```
| // n < 4759123141 chk = [2, 7, 61]
// n < 1122004669633 chk = [2, 13, 23, 1662803]
// n < 2^64 chk = [2, 325, 9375, 28178, 450775, 9780504,
      17952650227
 vector<long long> chk =
      {2,325,9375,28178,450775,9780504,1795265022};
 bool Check(long long a, long long u, long long n, int t) {
   a = fpow(a, u, n);
   if (a == 0 \mid \mid a == 1 \mid \mid a == n - 1) return true;
   for (int i = 0; i < t; ++i) {
     a = fmul(a, a, n);
     if (a == 1) return false;
     if (a == n - 1) return true;
   }
   return false:
bool IsPrime(long long n) {
   if (n < 2) return false;
   if (n % 2 == 0) return n == 2;
   long long u = n - 1; int t = 0;
   for (; !(u & 1); u >>= 1, ++t);
for (long long i : chk) {
     if (!Check(i, u, n, t)) return false;
   return true;
}
```

#### 6.14 Pollard's Rho

```
map<long long, int> cnt;
 void PollardRho(long long n) {
   if (n == 1) return;
   if (prime(n)) return ++cnt[n], void();
   if (n % 2 == 0) return PollardRho(n / 2), ++cnt[2], void(); long long x = 2, y = 2, d = 1, p = 1;
   auto f = [\&](auto x, auto n, int p) { return <math>(fmul(x, x, n) +
       p) % n; }
   while (true) {
     if (d != n && d != 1) {
       PollardRho(n / d);
       PollardRho(d);
       return;
     if (d == n) ++p;
     x = f(x, n, p); y = f(f(y, n, p), n, p);
     d = \_gcd(abs(x - y), n);
}
```

# 6.15 Meissel-Lehmer Algorithm

```
int64_t PrimeCount(int64_t n) {
   if (n <= 1) return 0;
   const int v = sqrt(n);
   vector<int> smalls(v + 1);
   for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;
   int s = (v + 1) / 2;
   vector<int> roughs(s);
   for (int i = 0; i < s; ++i) roughs[i] = 2 * i + 1;
   vector<int64_t> larges(s);
   for (int i = 0; i < s; ++i) larges[i] = (n / (2 * i + 1) + 1) / 2;
   vector<bool> skip(v + 1);
   int pc = 0;
   for (int p = 3; p <= v; ++p) {
      if (smalls[p] > smalls[p - 1]) {
       int q = p * p;
       pc++;
      if (1LL * q * q > n) break;
   }
}
```

```
skip[p] = true:
                   for (int i = q; i <= v; i += 2 * p) skip[i] = true;</pre>
                    int ns = 0;
                    for (int k = 0; k < s; ++k) {
                            int i = roughs[k];
                            if (skip[i]) continue;
                            int64_t d = 1LL * i * p;
                            larges[ns] = larges[k] - (d \ll v ? larges[smalls[d] - v ? ] larges[smalls[d] - v ? larges[smalls[d] - v ? larges[smalls[d] - v ] larges[
             pc] : smalls[n / d]) + pc;
                            roughs[ns++] = i;
                  }
                   s = ns:
                  for (int i = j * p, e = min(i + p, v + 1); i < e; ++i)
             smalls[i] -= c;
        }
for (int k = 1; k < s; ++k) {
  const int64_t m = n / roughs[k];</pre>
          int64_t = larges[k] - (pc + k - 1);
         for (int l = 1; l < k; ++l) {
                 int p = roughs[l];
if (1LL * p * p > m) break;
                  s = smalls[m / p] - (pc + l - 1);
         larges[0] -= s;
}
 return larges[0];
```

### 6.16 Discrete Logarithm

Description: to find x such that  $x^a \equiv b \pmod{p}$ , let g be the primitive root of p, find k such that  $g^k \equiv b \pmod{p}$  and x can be found by  $g^d$  where  $ad \equiv k \pmod{p-1}$ .

```
int DiscreteLog(int s, int x, int y, int m) {
  constexpr int kStep = 32000;
  unordered_map<int, int> p;
  int b = 1:
  for (int i = 0; i < kStep; ++i) {
    p[y] = i;
y = 1LL * y * x % m;
b = 1LL * b * x % m;
  for (int i = 0; i < m + 10; i += kStep) {
   s = 1LL * s * b % m;</pre>
    if (p.find(s) != p.end()) return i + kStep - p[s];
  return -1:
int DiscreteLog(int x, int y, int m) {
  if (m == 1) return 0;
  int s = 1;
  for (int i = 0; i < 100; ++i) {
    if (s == y) return i;
s = 1LL * s * x % m;
  if (s == y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
```

#### 6.17 Quadratic Residue

```
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
   a %= m;
    if (a == 0) return 0;
    const int r = __builtin_ctz(a);
    if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
    if (a \& m \& 2) s = -s;
    swap(a, m);
  return s:
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
  const int jc = Jacobi(a, p);
  if (jc == 0 || jc == -1) return jc;
  int b, d;
 for (; ; ) {
    b = rand() \% p;
```

```
d = (1LL * b * b + p - a) % p;
if (Jacobi(d, p) == -1) break;
}
int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
for (int e = (p + 1) >> 1; e; e >>= 1) {
   if (e & 1) {
      tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 % p)) % p;
      g0 = tmp;
   }
tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) % p;
f1 = (2LL * f0 * f1) % p;
f0 = tmp;
}
return g0;
}
```

#### 6.18 Gaussian Elimination

```
| double Gauss(vector<vector<double>> &d) {
   int n = d.size(), m = d[0].size();
double det = 1;
   for (int i = 0; i < m; ++i) {
     int p = -1;
     for (int j = i; j < n; ++j) {
       if (fabs(d[j][i]) < kEps) continue;</pre>
       if (p == -1 \mid | fabs(d[j][i]) > fabs(d[p][i])) p = j;
     if (p == -1) continue;
     if (p != i) det *= -1;
     for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
     for (int j = 0; j < n; ++j) {
       if (i == j) continue;
       double z = d[j][i] / d[i][i];
       for (int k = 0; k < m; ++k) d[j][k] -= z * d[i][k];
   for (int i = 0; i < n; ++i) det *= d[i][i];</pre>
   return det;
13
```

#### 6.19 Characteristic Polynomial

```
vector<vector<int>> Hessenberg(const vector<vector<int>> &A) {
  int N = A.size();
  vector<vector<int>> H = A;
  for (int i = 0; i < N - 2; ++i) {
    if (!H[i + 1][i]) {
      for (int j = i + 2; j < N; ++j) {
        if (H[j][i]) {
           for (int k = i; k < N; ++k) swap(H[i + 1][k], H[j][k]
     ]);
          for (int k = 0; k < N; ++k) swap(H[k][i + 1], H[k][j
     ]);
          break;
      }
    }
    if (!H[i + 1][i]) continue;
    int val = fpow(H[i + 1][i], kP - 2);
    for (int j = i + 2; j < N; ++j) {
  int coef = 1LL * val * H[j][i] % kP;</pre>
     for (int k = i; k < N; ++k) H[j][k] = (H[j][k] + 1LL * H[i + 1][k] * (kP - coef)) % kP;
      for (int k = 0; k < N; ++k) H[k][i + 1] = (H[k][i + 1] +
     1LL * H[k][j] * coef) % kP;
    }
  }
  return H;
}
vector<int> CharacteristicPoly(const vector<vector<int>> &A) {
  int N = A.size();
  auto H = Hessenberg(A);
  for (int i = 0; i < N; ++i) {
    for (int j = 0; j < N; ++j) H[i][j] = kP - H[i][j];
  vector<vector<int>>> P(N + 1, vector<int>(N + 1));
  P[0][0] = 1;
  for (int i = 1; i <= N; ++i) {
    P[i][0] = 0;
    for (int j = 1; j \le i; ++j) P[i][j] = P[i - 1][j - 1];
    int val = 1;
    for (int j = i - 1; j >= 0; --j) {
      int coef = 1LL * val * H[j][i - 1] % kP;
      for (int k = 0; k \le j; ++k) P[i][k] = (P[i][k] + 1LL * P
     [j][k] * coef) % kP;
```

```
if (j) val = 1LL * val * (kP - H[j][j - 1]) % kP;
}
if (N & 1) {
    for (int i = 0; i <= N; ++i) P[N][i] = kP - P[N][i];
}
return P[N];
}</pre>
```

