

Competitive Programming Notebook

IllegalSkillsException

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1 Introduction

This document was written to be used in programming competitions by my team: *HammerHappy*. Conciseness (not clarity) was the priority.

2 Strings

2.1 Knuth-Morris-Pratt (KMP)

```
std::vector<int> compute_prefix(
    const string& p) {
    int m = p.size();
    std::vector<int> pi(m);
    pi[0] = 0;

    int k = 0;
    for (int q = 1; q < m; q++) {
        while (k > 0 && p[k] != p[q])
            k = pi[k-1];
        if (p[k] == p[q])
            k++;
        pi[q] = k;
    }
    return pi;
}
```

```
void kmp_match(const string& s,
               const string& p) {
    std::vector<int> pi = compute_prefix(p);
    int q = 0;
    int n = s.size();
    int m = p.size();

    for (int i = 0; i < n; i++) {
        while (q > 0 && p[q] != s[i])
            q = pi[q-1];
        if (p[q] == s[i])
            q++;
        if (q == m) {
            std::cout
```

```
<< "Match at pos: "
<< (i - m + 1)
<< std::endl;
```

```
}
}
```

2.2 Range Minimum Query

```
int main() {
    int N, Q, i, j, k;

    scanf("%d%d", &N, &Q);

    for (i = 0; i < N; i++)
        scanf("%d", &n[i]);

    for (i = 0; i < N; i++)
        m[i][0] = M[i][0] = n[i];

    for (i=1; (1 << i) <= N; i++) {
        for (j = 0; j + (1 << i) - 1 < N; j++) {
            m[j][i] = min(m[j][i-1], m[j +
                (1 << (i-1))][i-1]);
            M[j][i] = max(M[j][i-1], M[j +
                (1 << (i-1))][i-1]);
        }
    }

    for (k = 0; k < Q; k++) {
        scanf("%d%d", &i, &j);
        i--;
        j--;
        int t, p;
        t = (int)(log(j - i + 1) / log(2));
        p = 1 << t;
        printf("%d\n", max(M[i][t],
            M[j - p + 1][t])
            - min(m[i][t], m[j - p + 1][t]));
    }
    return 0;
}
```

$$M[i][j] = \max(M[i][j-1], M[i+2^{j-1}][j-1])$$

$$RMQ_A(i, j) = \max(M[i][k], M[j-2^k+1][k])$$

2.3 Nth Permutation

```
/**
 * Computes kth (0 to s.size()-1) permutation
 * of string s
```

```

*/
std::string nth_permutation(uint64_t k,
                           const std::string &s) {
    uint64_t factorial = 1;
    for (uint64_t i = 1; i <= s.size(); ++i) {
        factorial *= i;
    }

    std::string s_copy = s;
    std::string res;
    for (uint64_t j = 0; j < s.size(); ++j) {
        // compute how many permutations
        // on the rest of the
        // string s[j + 1 .. s.size() - 1]
        factorial /= s.size() - j;

        // store character
        uint64_t l = k / factorial;
        res += s_copy[l];

        // remove already used character
        s_copy.erase(s_copy.begin() + l);

        // compute new value of k
        k = k % factorial;
    }
    return res;
}

```

3 Dynamic Programming

3.1 Longest Common Subsequence (LCS)

```

int L[MAX][MAX] = {{0}};
int LCS(char A[], char B[]) {
    // m = strlen(A)
    // n = strlen(B)
    for (int i = m; i >= 0; i--) {
        for (int j = n; j >= 0; j--) {
            if (!A[i] || !B[j])
                L[i][j] = 0;
            else if (A[i] == B[j])
                L[i][j] = 1 + L[i + 1][j + 1];
            else L[i][j] = max(L[i + 1][j],
                               L[i][j + 1]);
        }
    }
    return L[0][0];
}

```

```

int LCSString(int L[MAX][MAX]) {
    int i, j;
    i = j = 0;
    while (i < m && j < n) {
        if (A[i] == B[j]) {
            // put A[i] at the end
            // of solution string
            i++; j++;
        }
        if (L[i + 1][j] >= L[i][j + 1]) i++;
        else j++;
    }
}

```

3.2 Longest Increasing Subsequence (LIS)

3.2.1 $O(n^2)$ version

```

int pred[MAX_SIZE], lasti;
int LIS(int C[], int n) {
    int s[MAX_SIZE], max = INT_MIN;
    for (int i = 1; i < n; i++) {
        for (int j = 0; j < i; j++) {
            if (C[i] > C[j] && s[i] <= s[j]) {
                pred[i] = j;
                if ((s[i] = s[j] + 1) > max)
                    lasti = i;
                max = s[i];
            }
        }
    }
    return max;
}

```

```

void PrintLIS() {
    int i, j, aux[MAX_SIZE];
    for (j = max - 1, i = lasti; j >= 0; j--) {
        aux[j] = C[i];
        i = pred[i];
    }

    for (j = 0; j < max; j++)
        printf("%d\n", aux[j]);
}

```

3.2.2 $O(n \log n)$ version

```

int a, num[120000], n, ans[120000], sz;

while (scanf("%d", &n) == 1) {
    for (a = 0; a < n; a++)

```

```

    scanf("%d", &num[a]);
    sz = 0;
    for (a = 0; a < n; a++) {
        int* it = lower_bound(
            ans, ans + sz, num[a]);
        if (it != ans + sz) *it = num[a];
        else ans[sz++] = num[a];
    }
    printf("%d\n", sz);
}

```

3.3 MCM (Matrix Chain Multiplication)

```

void mcm() {
    int i, j, n = 3;
    for (i = 0; i < n; i++)
        m[i][i] = 0;

    for (i = n - 1; i >= 0; i--)
        for (j = i + 1; j <= n; j++)
            m[i][j] = calc(i, j);
}

int calc(int i, int j) {
    int res = INT_MAX;
    for (k = i; k < j; k++) {
        tmp = m[i][k] + m[k + 1][j] +
            Line[i] * Col[k] * Col[j];
        if (tmp < res) {
            res = tmp;
            s[i][j] = k;
        }
    }
    return res;
}

```

//printMCM(0,N-1);

```

void printMCM(int i, int j) {
    if (i == j) printf("A%d", i);
    else {
        putchar('(');
        printMCM(i, s[i][j]);
        putchar('*');
        printMCM(s[i][j] + 1, j);
        putchar(')');
    }
}

```

3.4 Knapsack

```
int n[WSIZE][ISIZE] = {{0}}
```

```

// put one zero in weight and value;
// e.g.
// weight={>0<,3,4,5}
// value={>0<,3,4,5,6};
int knapsack(int items, int W,
    int value[], int weight[]){
    for (int i = 1; i <= items; i++) {
        for (int j = 0; j <= W; j++) {
            if (weight[i] <= j) {
                if (value[i] + n[i-1][j-weight[i]]
                    > n[i-1][j]) {
                    n[i][j] = value[i] +
                        n[i-1][j-weight[i]];
                } else {
                    n[i][j] = n[i-1][j];
                }
            } else n[i][j] = n[i-1][j];
        }
    }
    return n[items][W];
}

```

```

void print_sequence(int items, int W, int weight[]) {
    int i = items, k = W;
    while (i > 0 && k > 0) {
        if (n[i][k] != n[i-1][k]) {
            printf("item %d is in\n", i);
            k = k - weight[i-1];
        }
        i--;
    }
}

```

3.5 Counting Change

```

int coins[] = {50,25,10,5,1};
int coin_change(int n) {
    table[0] = 1;
    for (int i = 0; i < 5; i++) {
        c = coins[i];
        for (int j = c; j <= n; j++)
            table[j] += table[j - c];
    }
    return table[n];
}

```

3.6 Coin Changing

```

int n[10000], i, N;
int coins[]={50,25,10,5,1},k;

scanf("%d", &N);
for (int i = 0; i <= N; i++)

```

```

    n[i] = INT_MAX;
n[0] = 0;
for (int i = 0; i < 5; i++) {
    for (k = 0; k <= N - coins[i]; k++) {
        n[k + coins[i]] =
            min(n[k] + 1, n[k + coins[i]]);
    }
}
printf("%d\n", n[N]);

```

3.7 Biggest Sum

```

#define SIZE 20000
int n[SIZE];

int biggest_sum() {
    int k, s, b;
    int xl, xr, best, prevx;

    cin >> k;
    for (int i = 1; i <= k; i++) {
        xr = xl = 0;

        cin >> s;
        for (int j = 0; j < s - 1; j++)
            cin >> n[j];

        prevx = xl = xr = 0;
        best = b = n[0];
        for (int j = 1; j < s - 1; j++) {
            if (b < 0)
                prevx = j;
            b = n[j] + max(0, b);
            if (b > best ||
                (b == best &&
                 j - prevx > xr - xl)) {
                xl = prevx;
                xr = j;
                best = b;
            }
        }
        if (best > 0)
            cout << "Biggest_sum" << i
                 << " is between "<< xl + 1
                 << "and_" << xr + 2
                 << endl;
    }
    return 0;
}

```

3.8 Edit Distance

Possible actions:

1. Delete a character
2. Insert a new character
3. Replace a character

```

int edit_distance(char *str1, char *str2) {
    int n[SIZE][SIZE];
    int i, j, value;

    for (i = 0; i <= str1_len; i++) n[i][0] = i;
    for (j = 0; j <= str2_len; j++) n[0][j] = j;

    for (i = 1; i <= str1_len; i++) {
        for (j = 1; j <= str2_len; j++) {
            value = (str1[i - 1] != str2[j - 1]);

            n[i][j] = min(n[i - 1][j - 1] + value,
                          n[i - 1][j] + 1,
                          n[i][j - 1] + 1);
        }
    }
    return n[str1_len][str2_len];
}

T(i, j) = min(C_d + T(i - 1, j),
              T(i, j - 1) + C_i,
              T(i - 1, j - 1) + (A[i] == B[j] ? 0 : C_r))

```

3.9 Integer Partitions

$P(n)$ represents the number of possible partitions of a natural number n . $P(4) = 5, 4, 3 + 1, 2 + 2, 2 + 1 + 1, 1 + 1 + 1 + 1$
 $P(0) = 1$
 $P(n) = 0, n < 0$
 $P(n) = p(1, n)$
 $p(k, n) = p(k + 1, n) + p(k, n - k)$
 $p(k, n) = 0$ if $k > n$
 $p(k, n) = 1$ if $k = n$

3.10 Box Stacking

A set of boxes is given. $Box_i = h_i, w_i, d_i$.

We can only stack box i on box j if $w_i < w_j$ and $d_i < d_j$.

To consider all the orientations of the boxes, replace each box with 3 boxes such that $w_i \leq d_i$ and $box_1[0] = h_i, box_2[0] = w_i, box_3[0] = d_i$.

Then, sort the boxes by decreasing area($w_i * d_i$).

$H(j)$ = tallest stack of boxes with box j on top.

$H(j) = \max_{i < j \& w_i > w_j \& d_i > d_j} (H(i)) + h_j$
Check $H(j)$ for all values of j .

3.11 Building Bridges

Maximize number of non-crossing bridges. Ex:

bridge1: 2, 5, 1, n , \dots , 4, 3

bridge2: 1, 2, 3, 4, \dots , n

Let $X(i)$ be the index of the corresponding city on northern bank. $X(1) = 3, X(2) = 1, \dots$

Find longest increasing subsequence of $X(1), \dots, X(n)$.

3.12 Partition Problem

Input: A given arrangement S of non-negative numbers s_1, \dots, s_n and an integer k .

Output: Partition S into k ranges, so as to minimize the maximum sum over all the ranges.

```
int M[1000][100], D[1000][100];
void partition_i(vector<int> &v, int k) {
    int p[1000], n = v.size();
    v.insert(v.begin(), 0);
    p[0] = 0;
    for(int i = 1; i < v.size(); i++)
        p[i] = p[i - 1] + v[i];

    for (int i = 1; i <= n; i++)
        M[i][1] = p[i];
    for (int i = 1; i <= k; i++)
        M[1][i] = v[1];
    for (int i = 2; i <= n; i++) {
        for (int j = 2; j <= k; j++) {
            M[i][j] = INT_MAX << 1 - 1;
            int s = 0;
            for (int x = 1; x <= i - 1; x++) {
                s = max(M[x][j - 1], p[i] - p[x]);
                if (M[i][j] > s) {
                    M[i][j] = s;
                    D[i][j] = x;
                }
            }
        }
    }
    printf("%d\n", M[n][k]);
}
```