# 6.20 $\mu$ function

```
int mu[kC], dv[kC];
vector<int> prime;
void Sieve() {
    mu[1] = dv[1] = 1;
    for (int i = 2; i < kC; ++i) {
        if (!dv[i]) {
            dv[i] = i, mu[i] = -1;
            prime.push_back(i);
        }
    for (int j = 0; i * prime[j] < kC; ++j) {
            dv[i * prime[j]] = prime[j];
            mu[i * prime[j]] = -mu[i];
            if (i % prime[j]] == 0) {
                mu[i * prime[j]] = 0;
                 break;
            }
        }
        }
    }
}</pre>
```

### 6.21 Partition Function

# 6.22 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_i = \lfloor \frac{n}{\lfloor \frac{n}{T_{i-1}+1} \rfloor} \rfloor
```

#### 6.23 De Bruijn Sequence

```
int res[kN], aux[kN], a[kN], sz;
 void Rec(int t, int p, int n, int k) {
  if (t > n) {
     if (n \% p == 0)
       for (int i = 1; i <= p; ++i) res[sz++] = aux[i];</pre>
  } else {
     aux[t] = aux[t - p];
     Rec(t + 1, p, n, k);
     for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t]) Rec(t +
      1, t, n, k);
  }
 int DeBruijn(int k, int n) {
  // return cyclic string of length k^n such that every string
     of length n using k character appears as a substring.
   if (k == 1) return res[0] = 0, 1;
  fill(aux, aux + k * n, 0);
   return sz = 0, Rec(1, 1, n, k), sz;
j }
```

# 6.24 Extended GCD

```
template <typename T> tuple<T, T, T> extgcd(T a, T b) {
   if (!b) return make_tuple(a, 1, 0);
   T d, x, y;
   tie(d, x, y) = extgcd(b, a % b);
   return make_tuple(d, y, x - (a / b) * y);
}
```

# 6.25 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \\ &= \begin{cases} \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ -2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

#### 6.26 Chinese Remainder Theorem

```
long long crt(vector<int> mod, vector<int> a) {
   long long mult = mod[0];
   int n = (int)mod.size();
   long long res = a[0];
   for (int i = 1; i < n; ++i) {
      long long d, x, y;
      tie(d, x, y) = extgcd(mult, mod[i] * 1ll);
      if ((a[i] - res) % d) return -1;
      long long new_mult = mult / __gcd(mult, 1ll * mod[i]) * mod
      [i];
      res += x * ((a[i] - res) / d) % new_mult * mult % new_mult;
      mult = new_mult;
      ((res %= mult) += mult) %= mult;
   }
   return res;
}</pre>
```

#### 6.27 Theorem

#### 6.27.1 Kirchhoff's Theorem

Denote L be a  $n\times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii}=d(i),\,L_{ij}=-c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

#### 6.27.2 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

#### 6.27.3 Cayley's Formula

- Given a degree sequence  $d_1, d_2, \dots, d_n$  for each labeled vertices, there
- are  $\frac{(n-2)!^k}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.

   Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### 6.27.4 Erdős–Gallai Theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all  $1 \le k \le n$ .

#### 6.28 Primes

```
\begin{array}{l} 97,101,131,487,593,877,1087,1187,1487,1787,3187,12721,\\ 13331,14341,75577,123457,222557,556679,999983,\\ 1097774749,1076767633,100102021,999997771,\\ 1001010013,1000512343,987654361,999991231,\\ 999888733,98789101,987777733,999991921,1000000007,\\ 1000000087,1000000123,1010101333,1010102101,\\ 100000000039,10000000000037,2305843009213693951,\\ 4611686018427387847,9223372036854775783,\\ 18446744073709551557\end{array}
```

# 7 Dynamic Programming

# 7.1 Dynamic Convex Hull

```
struct Line {
  mutable int64_t a, b, p;
  bool operator<(const Line &rhs) const { return a < rhs.a; }</pre>
  bool operator<(int64_t x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>>> {
  static const int64_t kInf = 1e18;
  int64_t Div(int64_t a, int64_t b) { return a / b - ((a \land b) <
      0 && a % b); }
  bool Isect(iterator x, iterator y) {
    if (y == end()) { x->p = kInf; return false; }
    if (x->a == y->a) x->p = x->b > y->b? kInf : -kInf;
    else x->p = Div(y->b - x->b, x->a - y->a);
    return x \rightarrow p >= y \rightarrow p;
  void Insert(int64_t a, int64_t b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (Isect(y, z)) z = erase(z);
    if (x != begin() \&\& Isect(--x, y)) Isect(x, y = erase(y));
    while ((y = x) != begin() \&\& (--x)->p >= y->p) Isect(x,
     erase(y));
  int64_t Query(int64_t x) {
    auto 1 = *lower_bound(x);
     return l.a * x + l.b;
|};
```

# 7.2 1D/1D Convex Optimization

```
struct segment {
   segment(int a, int b, int c): i(a), l(b), r(c) {}
 inline long long f(int l, int r) { return dp[l] + w(l + 1, r);
 void solve() {
   dp[0] = 011;
   deque<segment> deq; deq.push_back(segment(0, 1, n));
   for (int i = 1; i \le n; ++i) {
     dp[i] = f(deq.front().i, i);
     while (deq.size() && deq.front().r < i + 1) deq.pop_front()</pre>
     deq.front().l = i + 1;
     segment seg = segment(i, i + 1, n);
while (deq.size() && f(i, deq.back().l) < f(deq.back().i,</pre>
      deq.back().1)) deq.pop_back();
     if (deq.size()) {
       int d = 1048576, c = deq.back().1;
       while (d \gg 1) if (c + d \ll deq.back().r) {
         if (f(i, c + d) > f(deq.back().i, c + d)) c += d;
       deq.back().r = c; seg.l = c + 1;
     if (seg.l <= n) deq.push_back(seg);</pre>
  }
| }
```

#### 7.3 Condition

#### 7.3.1 Totally Monotone (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

#### 7.3.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j', \ B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', \ B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

#### 7.3.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \geq B[i][j+1] + B[i+1][j] then H_{i,j-1} \leq H_{i,j} \leq H_{i+1,j}
```