//n = number of elements of the initial set

```
void reconstruct_partition(
    const vector<int> &S, int n, int k) {
    if (k == 1) {
        for (int i = 1; i <= n; i++)
            printf("%d_", &S[i]);
        putchar('\n');
    } else {
        reconstruct_partition(S, D[n][k], k - 1);
        for (int i = D[n][k] + 1; i <= n; i++)
            printf("%d_", S[i]);
        putchar('\n');
    }
}
```

3.13 Balanced Partition

```
enum {DONT_GET, GET};
char **sol, **P;
```

```
// return 1 if there is a subset
// of v0...vi with sum j
// 0 otherwise
int calcP(int i, int j, const vi &v) {
    if (i < 0 || j < 0) return 0;
    if (P[i][j] != -1) return P[i][j];

    if (j == 0) { // trivial case
        sol[i][j] = DONT_GET;
        return P[i][j] = 1;
    }
    if (v[i] == j) {
        sol[i][j] = GET;
        return P[i][j] = 1;
    }

    int res1 = calcP(i - 1, j, v);
    int res2 = calcP(i - 1, j - v[i], v);
    if (res1 >= res2)
        P[i][j] = res1, sol[i][j] = DONT_GET;
    else P[i][j] = res2, sol[i][j] = GET;
    return P[i][j];
}
```

```
// v is the vector of values
// k is the maximum value in v
// sum is the sum of all elements in v
void balanced_partition(vi &v,
```

```

        int k, int sum) {
P = new char*[v.size()];
sol = new char*[v.size()];
for (int i = 0; i < v.size(); i++) {
    P[i] = new char[k * v.size() + 1];
    sol[i] = new char[k * v.size() + 1];
    for (int j = 0;
        j < k * v.size() + 1; j++)
        P[i][j] = -1, sol[i][j] = DONT_GET;
}
for (int i = 0; i < v.size(); i++)
    for (int j = 0;
        j < v.size() * k + 1; j++)
        calcP(i, j, v);
//calcP(v.size() - 1, sum/2, v);

int S = sum / 2;
if (sum & 1 || !P[v.size() - 1][S])
    cout << "ERROR" << endl;
else cout << "SUCCESS" << endl;
}

void free_mem(vi& v) {
    for (int i = 0; i < v.size(); i++) {
        delete P[i]; delete sol[i];
    }
    delete[] P;
    delete[] sol;
}

// get_solution(v.size() - 1,
// accumulate(v.begin(), v.end(), 0) / 2,
// v1, v2, v);
void get_solution(int i, int j,
    vi &S1, vi &S2, vi &v) {
    if (j < 0 || i < 0) return;
    if (sol[i][j] == GET) {
        S1.push_back(v[i]);
        return get_solution(i - 1, j - v[i],
            S1, S2, v);
    } else {
        S2.push_back(v[i]);
        return get_solution(i - 1, j,
            S1, S2, v);
    }
}

```

4 Graphs

4.1 Heap

```

#define LEFT(i) (2 * (i + 1) - 1)
#define RIGHT(i) (2 * (i + 1))
#define PARENT(i) (((i) + 1) / 2 - 1)

int *min_heap, *heap_place;
long int *keys;
int heap_size=0;

#define update_place(i) \
    heap_place[min_heap[(i)]]=(i)

void init_heap(int nelems) {
    min_heap = new int[nelems];
    keys = new long int[nelems];
    heap_place = new int[nelems];
    heap_size = nelems;
    for (int i = 0; i < nelems; i++) {
        min_heap[i] = i;
        heap_place[i] = i;
        keys[i] = LONG_MAX;
    }
}

void heap_min_heapify(int i) {
    int smallest, temp;
    int l = LEFT(i);
    int r = RIGHT(i);

    if (l < heap_size
        && keys[min_heap[l]]
            < keys[min_heap[i]])
        smallest = l;
    else smallest = i;

    if (r < heap_size &&
        keys[min_heap[r]]
            < keys[min_heap[smallest]])
        smallest = r;

    if (smallest != i) {
        temp = min_heap[i];
        min_heap[i] = min_heap[smallest];
        min_heap[smallest] = temp;
        update_place(smallest);
        update_place(i);
        heap_min_heapify(smallest);
    }
}

```

```

    }
}

int heap_extract_min() {
    if (heap_size < 1)
        return -1;
    int res = min_heap[0];
    heap_size--;
    min_heap[0] = min_heap[heap_size];
    update_place(0);
    heap_min_heapify(0);
    return res;
}

void heap_decrease_key(int elem,
                      long int key) {
    int i = heap_place[elem];

    keys[min_heap[i]] = key;

    while (i > 0
           && keys[min_heap[PARENT(i)]] >
           keys[min_heap[i]]) {
        int temp = min_heap[i];
        min_heap[i] = min_heap[PARENT(i)];
        min_heap[PARENT(i)] = temp;
        update_place(i);
        update_place(PARENT(i));
        i = PARENT(i);
    }
}

```

4.2 Find an Eulerian Path

```

stack<int> s;
vector<list<int>> adj;

void remove_edge(int u, int v) {
    for (list<int>::iterator it = adj[u].begin();
         it != adj[u].end(); it++) {
        if (*it == v) {
            it = adj[u].erase(it);
            return;
        }
    }
}

```

```

int path(int v) {
    int w;
    for (; adj[v].size(); v = w) {
        s.push(v);
        list<int>::iterator it = adj[v].begin();
        w = *it;
        remove_edge(v, w);
        remove_edge(w, v);
        edges--;
    }
    return v;
}

//u - source, v-destiny
int eulerian_path(int u, int v) {
    printf("%d\n", v);
    while (path(u) == u && !s.empty()) {
        printf("%d", u = s.top());
        s.pop();
    }
    return edges == 0;
}

```

4.3 Breadth First Search

```

bool adj[N][N];
int colour[N], d[N], p[N];
void bfs() {
    queue<int> q;
    int source = 0;

    for (int i = 0; i < N; i++) {
        d[i] = INF;
        p[i] = -1;
        colour[i] = WHITE;
    }

    d[source] = 0;
    colour[source] = GRAY;
    q.push(source);
    while (!q.empty()) {
        int u = q.front();
        q.pop();
        for (int v = 0; v < N; v++) {
            if (colour[v] == WHITE
                && adj[u][v]) {
                colour[v] = GRAY;
                d[v] = d[u] + 1;
                p[v] = u;
            }
        }
    }
}