# 8 Geometry

#### 8.1 Basic

```
bool same(double a, double b) { return abs(a - b) < eps; }</pre>
struct P {
  double x,
  P() : x(0), y(0) {}
  P(double x, double y) : x(x), y(y) {}
  P operator + (P b) { return P(x + b.x, y + b.y); }
P operator - (P b) { return P(x - b.x, y - b.y); }
P operator * (double b) { return P(x * b, y * b); }
P operator / (double b) { return P(x * b, y * b); }
  double operator * (P b) { return x * b.x + y * b.y; }
double operator ^ (P b) { return x * b.y - y * b.x; }
  double abs() { return hypot(x, y); }
  P unit() { return *this / abs(); }
  P spin(double o) {
     double c = cos(o), s = sin(o);
     return P(c * x - s * y, s * x + c * y);
  double angle() { return atan2(y, x); }
struct L {
   // ax + by + c = 0
  double a, b, c, o;
  P pa, pb;
  L(): a(0), b(0), c(0), o(0), pa(), pb() {}
  L(P pa, P pb) : a(pa.y - pb.y), b(pb.x - pa.x), c(pa \land pb), o
     (atan2(-a, b)), pa(pa), pb(pb) {}
    project(P p) { return pa + (pb - pa).unit() * ((pb - pa) *
     (p - pa) / (pb - pa).abs()); }
  P reflect(P p) { return p + (project(p) - p) * 2; }
  double get_ratio(P p) { return (p - pa) * (pb - pa) / ((pb -
     pa).abs() * (pb - pa).abs()); }
bool SegmentIntersect(P p1, P p2, P p3, P p4) {
  if (max(p1.x, p2.x) < min(p3.x, p4.x) | | max(p3.x, p4.x) <
     min(p1.x, p2.x)) return false
  if (max(p1.y, p2.y) < min(p3.y, p4.y) || max(p3.y, p4.y) <
     min(p1.y, p2.y)) return false
  return sign((p3 - p1) ^ (p4 - p1)) * sign((p3 - p2) ^ (p4 -
     p2)) <= 0 &&
       sign((p1 - p3) \land (p2 - p3)) * sign((p1 - p4) \land (p2 - p4))
       <= 0:
bool parallel(L x, L y) { return same(x.a * y.b, x.b * y.a); }
P Intersect(L x, L y) { return P(-x.b * y.c + x.c * y.b, x.a *
     y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a); }
```

# 8.2 KD Tree

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn], yl[maxn], yr[
     maxn];
point p[maxn];
int build(int 1, int r, int dep = 0) {
   if (l == r) return -1;
  function<bool(const point &, const point &)> f = [dep](const
     point &a, const point &b) {
if (dep & 1) return a.x < b.x;
else return a.y < b.y;</pre>
  int m = (l + r) >> 1;
nth_element(p + l, p + m, p + r, f);
  xl[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
lc[m] = build(l, m, dep + 1);
  if (~lc[m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
     xr[m] = max(xr[m], xr[lc[m]]);
    yl[m] = min(yl[m], yl[lc[m]]);
    yr[m] = max(yr[m], yr[lc[m]]);
```

double side(P &a, P &b, P &p) { return (b - a) ^ (p - a); }

struct Tri:

struct Edge {

```
Tri *tri;
  rc[m] = build(m + 1, r, dep + 1);
                                                                               int side;
                                                                               Edge() : tri(0), side(0) {}
  if (~rc[m]) {
    xl[m] = min(xl[m], xl[rc[m]]);
                                                                              Edge(Tri *_tri, int _side) : tri(_tri), side(_side) {}
    xr[m] = max(xr[m], xr[rc[m]]);
                                                                             };
    yl[m] = min(yl[m], yl[rc[m]]);
                                                                             struct Tri {
                                                                              P p[3];
    yr[m] = max(yr[m], yr[rc[m]]);
                                                                              Edge edge[3];
Tri *ch[3];
  return m;
                                                                              Tri() {}
bool bound(const point &q, int o, long long d) {
                                                                              Tri(P p0, P p1, P p2) {
p[0] = p0; p[1] = p1; p[2] = p2;
  double ds = sqrt(d + 1.0);
  if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
                                                                               ch[0] = ch[1] = ch[2] = 0;
  q.y < yl[o] - ds || q.y > yr[o] + ds) return false;
return true;
                                                                               bool has_ch() { return ch[0] != 0; }
                                                                               int num_ch() {
long long dist(const point &a, const point &b) {
                                                                               return ch[0] == 0 ? 0 : ch[1] == 0 ? 1 : ch[2] == 0 ? 2 : 3;
  return (a.x - b.x) * 111 * (a.x - b.x) +
      (a.y - b.y) * 111 * (a.y - b.y);
                                                                               bool contains(P &q) {
                                                                                for (int i = 0; i < 3; ++i)
if (side(p[i], p[(i + 1) % 3], q) < -eps) return false;</pre>
void dfs(const point &q, long long &d, int o, int dep = 0) {
  if (!bound(q, o, d)) return;
                                                                                return true
  long long cd = dist(p[o], q);
  if (cd != 0) d = min(d, cd);
                                                                             } pool[maxn * 10], *tris;
  if ((dep & 1) && q.x < p[o].x || !(dep & 1) && <math>q.y < p[o].y)
                                                                             void edge(Edge a, Edge b) {
                                                                              if (a.tri) a.tri->edge[a.side] = b;
    if (~lc[o]) dfs(q, d, lc[o], dep + 1);
                                                                              if (b.tri) b.tri->edge[b.side] = a;
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  } else {
                                                                             struct Trig {
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
                                                                              Trig() {
                                                                               the_root = new (tris++) Tri(P(-inf, -inf), P(inf * 2, -inf),
   P(-inf, inf * 2));
  }
                                                                               } // all p should in
                                                                              Tri *find(P p) { return find(the_root, p); }
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
                                                                               void add_point(P &p) { add_point(find(the_root, p), p); }
  root = build(0, v.size());
                                                                              Tri *the_root;
                                                                              static Tri *find(Tri *root, P &p) {
long long nearest(const point &q) {
                                                                               while (true) {
  long long res = 1e18;
                                                                                 if (!root->has_ch()) return root;
  dfs(q, res, root);
return res;
                                                                                 for (int i = 0; i < 3 \&\& root->ch[i]; ++i)
                                                                                  if (root->ch[i]->contains(p)) {
                                                                                   root = root->ch[i];
                                                                                   break:
8.3 Delaunay Triangulation
                                                                                  }
                                                                               }
/* Delaunay Triangulation:
                                                                               assert(false); // "point not found"
  Given a sets of points on 2D plane, find a
  triangulation such that no points will strictly
                                                                               void add_point(Tri *root, P &p) {
  inside circumcircle of any triangle.
                                                                               Tri *tab, *tbc, *tca;
find : return a triangle contain given point
                                                                               tab = new (tris++) Tri(root->p[0], root->p[1], p);
tbc = new (tris++) Tri(root->p[1], root->p[2], p);
add_point : add a point into triangulation
A Triangle is in triangulation iff. its has_chd is 0.
                                                                                tca = new (tris++) Tri(root->p[2], root->p[0], p);
Region of triangle u: iterate each u.edge[i].tri, each points are u.p[(i+1)%3], u.p[(i+2)%3]
                                                                                edge(Edge(tab, 0), Edge(tbc, 1));
                                                                                edge(Edge(tbc, 0), Edge(tca, 1));
calculation involves O(IVI^6) */
                                                                                edge(Edge(tca, 0), Edge(tab, 1));
const double inf = 1e9;
                                                                                edge(Edge(tab, 2), root->edge[2]);
double eps = 1e-6; // 0 when integer
                                                                                edge(Edge(tbc, 2), root->edge[0]);
// return p4 is in circumcircle of tri(p1,p2,p3)
                                                                                edge(Edge(tca, 2), root->edge[1]);
root->ch[0] = tab; root->ch[1] = tbc; root->ch[2] = tca;
bool in_cc(P &p1, P &p2, P &p3, P &p4) {
  int o1 = (abs(p1.x) >= inf * 0.99) || abs(p1.y) >= inf * 0.99);

                                                                                flip(tab, 2); flip(tbc, 2); flip(tca, 2);
 int o2 = (abs(p2.x) >= inf * 0.99 | | abs(p2.y) >= inf * 0.99);
 int o3 = (abs(p3.x) >= inf * 0.99 | | abs(p3.y) >= inf * 0.99);
                                                                               void flip(Tri *tri, int pi) {
 int rtrue = 01 + 02 + 03;
                                                                               Tri *trj = tri->edge[pi].tri;
int pj = tri->edge[pi].side;
 int rfalse = abs(p4.x) >= inf * 0.99 | | abs(p4.y) >= inf *
     0.99;
                                                                                if (!trj) return;
 if (rtrue == 3) return true;
                                                                                if (!in_cc(tri->p[0], tri->p[1], tri->p[2], trj->p[pj]))
 if (rtrue) {
                                                                                   return
  P in(0, 0), out(0, 0);
                                                                                /* flip edge between tri,trj */
  if (o1) out = out + p1; else in = in + p1;
                                                                                Tri *trk = new (tris++) Tri(tri->p[(pi + 1) % 3], trj->p[pj],
  if (o2) out = out + p2; else in = in + p2;
                                                                                    tri->p[pi]);
  if (o3) out = out + p3; else in = in + p3;
return (p4 - in) * (out - in) > 0;
                                                                                Tri *trl = new (tris++) Tri(trj->p[(pj + 1) % 3], tri->p[pi],
                                                                                    trj->p[pj]);
                                                                                edge(Edge(trk, 0), Edge(trl, 0));
 if (rfalse) return false;
                                                                                edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
                                                                               edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
 double u11 = p1.x - p4.x, u12 = p1.y - p4.y;
double u21 = p2.x - p4.x, u22 = p2.y - p4.y;
double u31 = p3.x - p4.x, u32 = p3.y - p4.y;
                                                                                edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
 double u13 = sq(p1.x) - sq(p4.x) + sq(p1.y) - sq(p4.y);
                                                                                tri - ch[0] = trk; tri - ch[1] = trl; tri - ch[2] = 0;
 double u23 = sq(p2.x) - sq(p4.x) + sq(p2.y) - sq(p4.y);
                                                                                trj->ch[0] = trk; trj->ch[1] = trl; trj->ch[2] = 0;
 double u33 = sq(p3.x) - sq(p4.x) + sq(p3.y) - sq(p4.y);
double det = -u13 * u22 * u31 + u12 * u23 * u31 + u13 * u21 *
u32 - u11 * u23 * u32 - u12 * u21 * u33 + u11 * u22 * u33;
                                                                               flip(trk, 1); flip(trk, 2);
flip(trl, 1); flip(trl, 2);
 return det > eps;
```

vector<Tri \*> triang;

void go(Tri \*now) {

set<Tri \*> vst;

```
if (vst.find(now) != vst.end()) return;
  vst.insert(now);
  if (!now->has_ch()) {
    triang.push_back(now);
    return;
  }
  for (int i = 0; i < now->num_ch(); ++i) go(now->ch[i]);
  }
  void build(int n, P *ps) {
    tris = pool;
    random_shuffle(ps, ps + n);
    Trig tri;
  for (int i = 0; i < n; ++i) tri.add_point(ps[i]);
    go(tri.the_root);
}</pre>
```

### 8.4 Voronoi Diagram

```
| int gid(P &p) {
   auto it = ptoid.find(p);
   if (it == ptoid.end()) return -1;
   return it->second;
 L make_line(P p, L l) {
   P d = 1.pb - 1.pa; d = d.spin(pi / 2);
   P m = (1.pa + 1.pb) / 2;
   l = L(m, m + d);
if (((l.pb - l.pa) ^ (p - l.pa)) < 0) l = L(m + d, m);
   return 1;
 double calc_ans(int i) {
   vector<P> ps = HPI(ls[i]);
   double rt = 0;
   for (int i = 0; i < (int)ps.size(); ++i) {
     rt += (ps[i] ^ ps[(i + 1) % ps.size()]);
   return abs(rt) / 2;
 void solve() {
   for (int i = 0; i < n; ++i) ops[i] = ps[i], ptoid[ops[i]] = i
   random_shuffle(ps, ps + n);
   build(n, ps);
   for (auto *t : triang) {
     int z[3] = \{gid(t->p[0]), gid(t->p[1]), gid(t->p[2])\};
for (int i = 0; i < 3; ++i) for (int j = 0; j < 3; ++j) if
      (i != j \&\& z[i] != -1 \&\& z[j] != -1) {
       L l(t->p[i], t->p[j]);
       ls[z[i]].push_back(make_line(t->p[i], l));
     }
   }
   vector<P> tb = convex(vector<P>(ps, ps + n));
   for (auto &p : tb) isinf[gid(p)] = true;
   for (int i = 0; i < n; ++i) {
     if (isinf[i]) cout << -1 << '\n';</pre>
     else cout << fixed << setprecision(12) << calc_ans(i) << '\</pre>
   }
| }
```

# 8.5 Sector Area

```
// calc area of sector which include a, b
|double SectorArea(P a, P b, double r) {
| double o = atan2(a.y, a.x) - atan2(b.y, b.x);
| while (o <= 0) o += 2 * pi;
| while (o >= 2 * pi) o -= 2 * pi;
| o = min(o, 2 * pi - o);
| return r * r * o / 2;
|}
```

# 8.6 Half Plane Intersection

```
bool jizz(L 11,L 12,L 13){
   P p=Intersect(12,13);
   return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;
}

bool cmp(const L &a,const L &b){
   return same(a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;
}

// availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
vector<P> HPI(vector<L> &ls){
   sort(ls.begin(),ls.end(),cmp);
   vector<L> pls(1,ls[0]);
```

```
for(int i=0;i<(int)ls.size();++i)if(!same(ls[i].o,pls.back().</pre>
     o))pls.push_back(ls[i]);
   deque<int> dq; dq.push_back(0); dq.push_back(1);
 #define meow(a,b,c) while(dq.size()>1u && jizz(pls[a],pls[b],
     pls[c]))
   for(int i=2;i<(int)pls.size();++i){</pre>
     meow(i,dq.back(),dq[dq.size()-2])dq.pop_back();
     meow(i,dq[0],dq[1])dq.pop_front();
     dq.push_back(i);
  meow(dq.front(),dq.back(),dq[dq.size()-2])dq.pop_back();
  meow(dq.back(),dq[0],dq[1])dq.pop_front();
   if(dq.size()<3u)return vector<P>(); // no solution or
      solution is not a convex
   vector<P> rt;
   for(int i=0;i<(int)dq.size();++i)rt.push_back(Intersect(pls[</pre>
      dq[i]],pls[dq[(i+1)%dq.size()]]));
1}
```

#### 8.7 Rotating Sweep Line

```
void rotatingSweepLine(vector<pair<int,int>> &ps){
  int n=int(ps.size());
   vector<int> id(n),pos(n);
   vector<pair<int,int>> line(n*(n-1)/2);
   int m=-1:
   for(int i=0;i<n;++i)for(int j=i+1;j<n;++j)line[++m]=make_pair</pre>
     (i,j); ++m;
   sort(line.begin(),line.end(),[&](const pair<int,int> &a,const
       pair<int,int> &b)->bool{
     if(ps[a.first].first==ps[a.second].first)return 0;
     if(ps[b.first].first==ps[b.second].first)return 1;
     return (double)(ps[a.first].second-ps[a.second].second)/(ps
      [a.first].first-ps[a.second].first) < (double)(ps[b.first</pre>
      ].second-ps[b.second].second)/(ps[b.first].first-ps[b.
      second].first);
   });
   for(int i=0;i<n;++i)id[i]=i;</pre>
   sort(id.begin(),id.end(),[&](const int &a,const int &b){
     return ps[a]<ps[b]; });</pre>
   for(int i=0;i<n;++i)pos[id[i]]=i;</pre>
   for(int i=0;i<m;++i){</pre>
     auto l=line[i];
     tie(pos[l.first],pos[l.second],id[pos[l.first]],id[pos[l.
     second]])=make_tuple(pos[l.second],pos[l.first],l.second,l
      .first):
}
```

#### 8.8 Triangle Center

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
   Point res;
   double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
   double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
   double ax = (a.x + b.x) / 2;
   double ay = (a.y + b.y) / 2;
   double bx = (c.x + b.x) / 2;
   double by = (c.y + b.y) / 2;
   double r1 = (sin(a2) * (ax - bx) + cos(a2) * (by - ay)) / (
    sin(a1) * cos(a2) - sin(a2) * cos(a1));
   return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
}
Point TriangleMassCenter(Point a, Point b, Point c) {
   return (a + b + c) / 3.0;
}
Point TriangleOrthoCenter(Point a, Point b, Point c) {
   return TriangleMassCenter(a, b, c) * 3.0 -
      TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
   Point res;
   double la = len(b - c);
   double lb = len(a - c);
   double lc = len(a - b);

res.x = (la * a.x + lb * b.x + lc * c.x) / (la + lb + lc);
   res.y = (la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
   return res;
13
```