```



```

        q.push(v);
    }
}
colour[u] = BLACK;
}
}

```

4.4 DFS/TopSort

$O(V + E)$

Recursive:

```

void dfs(int u) {
    colour[u] = GRAY;
    for (int v = 0; v < N; v++) {
        if (colour[v] == WHITE && adj[u][v]) {
            p[v] = u;
            dfs(v);
        }
    }
    colour[u] = BLACK;
    //put node in front of a list if topsort
}

```

Iterative:

```

typedef enum {WHITE, GRAY, BLACK} color_t;
vector<color_t> color;
// SCC: vector<int> close_time;

```

```

stack<int> dfs;
color = vector<color_t>(N, WHITE);
for (int i = 0; i < N; i++) {
    if (color[i] != WHITE)
        continue;
    dfs.push(i);
    // SCC2: for (int i = N - 1; i >= 0; i--) {
    // SCC2: if (color[close_time[i]] != WHITE)
    //         continue;
    // SCC2: dfs.push(close_time[i]);
    while (!dfs.empty()) {
        int u = dfs.top();
        switch(color[u]) {
            case WHITE:
                color[u] = GRAY;
                for (v in adj[u]) {
                    // SCC2: for (v in adj-t[u]) {
                    if (color[v] == WHITE) {
                        dfs.push(v);
                    }
                }
            }
        }
        break;
    }
}

```

```

case GRAY:
    color[u] = BLACK;
    dfs.pop();
    // put node in front of a list
    // if topsort
    // SCC1: close_time.push_back(u);
    break;
case BLACK:
    dfs.pop();
    break;
}
}
}

```

Maximum Spanning Tree:

Negate all the edge weights and determine the minimum spanning tree.

Minimum Product Spanning Tree:

Replace all the edge weights with their logarithm

Strongly Connected Components:

1. Run DFS: Save closing times of all vertexes.
2. Compute adj_t.
3. Run DFS: Reverse order of closing times. In adj_t.
4. Each resulting tree is a SCC.

4.5 Prim's Algorithm

4.5.1 Naive version

```

double Prim(int start, int nvert) {
    bool in[N];
    double dist[N];
    int p[N], v;

    for (int i = 0; i < nvert; i++) {
        in[i] = false;
        dist[i] = INT_MAX;
        p[i] = -1;
    }

    dist[start] = 0;
    v = start;
    while (!in[v]) {
        in[v] = true;
        for (int i = 0; i < nvert; i++) {
            if (adj[v][i] && !in[i]) {
                if (dist[i] > adj[v][i]) {
                    dist[i] = adj[v][i];

```

```

        p[i] = v;
    }
}

double d = FLT_MAX;
for (int i = 0; i < nvert; i++) {
    if (!in[i] && d > dist[i]) {
        v = i;
        d = dist[i];
    }
}
double res = 0;
for (int i = 0; i < nvert; i++)
    res += dist[i];
return res;
}

```

4.5.2 Set version

```
std::vector<std::pair<int, int>> graph[SIZE];
```

```

struct cmp_fn {
    bool operator() (const int&a,
                     const int &b) const {
        return dist[a] < dist[b] ||
            (dist[a] == dist[b] && a < b);
    }
};

int Prim(int start, int nvert) {
    set<int, cmp_fn> s;
    dist[start] = 0;
    s.insert(start);

    while (s.size()) {
        int v = *(s.begin());
        s.erase(s.begin());

        in[v] = true;
        for (unsigned int i = 0;
             i < graph[v].size(); i++) {
            int node = graph[v][i].first;
            int length = graph[v][i].second;
            if (!in[node]) {
                if (dist[node] > length) {
                    if (s.find(node) != s.end())
                        s.erase(s.find(node));
                    dist[node] = length;
                }
            }
        }
    }
}

```

```

        s.insert(node);
    }
}

ull res = 0;
for (int i = 0; i < nvert; i++)
    res += dist[i];
return res;
}

```

4.6 Dijkstra

4.6.1 Naive version

```

int dijkstra(int source, int dest,
             int nvert, int d[], int p[]) {
    bool in[N];
    int u;

    for (int i = 0; i < nvert; i++) {
        in[i] = false;
        d[i] = INF;
        p[i] = -1;
    }
    d[source] = 0;
    u = source;
    while (!in[u]) {
        in[u] = true;
        for (int v = 0; v < nvert; v++) {
            if (adj[u][v]
                && d[v] > d[u] + adj[u][v]) {
                p[v] = u;
                d[v] = d[u] + adj[u][v];
            }
        }

        int dist = INT_MAX;
        for (int i = 0; i < nvert; i++) {
            if (!in[i] && d[i] < dist) {
                u = i;
                dist = d[i];
            }
        }
    }
    return d[dest];
}

```

4.6.2 Set version

```
#define VPI std::vector<std::pair<int, int>>
```

```

int dijkstra(const VPI graph[SIZE],
             int S, int T) {
    dist[S] = 0;

    set<Node> s;
    s.insert(Node(S));

    int u = S;
    while (s.size()) {
        Node n = *(s.begin());
        s.erase(s.begin());
        u = n.x;
        seen[u] = true;

        unsigned int i;
        for (i = 0; i < graph[u].size(); ++i) {
            int node = graph[u][i].first;
            int lat = graph[u][i].second;
            if (!seen[node]
                && dist[node] > dist[u] + lat) {
                if (s.find(Node(node))
                    != s.end()) {
                    s.erase(s.find(Node(node)));
                }
                dist[node] = dist[u] + lat;
                s.insert(Node(node));
            }
        }
    }

    return dist[T];
}

```

4.7 Kth Shortest Paths $O(Km)$

```

/*
 * u - source node
 * p - predecessor vector
 * h - vector of transformation
 * v - result vector
 */
#define vvi vector<vector<int>>
void path(int u, const vector<int> &p,
         const vector<int> &h,
         vector<int> &v) {
    if (u != -1) {
        path(p[u], p, h, v);
        v.push_back(h[u]);
    }
}

```

```

}

vvi dijkstra(int source, int dest, int K) {
    vector<int> count(SIZE), d(10000),
                p(10000), h(10000), X;
    vvi res;

    for (int i = 0; i < N; i++)
        p[i] = -1;
    int elm = 1;
    h[elm] = source;
    d[elm] = 0;
    X.push_back(elm);

    while (count[dest] < K && !X.empty()) {
        int ind = 0;
        for (unsigned int i = 1;
            i < X.size(); i++) {
            if (d[X[i]] < d[X[ind]])
                ind = i;
        }
        int k = X[ind];
        X.erase(X.begin() + ind);
        int i = h[k];

        count[i]++;
        if (i == dest) {
            vector<int> v;
            path(k, p, h, v);
            res.push_back(v);
        }

        if (count[i] <= K) {
            for (int j = 0;
                j < SIZE; j++) {
                if (adj[i][j]) {
                    elm++;
                    d[elm] = d[k] + adj[i][j];
                    p[elm] = k;
                    h[elm] = j;
                    X.push_back(elm);
                }
            }
        }
    }
    return res;
}