# 8.9 Polygon Center

```
| Point BaryCenter(vector<Point> &p, int n) {
    Point res(0, 0);
    double s = 0.0, t;
    for (int i = 1; i < p.size() - 1; i++) {
        t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
        s += t;
        res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
        res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
    }
    res.x /= (3 * s);
    res.y /= (3 * s);
    return res;
}
```

# 8.10 Maximum Triangle

```
double ConvexHullMaxTriangleArea(Point p[], int res[], int
     chnum) {
  double area = 0, tmp;
  res[chnum] = res[0];
  for (int i = 0, j = 1, k = 2; i < chnum; i++) {
  while (fabs(Cross(p[res[j]] - p[res[i]], p[res[(k + 1) %</pre>
     chnum]] - p[res[i]])) > fabs(Cross(p[res[j]] - p[res[i]],
     p[res[k]] - p[res[i]])) k = (k + 1) % chnum;
     tmp = fabs(Cross(p[res[j]] - p[res[i]], p[res[k]] - p[res[i
     ]]));
    if (tmp > area) area = tmp;
    while (fabs(Cross(p[res[(j + 1) % chnum]] - p[res[i]], p[
     res[k]] - p[res[i]])) > fabs(Cross(p[res[j]] - p[res[i]],
    p[res[k]] - p[res[i]]))) j = (j + 1) % chnum;
tmp = fabs(Cross(p[res[j]] - p[res[i]], p[res[k]] - p[res[i
     ]]));
    if (tmp > area) area = tmp;
  return area / 2;
```

# 8.11 Point in Polygon

```
int pip(vector<P> ps, P p) {
  int c = 0;
  for (int i = 0; i < ps.size(); ++i) {
    int a = i, b = (i + 1) % ps.size();
    L l(ps[a], ps[b]);
    P q = l.project(p);
    if ((p - q).abs() < eps && l.inside(q)) return 1;
    if (same(ps[a].y, ps[b].y) && same(ps[a].y, p.y)) continue;
    if (ps[a].y > ps[b].y) swap(a, b);
    if (ps[a].y <= p.y && p.y < ps[b].y && p.x <= ps[a].x + (ps
        [b].x - ps[a].x) / (ps[b].y - ps[a].y) * (p.y - ps[a].y))
    ++c;
    }
    return (c & 1) * 2;
}</pre>
```

#### 8.12 Circle

```
struct C {
 P c;
double r;
  C(P \ c = P(0, 0), double \ r = 0) : c(c), r(r) \{\}
vector<P> Intersect(C a, C b) {
 if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  vector<P> p;
  if (same(a.r + b.r, d)) p.push_back(a.c + (b.c - a.c).unit()
     * a.r);
  else if (a.r + b.r > d \&\& d + a.r >= b.r) {
    double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d)
    P i = (b.c - a.c).unit();
    p.push_back(a.c + i.spin(o) * a.r);
    p.push_back(a.c + i.spin(-o) * a.r);
  return p;
double IntersectArea(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  if (d \ge a.r + b.r - eps) return 0;
  if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
  double p = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
 double q = a\cos((sq(b.r) + sq(d) - sq(a.r)) / (2 * b.r * d));

return p * sq(a.r) + q * sq(b.r) - a.r * d * sin(p);
```

```
// remove second level if to get points for line (defalut:
     segment)
vector<P> CircleCrossLine(P a, P b, P o, double r) {
  double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y);
  double C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
  vector<P> t;
  if (d >= -eps) {
    d = max(0., d);
    double i = (-B - sqrt(d)) / (2 * A);
double j = (-B + sqrt(d)) / (2 * A);
    if (i - 1.0 \le eps \&\& i \ge -eps) t.emplace_back(a.x + i * x)
      a.y + i * y);
    if (j - 1.0 \le eps \&\& j \ge -eps) t.emplace_back(a.x + j * x)
     , a.y + j * y);
  return t;
}
// calc area intersect by circle with radius r and triangle OAB
double AreaOfCircleTriangle(P a, P b, double r) {
  bool ina = a.abs() < r, inb = b.abs() < r;
  auto p = CircleCrossLine(a, b, P(0, 0), r);
  if (ina) {
    if (inb) return abs(a ^ b) / 2;
    return SectorArea(b, p[0], r) + abs(a ^p[0]) / 2;
  if (inb) return SectorArea(p[0], a, r) + abs(p[0] \land b) / 2;
  if (p.size() == 2u) return SectorArea(a, p[0], r) +
    SectorArea(p[1], b, r) + abs(p[0] ^ p[1]) / 2;
  else return SectorArea(a, b, r);
// for any triangle
double AreaOfCircleTriangle(vector<P> ps, double r) {
  double ans = 0;
  for (int i = 0; i < 3; ++i) {
    int j = (i + 1) \% 3;
    double o = atan2(ps[i].y, ps[i].x) - atan2(ps[j].y, ps[j].x
    if (o >= pi) o = o - 2 * pi;
    if (o <= -pi) o = o + 2 * pi;
    ans += AreaOfCircleTriangle(ps[i], ps[j], r) * (o >= 0 ? 1
     : -1);
  }
  return abs(ans);
```

#### 8.13 Tangent of Circles and Points to Circle

```
vector<L> tangent(C a, C b) {
#define Pij \
  P i = (b.c - a.c).unit() * a.r, j = P(i.y, -i.x); 
  z.emplace_back(a.c + i, a.c + i + j);
#define deo(I,J) \
  double d = (a.c - b.c).abs(), e = a.r I b.r, o = acos(e / d)
  P i = (b.c - a.c).unit(), j = i.spin(o), k = i.spin(-o);\
z.emplace_back(a.c + j * a.r, b.c J j * b.r);\
z.emplace_back(a.c + k * a.r, b.c J k * b.r);
  if (a.r < b.r) swap(a, b);</pre>
  vector<L> z;
  if ((a.c - b.c).abs() + b.r < a.r) return z;</pre>
  else if (same((a.c - b.c).abs() + b.r, a.r)) { Pij; }
  else {
    deo(-,+):
    if (same(d, a.r + b.r)) { Pij; }
    else if (d > a.r + b.r) \{ deo(+,-); \}
  return z;
}
vector<L> tangent(C c, P p) {
  vector<L> z;
  double d = (p - c.c).abs();
  if (same(d, c.r)) {
    P i = (p - c.c).spin(pi / 2);
    z.emplace_back(p, p + i);
  } else if (d > c.r) {
    double o = acos(c.r / d);
    P i = (p - c.c).unit(), j = i.spin(o) * c.r, k = i.spin(-o) * c.r;
    z.emplace_back(c.c + j, p);
    z.emplace_back(c.c + k, p);
  return z;
```