```

4.8 Floyd-Warshall $O(n^3)$

```

void floyd(int adj[NVERT][NVERT]) {
    for (int k = 1; k <= NVERT; k++) {
        for (int i = 1; i <= NVERT; i++) {
            for (int j = 1; j <= NVERT; j++) {
                int through_k = adj[i][k]
                    + adj[k][j];
                if (through_k < adj[i][j])
                    adj[i][j] = through_k;
            }
        }
    }
}

```

4.9 Bellman-Ford

```

typedef struct {
    int source;
    int dest;
    int weight;
} Edge;

```

```

void BellmanFord(Edge edges[], int edgecount,
    int nodecount, int source) {
    int *distance = new int[nodecount];
    for (int i=0; i < nodecount; i++)
        distance[i] = INT_MAX;

    // source node distance is set to zero
    distance[source] = 0;

    for (int i = 0; i < nodecount; i++) {
        for (int j = 0; j < edgecount; j++) {
            if (distance[edges[j].source]
                != INT_MAX) {
                int new_distance =
                    distance[edges[j].source] +
                    edges[j].weight;

                if (new_distance <
                    distance[edges[j].dest])
                    distance[edges[j].dest] =
                        new_distance;
            }
        }
    }

    for (int i = 0; i < edgecount; i++) {
        if (distance[edges[i].dest] >
            distance[edges[i].source] +
            edges[i].weight) {

```

```

                puts("Negative edge weight
                .....cycles detected!");
                free(distance);
                return;
            }
        }

        for (int i = 0; i < nodecount; i++) {
            printf("The shortest distance between
            .....nodes %d and %d is %d\n",
                source, i, distance[i]);
        }
        delete[] distance;
    }
}

```

4.10 Detecting Bridges

```

int dfs(int u, int p) {
    colour[u] = 1;
    dfsNum[u] = num++;
    int leastAncestor = num;
    for (int v = 0; v < N; v++) {
        if (M[u][v] && v!=p) {
            if (colour[v] == 0) {
                int rec = dfs(v, u);
                if (rec > dfsNum[u])
                    cout << "Bridge:_"
                        << u << "_" << v
                        << endl;
                leastAncestor =
                    min(leastAncestor, rec);
            }
            else {
                leastAncestor = min(leastAncestor,
                    dfsNum[v]);
            }
        }
    }
    colour[u] = 2;
    return leastAncestor;
}

```

4.11 Finding a Loop in a Linked List $O(n)$

```
function boolean hasLoop(Node startNode) {
    Node slowNode, fastNode1, fastNode2;
    slowNode = fastNode1 = fastNode2 = startNode;
    while (slowNode && fastNode1 = fastNode2.next()
           && fastNode2 = fastNode1.next()) {
        if (slowNode == fastNode1 ||
            slowNode == fastNode2)
            return true;
        slowNode = slowNode.next();
    }
    return false;
}
```

4.12 Tree diameter

Pick a root and start a DFS from it which returns both the diameter of the subtree and its maximum height. The diameter is the maximum of (left diameter, right diameter, left height + right height).

4.13 Union Find

```
int Rank[SIZE];
int P[SIZE];

void create_set(int x) {
    P[x] = x;
    Rank[x] = 0;
}

void merge_sets(int x, int y) {
    int px = find_set(x);
    int py = find_set(y);
    if (Rank[px] > Rank[py])
        P[py] = px;
    else P[px] = py;

    if (Rank[px] == Rank[py])
        Rank[py]++;
}

int find_set(int x) {
    if (x != P[px])
        P[x] = find_set(P[x]);
    return P[x];
}

void connected_components() {
    for each vertex i
```

```
do create_set(i);
```

```
for each edge (u,v)
    if (find_set(u) != find_set(v))
        merge_sets(u,v);
}

bool same_components(int u, int v) {
    if (find_set(u) == find_set(v))
        return true;
    else return false;
}
```

4.14 Edmonds Karp

```
struct edge {
    int dest;
    int max_weight;
    int flow;
    edge * residual;
}

int nnodes;

typedef map<int, edge*> node;
typedef node** graph;
graph grafo;

void create_edge(int source,
                 int dest, int weight) {
    if ((*grafo[source]).find(dest) ==
        (*grafo[source]).end()) {
        edge* e = new edge;
        edge* res = new edge;
        e->dest = dest;
        res->dest = source;
        e->max_weight = weight;
        res->max_weight = weight;
        e->flow = 0;
        res->flow = 0;
        e->residual = res;
        res->residual = e;
        (*grafo[source])[dest] = e;
        (*grafo[dest])[source] = res;
        return;
    }

    edge* e = (*grafo[source])[dest];
    edge* res = e->residual;
    e->max_weight += weight;
```

```

    res->max_weight += weight;
    res->flow += weight;
}

int update_path(int flowsource,
                int flowdest) {
    int flow = INT_MAX;
    int noded = flowdest;
    while (noded != flowsource) {
        int source=p[noded];
        edge* e=(*grafo[source])[noded];
        if (flow>e->max_weight-e->flow) {
            flow=e->max_weight-e->flow;
        }
        noded=source;
    }

    noded=flowdest;
    while (noded != flowsource) {
        int source = p[noded];
        edge* e = (*grafo[source])[noded];
        e->flow+=flow;
        e->residual->flow-=flow;
        noded=source;
    }
    return flow;
}

int edmonds_karp(int source, int dest) {
    int res = 0;
    while (1) {
        bfs(source);
        if (colour[dest] == WHITE) {
            return res;
        }
        res += update_path(source, dest);
    }
    return res;
}

```

4.15 Ford Fulkerson

```

#define V 110
int graph[V][V];

bool bfs(int rGraph[V][V], int s,
        int t, int parent[]) {
    // Create a visited array and
    // mark all vertices as not visited
    bool visited[V];

```

```

    memset(visited, 0, sizeof(visited));

    // Create a queue
    // enqueue source vertex and
    // mark source vertex as visited
    std::queue<int> q;
    q.push(s);
    visited[s] = true;
    parent[s] = -1;

    // Standard BFS Loop
    while (!q.empty()) {
        int u = q.front();
        q.pop();

        for (int v = 0; v < V; v++) {
            if (visited[v] == false
                && rGraph[u][v] > 0) {
                q.push(v);
                parent[v] = u;
                visited[v] = true;
            }
        }
    }

    // If we reached sink in BFS starting from
    // source, then return true, else false
    return (visited[t] == true);
}

```

4.16 Widest path problem

In an undirected graph, a widest path may be found as the path between the two vertices in the maximum spanning tree of the graph

5 Geometrical Algorithms

5.1 Circle

Formula is given by

$$x^2 + y^2 = r^2$$

5.2 Triangle's medians

Any triangle's area T can be expressed in terms of its medians m_a, m_b, m_c as follows. Denoting their semi-sum $(m_a + m_b + m_c)/2$ as s , we have

$$A = \frac{4}{3} \sqrt{s(s-m_a)(s-m_b)(s-m_c)}$$

The sides of the triangle are given, from the medians:

$$\begin{aligned} a &= \frac{2}{3} \sqrt{-m_a^2 + 2m_b^2 + 2m_c^2} \\ b &= \frac{2}{3} \sqrt{-m_b^2 + 2m_a^2 + 2m_c^2} \\ c &= \frac{2}{3} \sqrt{-m_c^2 + 2m_b^2 + 2m_a^2} \end{aligned}$$

5.3 Heron's formula

$$s = \frac{a+b+c}{2}$$

Area is given by

$$A = \sqrt{s(s-a)(s-b)(s-c)}$$

5.4 Dot Product

```
int dot(int [] A, int [] B, int [] C) {
    int AB[2], BC[2];
    AB[0] = B[0]-A[0];
    AB[1] = B[1]-A[1];
    BC[0] = C[0]-B[0];
    BC[1] = C[1]-B[1];
    int dot = AB[0] * BC[0] + AB[1] * BC[1];
    return dot;
}
```

5.5 Cross Product

```
int cross(int [] A, int [] B, int [] C) {
    int AB[2], AC[2];
    AB[0] = B[0]-A[0];
    AB[1] = B[1]-A[1];
    AC[0] = C[0]-A[0];
    AC[1] = C[1]-A[1];
    int cross = AB[0] * AC[1] - AB[1] * AC[0];
    return cross;
}
```

5.6 Point on segment

A point is on a segment if its distance to the segment is 0.

Given two different points (x_1, y_1) and (x_2, y_2) the values of A, B , and C for $Ax + By + C = 0$ are given by

$$\begin{aligned} A &= y_2 - y_1 \\ B &= x_1 - x_2 \\ C &= A * x_1 + B * y_1 \end{aligned}$$

5.7 Intersection of segments

```
double det = A1*B2 - A2*B1
if (det == 0) {
    //Lines are parallel
} else {
    double x = -(A1*C2 - A2*C1) / det
    double y = -(B1*C2 - B2*C1) / det
}
```

5.8 Position of point in relation to line

```
//Input: three points P0, P1, and P2
//Return: >0 for P2 left of the line through P0 and P1
//        = 0 for P2 on the line
//        < 0 for P2 right of the line
int isLeft( Point P0, Point P1, Point P2 ) {
    return ( (P1.x - P0.x) * (P2.y - P0.y)
            - (P2.x - P0.x) * (P1.y - P0.y) );
}
```

5.9 Distance between point and line/segment

If the line is in the form $Ax + By + C = 0$:

$$d = \frac{|Ax_0 + By_0 + C|}{\sqrt{A^2 + B^2}}$$

```
//Compute the dist. from AB to C
//if isSegment=true, AB is a seg., not a line.
double linePointDist(int [] A, int [] B,
    int [] C, boolean isSegment) {
    double dist = cross(A,B,C) / distance(A,B);
    if (isSegment) {
        int dot1 = dot(A,B,C);
        if (dot1 > 0)
            return distance(B,C);
        int dot2 = dot(B,A,C);
        if (dot2 > 0)
            return distance(A,C);
    }
    return abs(dist);
}
```

5.10 Polygon Area

```

int area = 0;
/*int N = lengthof(p);*/

for (int i = 1; i + 1 < N; i++) {
    int x1 = p[i][0] - p[0][0];
    int y1 = p[i][1] - p[0][1];
    int x2 = p[i+1][0] - p[0][0];
    int y2 = p[i+1][1] - p[0][1];
    int cross = x1*y2 - x2*y1;
    area += cross;
}
return fabs(area/2.0);

```

5.11 Convex Hull

```

#include <vector>
vector<point> ConvexHull(vector<point> P) {
    int n = P.size(), k = 0;
    vector<point> H(2*n);

    // Sort points lexicographically
    sort(P.begin(), P.end());

    // Build lower hull
    for (int i = 0; i < n; i++) {
        while (k >= 2
            && cross(H[k-2], H[k-1], P[i]) <= 0)
            k--;
        H[k++] = P[i];
    }

    // Build upper hull
    for (int i = n-2, t = k+1; i >= 0; i--) {
        while (k >= t
            && cross(H[k-2], H[k-1], P[i]) <= 0)
            k--;
        H[k++] = P[i];
    }

    H.resize(k);
    return H;
}

```

5.12 Closest pair of points

```

double delta_m(vp &ql, vp &qr, double delta) {
    uint64_t j = 0;
    double dm = delta;
}

```

```

for (uint64_t i = 0; i < ql.size(); i++) {
    point p = ql[i];

    while (j < qr.size()
        && qr[j].y < p.y - delta)
        j++;

    uint64_t k = j;
    while (k < qr.size()
        && qr[k].y <= p.y + delta) {
        dm = min(dm, dist(p, qr[k]));
        k++;
    }
}
return dm;
}

```

```

vp select_candidates(vp &p, int l, int r,
    double delta, double midx) {
    vp n;
    for (int i = l; i <= r; i++) {
        if (abs(p[i].x - midx) <= delta)
            n.push_back(p[i]);
    }
    return n;
}

```

```

double closest_pair(vp &p, int l, int r) {
    if (r - l + 1 < 2) return INT_MAX;
    int mid = (l + r) / 2;
    double midx = p[mid].x;
    double dl = closest_pair(p, l, mid);
    double dr = closest_pair(p, mid + 1, r);
    double delta = min(dl, dr);
}