#### 8.14 Area of Union of Circles

```
vector<pair<double, double>> CoverSegment(C &a, C &b) {
  double d = (a.c - b.c).abs();
        vector<pair<double, double>> res;
        if (same(a.r + b.r, d));
        else if (d \leftarrow abs(a.r - b.r) + eps) {
              if (a.r < b.r) res.emplace_back(0, 2 * pi);</pre>
        } else if (d < abs(a.r + b.r) - eps) {
              double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d)
                 ), z = (b.c - a.c).angle();
              if (z < 0) z += 2 * pi;
double l = z - o, r = z + o;
if (l < 0) l += 2 * pi;</pre>
               if (r > 2 * pi) r = 2 * pi;
              if (l > r) res.emplace_back(l, 2 * pi), res.emplace_back(0,
              else res.emplace_back(l, r);
        return res;
 }
  double CircleUnionArea(vector<C> c) { // circle should be
                 identical
        int n = c.size();
        double a = 0, w;
        for (int i = 0; w = 0, i < n; ++i) {
              vector<pair<double, double>> s = \{\{2 * pi, 9\}\}, z; for (int j = 0; j < n; ++j) if (i != j) {
                     z = CoverSegment(c[i], c[j]);
                     for (auto &e : z) s.push_back(e);
              sort(s.begin(), s.end());
auto F = [&] (double t) { return c[i].r * (c[i].r * t + c[i].r * 
                 ].c.x * sin(t) - c[i].c.y * cos(t)); };
               for (auto &e : s) {
                    if (e.first > w) a += F(e.first) - F(w);
                     w = max(w, e.second);
              }
        return a * 0.5;
į }
```

#### 8.15 Minimum Distance of 2 Polygons

```
// p, q is convex
 double TwoConvexHullMinDist(Point P[], Point Q[], int n, int m)
   int YMinP = 0, YMaxQ = 0;
   double tmp, ans = 999999999;
for (i = 0; i < n; ++i) if(P[i].y < P[YMinP].y) YMinP = i;
   for (i = 0; i < m; ++i) if(Q[i].y > Q[YMaxQ].y) YMaxQ = i;
  P[n] = P[0], Q[m] = Q[0];
for (int i = 0; i < n; ++i) {
     while (tmp = Cross(Q[YMaxQ + 1] - P[YMinP + 1], P[YMinP] -
      P[YMinP + 1]) > Cross(Q[YMaxQ] - P[YMinP + 1], P[YMinP] -
      P[YMinP + 1])) YMaxQ = (YMaxQ + 1) % m;
     if (tmp < 0) ans = min(ans, PointToSegDist(P[YMinP], P[</pre>
      YMinP + 1], Q[YMaxQ]));
     else ans = min(ans, TwoSegMinDist(P[YMinP], P[YMinP + 1], Q
      [YMaxQ], Q[YMaxQ + 1]));
     YMinP = (YMinP + 1) \% n;
   return ans;
ĺ}
```

#### 8.16 2D Convex Hull

```
bool operator < (const P &a, const P &b) { return same(a.x, b.x</pre>
    ) ? a.y < b.y : a.x < b.x; }
bool operator > (const P &a, const P &b) { return same(a.x, b.x
    ) ? a.y > b.y : a.x > b.x; }
#define crx(a, b, c) ((b - a) \wedge (c - a))
vector<P> convex(vector<P> ps) {
  vector<P> p;
  sort(ps.begin(), ps.end(), [&] (P a, P b) { return same(a.x,
    b.x) ? a.y < b.y : a.x < b.x; });
  for (int i = 0; i < ps.size(); ++i) {</pre>
   while (p.size() \ge 2 \& crx(p[p.size() - 2], ps[i], p[p.
    size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  int t = p.size();
  for (int i = (int)ps.size() - 2; i >= 0; --i) {
    while (p.size() > t && crx(p[p.size() - 2], ps[i], p[p.size
    () - 1]) >= 0) p.pop_back();
```

```
p.push_back(ps[i]);
  p.pop_back();
  return p;
}
int sgn(double x) { return same(x, 0) ? 0 : x > 0 ? 1 : -1; }
P isLL(P p1, P p2, P q1, P q2) {
 double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
struct CH {
  int n;
  vector<P> p, u, d;
  CH() {}
  CH(vector<P> ps) : p(ps) {
    n = ps.size();
    rotate(p.begin(), min_element(p.begin(), p.end()), p.end())
    auto t = max_element(p.begin(), p.end());
    d = vector<P>(p.begin(), next(t));
    u = vector < P > (t, p.end()); u.push_back(p[0]);
  int find(vector<P> &v, P d) {
    int l = 0, r = v.size();
    while (l + 5 < r) {
      int \hat{L} = (1 * 2 + r) / 3, R = (1 + r * 2) / 3;
      if (v[L] * d > v[R] * d) r = R;
      else l = L;
    int x = 1:
    for (int i = l + 1; i < r; ++i) if (v[i] * d > v[x] * d) x
     = i;
    return x;
  int findFarest(P v) {
    if (v.y > 0 \mid | v.y == 0 \& v.x > 0) return ((int)d.size() -
      1 + find(u, v)) % p.size();
    return find(d, v);
  P get(int 1, int r, P a, P b) {
    int s = sgn(crx(a, b, p[l % n]));
    while (l + 1 < r) {
      int m = (l + r) >> 1;
      if (sgn(crx(a, b, p[m % n])) == s) l = m;
      else r = m:
    return isLL(a, b, p[l % n], p[(l + 1) % n]);
  vector<P> getIS(P a, P b) {
    int X = findFarest((b - a).spin(pi / 2));
    int Y = findFarest((a - b).spin(pi / 2));
    if (X > Y) swap(X, Y)
    if (sgn(crx(a, b, p[X])) * sgn(crx(a, b, p[Y])) < 0) return
      \{get(X, Y, a, b), get(Y, X + n, a, b)\};
    return {};
  void update_tangent(P q, int i, int &a, int &b) {
    if (sgn(crx(q, p[a], p[i])) > 0) a = i;
    if (sgn(crx(q, p[b], p[i])) < 0) b = i;
  void bs(int 1, int r, P q, int &a, int &b) {
    if (l == r) return;
    update_tangent(q, 1 % n, a, b);
    int s = sgn(crx(q, p[l % n], p[(l + 1) % n]));
    while (l + 1 < r) {
      int m = (l + r) >> 1;
      if (sgn(crx(q, p[m % n], p[(m + 1) % n])) == s) l = m;
    update_tangent(q, r % n, a, b);
  bool contain(P p) {
    if (p.x < d[0].x | l p.x > d.back().x) return 0;
    auto it = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
    if (it->x == p.x) {
    if (it->y > p.y) return 0;
} else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
    it = lower_bound(u.begin(), u.end(), P(p.x, 1e12), greater<
     P>());
    if (it->x == p.x) {
      if (it->y < p.y) return 0;</pre>
    } else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
    return 1:
```

```
bool get_tangent(P p, int &a, int &b) { // b -> a
    if (contain(p)) return 0;
    a = b = 0;
    int i = lower_bound(d.begin(), d.end(), p) - d.begin();
    bs(0, i, p, a, b);
    bs(i, d.size(), p, a, b);
    i = lower_bound(u.begin(), u.end(), p, greater<P>()) - u.
    begin();
    bs((int)d.size() - 1, (int)d.size() - 1 + i, p, a, b);
    bs((int)d.size() - 1 + i, (int)d.size() - 1 + u.size(), p,
    a, b);
    return 1;
}
```

#### 8.17 3D Convex Hull

```
double absvol(const P a,const P b,const P c,const P d) {
  return abs(((b-a)^{(c-a)})*(d-a))/6;
}
struct convex3D {
static const int maxn=1010;
struct T{
  int a,b,c;
  bool res;
  T(){}
  T(int a, int b, int c, bool res=1):a(a),b(b),c(c),res(res){}
int n,m;
P p[maxn];
T f[maxn*8];
int id[maxn][maxn];
bool on(T &t,P &q){
  return ((p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>eps;
void meow(int q,int a,int b){
  int g=id[a][b];
  if(f[g].res){
    if(on(f[g],p[q]))dfs(q,g);
      id[q][b]=id[a][q]=id[b][a]=m;
      f[m++]=T(b,a,q,1);
  }
}
void dfs(int p,int i){
  f[i].res=0;
  meow(p,f[i].b,f[i].a);
  meow(p,f[i].c,f[i].b);
  meow(p,f[i].a,f[i].c);
void operator()(){
  if(n<4)return;</pre>
  if([&](){
    for(int i=1;i<n;++i)if(abs(p[0]-p[i])>eps)return swap(p[1],
    p[i]),0;
return 1;
  }() || [&](){
    for(int i=2;i<n;++i)if(abs((p[0]-p[i])^(p[1]-p[i]))>eps)
     return swap(p[2],p[i]),0;
    return 1
  }() || [&](){
    for(int i=3; i<n; ++i)if(abs(((p[1]-p[0])^(p[2]-p[0]))*(p[i]-
     p[0]))>eps)return swap(p[3],p[i]),0;
     eturn 1;
  }())return;
  for(int i=0;i<4;++i){</pre>
    T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
    if(on(t,p[i]))swap(t.b,t.c);
    id[t.a][t.b]=id[t.b][t.c]=id[t.c][t.a]=m;
    f[m++]=t;
  for(int i=4;i<n;++i)for(int j=0;j<m;++j)if(f[j].res && on(f[j</pre>
     ],p[i])){
    dfs(i,j);
    break;
  for(int i=0;i<mm;++i)if(f[i].res)f[m++]=f[i];</pre>
bool same(int i,int j){
  return !(absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].a])>eps
     | absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].b])>eps | |
     absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].c])>eps);
int faces(){
```

```
int r=0;
for(int i=0;i<m;++i){
  int iden=1;
  for(int j=0;j<i;++j)if(same(i,j))iden=0;
  r+=iden;
}
return r;
}
} tb;</pre>
```

#### 8.18 Minimum Enclosing Circle

```
pt center(const pt &a, const pt &b, const pt &c) { pt p0 = b - a, p1 = c - a;
   double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0 \land p1;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
circle min_enclosing(vector<pt> &p) {
  random_shuffle(p.begin(), p.end());
   double r = 0.0;
  pt cent;
   for (int i = 0; i < p.size(); ++i) {</pre>
     if (norm2(cent - p[i]) <= r) continue;</pre>
     cent = p[i];
     r = 0.0;
     for (int j = 0; j < i; ++j) {
  if (norm2(cent - p[j]) <= r) continue;</pre>
        cent = (p[i] + p[j]) / 2;
r = norm2(p[j] - cent);
        for (int k = 0; k < j; ++k) {
          if (norm2(cent - p[k]) <= r) continue;</pre>
          cent = center(p[i], p[j], p[k]);
          r = norm2(p[k] - cent);
     }
  return circle(cent, sqrt(r));
```

#### 8.19 Closest Pair

```
double closest_pair(int l, int r) {
   // p should be sorted increasingly according to the x-
     coordinates.
   if (l == r) return 1e9;
  if (r - l == 1) return dist(p[l], p[r]);
  int m = (l + r) >> 1;
  double d = min(closest_pair(l, m), closest_pair(m + 1, r));
   vector<int> vec;
   for (int i = m; i >= l \&\& fabs(p[m].x - p[i].x) < d; --i) vec
      .push_back(i);
   for (int i = m + 1; i \le r \&\& fabs(p[m].x - p[i].x) < d; ++i)
      vec.push_back(i);
   sort(vec.begin(), vec.end(), [&](int a, int b) { return p[a].
     y < p[b].y; \});
   for (int i = 0; i < vec.size(); ++i) {</pre>
     for (int j = i + 1; j < vec.size() && fabs(p[vec[j]].y - p[
     vec[i]].y) < d; ++j) {
       d = min(d, dist(p[vec[i]], p[vec[j]]));
    }
  return d:
}
```

# 9 Miscellaneous

#### 9.1 Bitwise Hack

```
| long long next_perm(long long v) {
| long long t = v | (v - 1);
| return (t + 1) | (((~t & -~t) - 1) >> (__builtin_ctz(v) + 1))
| ;
| }
| void subset(long long s) {
| long long sub = s;
| while (sub) sub = (sub - 1) & s;
| }
```

#### 9.2 Hilbert's Curve (faster Mo's algorithm)

```
long long hilbert(int n, int x, int y) {
  long long res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x \& s) > 0;
    int ry = (y & s) > 0;
res += s * 1ll * s * ((3 * rx) ^ ry);
    if (ry == 0) {
      if (rx == 1) x = s - 1 - x, y = s - 1 - y;
       swap(x, y);
    }
  return res;
}
```

#### 9.3Mo's Algorithm on Tree

```
void MoAlgoOnTree() {
   Dfs(0, -1);
   vector<int> euler(tk);
   for (int i = 0; i < n; ++i) {
  euler[tin[i]] = i;</pre>
     euler[tout[i]] = i;
   vector<int> l(q), r(q), qr(q), sp(q, -1);
   for (int i = 0; i < q; ++i) {
   if (tin[u[i]] > tin[v[i]]) swap(u[i], v[i]);
     int z = GetLCA(u[i], v[i]);
     sp[i] = z[i];
     if (z == u) l[i] = tin[u[i]], r[i] = tin[v[i]];
     else l[i] = tout[u[i]], r[i] = tin[v[i]];
     qr[i] = i;
   sort(qr.begin(), qr.end(), [&](int i, int j) {
  if (l[i] / kB == l[j] / kB) return r[i] < r[j];</pre>
     return l[i] / kB < l[j] / kB;</pre>
   vector<bool> used(n);
   // Add(v): add/remove v to/from the path based on used[v]
   for (int i = 0, tl = 0, tr = -1; i < q; ++i) {
  while (tl < l[qr[i]]) Add(euler[tl++]);</pre>
     while (tl > l[qr[i]]) Add(euler[--tl]);
     while (tr > r[qr[i]]) Add(euler[tr--]);
     while (tr < r[qr[i]]) Add(euler[++tr]);</pre>
      // add/remove LCA(u, v) if necessary
   }
į }
```

#### 9.4 Java

```
import java.io.*
import java.util.*
import java.lang.*
import iava.math.*:
public class filename{
 static Scanner in = new Scanner(System.in);
  public static void main(String[] args) throws Exception {
    Scanner fin = new Scanner(new File("infile"));
PrintWriter fout = new PrintWriter("outfile", "UTF-8");
    fout.println(fin.nextLine());
    fout.close();
    while (in.hasNext()) {
      String str = in.nextLine(); // getline
      String stu = in.next(); // string
    System.out.println("Case #" + t);
    System.out.printf("%d\n", 7122);\\
    int[][] d = {{7,1,2,2},{8,7}};
    int g = Integer.parseInt("-123");
    long f = (long)d[0][2];
    List<Integer> l = new ArrayList<>();
    Random rg = new Random();
    for (int i = 9; i >= 0; --i) {
      l.add(Integer.valueOf(rg.nextInt(100) + 1));
      l.add(Integer.valueOf((int)(Math.random() * 100) + 1));
    Collections.sort(l, new Comparator<Integer>() {
      public int compare(Integer a, Integer b) { return a - b;
    });
    for (int i = 0; i < l.size(); ++i)</pre>
      System.out.print(l.get(i));
    Set<String> s = new HashSet<String>(); // TreeSet
    s.add("jizz");
```

```
System.out.println(s):
     System.out.println(s.contains("jizz"));
     Map<String, Integer> m = new HashMap<String, Integer>();
     m.put("lol", 7122);
     System.out.println(m);
     for(String key: m.keySet())
  System.out.println(key + " : " + m.get(key));
     System.out.println(m.containsKey("lol"));
     System.out.println(m.containsValue(7122));
     System.out.println(Math.PI);
     System.out.println(Math.acos(-1));
     BigInteger bi = in.nextBigInteger(), bj = new BigInteger("
     -7122"), bk = BigInteger.valueOf(17171);
int sgn = bi.signum(); // sign(bi)
     bi = bi.subtract(BigInteger.ONE).multiply(bj).divide(bj).
      and(bj).gcd(bj).max(bj).pow(87);
     int meow = bi.compareTo(bj); // -1 0 1
String stz = "f5abd69150";
     BigInteger b16 = new BigInteger(stz, 16);
     System.out.println(b16.toString(2));
}
```