```

```

vp ql, qr;
ql = select_candidates(p, l,
    mid, delta, midx);
qr = select_candidates(p, mid + 1,
    r, delta, midx);

```

```

double dm = delta_m(ql, qr, delta);

```

```

vp res;
merge(p.begin() + l, p.begin() + mid + 1,
    p.begin() + mid + 1, p.begin() + r + 1,
    back_inserter(res), cmp);
copy(res.begin(), res.end(), p.begin() + l);
return min(dm, min(dr, dm));
}

```


5.13 Test if point is inside a polygon

```

int wn_PnPoly(Point P, Point* V, int n) {
    int wn = 0; // the winding number counter

    // loop through all edges of the polygon
    for (int i = 0; i < n; i++) {
        if (V[i].y <= P.y) {
            if (V[i + 1].y > P.y)
                if (isLeft(V[i],
                    V[i + 1], P) > 0)
                    ++wn;
        } else {
            if (V[i + 1].y <= P.y)
                if (isLeft(V[i],
                    V[i + 1], P) < 0)
                    --wn;
        }
    }
    return wn;
}

```

5.14 Circle from 3 points

```

int main() {
    double ax, ay, bx, by, cx, cy, xres, yres;
    double xmid, ymid, A1, B1, C1, A2, C2, B2, dist;

    while (scanf("%lf %lf %lf %lf %lf %lf",
        &ax, &ay, &bx, &by, &cx, &cy) == 6) {
        A1 = by - ay;
        B1 = ax - bx;
        xmid = min(ax, bx) + (max(ax, bx)
            - min(ax, bx)) / 2.0;
        ymid = min(ay, by) + (max(ay, by)
            - min(ay, by)) / 2.0;
        C1 = -B1 * xmid + A1 * ymid;

        B2 = bx - cx;
        A2 = cy - by;
        xmid = min(bx, cx) + (max(bx, cx)
            - min(bx, cx)) / 2.0;
        ymid = min(by, cy) + (max(by, cy)
            - min(by, cy)) / 2.0;
        C2 = -B2 * xmid + A2 * ymid;

        //intersection of segments
        intersection(A1, B1, C1, A2,
            B2, C2, &xres, &yres);
        dist = sqrt(pow(xres - bx, 2)

```

```

        + pow(yres - by, 2));
    }
}

```

```

return 0;
}

```

6 Numerical

6.1 Check if float is an integer

```

#define EQ(a,b) (fabs((a) - (b)) < EPS)
#define IS_INT(a) ( EQ((a), ceil(a)) || \
    EQ((a), floor(a)) )

```

6.1.1 Big Mod

```

(BP)%M

```

```

typedef long long int lli;

```

```

long int bigmod(long long int B,
    long long int P, long long int M) {
    if (P == 0)
        return 1;
    else if (P & 1) {
        lli tmp =
            bigmod(B, (P - 1) >> 1, M) % M;
        tmp = (tmp * tmp * B) % M;
        return tmp;
    } else {
        lli tmp = bigmod(B, P >> 1, M) % M;
        return (tmp * tmp) % M;
    }
}

```

6.2 Triangle area

$$A = \frac{1}{2} * a * b * \sin(C)$$

6.3 Heron's formula

Let $s = \frac{1}{2}(a + b + c)$ then

$$A = \sqrt{s(s-a)(s-b)(s-c)}$$

6.4 Choose

```

(nk)

```

```

long long memo[SIZE][SIZE]; //initialized to -1
long long binom(int n, int k){
    if (memo[n][k] != -1) return memo[n][k];
    if (n < k) return 0;
    if (n == k) return 1;
    if (k == 0) return 1;
    return memo[n][k] = binom(n - 1, k)
        + binom(n - 1, k - 1);
}

```

6.5 Modulo:

```

int mod(int a, int n) {
    return (a%n + n)%n;
}

```

6.6 LCM / GCD

$$gcd(a, b) * lcm(a, b) = a * b$$

```

int gcd(int a, int b){
    if (!b)
        return a;
    else return gcd(b, a % b);
}

```

```

struct triple{
    int gcd, x, y;
    int triple(int g = 0, int a = 0, int b = 0){
        gcd(g), x(a), y(b) {}
    }
};

```

```

triple ExtendedEuclid(int a, int b){
    if (!b)
        return triple(a, 1, 0);

    triple t = ExtendedEuclid(b, a % b);
    return triple(t.gcd, t.y,
        t.x - (a / b) * t.y);
}

```

```

int LCM(int a, int b){
    return a * b / gcd(a, b);
}

```

6.7 Base conversion

```

void base(char *res, int num, int base){
    char tmp[100];

```

```

    int i, j;
    for (i = 0; num; i++) {
        tmp[i] =
            "0123456789ABCDEFGHIJKLM"[num % base];
        num /= base;
    }
    tmp[i] = 0;
    for (i--, j = 0; i >= 0; i--, j++)
        res[j] = tmp[i];
    res[j] = 0;
}

```

6.8 Horner's Rule

$$P(x) = \sum_{k=0}^n a_k x^k = a_0 + x(a_1 + x(a_2 + \dots + (a_{n-1} + x1_n)))$$

```

double Horner(double coef[], int degree, int x) {
    double res = 0;

    for (int i = degree; i >= 0; i--)
        res = coef[i] + x * res;
    return res;
}

```

6.9 Matrix Multiplication

```

void Matrix_Multiply(int A[N][P],
    int B[P][M], int N){
    int C[N][M], i, j, k;
    for (i = 0; i < N; i++){
        for (j = 0; j < P; j++){
            for (k = 0; k < P; k++){
                C[i][j] += A[i][k] * B[k][j];
            }
        }
    }
}

```

6.10 Long Arithmetic

Take care of leading zeroes.

Addition:

```

// make sure num1 and num2 are
// filled with '\0' after digits
void add(char *num1, char *num2, char *res){
    int i, carry=0;
    reverse(num1, num1 + strlen(num1));
    reverse(num2, num2 + strlen(num2));

```

```

for (i = 0; num1[i] || num2[i]; i++){
    res[i] = num1[i] + num2[i]
        - '0' + carry;
    if (!num1[i] || !num2[i])
        res[i] += '0';
    if (res[i] > '9'){
        carry = 1;
        res[i] -= 10;
    } else carry = 0;
}
if (carry) res[i] = '1';
reverse(res, res + strlen(res));
}

```

Multiplication

```

void mul(char *num1, char *num2, char *str) {
    int i, j, res[2*SIZE] = {0}, carry = 0;

    reverse(num1, num1 + strlen(num1));
    reverse(num2, num2 + strlen(num2));
    for (i = 0; num1[i]; i++)
        for (j = 0; num2[j]; j++)
            res[i + j] += (num1[i] - '0')
                * (num2[j] - '0');
    for (i = 2 * SIZE - 1;
         i >= 0 && !res[i]; i--);

    if (i < 0) {
        strcpy(str, "0");
        return;
    }
    for (j = 0; i >= 0; i--, j++){
        str[j] = res[i] + carry;
        carry = str[j] / 10;
        str[j] %= 10;
        str[j] += '0';
    }
    if (carry)
        str[j] = carry + '0';
}

```

6.11 Infix para Postfix

```

#define oper(a) ((a) == '+' || (a) == '-' \
    || (a) == '*' || (a) == '/')

```

```

// true if either: !!
// b is left associative and

```

```

// its precedence is <= than a
//
// b is right associative and
// its prec is < than a
bool be_prec(char a, char b) {
    int p[300];
    p['+'] = p['-'] = 1;
    p['*'] = p['/'] = 2;
    return p[a] >= p[b];
}

```

```

string shunting_yard(string exp) {
    int i = 0;
    string res;
    stack<char> s; //operators (1 char!)

```

```

    while (i < exp.size()) {
        // if it's a function token
        // push it onto the stack

        // If it is a func arg
        // separator (e.g., a comma):
        // Until the topmost
        // elem of the stack is '('
        // pop the elem from the stack and
        // append it to res.
        // If no '(' -> error
        // do not pop '('

        if (isdigit(exp[i]) || exp[i] == 'x') {
            //number. add isalpha() for vars
            for (; i < exp.size()
                && (isdigit(exp[i])
                    || exp[i] == 'x');
                i++) {
                res.push_back(exp[i]);
            }
            res.push_back(' ');
            i--; //there's a i++ down there
        } else if (exp[i] == '(') {
            s.push('(');
        } else if (exp[i] == ')') {
            while (!s.empty()
                && s.top() != '(') {
                res += s.top() + string(" ");
                s.pop();
            }
            if (s.top() != '(') ;//error
            else s.pop();
        } else if (oper(exp[i])) { //operator

```

```

    while (!s.empty()
           && oper(s.top())
           && be_prec(s.top(), exp[i])) {
        res += (s.top() + string("_"));
        s.pop();
    }
    s.push(exp[i]);
}
i++;
}
while (!s.empty()) {
    if (s.top() == '('
        || s.top() == ')')
        cout << "Error" << endl;
    res += (s.top() + string("_"));
    s.pop();
}
if (*(res.end() - 1) == '_')
    res.erase(res.end() - 1);
return res;
}

```

6.12 Calculate Postfix expression

```

// exp is in postfix
double calc(string exp) {
    stack<double> s;
    istringstream iss(exp);
    string op;

    while (iss >> op) {
        // ATTENTION TO THIS
        if (op.size() == 1 && oper(op[0])) {
            if (s.size() < 2)
                exit(-1); // error
            double a = s.top(); s.pop();
            double b = s.top(); s.pop();
            switch (op[0]) {
                case '+': s.push(b + a); break;
                case '-': s.push(b - a); break;
                case '*': s.push(b * a); break;
                case '/': s.push(b / a); break;
            }
        } else {
            istringstream iss2(op);
            double tmp;
            iss2 >> tmp;
            s.push(tmp);
        }
    }
}

```

```

    return s.top();
}

```

6.13 Postfix to Infix

```

/*
 * Pass a stack with the expression
 * to rpn2infix.
 * Ex: (bottom) 3 4 5 * + (top)
 */
string rpn2infix(stack<string> &s) {
    string x = s.top();
    s.pop();
    if (isdigit(x[0])) return x;
    else return string("(") +
        rpn2infix(s) + x +
        rpn2infix(s) + string(")");
}

```

6.14 Matrix Multiplication

$$C_{ij} = \sum_{k=1}^n a_{ik} \cdot b_{kj}$$

```

void matrix_mul(int A[N][P], int B[P][M]) {
    int C[N][M], i, j, k;
    for (i = 0; i < N; i++) {
        for (j = 0; j < P; j++) {
            C[i][j] = 0;
            for (k = 0; k < P; k++)
                C[i][j] += A[i][k] * B[k][j];
        }
    }
}

```

6.15 Catalan Numbers

$$C_n = \frac{(2n)!}{(n+1)!n!}$$

- C_n counts the number of expressions containing n pairs of parentheses which are correctly matched
- C_n is the number of different ways a convex polygon with $n+2$ sides can be cut into triangles by connecting vertices with straight lines.

6.16 Fibonnaci

```

long fib(long n){
    long matrix[2][2] = {{1, 1}, {1, 0}};
    long res[2][2] = {{1, 1}, {1, 0}};
    while (n) {
        if (n & 1) {
            matrix_mul(matrix, res, res);
        }
        matrix_mul(matrix, matrix, matrix);
        n /= 2;
    }
    return res[1][1];
}

```

7 Sorting / Search

7.1 Counting Sort

```

int count[SIZE] = {0};
int output[SIZE] = {0};
void linear_sort(int v[SIZE], int N) {
    int max = 0;
    for (int i = 0; i < N; ++i) {
        if (v[i] > max)
            max = v[i];
        count[v[i]]++;
    }

    for (int i = 1; i <= max; ++i) {
        count[i] += count[i-1];
    }

    for (int i = 0; i < N; ++i) {
        output[count[v[i]]-1] = v[i];
        count[v[i]]--;
    }
}

```

7.2 Binary Search - Lower bound

```

int lower_bound(int l, int r, ull q) {
    while (l < r) {
        ull mid = (l+r) / 2;

        if (v[mid] < q) {
            l = mid + 1;
        } else if (v[mid] > q) {
            r = mid - 1;
        } else {
            r = mid;
        }
    }
}

```

```

    }

    return l;
}

```

7.3 Binary Search - Upper bound

```

int upper_bound(int l, int r, ull q) {
    while (l < r) {
        int mid = (l+r) / 2;

        if (v[mid] < q) {
            l = mid + 1;
        } else if (v[mid] > q) {
            r = mid - 1;
        } else {
            l = mid+1;
        }
    }

    return l;
}

```