# 9.5 Dancing Links

```
namespace dlx {
int lt[maxn], rg[maxn], up[maxn], dn[maxn], cl[maxn], rw[maxn],
      bt[maxn], s[maxn], head, sz, ans;
void init(int c) {
  for (int i = 0; i < c; ++i) {</pre>
    up[i] = dn[i] = bt[i] = i;
    lt[i] = i == 0 ? c : i - 1;
    rg[i] = i == c - 1 ? c : i + 1;
    s[i] = 0;
  rg[c] = 0, lt[c] = c - 1;
 up[c] = dn[c] = -1;
head = c, sz = c + 1;
void insert(int r, const vector<int> &col) {
  if (col.empty()) return;
  int f = sz;
  for (int i = 0; i < (int)col.size(); ++i) {</pre>
    int c = col[i], v = sz++;
    dn[bt[c]] = v;
    up[v] = bt[c], bt[c] = v;
    rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
    rw[v] = r, cl[v] = c;
    ++s[c];
    if (i > 0) lt[v] = v - 1;
  lt[f] = sz - 1;
}
void remove(int c) {
  lt[rg[c]] = lt[c], rg[lt[c]] = rg[c];
  up[dn[j]] = up[j], dn[up[j]] = dn[j], --s[cl[j]];
 }
}
void restore(int c) {
  for (int i = up[c]; i != c; i = up[i]) {
    for (int j = lt[i]; j != i; j = lt[j])

      ++s[cl[j]], up[dn[j]] = j, dn[up[j]] = j;
  lt[rg[c]] = c, rg[lt[c]] = c;
}
// Call dlx::make after inserting all rows.
void make(int c) {
  for (int i = 0; i < c; ++i)
    dn[bt[i]] = i, up[i] = bt[i];
void dfs(int dep) {
  if (dep >= ans) return;
  if (rg[head] == head) return ans = dep, void();
  if (dn[rg[head]] == rg[head]) return;
  int c = rg[head];
  int w = c;
  for (int x = c; x != head; x = rg[x]) if (s[x] < s[w]) w = x;
  remove(w);
  for (int i = dn[w]; i != w; i = dn[i]) {
    for (int j = rg[i]; j != i; j = rg[j]) remove(cl[j]);
    dfs(dep + 1);
    for (int j = lt[i]; j != i; j = lt[j]) restore(cl[j]);
```

```
| }
| restore(w);
|}
|int solve() {
| ans = 1e9, dfs(0);
| return ans;
|}
```

#### 9.6 Offline Dynamic MST

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second =
     weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i such
     that cnt[i] == 0
void contract(int l, int r, vector<int> v, vector<int> &x,
     vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
    if (cost[i] == cost[j]) return i < j;</pre>
    return cost[i] < cost[j];</pre>
  djs.save();
  for (int i = l; i <= r; ++i) djs.merge(st[qr[i].first], ed[qr</pre>
     [i].first]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
  dis.undo():
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[x[i]],</pre>
     ed[x[i]]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
    }
  djs.undo();
}
void solve(int l, int r, vector<int> v, long long c) {
    cost[qr[l].first] = qr[l].second;
    if (st[qr[1].first] == ed[qr[1].first]) {
    printf("%lld\n", c);
    return.
      return:
    int minv = qr[l].second;
    for (int i = 0; i < (int)v.size(); ++i) minv = min(minv,</pre>
     cost[v[i]]);
    printf("%lld\n", c + minv);
    return;
  }
  int m = (l + r) >> 1;
  vector<int> lv = v, rv = v;
  vector<int> x, y;
  for (int i = m + 1; i \le r; ++i) {
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first);
  contract(l, m, lv, x, y);
long long lc = c, rc = c;
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    lc += cost[x[i]];
    \label{eq:djs.merge} \mbox{djs.merge(st[x[i]], ed[x[i]]);}
  solve(l, m, y, lc);
  djs.undo();
  x.clear(), y.clear();
  for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
  for (int i = l; i <= m; ++i) {
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first);
  contract(m + 1, r, rv, x, y);
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    rc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
```

```
| solve(m + 1, r, y, rc);
| djs.undo();
| for (int i = l; i <= m; ++i) cnt[qr[i].first]++;
|}</pre>
```

#### 9.7 Manhattan Distance MST

```
void solve(int n) {
   init();
   vector<int> v(n), ds;
   for (int i = 0; i < n; ++i) {
    v[i] = i;
     ds.push_back(x[i] - y[i]);
   sort(ds.beain(), ds.end());
   ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
   sort(v.begin(), v.end(), [\&](int i, int j) { return x[i] == x}
      [j] ? y[i] > y[j] : x[i] > x[j]; y[j] : x[i] > x[j]; y[i]
   for (int i = 0; i < n; ++i) {
     int p = lower_bound(ds.begin(), ds.end(), x[v[i]] - y[v[i]]
     ]]) - ds.begin() + 1;
     pair<int, int> q = query(p);
     // query return prefix minimum
     if (~q.second) add_edge(v[i], q.second);
     add(p, make\_pair(x[v[i]] + y[v[i]], v[i]));
}
void make_graph() {
   solve(n);
   for (int i = 0; i < n; ++i) swap(x[i], y[i]);
   solve(n);
   for (int i = 0; i < n; ++i) x[i] = -x[i];
   solve(n);
   for (int i = 0; i < n; ++i) swap(x[i], y[i]);
   solve(n);
}
```

#### 9.8 IOI 2016 Aliens Trick

```
long long Alien() {
  long long c = kInf;
  for (int d = 60; d >= 0; --d) {
    // cost can be negative as well, depending on the problem.
    if (c - (1LL << d) < 0) continue;
    long long ck = c - (1LL << d);
    pair<long long, int> r = check(ck);
    if (r.second == k) return r.first - ck * k;
    if (r.second < k) c = ck;
}
pair<long long, int> r = check(c);
return r.first - c * k;
}
```

### 9.9 Matroid Intersection

Start from  $S = \emptyset$ . In each iteration, let

- $Y_1 = \{x \notin S \mid S \cup \{x\} \in I_1\}$
- $Y_2 = \{x \notin S \mid S \cup \{x\} \in I_2\}$

If there exists  $x \in Y_1 \cap Y_2$ , insert x into S. Otherwise for each  $x \in S, y \not\in S,$  create edges

•  $x \to y \text{ if } S - \{x\} \cup \{y\} \in I_1.$ •  $y \to x \text{ if } S - \{x\} \cup \{y\} \in I_2.$ 

Find a shortest path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x) if  $x \notin S$ . Find the path with the minimum number of edges among all minimum length paths and alternate it.